



**RT-11**  
**Software Support Manual**

Order No. DEC-11-ORPGA-B-D, DN1

digital



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## PREFACE

The RT-11 Software Support Manual covers the internal description of the RT-11 software system. Chapter 1 presents an overview of the system and discusses conventions used throughout the manual. Chapters 2 through 6 describe in detail various aspects of the monitor and system structure, including memory layout, monitor tables, file structures, file formats, system device structure, bootstrap operation, I/O queuing system, device handlers and F/B monitor description. Chapter 7 discusses the operation of the BATCH compiler and run-time handler.

The appendixes provide example handler listings, including a foreground terminal handler (Appendix B) and a sample foreground program (Appendix D). Complete flowcharts of both the Single-Job and Foreground/Background Monitors are shown in Appendix E.

The reader should be thoroughly familiar with the RT-11 system. Although the information in this manual is aimed at V02B and V02C users, it should be adequate for Version 2 users also; excluding a few minor alterations (to permit the addition of the new V02B devices), the construction of the monitors has changed very little between the two versions. A comprehensive list of differences between the V02B and V02C and between V2 and V02B systems is included in RT-11 System Release Notes (V02C), (DEC-11-ORNRA-A-D).

It is assumed that the user has read the RT-11 System Reference Manual (DEC-11-ORUGA-B-D) or (DEC-11-ORUGA-C-D) and all other documentation included in the RT-11 kit, and is an experienced PDP-11 programmer. It is recommended that RT-11 monitor source listings be available for reference.



## CHAPTER 1

### RT-11 OVERVIEW

#### 1.1 INTRODUCTION

RT-11 is a single-user programming and operating system designed for the PDP-11 series of computers. It permits the use of a wide range of peripherals and up to 28K of either solid state or core memory (hereafter referred to as memory).

RT-11 provides two operating environments: Single-Job (S/J) operation, and a powerful Foreground/Background (F/B) capability. Either environment is controlled by a single user from the console terminal keyboard by means of the appropriate monitor--S/J or F/B. The monitors are upwards compatible; features that are used only in a F/B environment are treated as no-ops under the S/J Monitor.

A feature common to both operating environments is the inclusion of a full complement of system development and utility programs to aid the programmer in the development of his own applications.

The normal use and operation of the monitors and system programs is discussed in detail in the RT-11 System Reference Manual. Concepts and applications that are specialized and useful to the more experienced programmer are included in this manual.

#### 1.2 SYSTEM CONCEPTS AND TERMINOLOGY

The basic concepts necessary to use RT-11 effectively are defined in the RT-11 System Reference Manual. The user should be familiar with those concepts before proceeding to use this manual.

Abbreviations used throughout this document are:

<u>TERM</u>	<u>MEANING</u>
KMON	Keyboard Monitor The console terminal interface to RT-11. KMON runs as a background job and allows the user to run programs, assign device names, and generally control the system.
USR	User Service Routines The nonresident (swapping) part of RT-11. The USR performs file-oriented operations.
CSI	Command String Interpreter The CSI is part of the USR. It accepts a string of characters from memory or from the console and performs specified file operations, or syntactically analyzes a command string and constructs a table from the information supplied.
RMON	Resident Monitor RT-11 provides a choice of two Resident Monitors: a Single-Job Monitor and a Foreground/Background Monitor. RMON specifically provides the following services:  EMT dispatcher Keyboard (console) interrupt service TT: resident device handler (F/B only) Read/Write processor USR swap routines I/O queuing routines System device handler System I/O tables Message handler (F/B only) Job scheduler (F/B only)
CSW	Channel Status Word Each bit in the CSW contains information relevant to the status of a channel; see Chapter 9 (.SAVESTATUS) of the <u>RT-11 System Reference Manual</u> .

<u>TERM</u>	<u>MEANING</u>
JSW	Job Status Word The JSW contains information in bytes 44 and 45 about the job currently in memory.
F/B	The Foreground/Background version of the monitor
S/J	The Single-Job version of the monitor
B/G	The background job
F/G	The foreground job
<CR>	Carriage Return
<LF>	Line Feed

Various mnemonic names (e.g., BLIMIT, SYSLOW), referred to from within the text and in diagrams and flowcharts, represent the actual symbolic names as they appear in the monitor source listings.

To avoid confusion, underlining is used in most examples to designate computer printout; square brackets, [ and ], are used to enclose comments. Values for symbolic names used in examples can be found in Table 2 of RT-11 System Release Notes.





CHAPTER 2  
MEMORY LAYOUT

RT-11 operates properly in any configuration between 8K and 28K (words) of memory (16K to 28K for the F/B Monitor). No user intervention is required when programs are moved to a different size machine; i.e., programs correctly developed in one environment will work in any size environment (providing there is sufficient memory) with no relinking necessary.

Figure 2-1 shows a general diagram of the memory layout in an RT-11 system.

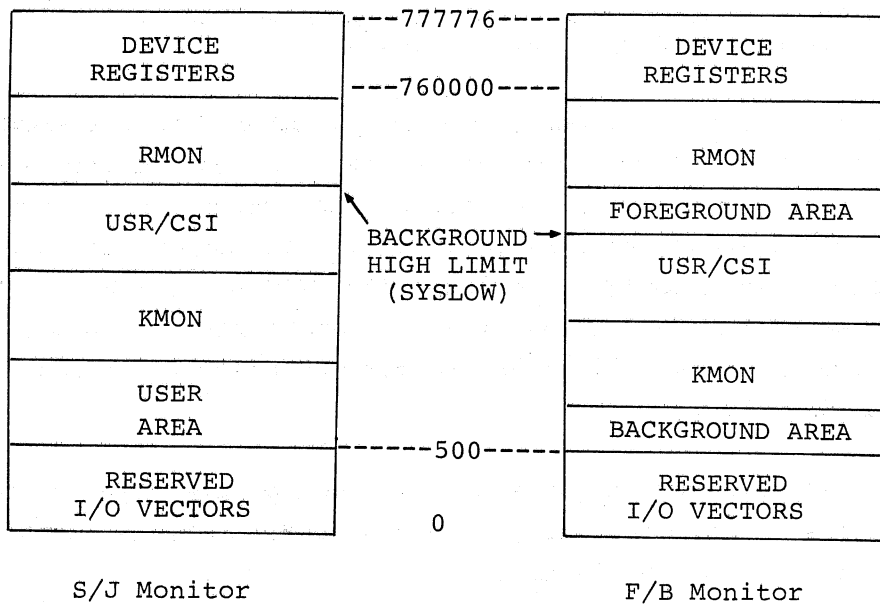


Figure 2-1  
Monitor Memory Layout

The memory area diagrammed is arranged as follows:

<u>Memory Area</u>	<u>Use</u>
0-477	Reserved for I/O vectors, RT-11 system communication area.
500-SYSLOW	Space available for user (background) programs. (The high limit of memory for the background is contained in SYSLOW, a location in the monitor data base.)

Space for foreground programs and LOAded handlers is allocated as needed, reducing the amount of space available for a background job.

The areas marked KMON and USR/CSI are the areas that these units normally occupy when they are in memory. The amount of memory that a user program occupies is determined by:

1. The initial size of the program, or
2. The amount of memory the user program requests via a .SETTOP programmed request.

When a user program (background job) is executed (via the KMON commands R, RUN, or GET and START), the top of memory is set to correspond to the size of the program. If the top of user memory never exceeds KMON, both KMON and USR/CSI are resident. If all of memory (up to SYSLOW) is requested (via a .SETTOP), neither the KMON nor the USR is resident and swapping of the USR is required. Programs performing many file-oriented operations gain from having the USR resident, since no time is spent swapping the USR.

The KMON, USR, and RMON modules normally occupy the upper segment of memory. This implies that larger memory configurations automatically have more free memory available.

The area marked DEVICE REGISTERS is the top 4K of memory in any PDP-11 computer. This area is reserved for the status and control registers of peripheral devices.

## 2.1 FOREGROUND JOB AREA LAYOUT

The foreground job area is located above the KMON/USR, as shown in Figure 2-1, and is allocated by the FRUN command. The actual layout of the job within the foreground area is shown in Figure 2-2. The *impure area* (described in Section 2.5.3) occupies the lowest 207 words of the job area and contains terminal ring buffers, I/O channels, and other job-specific information.

The foreground stack is located immediately above the impure area with a default size of 128 words; this may be changed using the FRUN /S switch. The program may specify a different location for the stack by using an .ASECT into location 42, in which case the /S switch is ignored and the program itself must allocate stack space. Wherever the stack is located, stack overflow will most probably cause program malfunction before penetrating the task area boundary, since either the program itself or the impure area will be corrupted.

### NOTE

Users must not use a relocatable symbol as the contents of location 42 when resetting the initial stack pointer via an .ASECT in a foreground job; such a symbol is not relocated when it occurs in an .ASECT in a foreground job. To set the stack to relative location 1000 in a foreground job, use:

```
.ASECT  
.=42  
.WORD 1000
```

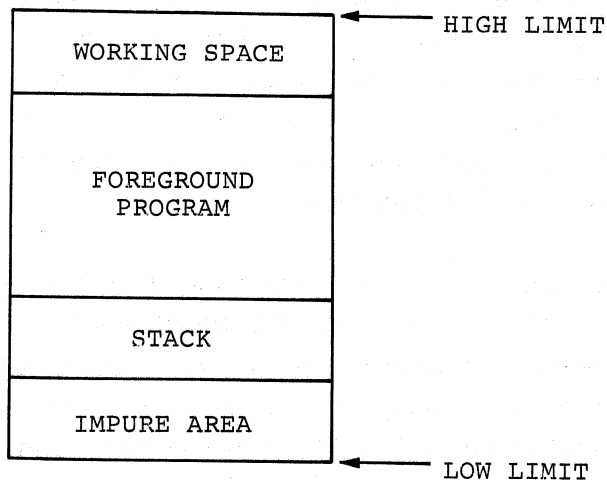


Figure 2-2  
Foreground Job Area Layout

The space allocated for the foreground program is sufficient to contain the program code itself, as indicated by location 50 (in block 0 of the file); location 50 is set by the Linker and designates the program's high limit. If the foreground job requires working space, this space must either be reserved from within the program (e.g., using .BLKW) or allocated at run-time using the FRUN /N switch. Space allocated with the /N switch is located above the program as shown in Figure 2-2. Location 50 will point to the top of the program area and a .SETTOP will permit access to any working space.

## 2.2 JOB BOUNDARIES IN F/B

The actual job boundaries are stored (in RMON) in limit tables for both foreground and background jobs. The FLIMIT table contains high and low boundaries for the foreground, and the BLIMIT table contains boundaries for the background. .SETTOPs are permitted for any job up to its high limit. The SYSLOW pointer mentioned earlier is equivalent to the background BLIMIT high pointer entry. This is shown in Figure 2-3.

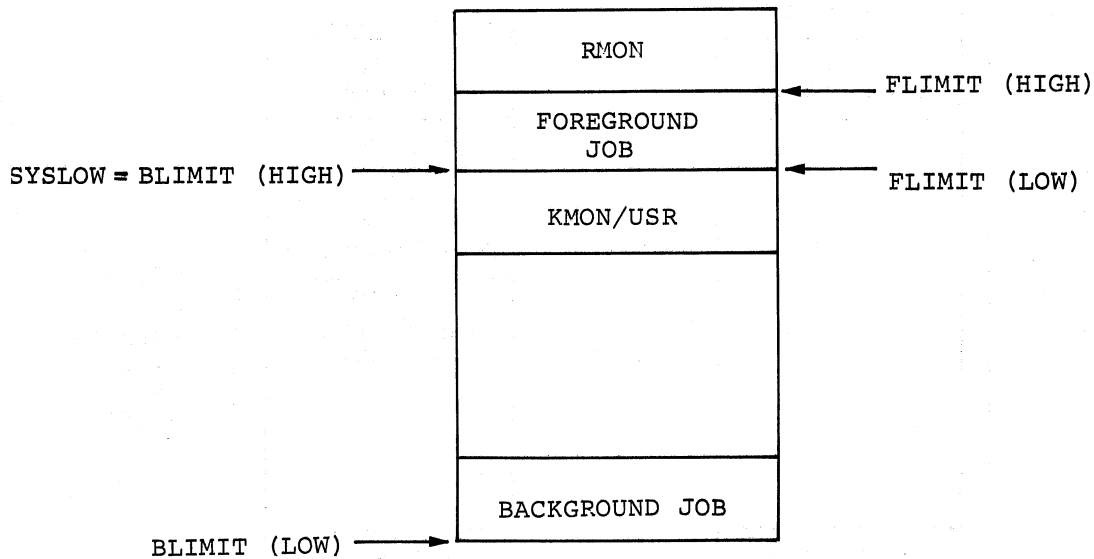


Figure 2-3  
Job Limits

The limit pointers for a foreground job are fixed once the job has been loaded into memory. A program that requires working space and uses a .SETTOP will fail if the space is not allocated with the /N switch (a FORTRAN program is a typical case; see Appendix G, Section G.1, of the RT-11 System Reference Manual). The high limit pointer (SYSLOW) for the background, however, is not fixed and will change as space is allocated for LOAded handlers, the text scroller, and foreground jobs. In addition, if the USR is made permanently resident (using the SET USR NOSWAP command), SYSLOW (BLIMIT HIGH) will again change. This is shown in Figure 2-4.

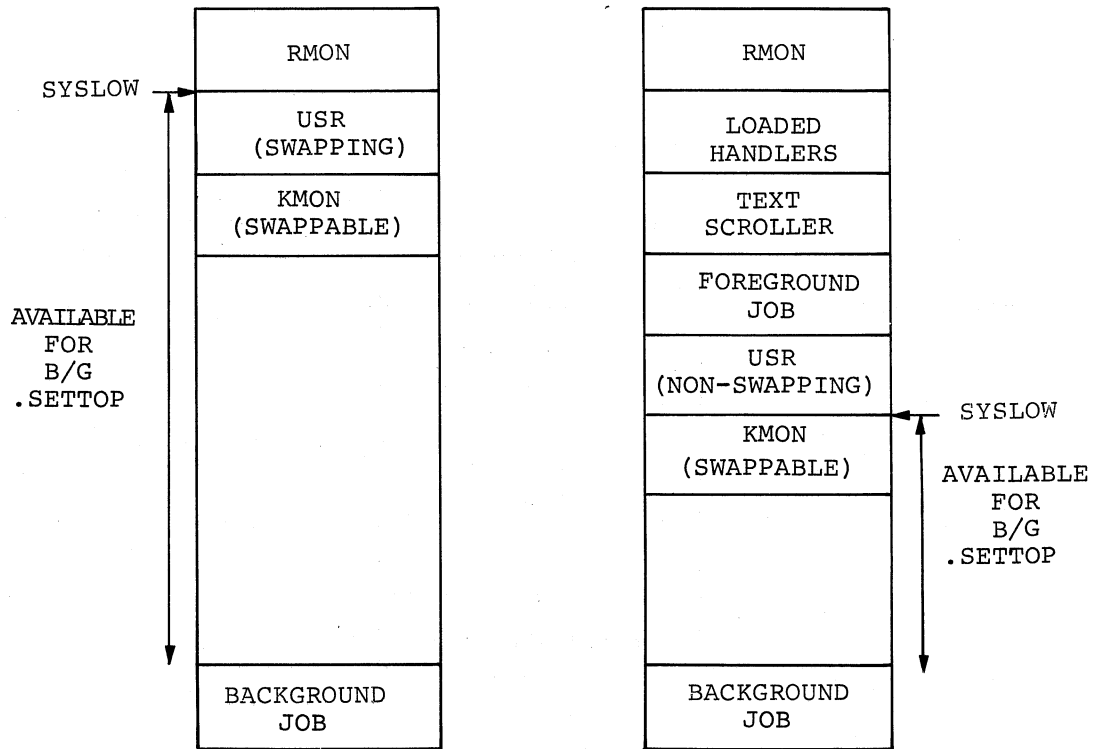


Figure 2-4  
Background SYSLOW Examples

### 2.3 'FLOATING' USR POSITION

The RT-11 USR is normally located in the memory area directly below that pointed to by SYSLOW. For the Version 1 monitor, this was directly below the RMON. For the Version 2 and 2B monitors, the USR position varies as handlers, the scroller, and foreground jobs (in F/B) are loaded into memory; the SYSLOW pointer is corrected for each change in memory configuration. In any case, the SYSLOW position is considered the normal USR swapping position.

It is possible, however, to cause the USR to swap into another location in memory. This is done by setting location 46 (in the system communication area) to the address at which the USR is to swap; if the contents of location 46 are nonzero and even, the monitor loads the USR at the new address. Note, however, that if no swapping is required, the USR is *not* loaded at the address indicated in location 46. Location 46 is cleared by an exit to the Keyboard Monitor (via an .EXIT, .HRESET, .SRESET, or CTRL C).

It is possible to make the USR permanently resident (i.e., non-swapping). Using the SET USR NOSWAP Keyboard Monitor command makes the USR permanently resident at its *normal* position, that is, below the memory area pointed to by SYSLOW.

#### 2.4 MONITOR MEMORY ALLOCATION

RT-11 uses a dynamic memory allocation scheme to provide memory space for LOAded handlers, foreground jobs (F/B Monitor only) and the display text scroller. Memory is allocated in the region above the KMON/USR and below RMON. If there is insufficient memory in this region (initially, after the system is bootstrapped, there is none), memory is taken from the background region by "sliding down" the KMON/USR the required number of words.

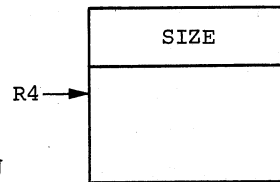
When memory allocated in this manner is released, the memory block is returned to a singly-linked free memory list, the list head of which is in RMON. Any contiguous blocks are concatenated into a single larger block. A block found to be contiguous with the KMON/USR is reclaimed by "sliding up" the KMON/USR, removing the block from the list.

Memory allocation and release is achieved by calls to the GETBLK and PUTBLK routines located in the KMON overlays (the GETBLK and PUTBLK routines are flowcharted in Appendix E). The requested number of words is passed to GETBLK in R0, and the address of the block is returned in R4. An extra word of memory is allocated by GETBLK, which then stores the size of the block in that word. R4 points to the first available word in the block (see Figure 2-5a). When releasing memory, R4 must point to the first available word, the same address returned by GETBLK during allocation (as shown in Figure 2-5b). The block will be linked into the free memory list (shown in Figure 2-5c).

a) Allocating a memory block

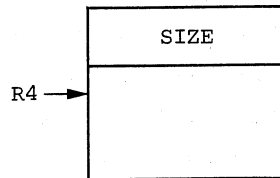
Call sequence:  
 $R0 = SIZE$   
 JSR PC, GETBLK

(returns with R4 pointing to the allocated block)



b) Releasing a memory block

Call sequence:  
 $R4 \rightarrow BLOCK$   
 JSR PC, PUTBLK



c) Free memory list

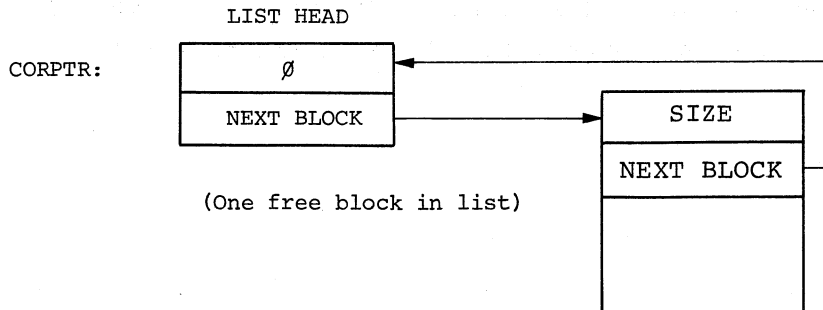


Figure 2-5  
 Memory Allocation



When a block of memory of sufficient size is not available, GETBLK must create a hole in memory by sliding down the KMON/USR. This is achieved by a call to KUMOVE, a small routine located physically at the front of the KMON. KUMOVE does the actual work of moving the KMON/USR up in memory. For moves downward, an auxiliary subroutine, MOVEDN, located at the top of the USR, is used.

Whenever a request is made for a block of a certain number of words, the memory allocator searches memory for the first highest block that is large enough to satisfy the request (that is, equal to or larger than the requested number). The goal of the memory allocator is to minimize the amount of free (unused) memory in the foreground region, making the maximum amount of memory available to the background. Contiguous blocks of free memory are merged and reclaimed whenever possible. The search time of the singly-linked list is not a factor, since at any time there will be few nodes (free memory areas) in the list, and the allocator minimizes the number.

## 2.5 MEMORY AREAS OF INTEREST

This section describes memory areas of particular interest and indicates the contents of those locations. The areas covered are:

1. Monitor Fixed Offsets (F/B & S/J)
2. F/B Impure Area
3. Resident Bitmap (F/B & S/J)
4. Tables

### 2.5.1 Monitor Fixed Offsets

Certain values are maintained at fixed locations from the start of the Resident Monitor in both F/B and S/J; these quantities (listed in Table 2-1) may be accessed by user programs. The technique used to access these offsets is as follows:

OFFSET = the byte offset to the word desired  
RMON = 54

```
MOV @#RMON,Rn          ;ANY GENERAL REGISTER  
MOV OFFSET(Rn),Rn
```

Rn now contains the desired quantity. If a byte quantity is desired, a better method is:

```
CLR Rm
MOV @#RMON, Rn
BISB OFFSET(Rn), Rm
```

This ensures that the high-order bits of the register are not set by a MOV<sub>B</sub> into the register.

Table 2-1  
Fixed Offsets

Offset (from Start of RMON)		Tag	Byte Length	Description
Octal	Decimal			
0		-	4	Serves as a link to interrupt entry code.
4		\$CSW	160 <sub>10</sub>	Default I/O channels for the background (16 <sub>10</sub> @ 5 words each).
244	164	\$SYSCH	10 <sub>10</sub>	Internal I/O channel used for system functions.
256	174	BLKEY	2	Segment number of the directory now in memory. 0 implies no directory is there.
260	176	CHKEY	2	Device index and unit number of the device whose directory is in memory. Bits 1-5 are the device index, bits 8-10 are the unit number.
262	178	\$DATE	2	Current date value. (The format is shown in Chapter 3, section 3.1.2.5.)
264	180	DFLG	2	"Directory operation in progress" flag. Used to inhibit ^C from aborting a job until directory operation is finished.
266	182	\$USRLC	2	Normal location of USR.
270	184	QCOMP	2	Address of I/O completion manager, COMPLT.
272	186	SPUSR	2	Flag word used by MT/CT. If a USR function performed by MT or CT fails, this word is made non-zero.

(continued on next page)

Table 2-1 (Cont.)  
Fixed Offsets

Offset (from Start of RMON)		Tag	Byte Length	Description																						
Octal	Decimal																									
274	188	SYUNIT	2	High-order byte contains the unit number of the current system device.																						
276	190	SYSVER	1	Monitor version number (2 in Versions 2, 2B, and 2C).																						
277	191	SYSUPD	1	Version release number (1 for V02, 2 for V02B, etc.)																						
300	192	CONFIG	2	System configuration word. A 16-bit series of flags whose meanings are:  <table border="0" style="margin-left: 40px;"> <thead> <tr> <th style="text-align: left;"><u>Bit #</u></th> <th style="text-align: left;"><u>Meaning</u></th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0 → S/J Monitor 1 → F/B Monitor</td> </tr> <tr> <td>2</td> <td>1 → VT11 hardware exists</td> </tr> <tr> <td>3</td> <td>1 → RT-11 BATCH controls the background</td> </tr> <tr> <td>5</td> <td>0 → 60-cycle KW11L clock 1 → 50-cycle clock</td> </tr> <tr> <td>6</td> <td>1 → 11/45 FPP present</td> </tr> <tr> <td>7</td> <td>0 → No foreground job present 1 → Foreground job is in memory</td> </tr> <tr> <td>8</td> <td>1 → User is linked to VT11 scroller</td> </tr> <tr> <td>9</td> <td>1 → USR is resident via SET USR</td> </tr> <tr> <td>11</td> <td>0 → No PDP-11/03 processor 1 → PDP-11/03 processor</td> </tr> <tr> <td>15</td> <td>1 → KW11 clock is present (always set if bit 11 is 1)</td> </tr> </tbody> </table> <p>Any bits not currently assigned are reserved by DIGITAL for future use and should not be used arbitrarily by user programs.</p>	<u>Bit #</u>	<u>Meaning</u>	0	0 → S/J Monitor 1 → F/B Monitor	2	1 → VT11 hardware exists	3	1 → RT-11 BATCH controls the background	5	0 → 60-cycle KW11L clock 1 → 50-cycle clock	6	1 → 11/45 FPP present	7	0 → No foreground job present 1 → Foreground job is in memory	8	1 → User is linked to VT11 scroller	9	1 → USR is resident via SET USR	11	0 → No PDP-11/03 processor 1 → PDP-11/03 processor	15	1 → KW11 clock is present (always set if bit 11 is 1)
<u>Bit #</u>	<u>Meaning</u>																									
0	0 → S/J Monitor 1 → F/B Monitor																									
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3	1 → RT-11 BATCH controls the background																									
5	0 → 60-cycle KW11L clock 1 → 50-cycle clock																									
6	1 → 11/45 FPP present																									
7	0 → No foreground job present 1 → Foreground job is in memory																									
8	1 → User is linked to VT11 scroller																									
9	1 → USR is resident via SET USR																									
11	0 → No PDP-11/03 processor 1 → PDP-11/03 processor																									
15	1 → KW11 clock is present (always set if bit 11 is 1)																									
302	194	SCROLL	2	Address of the VT11 scroller.																						
304	196	TTKS	2	Address of console keyboard status.																						
306	198	TTKB	2	Address of console keyboard buffer.																						

(continued on next page)

Table 2-1 (Cont.)  
Fixed Offsets

Offset (from Start of RMON)		Tag	Byte Length	Description
Octal	Decimal			
310	200	TTPS	2	Address of console printer status.
312	202	TTPB	2	Address of console printer buffer.  (See Section 2.6, Using Auxiliary Terminals as the Console Terminal.)
314	204	MAXBLK	2	Largest output file permitted with an indefinite length request (initially defined as -1, which implies that no limit is defined).
316	206	E16LST	2	Offset from start of RMON to the dispatch table for EMT's 340-357. (This is used by the BATCH processor.)
320	208	CNTXT	2	Pointer to the impure area for the current executing job.
		(F/B)		
322	210	JOBNUM	2	Executing job's number (0 = B/G, 2 = F/G).
320	208	\$TIME	4	Two words of time of day in the S/J Monitor.
322	210	(S/J)		
324	212	SYNCH	2	Address of monitor routine to handle .SYNCH request.
326	214	LOWMAP	20 <sub>10</sub>	Start of low memory protection map. (This map protects vectors at locations 0-476.)
352	234	USRLOC	2	Pointer to current entry point of USR.
354	236	GTVECT	2	Pointer to VT11 vector. The vector is initially positioned at 320.
356	238	ERRCNT	1	Error count byte (for future use by system programs).
357	239	FUTURE	5	Reserved by DIGITAL for future use.

## 2.5.2 Table Descriptions

The monitor device tables discussed in this section include:

\$PNAME  
\$STAT  
\$ENTRY  
\$DVREC  
\$HSIZE  
\$DVSIZ  
\$UNAM1, \$UNAM2  
\$OWNER

The size of these tables is fixed and is governed by the \$SLOT assignment; the default value is 14<sub>10</sub> entries per table. To alter this, it is necessary to first edit a new value of \$SLOT into the monitor source program, then reassemble and relink new monitors.

2.5.2.1 \$PNAME (Permanent Name Table) - \$PNAME is the central table around which all the others are constructed. There is an entry in \$PNAME for each device in the system. Each entry consists of a single word that contains the .RAD50 code for the two-character permanent device name for that device; for example the entry for DECTape is .RAD50 /DT/. The position of devices in this table is non-critical, but their relative position determines the general device index used in various places in the monitor; thus, all other tables must be organized in the same order as \$PNAME (the index into \$PNAME serves as the index into all the other tables for the equivalent device).

2.5.2.2 \$STAT (Device Status Table) - Each device in the system must have a status entry in its corresponding slot in \$STAT. The status word is broken down into two bytes as follows:

Even byte - contains a device identifier. Each unique type of device in the system has an identifying integer. Those defined are:

0 = RK05 Disk  
1 = TC11 DECTape  
2 = Reserved  
3 = Line Printer (LP11, LS11, LV11)  
4 = Console Terminal (LT33/35, LA30/36,  
VT05, VT50)  
5,6 = Reserved  
7 = PC11 High-speed Reader  
10 = PC11 High-speed Punch  
11 = Magtape (TM11, TU10)  
12 = RF11 Disk  
13 = TA11 Cassette

14 = Card Reader (CR11, CM11)  
 15 = Reserved  
 16 = RJS03/4 Fixed-head Disks  
 17 = Reserved  
 20 = TJU16 Magtape  
 21 = RP11/RP02/RP03 Disk  
 22 = RX11/RX01 Diskette

Odd byte - Bit flags with the following meanings:

Bit 15: 1 = Random-access device (disk, DEctape)  
           0 = Sequential-access device (line printer, papertape, card reader, magtape, cassette, terminal)  
 Bit 14: 1 = Read-only device (card reader, papertape reader)  
 Bit 13: 1 = Write-only device (line printer, papertape punch)  
 Bit 12: 1 = NonRT-11 directory-structured device (magtape, cassette)  
 Bit 11: 1 = Enter handler abort entry every time a job is aborted.  
           0 = Handler abort entry taken only if there is an active queue element belonging to aborted job.  
 Bit 10: 1 = Handler accepts .SPFUN requests (e.g., MT, CT, DX).  
           0 = .SPFUN requests are rejected as illegal.  
 Bits 9-8: Reserved

2.5.2.3 \$ENTRY (Handler Entry Point Table) - Whenever a handler is made resident, either by a .FETCH or with the LOAD command, the \$ENTRY slot for that device is made to point to the fourth word of the device handler. The entry is zeroed when the handler is .RELEASED or UNLOADED.

2.5.2.4 \$DVREC (Device Handler Block Table) - This table (filled in at system bootstrap time) reflects the absolute block position of each of the device handlers on the system device. Since handlers are treated as files under RT-11, their position on the system device is not necessarily fixed. Thus, each time the system is bootstrapped, the handlers are located and \$DVREC is updated with the value of the second block of the handler file. (Because the handlers are linked at 1000, the actual handler code starts in the second block of the file.) A zero entry in the \$DVREC table indicates that no handler for the device in that slot was found on the system device.

2.5.2.5 \$HSIZE (Handler Size Table) - This table contains the size, in bytes, of each device handler. The table is set up at assembly time with the correct values and is used when a .FETCH is executed to provide the size of the specified handler. This size is also returned to the user as one of the values returned in a .DSTAT request.

2.5.2.6 \$DVSIZ (Device Directory Size Table) - Entries in this table are non-zero for file-structured devices only and reflect the number of 256<sub>10</sub>-word blocks contained on the device. The current devices and their entries are:

<u>Device</u>	<u>Number of 256-Word Blocks</u>	<u>Device</u>	<u>Number of 256-Word Blocks</u>
RK11	11300 <sub>8</sub>	RP02	116300 <sub>8</sub>
TC11	1102 <sub>8</sub>	RJS03	2000 <sub>8</sub>
		RJS04	4000 <sub>8</sub>
RF11	2000 <sub>8</sub> (1 platter) 4000 <sub>8</sub> (2 platters) 6000 <sub>8</sub> (3 platters) 10000 <sub>8</sub> (4 platters)	RX01	752 <sub>8</sub>

The default for RF11 and RJS03/4 is one platter, or 2000<sub>8</sub> blocks. It is possible to alter the system to indicate the correct number of platters. Instructions are in Chapter 4 of the RT-11 System Generation Manual, (DEC-11-ORGMA-A-D).

2.5.2.7 \$UNAM1, \$UNAM2 (User Name Tables) - These tables are used in conjunction with ASSIGN keyboard functions. The form of the ASSIGN command is:

```
..ASSIGN pnam:unam<CR>
```

where:

pnam - a system device name/unit number

unam - a user-assigned device name

A typical example is:

```
..ASSIGN DT1:DK
```

The default device name, DK, is now directed to DECTape unit 1. The user-assigned name is stored in an available slot in \$UNAM2, while the device's permanent name/unit is stored in the corresponding slot in \$UNAM1. The system uses a common device name lookup routine that maps any user-assigned name in the \$UNAM2 table into a physical device name to be used in an operation. The total number of ASSIGNs permitted is limited by the value of \$SLOT.

The command:

```
.ASSIGN<CR>
```

zeroes \$UNAM2, thus removing all user assignments.

2.5.2.8 \$OWNER (Device Ownership Table) - This table is used only under F/B to arbitrate device ownership. The table is \$SLOT\*2 words in length and is divided into 2-word entries per device. Each 2-word entry is divided into eight 4-bit fields capable of holding a job number. Thus, each device is presumed to have up to eight units, each assigned independently of the others. However, if the device is nonfile-structured, the ownership is assigned to all units.

When a job attempts to access a particular unit of a device, the F/B Monitor checks to be sure the unit being accessed is either public or belongs to the requesting job. If the unit is owned by the other job, a fatal error is generated.

The device is assumed to be public if the 4-bit field is 0. If it is not public, the field contains a code equal to the job number plus one. Since job numbers are always even, the ownership code is odd. Bit 0 of the field being set is then used to indicate that the unit ownership is assigned to a job (1 for the background job and 3 for the foreground job).

2.5.2.9 DEVICE Macro - The DEVICE macro call is used in RMON to allow quick and easy insertion of new devices at assembly time. The form of the macro call is:

```
DEVICE NAME,SIZ,STAT,ENTRY
```



where:

- NAME - two characters of the permanent device name
- SIZ - the size of the device's directory in 256-word blocks; 0 means nonfile-structured or special
- STAT - the sum of all \$STAT table entries that apply for this device plus the device id (from section 2.5.2.2):

FILST\$ = 100000	Random-access device (disk, DEctape)
RONLY\$ = 40000	Read-only device
WONLY\$ = 20000	Write-only device
SPECL\$ = 10000	Non RT-11 directory-structured device (including MT and CT)
HNDLR\$ = 4000	Handler abort entry
SPFUN\$ = 2000	Special function requests

ENTRY - the 2-character device name with the SYS appended, if this is a system device.

Thus, a sample call is:

```
DEVICE TT,0,4
```

The SIZ entry is 0, since TT is a nonfile-structured device.

The entry for DEctape is:

```
DEVICE DT,1102,1+FILST$,DTSYS
```

The 1+FILST\$ indicates that the device code is 1 and FILST\$ is defined as 100000. The entry for DTSYS is present because DT can be a system device.

In addition to the DEVICE macro, another macro, HSIZE, is defined and sets the handler size for the \$HSIZE table. The format of the HSIZE macro call is:

```
HSIZE HAN,BYT,TYPE
```

where:

- HAN - the 2-letter device name
- BYT - the handler size in bytes
- TYPE - SYS if the device can be a system device; blank otherwise

Chapter 5 shows the use of HSIZE in adding a handler to the RT-11 system. The KMON portion of the monitor source listing should be consulted for greater detail.

### 2.5.3 F/B Impure Area

An impure area is defined here as that area of memory where the monitor stores all job-dependent data. Thus, the impure area contains all information that the monitor requires to effectively run two independent jobs, both of which are memory-resident. This section details the contents and location of each word (byte) in the impure area.

A table that points to the impure area for a particular job is in the F/B monitor's data base. This table is at \$IMPUR and currently consists of two words: the first is a pointer to the background's impure area (which is permanently resident in RMON at location BKGND), the second is the foreground's pointer. The \$IMPUR table is accessed by using IMPLOC, located at an offset of 422 into RMON. IMPLOC points beyond the end of \$IMPUR to \$IMPUR +4 to facilitate accessing the \$IMPUR table from the top down in order of decreasing priority.

Under RT-11, a background job is always running and will be the KMON if no other background job exists. However, the foreground impure area pointer may be 0 if no foreground job is in memory. When an FRUN command is given, a foreground impure area is created for the job and the \$IMPUR entry for the foreground pointer is updated to point to the impure area.

A foreground program can determine whether the KMON is resident by testing KMONIN, located at an offset of 424 into RMON. KMONIM is non-zero if the KMON is resident and zero if a background job is running. In addition, the file name of the running foreground or background job is located in the job's impure area at offset I.NAME (376). Note that for a background job, KMONIN must first be tested to determine whether the name belongs to an active job since the file descriptor is not cleared when KMON is entered.

Table 2-2 is a detailed breakdown of the contents of the impure area. The offset mentioned is the offset from the start of the impure area itself; thus, the first word in the area has a 0 offset.

Table 2-2  
Impure Area

Offset	Mnemonic	Octal Length (Bytes)	Contents
0	I.JSTA	2	Job status.
2	I.QHDR	2	I/O Queue Header.
4	I.CMPE	2	Last entry in completion queue. I/O completion routines are queued for execution. This is the pointer to the last routine to be entered.
6	I.CMPL	2	Completion queue header.
10	I.CHWT	2	Pointer to channel during I/O wait. When a job is waiting for I/O, the address of the channel area in use goes here.
12	I.PCHW	2	Saved channel pointer during execution of a completion routine. The contents of I.PCHW are put in R0 when a completion routine is entered.
14	I.PERR	2	Error byte 52 and 53 saved during completion routines.
16	I.PTTI	2	Previous TT input character.
20	I.TTLC	2	Terminal input ring buffer line count.
22	I.TID	2	Pointer to job ID area.
24	I.JNUM	2	Job number of job that owns this impure area.
26	I.CNUM	2	Number of I/O channels defined. 16 <sub>10</sub> is default, .CDFN can be used to define new ones.
30	I.CSW	2	Pointer to job's channel area.
32	I.IOCT	2	Count of total I/O operations outstanding.
34	I.SCTR	2	Suspension count. Zero means the number of .SPNDs = the number of .RSUMs.
36	I.SPLS	2	Address of the .DEVICE request list.

(continued on next page)

Table 2-2 (Cont.)  
Impure Area

Offset	Mnemonic	Octal Length (Bytes)	Contents
40	I.TRAP	2	Address of user trap routine. Set by .TRPSET.
42	I.FPP	2	Address of FPP exception routine. Set by .SFPA.
44	I.SWAP	4	Address and number of extra words to be included in the context switch operation. Set by .CNTXSW request.
50	I.SP	2	Saved stack pointer. When this job is made inactive, the active value of SP is saved here.
52	I.BITM	24	Low memory protection bitmap. This map reflects the user's .PROTECT requests.
(76 through 332 concern the console terminal)			
76	I.IRNG	2	Input ring buffer low limit.
100	I.IPUT	2	Input "PUT" pointer for inter- rupts.
102	I.ICTR	2	Input character counter.
104	I.IGET	2	Input "GET" pointer for .TTYIN.
106	I.ITOP	2	Input ring buffer high limit.
110		144	Input ring buffer.
254	I.OPUT	2	Output "PUT" pointer for interrupts.
256	I.OCTR	2	Output character counter.
260	I.OGET	2	Output "GET" pointer for interrupts.
262	I.OTOP	2	Output ring buffer high limit.
264		50	Output ring buffer.
334	I.QUE	20	Initial I/O queue element.

(continued on next page)

Table 2-2 (Cont.)  
Impure Area

Offset	Mnemonic	Octal Length (Bytes)	Contents
354	I.MSG	12	Message channel. Used by .RCVD and .SDAT. This channel is permanently open.
366		10	Job ID area. Contains (<CR><LF>)B(<CR><LF>) or (<CR><LF>)F(<CR><LF>) for terminal prompting. Space has been left for up to a 3-character job name.

#### 2.5.4 Low Memory Bitmap (LOWMAP)

RT-11 maintains a bitmap which reflects the protection status of low memory, locations 0-476. This map is required in order to avoid conflicts in the use of the vectors. In F/B, the .PROTECT request allows a program to gain exclusive control of a vector or a set of vectors. When a vector is protected, the bitmap is updated to indicate which words are protected. If a word in low memory is protected, it will not be destroyed when a new background program is run.

The bitmap is a 20<sub>10</sub> byte table which starts 326 bytes from the beginning of the Resident Monitor. Table 2-3 lists the offset from RMON and the corresponding locations represented by that byte:

Table 2-3  
Bitmap Byte Table

Offset	Locations (octal)	Offset	Locations (octal)
326	0-16	340	240-256
327	20-36	341	260-277
330	40-56	342	300-316
331	60-76	343	320-336
332	100-116	344	340-356
333	120-136	345	360-376
334	140-156	346	400-416
335	160-176	347	420-436
336	200-216	350	440-456
337	220-236	351	460-476

Each byte in the table reflects the status of  $16_{10}$  words of memory. The first byte in the table controls locations 0-16, the second byte controls locations 20-36, and so on. The bytes are read from left to right. Thus, if locations 0-3 are protected, the first byte of the table contains:

11000000

Note that only individual words are protected, not bytes. Thus, protecting location 0 always implies that the word at location 0 is protected, meaning both locations 0 and 1. If locations 24 and 26 are protected, the second byte of the table contains:

00110000

since the leftmost bit represents location 20 and the rightmost bit represents location 36. To protect locations 300-306, the leftmost 4 bits of byte 342 must be set:

11110000

resulting in a value of 360 for that byte.

2.5.4.1 S/J Restrictions - The S/J Monitor does not support the .PROTECT request. If users wish to protect vectors, the protection must be done in one of two ways:

1. Manually, with PATCH, or
2. Dynamically (from within the user's program)

To protect locations 300-306 dynamically, the following instructions are used:

```
MOV @#54,R0
BISB #†B11110000,342(R0)
```

Protecting locations with PATCH implies that the vector is permanently protected, even if the system is re-bootstrapped, while the second method provides a temporary measure and does not hold across bootstraps. However, users are cautioned that the second method involves storing directly into the monitor; for this reason it is recommended that S/J users use method 1.

## 2.6 USING AUXILIARY TERMINALS AS THE CONSOLE TERMINAL

This section describes how RT-11 can be modified to allow a terminal other than the standard console unit 0 to become the console terminal. This procedure is useful in cases where it is desirable to be able to use different console capabilities at different times (for example, at certain times the hard copy output of an LA30 is required, while at other times the speed of a VT05 is desirable). The only information required to make the alteration is:

- 1) the address of the auxiliary terminal's interrupt vectors, and
- 2) the I/O page addresses of the keyboard and printer status register and buffer.

RT-11 is designed so that all console references are done indirectly through centralized pointers. Thus, changing several system locations causes all operations to be transferred to a new terminal.

For this example, assume that the new terminal's interrupt vectors are at 300,302 and 304,306 and that its I/O page addresses are:

TKS at 177500  
TKB at 177502  
TPS at 177504  
TPB at 177506

Also assume that the new terminal is a parallel interface so that no fill characters are required.

.R PATCH <CR>

PATCH Version number

FILE NAME--

\*MONITR.SYS/M<CR>

\*BASE;ØR<CR>

\*6Ø/\_\_\_\_\_VECTIN<LF>

62/\_\_\_\_\_STATIN<LF>

64/\_\_\_\_\_VECTOUT<LF>

66/\_\_\_\_\_STATOUT<CR>

\*3ØØ/\_\_\_\_\_nnnnn VECTIN<LF>

3Ø2/\_\_\_\_\_nnnnn STATIN<LF>

3Ø4/\_\_\_\_\_nnnnn VECTOUT<LF>

3Ø6/\_\_\_\_\_nnnnn STATOUT<CR>

\*Ø,xx3Ø4/\_\_\_\_\_17756Ø 1775ØØ<LF>

Ø,xx3Ø6/\_\_\_\_\_177562 1775Ø2<LF>

Ø,xx31Ø/\_\_\_\_\_177564 1775Ø4<LF>

Ø,xx312/\_\_\_\_\_177566 1775Ø6<CR>

\*Ø,xx342\\_\_\_\_\_Ø 36Ø<CR>

\*E

.

[The current values for the BASE address and for the input/output vectors and status are in Table 2 of RT-11 System Release Notes. They must be copied into the new terminal's vectors.] [nnnnn are arbitrary numbers]

[xx = 16 for S/J, 17 for F/B. Modify monitor's central I/O page pointers]

[Protect new vectors]

The bootstrap must also be changed to relocate the new vector locations when the monitor is first loaded into memory. The bootstrap contains a list of items that must be relocated; the list is located at RELLST in the bootstrap code. The exact position of RELLST varies with each monitor and must be obtained from Table 2 of RT-11 System Release Notes (V02C). The patching procedure is:

.R PATCH <CR>

PATCH Version number

FILE NAME--

\*MONITR.SYS/M<CR>

\*RELLST+1Ø/ 6Ø 3ØØ<LF>

RELLST+12/ 64 3Ø4<CR>

\*E

.R PIP<CR>

\*A=MONITR.SYS/U<CR>

\*SY:/O<CR>

[Bootstrap must be rewritten. Rebootstrap; system will appear on new terminal.]



It is also possible to write a user program that would perform this procedure dynamically at run-time. Such a program would modify the monitor's protection map and the central I/O page pointers, then set up locations 300-306 and exit. If done dynamically, the monitor file itself is unchanged; thus when the system is bootstrapped, the console terminal reverts to the usual unit.

## 2.7 MAKING TTY SET OPTIONS PERMANENT IN F/B MONITOR

The F/B Monitor may be configured for different console terminal requirements by use of the TTY options of the SET command. These changes are not permanent and must be made each time the monitor is bootstrapped. By using the patching procedures in this section, the various options required for the installation may be made a permanent part of the F/B Monitor.

Table 2-4 is a description of the TTY options and their default functions in the F/B Monitor as distributed.

Table 2-4  
Default Functions for TTY Options

Option	Default	Description
TAB/NOTAB	NOTAB	Hardware tabs converted to spaces.
CRLF/NOCLRF	CRLF	<CR><LF> inserted if WIDTH reached.
FORM/NOFORM	NOFORM	Form Feed converted to Line Feeds.
FB/NOFB	FB	CTRL F/CTRL B cause context switch.
PAGE/NOPAGE	PAGE	CTRL S holds output, CTRL Q continues it.
SCOPE/NOSCOPE	NOSCOPE	VT05, VT50, VT11 is the console terminal (rubout produces backspace, space, backspace).
WIDTH	72(10)	Width of carriage.

The three options enabled are PAGE, CRLF, and FB. The carriage width is set to 72(10) characters (110 octal).

To permanently change these options, the words TTCNFG, TTWIDT and LISTFB in the F/B Monitor must be patched. The exact locations of these words and the BASE address are found in Table 2 of RT-11 System Release Notes (V02C). The numbers used in the following examples are for illustration purposes only and may not be correct for all systems.

### 2.7.1 Carriage Width

The carriage width is the line width at which the CTRL option generates a carriage return/line feed. This width is changed by patching the word TTWIDT, which for this example is assumed to be located at 21410. See Table 2 of RT-11 System Release Notes (V02C) for the exact locations of BASE and TTWIDT.

```
.R PATCH <CR>
```

```
PATCH Version number
```

```
FILE NAME--
```

```
*MONITR.SYS/M<CR>
```

```
*BASE;ØR<CR>
```

```
*Ø,2141Ø\ _____ 11Ø _____ 2Ø4<CR>
```

```
*E
```

```
:
```

[The /M is necessary; set relocation registers; open with backslash]

In this example, the width is changed from 72<sub>10</sub> to 132<sub>10</sub> (204<sub>8</sub>).

### 2.7.2 Other Options

Other options are changed by setting or clearing the appropriate bits in TTCNFG. To determine the new value to be inserted in TTCNFG, Table 2-5 is used. For each option, select the permanent value desired. Add together the octal bit patterns for each value selected to determine the new value of TTCNFG.

Table 2-5  
TTCNFG Option Bits

Option	Bit Pattern
TAB	000001
CRLF	000002
FORM	000004
FB	000010
PAGE	000200
SCOPE	100000
Any NO option	000000

For example, the monitor default is PAGE, CRLF and FB. Adding together the bit patterns for PAGE, CRLF and FB produces the octal value 212 (= 200 + 10 + 2).

To change this to SCOPE, PAGE, FB, add together the numbers 100000, 200 and 10 to get 100210, the new value of TTCNFG. Using the location of TTCNFG obtained from Table 2 of RT-11 System Release Notes is:

```
.R PATCH <CR>
PATCH Version number
FILE NAME--
*MONITR.SYS/M<CR>
*BASE;ØR<CR>
*Ø,TTCNFG/_____ 212 _____ 1ØØ21Ø<CR>
*E
:
```

If the FB option is changed, an additional step is necessary. Bit 15 of LISTFB must be changed to reflect the new FB option. Bit 15 must be 0 if the option is FB and must be 1 if the option is NOFB. For example, to change the monitor default to FORM, TAB, NOFB, the value of TTCNFG is 5 (4 + 1 + 0), and bit 15 of LISTFB must be a 1. The patch procedure is:

.R PATCH <CR>

PATCH Version number

FILE NAME--

\*MONITR.SYS/M<CR>

[The /M is necessary;

\*BASE;ØR<CR>

set relocation register;

\*Ø, TTCNFG/ 212 5<CR>

change TTCNFG;

\*Ø, LISTFB/ 3316 1Ø3316<CR>

set bit 15 in LISTFB.]

\*E

:

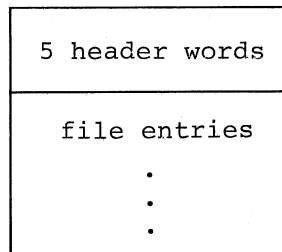
After making any of these patches, it is necessary to bootstrap the system to load the new version of the monitor.

CHAPTER 3  
FILE STRUCTURES AND FILE FORMATS

3.1 DEVICE DIRECTORY SEGMENTS

The device directory begins with physical block 6 of any directory-structured device and consists of a series of directory segments that contain the names and lengths of the files on that device. The directory area is variable in length, from 1 to 31 (decimal) directory segments. PIP allows specification of the number of segments when the directory is zeroed. The default value is four directory segments. Each directory segment is made up of two physical blocks; thus, a single directory segment is 512 words in length.

A directory segment has the following format:



3.1.1 Directory Header Format

Each directory segment contains a 5-word header block, leaving 507 (decimal) words for directory entries. The contents of the header words are described in Table 3-1.

Table 3-1  
Directory Header Words

Word	Contents
1	The number of segments available for entries. This number is specified in PIP when the device is zeroed and must be in the range $1 \leq N \leq 31_{10}$ .
2	Segment number of the next logical directory segment. The directory may, in certain cases, be a linked list. This word is the link word between logically contiguous segments; if equal to 0, there are no more segments in the list. Refer to Section 3.2.1, Directory Segment Extensions, for more details on the link word.
3	The highest segment currently open (each time a new segment is created, this number is incremented). This word is updated only in the first segment and is unused in any but the first segment.
4	The number of extra bytes per directory entry. This number can be specified when the device is zeroed with PIP. Currently, RT-11 does not allow direct manipulation of information in the extra bytes.
5	Block number where files in this segment begin.

### 3.1.2 Directory Entry Format

The remainder of the segment is filled with directory entries. An entry has the following format:

STATUS WORD	
NAME (CHARS 1-3)	
NAME (CHARS 4-6)	
EXTENSION	
TOTAL FILE LENGTH	
JOB #	CH #
DATE	
EXTRA WORDS	
.	
.	
.	

} IN RAD5Ø

} OPTIONAL

Figure 3-1  
Directory Entry Format

3.1.2.1 Status Word - The Status Word is broken down into two bytes of data:

Even byte: Reserved for future use.

Odd byte: Indicates the type of entry. Currently RT-11 recognizes the file types listed in Table 3-2:

Table 3-2  
File Types

Value	File Type
1	Tentative File, i.e., one that has been .ENTERed but not .CLOSEd. Files of this type are deleted if not eventually .CLOSEd and are listed by PIP as <UNUSED> files.
2	An empty file. The name, extension, and date fields are not used. PIP lists an empty file as <UNUSED> followed by the length of the unused area.

(continued on next page)

Table 3-2 (Cont.)

File Types

Value	File Type
4	A permanent entry. A tentative file that has been .CLOSEd is a permanent file. The name of a permanent file is unique; there can be only one file with a given name and extension. If another exists before the .CLOSE is done, it is deleted by the monitor as part of the .CLOSE operation.
10	End-of-segment marker. RT-11 uses this to determine when the end of the directory segment has been reached during a directory search.

3.1.2.2 Name and Extension - These three words (in .RAD50) represent the symbolic name and extension assigned to a file.

3.1.2.3 Total File Length - The file length consists of the number of blocks currently a part of the file. Attempts to read or write outside the limits of the file result in an End of File error.

3.1.2.4 Job Number and Channel Number - A tentative file is associated with a job in one of two ways:

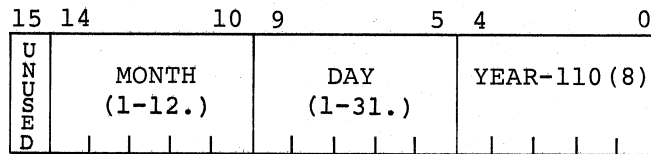
1. Under the S/J Monitor, the sixth word of the entry holds the channel number on which the file is open. This enables the monitor to locate the correct tentative entry for the channel when the .CLOSE is given. The channel number is loaded into the even byte of the sixth word.
2. In F/B, the channel number is put into the even byte of the sixth word; in addition, the number of the job that is opening the file is put into the odd byte of the word. This is required to uniquely identify the correct tentative file during the .CLOSE and is necessary because both jobs may have files open on their respective channels; the job number differentiates the tentative files.

NOTE

This sixth word is used only when the file is marked as tentative. Once it becomes permanent, the word becomes unused. Its function while permanent is reserved for future use.



3.1.2.5 Date - When a tentative file is created via .ENTER, the system date word is put into the creation date slot for the file. The date word is in the following format:



3.1.2.6 Extra Words - The number of extra words is determined by the number of extra bytes per entry in the header words. Although PIP provides for allocation and listing of extra words, RT-11 provides no direct facilities for manipulating this extra information. Any user program wishing to access these words must perform its own direct operations on the RT-11 directory.

Figure 3-2 shows a typical RT-11 directory segment:

HEADER BLOCK	4	FOUR SEGMENTS AVAILABLE
	0	NO NEXT SEGMENT
	1	HIGHEST OPEN IS #1
	0	NO EXTRA WORDS/ENTRY
	16	FILES START AT BLOCK $16_8$
FILE ENTRIES	2000	PERMANENT ENTRY
	51646	RAD5Ø FOR "MON"
	35562	RAD5Ø FOR "ITR"
	75273	RAD5Ø FOR "SYS"
	42	FILE IS $34_{10}$ ( $42_8$ ) BLOCKS LONG
	0	
	0	NO CREATION DATE
	1000	AN EMPTY ENTRY
	0	(THE NAME AND EXTENSION OF AN
	0	EMPTY IS NOT IMPORTANT
	0	
	100	$64_{1Ø}$ ( $1ØØ_8$ ) BLOCKS LONG
	0	
	0	
	2000	PERMANENT
62570	RAD5Ø FOR "PIP"	
0		
50553	RAD5Ø FOR "MAC"	
11	FILE IS $9_{10}$ ( $11_8$ ) BLOCKS LONG	
0		
0	NO CREATION DATE	
4ØØ	TENTATIVE FILE ON CHANNEL 1	
62570	RAD5Ø FOR "PIP"	
0		
50553	RAD5Ø FOR "MAC"	
20		
1	JOB #, CHANNEL #	
0		
1000	EVERY TENTATIVE MUST BE FOLLOWED BY	
0	AN EMPTY ENTRY	
0		
0		
1020	FILE IS $528_{1Ø}$ ( $1Ø2Ø_8$ ) BLOCKS LONG	
0		
0		
4000	END OF DIRECTORY SEGMENT	

Figure 3-2  
Directory Segment

When the tentative file PIP.MAC is .CLOSEd, the permanent file PIP.MAC is deleted.

To find the starting block of a particular file, first find the directory segment containing the entry for the desired file. Then take the starting block number given in the fifth word of that directory segment and add to it the length of each file in the directory before the desired file. For example, in Figure 3-2, the permanent file PIP.MAC will begin at block number 160 (octal).

### 3.2 SIZE AND NUMBER OF FILES

The number of files that can be stored on an RT-11 device depends on the number of segments in the device's directory and the number of extra words per entry. The maximum number of directory segments on any RT-11 device is  $31_{10}$ . This theoretically leaves room for a maximum of:

$$31 \times \left[ \frac{512-5}{7+N} \right]$$

directory entries, where N equals the number of extra information words per entry. If N=0, this indicates that the maximum is  $2232_{10}$  entries.

If files are added sequentially (that is, one immediately after another) without deleting any files, roughly one half the total number of entries will fit on the device before a directory overflow occurs. This results from the way filled directory segments are handled.

When a directory segment becomes full and it is necessary to open a new segment, approximately one half the entries of the filled segment are moved to the newly-opened segment (this process is illustrated in Section 3.2.1); thus, when the final segment is full, all previous segments have approximately one half their total capacity. If this process were not done and a file was deleted from a full segment, the space from the deleted file could not be reclaimed. Every tentative file must be followed by an empty entry (for recovering unused blocks when the file is made permanent). Though only one file is deleted, two entries (tentative and empty) are needed to reclaim the space.

If files are continuously added to a device, the maximum number of entries will be:

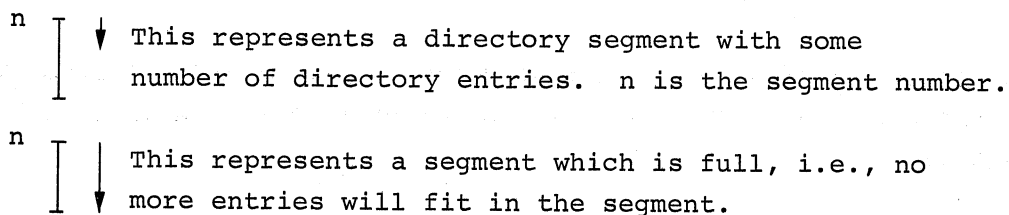
$$(M+1) \left[ \frac{507}{2(7+N)} \right]$$

where M equals the number of segments available on the device and N equals the number of extra words.

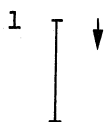
The theoretical total can be realized by compressing the device (using the PIP /S operation) when the directory fills up. PIP packs the directory segments as well as the physical device.

### 3.2.1 Directory Segment Extensions

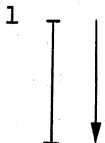
RT-11 allows a maximum of 31 (decimal) directory segments. This section covers the processing of a directory segment. For illustrative purposes, the following symbols are used:



Systems start out with entries entered into segment 1:



As entries are added, segment 1 fills:



When this occurs and an attempt is made to add another entry to the directory, the system must open another directory segment. If another segment is available, the following occurs:

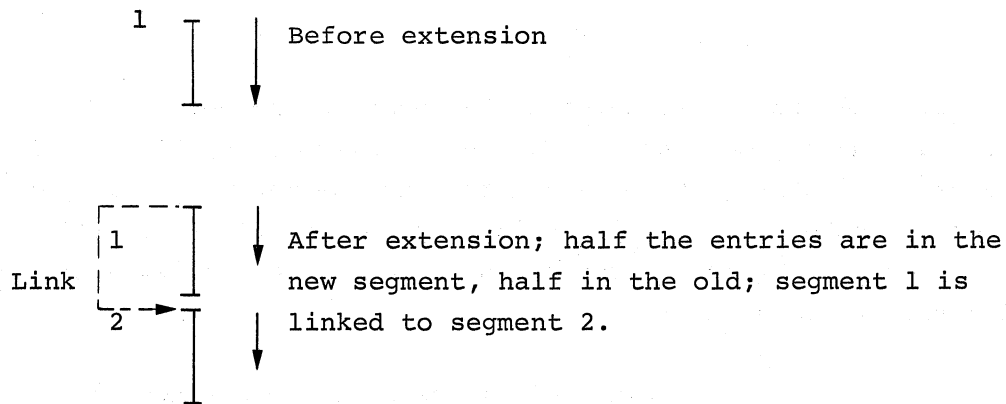
1. one half of the entries from the filled segment are put into the next available segment,
2. the shortened segment is re-written to the disk,
3. the directory segment links are set, and
4. the file is entered in the newly created segment.

NOTE

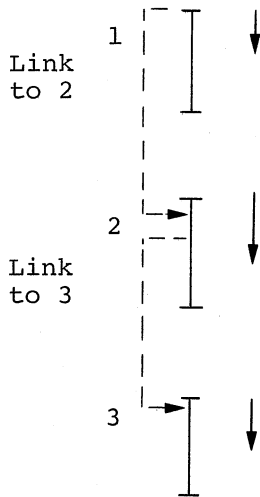
If the last segment becomes full and an attempt is made to enter another file, a fatal error occurs and an error message is generated:

?M-DIR OVFLO?

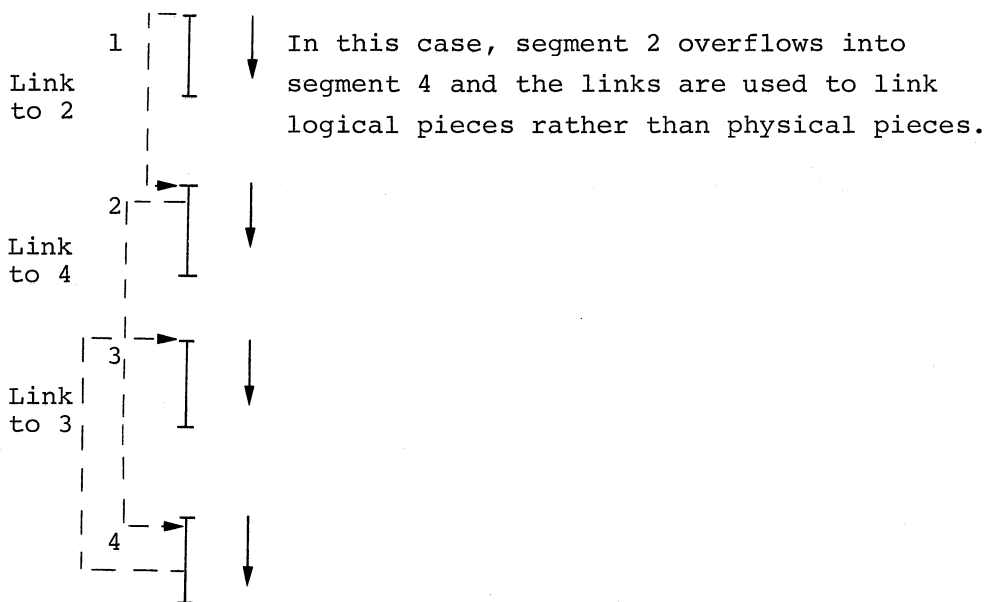
Thus, in the normal case, the segment appears as:



If many more files are entered, they fill up the second segment and overflow into the third segment, if it is available:



In this case, the links between the segments are not strictly necessary, as the segments are contiguous. However, the links do become necessary if a large file is deleted from segment 2 and many small files are entered, since it would then be possible to overflow segment 2 again. If this occurred and a fourth segment existed, the directory would appear:



### 3.3 MAGTAPE AND CASSETTE FILE STRUCTURE

#### 3.3.1 Magtape File Structure

This section covers the magtape file structure as implemented in RT-11, Versions 2B and 2C. The structure is slightly different from that of Version 2. However, RT-11 V02B and V02C can read magtapes written under Version 2.

RT-11 magtapes use a subset of the VOL1, HDR1, and EOF1 ANSI standard labels. Each magtape file has the format:

```
HDR1*---data---*EOF1*
```

where each asterisk represents a tape mark.

A volume containing a single file has the following format:

```
VOL1 HDR1*---data---*EOF1**
```

A volume containing two files has the following format:

```
VOL1 HDR1*---data---*EOF1*HDR1*---data---*EOF1**
```

A double tape mark following an EOF1 label indicates logical end of tape.

A zeroed magtape has the following format:

```
VOL1**
```

Each label occupies the first 80 bytes of a 256-word physical block, and each byte in the label contains an ASCII character (i.e., if the content of a byte is listed as '1', the byte contains the ASCII code for '1'). Table 3-3 shows the contents of the first 80 bytes in the three labels. Note that VOL1, HDR1, and EOF1 each occupy a full 256-word block, of which only the first 80 bytes are meaningful.

The meanings of the table headings are:

CP - character position in label  
Field Name - reference name of field  
L - length of field in bytes  
Content - content of field

Table 3-3  
ANSI MT Labels Under RT-11

Volume-Header Label (VOL1)			
CP	Field Name	L	Content
1-3	Label identifier	3	VOL
4	Label number	1	1
5-10	Volume identifier	6	RT1101
11	Accessibility	1	Blank
12-37	(Reserved)	26	Blanks
38-51	Owner identifier	14	DD%% {used to indicate an RT-11 MT to RSX-11D
52-79	(Reserved)	28	Blanks
80	Label-Standard Version	1	1
First File Header Label (HDR1)			
CP	Field Name	L	Content
1-3	Label identifier	3	HDR
4	Label number	1	1
5-21	File identifier	17	6-character ASCII file name, followed by '.', followed by 3-character ASCII file extension; left justified, remainder of field is blanks
22-27	File Set identifier	6	RT1101
28-31	File Section Number	4	0001
32-35	File Sequence Number	4	0001
36-39	Generation Number	4	0001
40-41	Generation Vsn Number	2	00
42-47	Creation Date	6	Blank then year*1000+day of year in ASCII ( $\Delta$ YYDDD); e.g., 2/1/75= $\Delta$ 75032
48-53	Expiration Date	6	blank then 00000
54	Accessibility	1	blank
55-60	Block Count	6	000000
61-73	System Code	13	RT11 left-justified followed by blanks
74-80	(Reserved)	7	blanks
<u>First End-of-File Label (EOF1)</u>			
Same as HDR1 except that the label identifier (CP 1-3) is <i>EOF</i> , not <i>HDR</i> , and the block count field (CP 55-60) contains the number of blocks in the file as a decimal value encoded in ASCII characters (for example, if the file was 12 blocks long, the block count field would be 00012).			



3.3.1.1 Bootable Magtape File Structure - An RT-11 bootable magtape is a multi-file volume that has the following format:

```
VOL1 BOOT HDR1*---data---*EOF1**
```

where BOOT is a 256-word physical block containing the magtape bootstrap loaders.

The format of the bootable magtape is not standard, because of the BOOT block, but other systems that will skip the BOOT block to HDR1 will be able to read RT-11 bootable magtapes if they can read regular RT-11 magtapes.



3.3.1.2 Moving MT to Other Industry-Compatible Environments - RT-11 V02C magtapes may be read by RSX-11D Version 6. RT-11 magtapes should be mounted, under RSX-11D, by using the /OVR switch of the MOUNT command, or by specifying a volume label of "RT1101". RSX-11D Version 6 will not allow the user to write on RT-11 V02B magtapes once they have been mounted. RT-11 V02C can read RSX-11D Version 6 magtapes, but RT-11 users should not attempt to write on tapes created by RSX-11D. Users should note that data structures differ between the two systems and these differences must be handled by the user.

RT-11 V02C magtapes may be read on IBM systems that support ANSI standard label processing. RT-11 V02C magtapes to be read by IBM systems should consist of single file volumes (one file per magtape). Important JCL parameters for reading RT-11 V02C tapes under an IBM OS system are as follows:

(In the DD statement of the Job Control Language)

```
DISP = OLD
LABEL = (01,AL,,IN)
VOL = (,RETAIN,SER=RT1101)
DSN = RTFILE.MAC
BLKSIZE = 512
DEN = 2      (for 800 bpi 7-track or 9-track tape)
```

The DSN parameter is the Data Set Name or the RT-11 filename and extension. Files to be moved to other systems should be created with full 6-character filenames and 3-character extensions; filenames less than 6 characters should be enclosed in quotes.

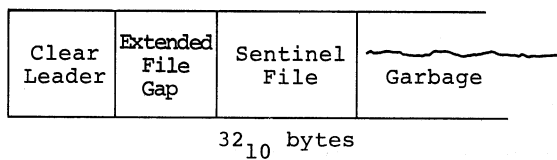
3.3.1.3 Recovering From Bad Tape Errors - When a bad tape error occurs on magtape, the magtape handler will retry the desired function, and, if the error persists, will attempt to save the tape's file structure. It does this on writes, for example, by retrying the write 10 times, using the write with extended file gap to space past the bad tape. If, after retrying, the error still exists, the file will be closed, containing all data written prior to the write on which the error occurred. The user should still be able to write additional files on the tape, since the bad portion of the tape will be within the area of the closed file.

If a bad tape error occurs when writing the file header during ENTER, and retry fails, the handler writes logical end of tape after the previous file on the tape. The remainder of the tape can be accessed only if the last complete file on the tape can be extended (or overwritten by a file of different length) so that the bad tape error does not occur on the file header when a subsequent file is ENTERed.

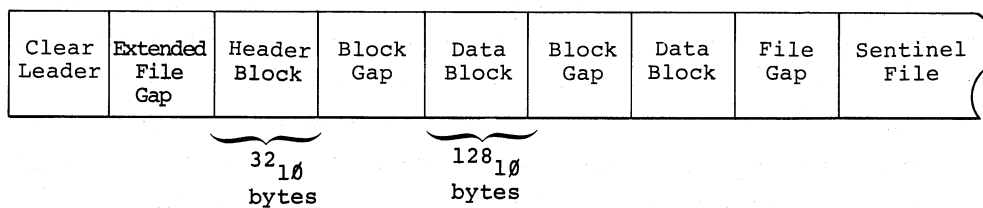
If a bad tape error occurs while writing the end of file label (EOF1) during CLOSE, the handler writes a triple tape mark to signify end of file and logical end of tape. Additional files can be added to the tape only if the last complete file can be extended (or overwritten by a file of different length) so that the bad tape error does not occur at the EOF1 label.

### 3.3.2 Cassette File Structure

A blank (newly initialized) cassette appears in the format:

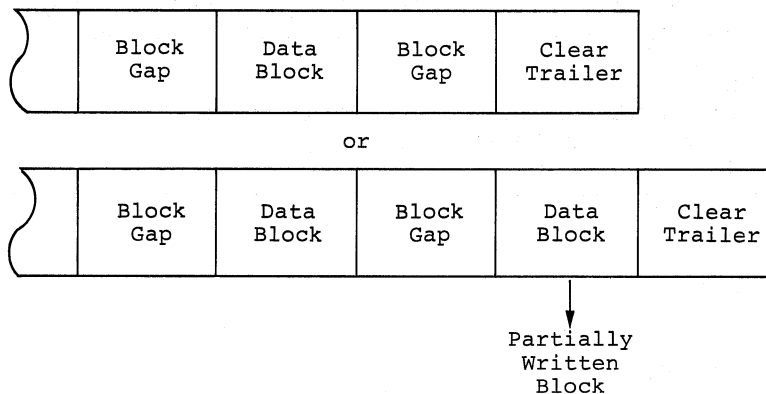


while a cassette with a file on it appears as:



Files normally have data written in  $128_{10}$ -byte blocks. This can be altered by writing cassettes while in *hardware* mode. (In hardware mode, the user program must handle the processing of any headers and sentinel files; in *software* mode the handler automatically does this. Refer to Appendix H of the RT-11 System Reference Manual.)

The preceding diagram shows a file terminated in the usual manner (by a sentinel file). However, the physical end of cassette may occur before the actual end of the file. This format appears as:



In the latter case, for multi-volume processing the partially written block must be the first data block of the next volume.

3.3.2.1 File Header - The File Header is a  $32_{10}$ -byte block that is the first block of any data file on a cassette. If the first byte of the header is null, the header is interpreted as a sentinel file, which is an indication of logical end of cassette. The format of the header is described in Table 3-4.

Table 3-4  
CT File Header Format

Byte Number	Contents
0-5	File name in ASCII characters (ASCII is assumed to imply a 7-bit code)
6-8	Extension in ASCII characters
9	Data type (0 for RT-11)
10,11	Block length of $128_{10}$ ( $200_8$ ); Note: byte 10=0 (high-order), byte 11= $200_8$ (low-order)
12	File sequence number. (0 for a single-volume file or the first volume of a multi-volume file; successive numbers are used for continuations)
13	Level 1; this byte is a 1
14-19	Date of file creation (6 ASCII digits representing day (01-31); month (01-12), and last two digits of the year; 0 or $40_8$ in first byte means no date present)
20,21	Zero
22	Record attributes (0 in RT-11 cassettes)
23-28	Reserved for future use
29-31	Reserved for user

### 3.4 RT-11 FILE FORMATS

#### 3.4.1 Object Format (.OBJ)

An object module is a file containing a program or routine in a binary, relocatable form; object files normally have an .OBJ extension. Object modules are produced by language processors (such as MACRO or FORTRAN) and are processed by the Linker to become a runnable program (in SAV, LDA, or REL format, discussed later). Object files may also be processed by the Librarian to produce library .OBJ files, which are then used by the Linker. Figure 3-3 illustrates this process.

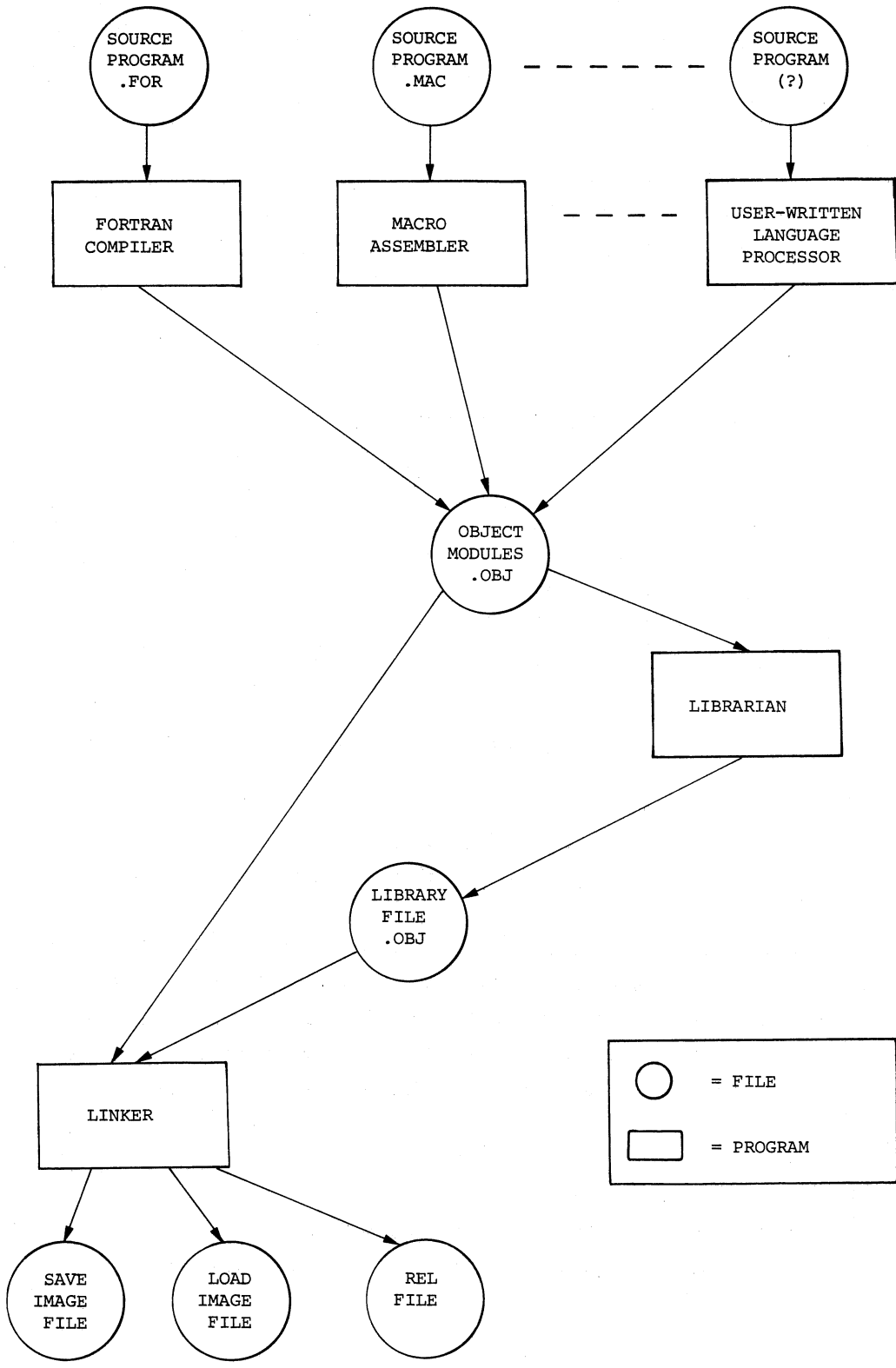


Figure 3-3  
Object Module Processing

Many different object modules may be combined to form one file; each object module remains complete and independent. However, object modules combined into a library by the Librarian are no longer independent -- they become part of the library's structure.

Object modules are made up of formatted binary blocks. A formatted binary block is a sequence of 8-bit bytes (stored in an RT-11 file, on paper tape, or by some other means) and is arranged as illustrated in Figure 3-4.

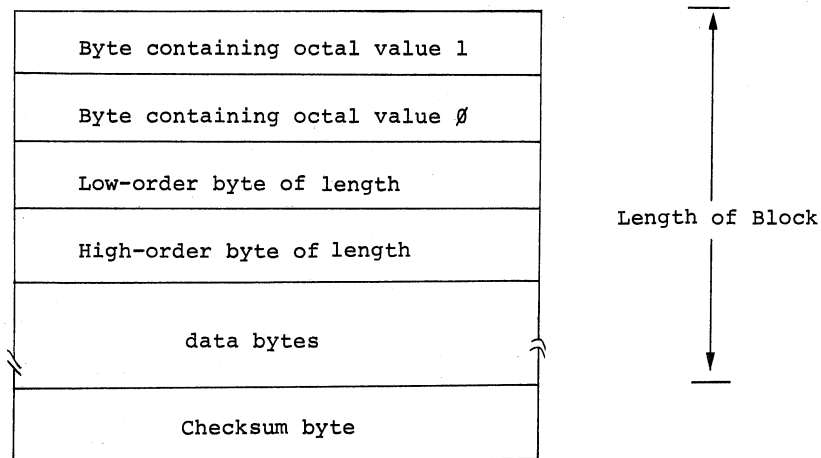


Figure 3-4  
Formatted Binary Block

Each formatted binary block has its length stored within it; the length includes all bytes of the block except the checksum byte. The data portion of each formatted binary block contains the actual object module information (described later). The checksum byte is computed such that the sum of all bytes in the formatted binary block, including the checksum byte, is zero when the sum is masked to 8 bits.

Formatted binary blocks are used to hold various kinds of information in an object module; this information is always contained completely in the data portion of the block, surrounded by the formatted binary block structure.



Eight types of data blocks may be present in an object module:

<u>Identification Code</u>	<u>Type of Block</u>	<u>Function</u>	
1	GSD blocks	hold the Global Symbol Directory information	
2	ENDGSD block	signals the end of GSD blocks in a module	
3	TXT blocks	hold the actual binary "text" of the program	
4	RLD blocks	hold Relocation Directory information	
5	ISD blocks	hold Internal Symbol Directory - not supported by RT-11	
6	ENDMOD block	signals end of the object module	
7	Librarian Header Block	17 words holding the status of the library file	} Library File Only
10	Librarian End Block	signals the end of the library file	

The structure of object modules produced by a language processor will be described first, followed by details specific only to Library .OBJ files.

The first block of an object module must be a GSD block, and all GSD blocks must appear before the ENDGSD block. The ENDMOD block must be the last block of the module. Except for these three restrictions, blocks may appear in any order within an object module.

When a 16-bit word is stored as part of the data in a block, it is always stored as two consecutive 8-bit bytes, with the low-order byte first.

The first word (data word) of each type of block mentioned above contains the identification code of that block type (1 = GSD block, etc.) with any information present following the identification word.

3.4.1.1 Global Symbol Directory - The object module's global symbol directory contains the following information:

- 0 - Module Name
- 1 - Program Section (CSECT) Definitions
- 2 - Internal Symbol Table Name (not supported by RT-11)
- 3 - Transfer (Start) Address
- 4 - Global Symbol Definitions or References

Each piece of information in the GSD is contained in a *GSD item*, formatted as shown in Figure 3-5:

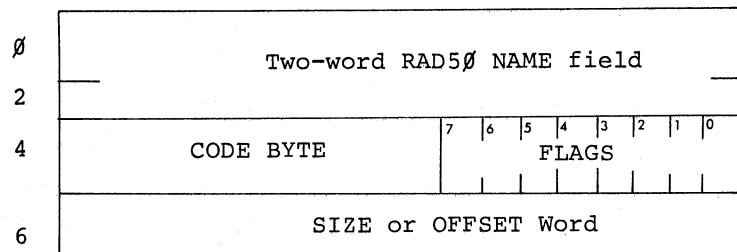


Figure 3-5  
GSD Structure

The code byte identifies the information contained in a GSD item according to the codes listed above (0 = Module Name, 1 = Program Section Definition, etc.). The first GSD item of an object module must contain the Module Name information (FLAGS, CODE, and SIZE = 0). There may be no more than five GSD items per GSD block (i.e., per formatted binary block). As many GSD blocks as necessary may be present, but all must appear before the ENDGSD block. GSD blocks need not be contiguous.

Flags are coded as follows:

- Bits 0,1,2,4,7      unused
- Bit 3:              0 = undefined, 1 = defined (used only with Global Symbols)
- Bit 5:              0 = absolute, 1 = relocatable
- Bit 6:              0 = internal, 1 = global

All program sections (CSECTs) defined in a module must be declared in GSD items (code byte = 1). The size word of each program section definition should contain the size in bytes to be reserved for the section. Program sections may be declared more than once, in which case the largest declared size of the section will be used. All global symbols that are defined in a given program section must appear in the GSD items immediately following the definition item of that program section.

A special program section named ".  ABS." (where    represents a space) is called the absolute section. The absolute section has the special attribute that it is always allocated by the Linker beginning at location 0 of memory. All global symbols that contain absolute (non-relocatable) values should be declared immediately after the GSD item that defines the absolute section. If it is not desired to allocate any memory space to the absolute section, its size word may be specified as zero, even if absolute global symbol definitions occur after it. Flag bit 5 of each absolute global symbol is always set to zero. GSD items that contain the definitions of global symbols (code byte = 4) must immediately follow the program section declaration into which they are to be defined. Flag bit 3 is set to 1 to indicate a symbol definition, bit 5 is set if and only if the symbol is relocatable, and bit 6 is set to indicate that the symbol being defined is a global. In addition, the offset word is set to contain the defined value of the global symbol, relative to the base of the program section in which the global is defined. At link time, the Linker assigned section base is added to get the final value of the global symbol.

Global symbols that are referenced but not defined in the current object module must also appear in GSD items. These *global references* may appear in any GSD item except the very first (which contains the module name). Global references are recognized by code byte 4 with flag bit 3=0, bit 5 is undetermined, and bit 6=1. All global symbols used in the RLD of the object module (described later) must appear in at least one Global Symbol or Program Section GSD item.

If RT-11 is to begin execution of a program within a particular object module of that program, then the information on where to start is given in a Transfer Address (code=3) GSD item. The first even transfer address encountered by the Linker will be passed to RT-11 as the program start address. Whenever the resulting program is run (using R or RUN

for SAV images, FRUN for REL files, or the absolute loader for LDA files), the start address is used to indicate the first executable instruction. If no transfer address is present or if all are odd, the resulting program will not self-start when run. In a Transfer Address GSD item, the name field is used to specify a program section (or global name) and the offset word is used to indicate the offset from the base of that program section (or global) to the starting point of the program. The program section or global name referenced need not be defined in the current object module, but must be defined in some object module included at link time.

NOTE

Program Section and Global names must begin with an alphabetic or numeric character, except for the names `._ABS.` and `._L.L.L.L.L.L.L.L.`

3.4.1.2 ENDGSD Block - The ENDGSD block contains a single data word, and that is the identification code of the ENDGSD block (2). All GSD blocks in an object module must precede the ENDGSD block.

3.4.1.3 TXT Blocks and RLD Blocks - The first TXT block (3) in an object module (if present) must be preceded by an RLD block (4).

TXT blocks contain the actual binary form of the programs and are formatted as shown in Figure 3-6:

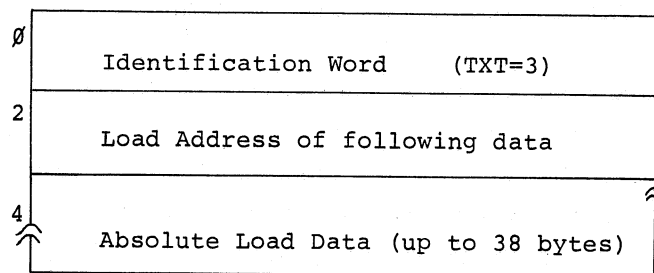


Figure 3-6  
TXT Block Format

The load address of a TXT block gives the relative address of the first byte of the absolute load data. The address is relative to the base of the last program section given in a Location Counter Definition RLD command (explained later).

The Absolute Load Data contains the actual bytes that will be loaded into memory when the program is run (except for relocations, described later).

RLD blocks contain variable length RLD commands, used to modify and complete the information contained in TXT blocks. Except for the Location Counter commands, RLD information must appear in an RLD block immediately following the TXT block to be modified.

Available RLD commands are:

1. Internal Relocation
2. Global Relocation
3. Internal Displaced Relocation
4. Global Displaced Relocation
5. Global Additive Relocation
6. Global Additive Displaced Relocation
7. Location Counter Definition
8. Location Counter Modification (not used by RT-11)
9. Set Program Limits

The location counter commands (numbers 7 and 8) are the only two RLD commands that must appear in an RLD block preceding the text blocks modified. The first RLD block must precede the first TXT block and must contain only a location counter definition command (7) in order to declare a program section for loading the first text block. (The location counter modification command (8) is included for compatibility with other systems, but is not used by RT-11.)

The data portion of an RLD block must not be larger than  $42_{10}$  bytes including the identification word (RLD=4) and all RLD commands.

All global names and program section names that appear in RLD commands must appear in GSD items in the same object module. Figure 3-7 shows the format of each RLD command (each part except the first word is optional and may not appear in some commands):

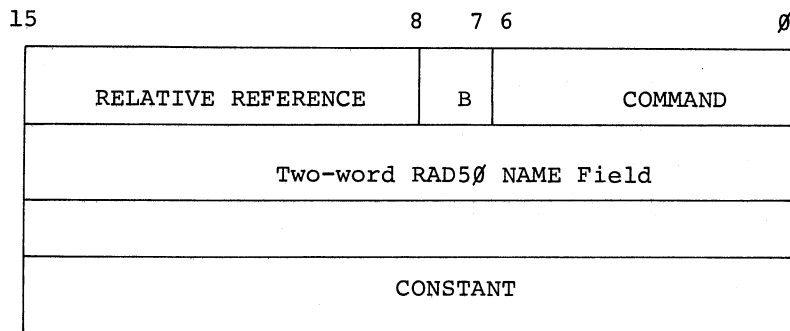


Figure 3-7  
RLD Format

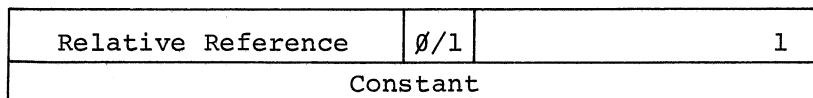
An RLD command may be 1, 2, 3, or 4 words long.

The Command Field contains the command code (1 = Internal Relocation, etc.). The Command Field occupies bits 0-6 of the first word of the command. The B field (bit 7) indicates a word command if 0 or a byte command if 1 (only valid for commands 1 through 6). The Relative Reference Field is a pointer into the preceding TXT block and is used with RLD commands that require text locations for modification (commands 1 through 6 and 9). This field specifies the displacement from the beginning of the preceding TXT block to the referenced text data byte (or word). The beginning of the TXT block is the identification word (the first word of the data portion of the block). Thus, the smallest relative reference will normally be 4 (the first byte (word) of the preceding TXT block).

The Name Field is used to hold a Global or Program Section name if the command requires it.

The Constant Field is used to hold a relative address or additive quantity if the command requires it. RLD commands are processed by the Linker as shown in the following situations:

1. Internal Relocation (code 1) - Add the current program section's base to the specified constant and place the result where indicated. This command relocates a direct pointer to an internal relocatable symbol.



Examples:

- a) .WORD LOCAL
- b) MOV #LOCAL,%Ø

2. Global Relocation (code 2) - Place the value of the specified global symbol where indicated. This command generates a direct pointer to an external symbol.

Relative Reference	Ø/1	2
Global Name		

Examples:

- a) .WORD GLOBAL
- b) MOV #GLOBAL,RØ

3. Internal Displaced Relocation (code 3) - Calculate the displacement from the position of the current location plus two to the specified absolute address, and store the result where indicated. This command occurs only when there is a reference to an absolute (non-relocatable) location from a relocatable section.

Relative Reference	Ø/1	3
Constant		

Examples:

- a) ABS=17755Ø
  - TST ABS
  - b) CLR 17755Ø
- } both addresses cause internal displaced relocation to occur

4. Global Displaced Relocation (code 4) - Calculate the displacement from the current location plus two to the specified global address, and store the result where indicated.

Relative Reference	Ø/1	4
Global Name		

Example:

```
.GLOBL GLOBAL
MOV GLOBAL,RØ
```

5. Global Additive Relocation (code 5) - Add the value of the specified global symbol to the specified constant, and store the result where indicated.

Relative Reference	Ø/1	5
Global Name		
Constant		

Example:

```
.GLOBL GLOBAL
CMP #GLOBAL+6,RØ
```

6. Global Additive Displaced Relocation (code 6) - Calculate the displacement from the current location plus two to the address specified by the sum of the global symbol value and the given constant, and place the result where indicated.

Relative Reference	Ø/1	6
Global Name		
Constant		

Example:

```
.GLOBL GLOBAL
CLR GLOBAL+6
```



7. Location Counter Definition (code 7) - This command is used to specify the program section into which the following TXT blocks are to be loaded.

7
Program Section Name
Constant

This command is generated whenever .ASECT or .CSECT is used to initiate or continue a program section. The constant word is effectively ignored by RT-11 and may be used for diagnostic purposes to indicate the relative point at which a program section is being entered.

8. Location Counter Modification (code 10<sub>8</sub>) - This command is used to enter the current program section at a different point. This command is effectively ignored by RT-11 and is used for diagnostic purposes only.

10
Constant

Examples:

- a) `.=100 ;IF WE ARE IN THE ASECT`  
 b) `.=.-20 ;IF WE ARE IN A RELOCATABLE SECTION`

9. Set Program Limits (code 11<sub>8</sub>) - This command (generated by the .LIMIT assembler directive) causes two words in the preceding TXT block to be modified. The first word is to be set to the lowest relocated address of the program. The second word is to be set to the address of the first free location following the relocated code. Note that both words to be modified must appear in the same TXT block.

Relative Reference	11
--------------------	----

In addition to the above commands, note that commands numbered 14<sub>8</sub>, 15<sub>8</sub>, and 16<sub>8</sub> can be generated by MACRO. These commands are identical to commands 4, 5, and 6 respectively, but are used when the *global* is really a program section name.

3.4.1.4 ISD Internal Symbol Directory - Not supported by RT-11.

3.4.1.5 ENDMOD Block - Every object module must end with an ENDMOD block. The ENDMOD block contains a single data word -- the identification code of the ENDMOD block (6).

3.4.1.6 Librarian Object Format - A library .OBJ file contains information additional to that previously defined. The object modules in a library file are preceded by a Library Header Block and Library Directory, and are followed by the Library End Block or trailer. This is illustrated in Figure 3-8.

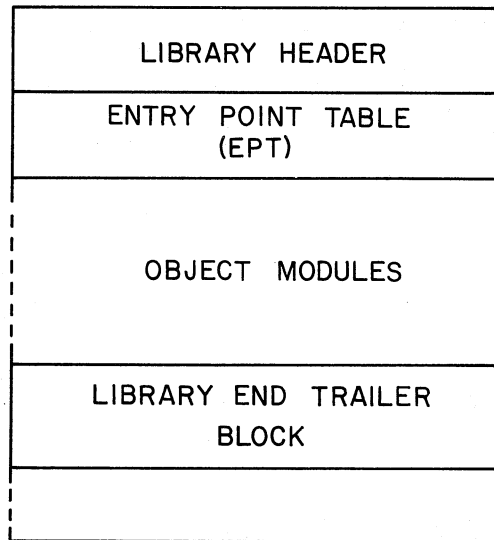


Figure 3-8  
Library File Format

Diagrams of each component in the library file structure are included here, but Chapter 7 of the RT-11 System Reference Manual should be consulted for details.

The library header is composed of 17<sub>10</sub> words describing the status of the file. The contents of the 17 words are shown in Figure 3-9.

1	} FORMATTED BINARY BLOCK HEADER
56 <sub>8</sub>	
7	LIBRARIAN CODE
X	VERSION NUMBER
0	RESERVED
X	MONTH-DAY-YEAR (OR Ø IF NO DATE)
0	} RESERVED
0	
0	
0	
0	
12 <sub>8</sub>	EPT RELATIVE START ADDRESS
X1	EPT ENTRIES ALLOCATED IN BYTES
0	EPT ENTRIES AVAILABLE (NOT USED IN VI)
X2	NEXT INSERT RELATIVE BLOCK NUMBER
X3	NEXT BYTE WITHIN BLOCK
0	NOT USED (MUST BE ZERO)

Figure 3-9  
Library Header Format

The Entry Point Table (EPT), Figure 3-10, is composed of four-word entries which contain information related to all object modules in the library file.

0	SYMBOL CHARS 1-3 (RAD5Ø)		
2	SYMBOL CHARS 4-6 (RAD5Ø)		
4		ADDRESS OF BLOCK	BIT 15=1-MODULE NAME Ø-CSECT OR ENTRY POINT NAME
6	# OF CSECTS IN OBJECT MODULE	RELATIVE BYTE IN BLOCK	RELATIVE BYTE MAXIMUM=777 <sub>8</sub> CSECTS MAXIMUM =177 <sub>8</sub>

Figure 3-10  
Entry Point Table Format

Object modules follow the Entry Point Table and consist of the types of data blocks already discussed: GSD, ENDGSD, TXT, RLD, and ENDMOD. The information in these blocks is used by the Linker during creation of the load module.

Following all object modules is a specially coded Library End Block (trailer), which signifies the end of the file, shown in Figure 3-11.

	1	FORMATTED BINARY HEADER
	10	FORMATTED BINARY LENGTH
	10	TYPE CODE
	0	NOT USED (MUST BE ZERO)
	357	CHECKSUM BYTE

Figure 3-11  
Library End Trailer

#### 3.4.2 Formatted Binary Format (.LDA)

The Linker /L switch produces output files in a paper tape compatible binary format.

Paper tape format, shown in Figure 3-12, is a sequence of formatted binary blocks (as explained in Section 3.4.1 and in Figure 3-4). Each formatted binary block represents the data to be loaded into a specific portion of memory. The data portion of each formatted binary block consists of the absolute load address of the block followed by the absolute data bytes to be loaded into memory beginning at the load address. There may be as many formatted binary blocks as necessary in an LDA file. The last formatted binary block of the file is special; it contains only the program start address in its data portion. If this address is even, the loader passes control to the loaded program at this address. If it is odd, the loader halts upon completion of loading. The final block of the LDA file is recognized by the fact that its length is 6.

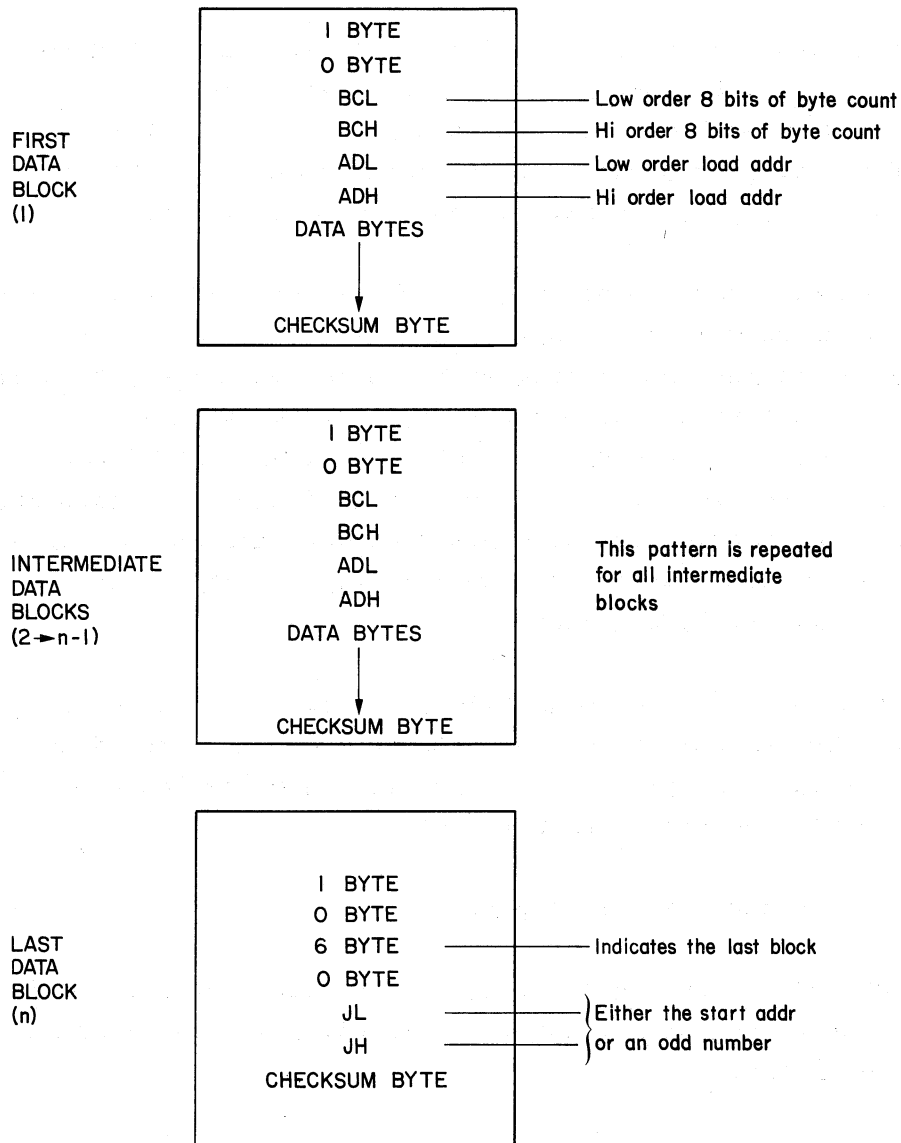


Figure 3-12  
Formatted Binary Format

The load module's binary blocks contain only absolute binary load data and absolute load addresses; all global references have been resolved and the appropriate relocation has been performed by the Linker.

### 3.4.3 Save Image Format (.SAV)

Save image format is used for programs that are to be run in the background. This format is essentially an image of the program as it would appear in memory (block 0 of the file corresponds to memory locations 0-776, block 1 to locations 1000-1776, and so forth).

Locations 360-377 in block 0 of the file are restricted for use by the system. The Linker stores the program memory usage bits in these eight words. Each bit represents one 256-word block of memory and is set if the program occupies that block of memory. This information is used by the R, RUN, and GET commands when loading the program.

When loading a save image program into memory, KMON reads block 0 of the file to extract the memory usage bits. These bits are used to determine whether the program will overlay either the KMON or the USR. If these portions of the monitor will not be overlaid, the entire program is loaded; if the USR and KMON must swap, KMON loads the resident portion of the program, up to the start of KMON. It then puts the portion of the program that overlays KMON/USR into the system swap blocks. When the program starts, the monitor swaps in the virtual portion of the program, overlaying KMON.

When block 0 of a save image file is loaded, each word is checked against the protection bit map (LOWMAP), which is resident in RMON. Locations that are protected in the map, such as location 54 and the system device vectors, are not loaded.

#### 3.4.4 Relocatable Format (.REL)

A foreground job is linked using the Linker /R switch. This causes the Linker to produce output in a linked, relocatable format, with a REL file extension.

The object modules used to create a REL file have been linked and all global references have been resolved. The REL file is not relocated, so it has an effective start address of 0, with relocation information included to be used at FRUN time. The relocation information in the file is used to determine which words in the program must be relocated when the job is installed in memory.

In order to determine if the code to be relocated (as indicated in the relocation information blocks) is to have positive or negative relocation (relative to the start address of the program), the following criteria from the text modification commands is used (R = relative address, G = global address, C = constant):

1. Internal Relocation (.WORD R) - always positive relocation (absolute)
2. Global Relocation (.WORD G) - positive relocation only if the global is not absolute

3. Internal Displaced Relocation - always negative relocation (MOV 54,R)
4. Global Displaced Relocation (MOV G,R) - negative relocation only where the global is defined as absolute elsewhere
5. Global Additive Relocation (.WORD G + C) - same as 2 above
6. Global Additive Displaced (MOV G + C,R) - same as 4 above
7. Program Counter Commands - not applicable
8. Set Program Limits - always positive relocation (requires 2 RELs; limit is two words)

There are two types of REL files to consider, those programs with overlay segments and those without.

3.4.4.1 Non-Overlay Programs - A REL file for a non-overlaid program appears as shown in Figure 3-13:

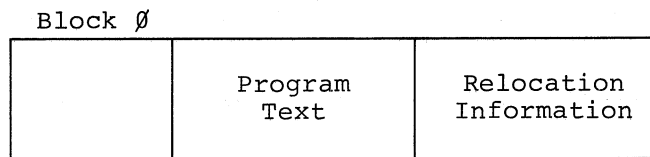


Figure 3-13  
REL File Without Overlays

Block 0 (relative to start of the file) contains certain information required by the FRUN processor:

<u>Offset from Beginning of Block 0</u>	<u>Contents</u>
52	Size of the program root segment in bytes
54	Size of the overlay region in words; 0 if no overlays
56	REL file identification word, which must contain the RAD50 value of the characters 'REL'
60	Relative block number of relocation information

In addition, the system communication locations (34-50) contain the following information:

<u>Offset from Beginning of Block 0</u>	<u>Contents</u>
34,36	TRAP vector
40	Start address of program
42	Initial setting of stack pointer
44	Job Status Word
46	USR swap address
50	Highest memory address in user's program

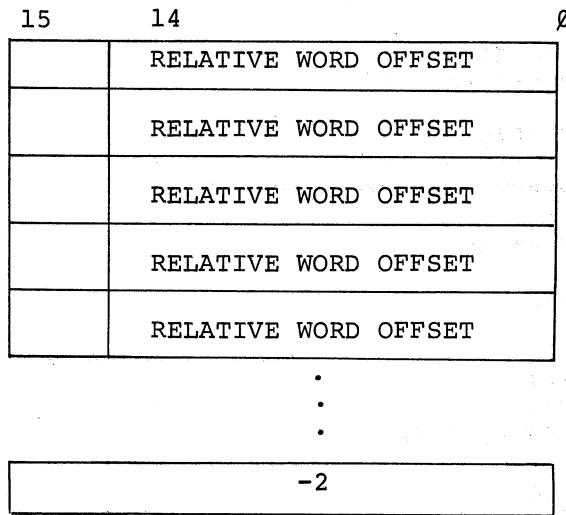
In the case of non-overlaid programs, the FRUN processor performs the following general steps to install a foreground job.

1. Block 0 of the file is read into an internal monitor buffer.
2. The amount of memory required for the job is obtained from location 52 of block 0 of the file, and the space is allocated.
3. The program text is read into the space just allocated for it.
4. The relocation information is read into an internal buffer.
5. The locations indicated in the relocation information area are relocated by adding the relocation quantity, which is the starting address the job occupies in memory.

The relocation information consists of a list of addresses relative to the start of the user's program. This list is scanned, and the appropriate locations in the user's program area are updated with a constant. The job is then ready to be started.



The relocation information is in the following format:



Bits 0-14 represent the relative address to relocate divided by two. This implies that relocation is always done on a word boundary, which is the case. Bit 15 is used to indicate the type of relocation to perform, positive or negative. The relocation constant (which is the load address of the program) is added to or subtracted from the indicated location depending on the sense of bit 15; 0 implies addition, 1 implies subtraction. 177776 terminates the list of relocation information.

Following is an example of a simple, non-overlaid program linked to produce a REL file. A dump of the file follows the program.

```

.TITLE FTFST
.MCALL ..V2...REGDEF,.LOOKUP,.READW,.QSET,.PRINT,.EXIT
..V2..
.REGDEF
ST: .QSET #QLIST,#7
    .LOOKUP #AREA,#0,#PTR
    RCC 1$
    .PRINT #LKFAIL
    .EXIT
1$: .READW #AREA,#0,#RUFF,#254.,#0
    .RCC 2$
    .PRINT #RDFAIL
    .EXIT
2$: .PRINT #OK
    .EXIT

QLIST: .BLKW 7*7
AREA: .BLKW 20.
PTR: .RAD50 /PR FILE12/
BUFF: .NLIST
    .REPT 254.
    .WORD 0
    .ENDR
    .LIST
LKFAIL: .ASCIZ /LOOKUP FAILED/
RDFAIL: .ASCIZ /READW FAILED/
OK: .ASCIZ /READW OK/
    .EVEN
    .NLIST
    .REPT <ST+1776-.>/2
    .WORD 0
    .ENDR
    .NLIST
.END ST

```

```

1 000000
2 000000
3 000000
4 000000
5 000000
6 000000
7 000000
8 000000
9 000000
10 000000
11 000000
12 000000
13 000000
14 000000
15 000000

          000007
          000144
          000306
          000001
          000002
          000356
          000004
          001364
          000306
          000010
          000002
          000364
          000400
          000000
          000000
          001402
          001417

          112700
          012746
          104353
          012700
          112760
          105010
          012760
          005060
          104375
          103004
          012700
          104351
          104350
          012700
          112760
          105010
          012760
          012760
          012760
          104375
          103004
          012700
          104351
          104350
          012700
          012760
          012760
          012760
          012760
          104375
          103004
          012700
          104351
          104350
          012700
          104351
          012700
          104351
          012700
          104351

          .MCALL
          .V2...
          .REGDEF
          .ST:
          .ITF NR <#7>,
          .GSET #QLIST,#7
          .MOV #7,%
          .FMT
          .LOOKUP #AREA,#0,#PTR
          .MOV #AREA,%
          .MOV #1,1(0)
          .CLRB (0)
          .MOV #PTR,2:(0)
          .CLP 4.(0)
          .FMT
          .ACC 1$
          .PRINT #LKFAIL
          .ITF NR <#LKFAIL>,
          .FMT
          .EXIT
          .READW #AREA,#0,#RUFF,#256.,#0
          .MOV #AREA,%
          .MOV #8.,1(0)
          .CLRB (0)
          .MOV #0,2.(0)
          .MOV #0,2.(0)
          .MOV #8UFF,4.(0)
          .MOV #0,8.(0)
          .FMT
          .ACC 2$
          .PRINT #RDFAIL
          .ITF NR <#RDFAIL>,
          .FMT
          .EXIT
          .PRINT #OK
          .ITF NR <#OK>,
          .MOV #OK,%
          .FMT
          #QLIST,=(6.)
          =0353
          #AREA,%
          #1,1(0)
          (0)
          #PTR,2:(0)
          4.(0)
          =0375
          MOV #LKFAIL,%
          =0350
          =0350
          #AREA,%
          #8.,1(0)
          (0)
          #0,2.(0)
          MOV #8UFF,4.(0)
          #0,8.(0)
          =0375
          MOV #RDFAIL,%
          =0350
          =0350
          #OK,%
          =0351

```

```

16 000142      104350      .EXIT      FMT      F0350
17 000142
18 000144
19 000306      .BLKW      7*7
20 000356      .BLKW      20.
21 001364      023320      023364      022070      .BLKW      /PR FILE12/
22 001367      114      117      120      .ASCIZ    /LOOKUP FATLFD/
23 001372      113      125      101
24 001375      040      106      105
25 001400      111      114      101
26 001402      104      000      101      .ASCIZ    /READW FAILED/
27 001405      122      105      040
28 001410      104      127      111
29 001413      106      101      104
30 001416      114      105      101
31 001417      000      105      101      .ASCIZ    /READW OK/

```

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```

001422      104      127      040
001425      117      113      000
                                .EVEN

```

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SYMBOL TABLE

```

AREA 000306R      RUFF      000364R      LKFATL 001364R      OK      001417R      PC      =X0000007
PTR 000356P      ALTST      000144R      PDFATL 001402R      R0      =X0000000
R2 =X000002      R3 =X000003      P4 =X000004      R5      =X0000001
ST 000000R      ...V2 = 000001
. ABS. 000000      000
      001776      001
ERRORS DETECTED: 0
FREE CORE: 15895. WORDS
,LP:/N:TTM/L:MF8=FTEST

```

```

BLOCK NUMBER 0000
000/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
020/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
040/ 000000 000000 000000 000000 001776 001776 000000 070524 * . . . . . TQ *
060/ 000003 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
100/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
120/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
140/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
160/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
200/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
220/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
240/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
260/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
300/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
320/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
340/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
360/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
400/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
420/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
440/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
460/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
500/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
520/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
540/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
560/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
600/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
620/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
640/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
660/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
700/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
720/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
740/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *
760/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . *

```

In block 0, word 50 shows the highest, non-relocated, memory address in the user program. Word 52 shows the program size in bytes. Word 54 shows the size of the overlay region. The value is non-zero only for programs with overlays. Word 60 contains a 3, indicating that the relocation information begins at block 3 of the file.

```

BLOCK NUMBER 0001
000/ 112700 000007 012746 000144 104353 012700 000306 112760 *.....F.N.K..F.P.*
020/ 000001 000001 105010 012760 000356 000002 005060 000004 *.....P.N...O...*
040/ 104375 103004 012700 001364 104351 104350 012700 000306 *J.....T.I.H..F.*
060/ 112760 000010 000001 105010 012760 000000 000002 012760 *P.....P.....P.*
100/ 000364 000004 012760 000400 000006 012760 000000 000010 *T.....P.....P...*
120/ 104375 103004 012700 001402 104351 104350 012700 001417 *J.....T.I.H..O...*
140/ 104351 104350 000000 000000 000000 000000 000000 000000 *I.H.....P.....*
160/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
200/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
220/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
240/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
260/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
300/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
320/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
340/ 000000 000000 000000 000000 000000 000000 000000 063320 *.....P.....PF*
360/ 023364 022070 000000 000000 000000 000000 000000 000000 *T&AS.....P.....*
400/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
420/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
440/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
460/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
500/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
520/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
540/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
560/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
600/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
620/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
640/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
660/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
700/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
720/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
740/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*
760/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....P.....P...*

```

This block corresponds to locations 0-776 in the assembly listing.

```

BLOCK NUMBER 0002
000/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
020/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
040/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
060/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
100/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
120/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
140/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
160/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
200/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
220/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
240/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
260/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
300/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
320/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
340/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
360/ 000000 000000 000000 047514 045517 050125 043040 044501 042514 *.....LOOKUP FAIL*
400/ 000104 042522 042101 020127 040506 046111 042105 051000 *D_READW FAILD,R*
420/ 040505 053504 047440 000113 000000 000000 000000 000000 *FADW OK.....*
440/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
460/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
500/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
520/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
540/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
560/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
600/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
620/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
640/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
660/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
700/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
720/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
740/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
760/ 000000 000000 000000 000000 000000 000000 000000 041056 *.....B*

```

This block corresponds to locations 1000-1776 in the assembly listing.

```

BLOCK NUMBER 0003
000/ 000003 000006 000014 000023 000027 000040 000053 000057 * . . . . . + / *
020/ 177776 000000 000000 000000 000000 000000 000000 000000 * . . . . .
040/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
060/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
100/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
120/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
140/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
160/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
200/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
220/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
240/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
260/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
300/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
320/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
340/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
360/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
400/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
420/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
440/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
460/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
500/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
520/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
540/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
560/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
600/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
620/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
640/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
660/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
700/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
720/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
740/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .
760/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . .

```

This block shows the root relocation information. The first word of block 3 is a 3; since this is positive, positive relocation is indicated. Locations 6, 14, 30, 46, 56, 100, 126, and 136 must all be positively relocated at FRUN time. (On examination of the assembly listing, those locations marked with a ' need to be relocated.) The 177776 terminates the list.

Had negative relocation been indicated at relative location 6, block 3 would have shown 100003, 6, 14, 23, 27, 40, 53, 57, 177776.

3.4.4.2 REL Files with Overlays - When overlays are included in a program, the file is similar to that of a non-overlaid program. However, the overlay segments must also be relocated. Since overlays are not permanently memory resident but are read in from the file as needed, they require an additional operation. Each overlay segment is relocated (by FRUN) and then rewritten into the file. Then, when the overlay is called in, it will be properly relocated. This process takes



place each time an overlaid file is run with FRUN. The relocation information for overlay files contains both the list of addresses to be modified and the original contents of each location. This allows the file to be FRUN after the first usage.

NOTE

.ASECTS are illegal above 1000<sub>8</sub> and restricted in an overlaid foreground job. Refer to Chapter 6 of the RT-11 System Reference Manual.

A REL file with overlays appears as shown in Figure 3-14:

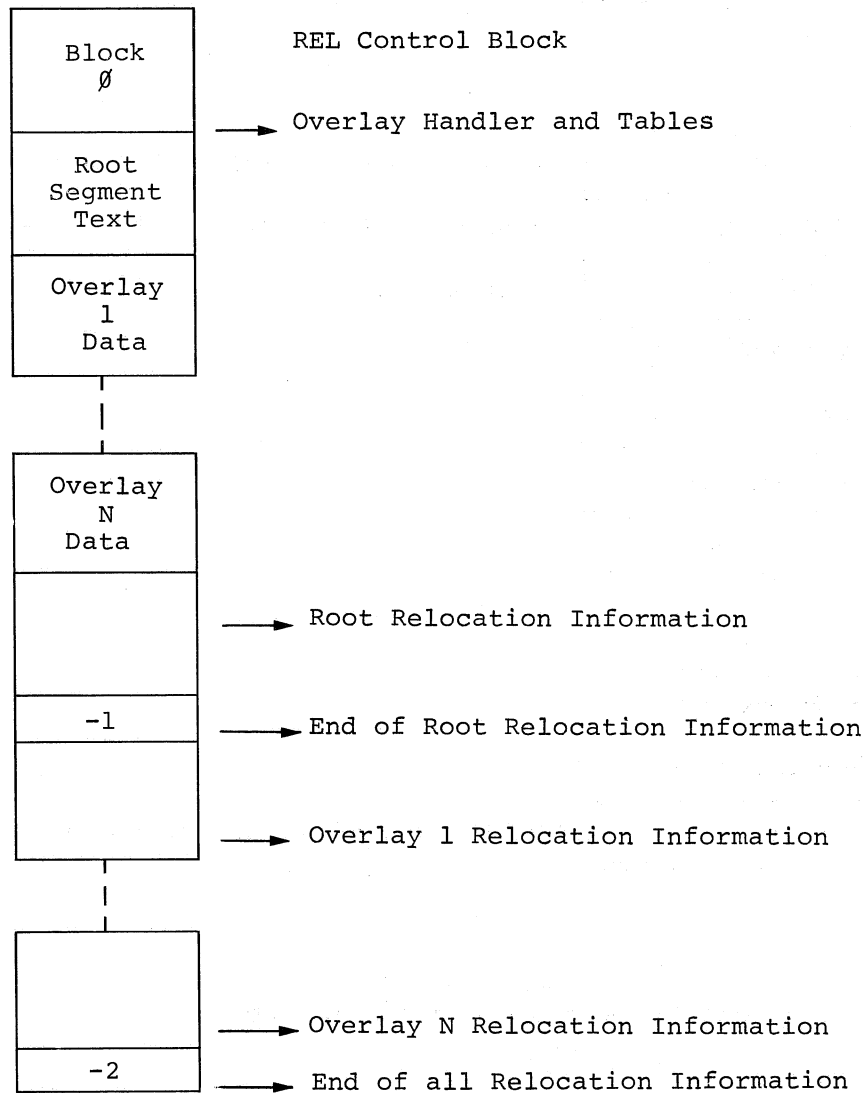


Figure 3-14  
REL File With Overlays

In this case, location 54 of block 0 of the REL file contains the size of the overlay region, in words. This is used to allocate space for the job when added to the size of the program base segment in location 52.

After the program base (root) code has been relocated, each existing overlay is read into the program overlay region in memory, relocated via the overlay relocation information, and then written back into the file.

The root relocation information section is terminated with a -1. This -1 is also an indication that an overlay segment relocation block follows. The overlay segment relocation block is shown in Figure 3-15:

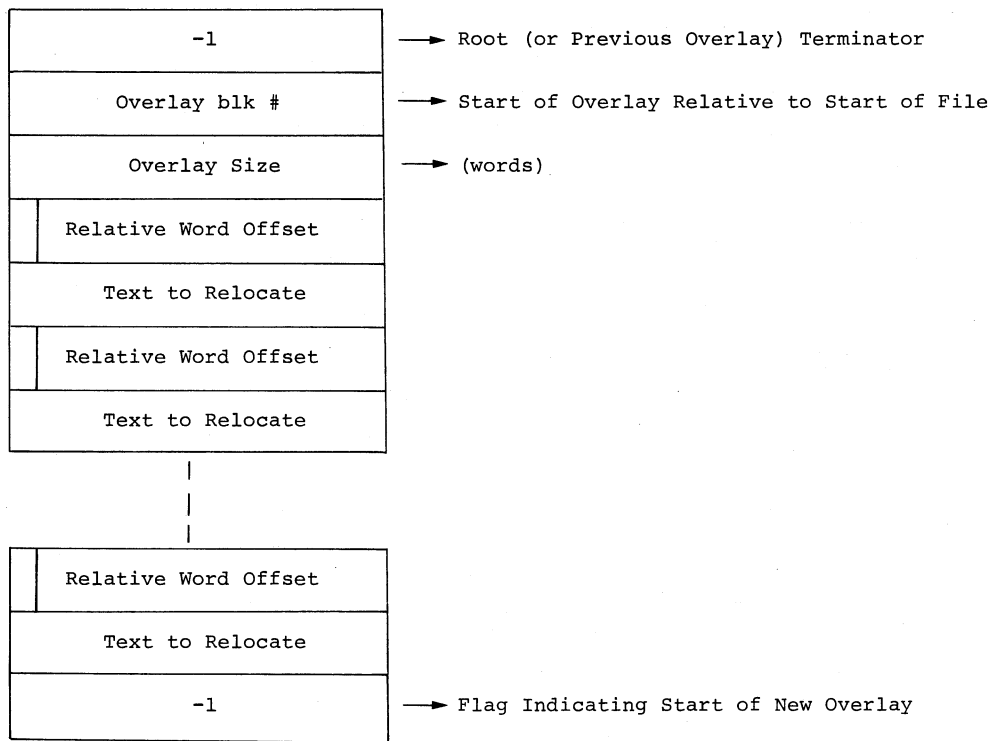


Figure 3-15  
Overlay Segment Relocation Block

The displacement is relative to the start of the program and is interpreted as in the nonoverlaid file (i.e., bit 15 indicates the type of relocation, and the displacement is the true displacement divided by two). Encountering -1 indicates that a new overlay region begins here. A -2 indicates the termination of all relocation information.



CHAPTER 4  
SYSTEM DEVICE

4.1 DETAILED STRUCTURE OF THE SYSTEM DEVICE

The RT-11 system device holds all the components of the system and is used by RT-11 to store device handlers and the monitor file. The layout of the system device is:

<u>Block #</u>	<u>Contents</u>
0	Bootstrap
1	Reserved for volume identification information
2	Bootstrap
3 to 5	Reserved for monitor or bootstrap expansion
6 to $(N*2)+5$	Directory segments; N is the number of directory segments
$(N*2)+6$ to end	File storage

All other system components, i.e., the monitor and device handlers, are files on the system device:

<u>File</u>	<u>Contains</u>
MONITR.SYS	The current RT-11 monitor; contains bootstrap, KMON, USR/CSI, RMON, KMON overlays, scratch blocks
SYSMAC.SML	System Macro Library
SYSMAC.S8K	8K System Macro Library
LP.SYS	Line printer handler

<u>File</u>	<u>Contains</u>
DT.SYS	DEctape handler
TT.SYS	Console handler (S/J only)
RK.SYS	RK disk handler
DS.SYS	RJS03/4 fixed-head disk handler
DX.SYS	RX01 flexible disk handler
DP.SYS	RP disk handler
PR.SYS	High-speed reader handler
PP.SYS	High-speed punch handler
CR.SYS	Card reader handler
RF.SYS	RF disk handler
CT.SYS	Cassette handler
MT.SYS	TM11 magtape handler
MM.SYS	TJUL6 magtape handler
BA.SYS	BATCH run-time handler

In general, files with the .SYS extension are parts of the monitor system. The bootstrap records the block numbers of the relevant areas in the monitor tables at bootstrap time. Thus, RT-11 is extremely flexible with respect to the interchange and construction of systems.

#### 4.2 CONTENTS OF MONITR.SYS

Following is the block layout of the RT-11 monitor file, MONITR.SYS. Block numbers are relative to the start of the file.

<u>F/B</u> <u>Monitor</u>	<u>Block # (decimal)</u>	<u>Contents</u>
	0-1	Copy of system bootstrap (blocks 0 and 2 of the system device)
	2-17	Swap blocks
	18-24	KMON (includes 2-block KMON overlay area)
	25-32	USR/CSI
	33-47	RMON
	48-57	KMON overlays

<u>S/J Monitor</u>	<u>Block # (decimal)</u>	<u>Contents</u>
	0-1	Copy of system bootstrap
	2-16	Swap blocks
	17-22	KMON (includes 1-block KMON overlay area)
	23-30	USR/CSI
	31-37	RMON
	38-44	KMON overlays

#### 4.3 KMON OVERLAYS

The KMON overlays are one block in size in the S/J Monitor and two blocks in size in the F/B Monitor. The contents of each overlay are described in this list:

<u>Overlay #</u>	<u>S/J</u>	<u>F/B</u>
0	DATE, TIME	DATE, TIME, SAVE, ASSIGN
1	SAVE, ASSIGN	LOAD, UNLOAD, SUSPEND, RESUME, CLOSE, FRUN (Part 1)
2	LOAD, UNLOAD, CLOSE	FRUN (Part 2)
3	GT ON/OFF	GT ON/OFF, SET
4	SET	

#### 4.4 DETAILED OPERATION OF THE BOOTSTRAP

Bootstrapping a system causes a fresh copy of that system to be installed in memory. In the RT-11 boot, certain system device resident tables are also updated. Following is a detailed description of the bootstrap.

<u>Action</u>	<u>Explanation</u>
1. User executes hardware bootstrap	On all system devices except diskette, this causes block 0 of the system device to be read into 0-777. Control then passes to location 0. On diskette, causes logical block 0 to be read into 0-777. Hardware bootstrap reads 64 words from track 1, sector 1. Control passes to location 0, where 64 words from each of sectors 3, 5, and 7 (track 1) are read.
2. Second part of bootstrap is read	The first part of the boot reads the second half into 1000-1777. On diskette, the first part of the boot reads logical block 2 (sectors 9, 11, 13, 15) into 1000-1777.
3. Determine how much memory is available	Boot sets a trap at location 4 and then starts addressing memory. When the trap is taken, illegal memory has been addressed.
4. Look for special devices	Boot sets a trap at location 10 and then tries to address the clock, FPU, and VT11 display processor. Their presence or absence is indicated in the CONFIG word in RMON. (If a PDP-11/03 processor is present, the bootstrap assumes that a clock is present.)
5. Check memory size	If memory is too small to read in the monitor, a message is printed and the boot halts.
6. Read in directory and find MONITR.SYS	The entire directory is searched. If MONITR.SYS is not found, a HALT occurs after the boot prints an error message.
7. Read the monitor into memory	The monitor file, MONITR.SYS, is read into the highest bank of memory.
8. Put pointers to monitor file blocks into RMON	RMON references the monitor swap blocks directly. Thus, the position of the swap blocks varies as the placement of MONITR.SYS varies. The real position of the blocks is updated for each boot operation.
9. Update position-dependent areas in RMON.	MONITR.SYS is initially linked at 8K. However, if more than 8K is available, RT-11 uses it. To do that, certain words must be updated to point to the actual areas of high memory where they will be. Boot contains a list of all words to be updated, located at RELLST in BSTRAP.MAC.



<u>Action</u>	<u>Explanation</u>
10. Update processor-dependent area in RMON	If processor is a PDP-11/03, any PS references in the monitor are changed to use the MFPS and MTPS instructions.
11. LOOKUP the device handlers in system and store their record numbers in \$DVREC	Boot looks at \$PNAME table to find the names of the devices in the system. The extension .SYS is appended. Thus, the PR handler is a file called PR.SYS. The location of the handler is then placed in \$DVREC. If the LOOKUP fails, the device gets a 0 in its \$DVREC entry. That implies that the device handler does not exist.
12. Print bootstrap header	Boot prints monitor identification message "RT-11" followed by monitor type ("FB" or "SJ") followed by version number.
13. Set up locations 0 and 2	Boot puts a "BIC R0,R0" in location zero and an .EXIT EMT in location 2.
14. Turn on KW11-L Clock	The bootstrap turns on the clock, if present in the configuration and processor is not a PDP-11/03.
15. Exit to Keyboard Monitor	

#### 4.5 FIXING THE SIZE OF A SYSTEM

RT-11 is designed to automatically operate from the top of the highest available 4K memory bank. However, it is possible to force the system to operate from a specified area that is not necessarily the highest. For instance, the following series of commands causes RT-11 to run in a 16K environment, even though the configuration actually has 28K of memory:

```
.R PATCH<CR> [Run RT-11 PATCH program.]
```

PATCH Version number

```
FILE NAME--
*MONITR.SYS/M<CR>
*BHALT/ 407 0<CR>
*E
.R PIP
*A=MONITR.SYS/U<CR>
*SY:/O
```

[Specifying MONITR.SYS/M indicates it is a monitor file. Change location "BHALT" from a 407 to a 0 (HALT). The correct address of BHALT can be found in Table 2 of RT-11 System Release Notes (V02C). E causes an exit to the monitor. Now run PIP to update the bootstrap and reboot the system.]

When the bootstrap is performed, the computer halts. The halt allows the user to enter the desired size in the switch register. With this patch installed, the V2 bootstrap uses the top five bits (bits 11-15) of the switch register to determine memory size. If the switch register contains the number 160000 or greater (e.g., if the register is unchanged after booting the system), a normal memory determination is performed. Otherwise, the top five bits are taken to be a number representing the number of 1K word blocks of memory. Each bit has the following value:

<u>Switch Register</u>	<u>Memory Size</u>
4000	1K
10000	2K
20000	4K
40000	8K
100000	16K

A combination of the bits will produce the range of system sizes from 8K through 28K, in 1K increments.

Examples:

1. To boot a system into 24K on a 28K configuration, use the combination:

$$140000 = 100000 (16K) + 40000 (8K)$$

2. To boot the S/J Monitor into 11K, use the combination:

$$54000 = 40000 (8K) + 10000 (2K) + 4000 (1K)$$

When the switch register is set properly, press the CONTInue switch and the bootstrap will be executed.

If the CONTInue switch is pressed immediately following the halt without changing the switch settings, a normal memory determination is done. To change the bootstrap back to its original (non-halting) form, execute the same commands as above, but change the 0 at BHALT back to a 407.

This procedure allows the user to 'protect' memory areas, since RT-11 never accesses memory outside the bounds within which it runs.

Another useful procedure, when desiring to always boot a system into a specific memory size or when the console switch register is not available, is to determine the bit combination corresponding to the choice of memory size, as explained above. Then enter the following commands, where xxxxx represents the bit pattern just determined:

```
.R PATCH<CR>
```

[Run RT-11 PATCH program.]

```
PATCH Version number
```

```
FILE NAME--
```

```
*MONITR.SYS/M<CR>
```

```
*BHALT/ 407 240<LF>
```

```
*BHALT+2/ 13702 12702<LF>
```

```
*BHALT+4/ 177570 xxxxxx<CR>
```

```
*E
```

```
.  
-
```

[NOP the branch at BHALT  
Change MOV @#SR,R2 to  
MOV #VAL,R2. Address of  
switch register is replaced  
with one of the bit combina-  
tions described previously.]

For the patch addresses for other system devices, and for the address of BHALT, consult Table 2 of RT-11 System Release Notes (V02C).



## CHAPTER 5

### I/O SYSTEM, QUEUES, AND HANDLERS

I/O transfers in RT-11 are handled by the monitor through routines known as device handlers. Device handlers are resident on the system mass storage device and can be called into memory at a location specified by the user (via a .FETCH handler request or KMON LOAD command). Only the device handlers distributed with the system in use (V2 or V02B) may be used; the system will malfunction otherwise.

This chapter describes how to write a new device handler and add it to the system. A summary of differences between Version 1 and Version 2 Device Handler requirements is included for the user who wishes to update old device handlers. Instructions and examples for making a device the system device and for writing a new bootstrap for the device are also included.

#### 5.1 QUEUED I/O IN RT-11

Once a device handler is in memory, any .READ/.WRITE requests for the corresponding devices are interpreted by the monitor and translated into a call to the I/O device handler. To facilitate overlapped I/O and computation, all I/O requests to RT-11 are done through an I/O queue. This section details the structure of the I/O queueing system.

##### 5.1.1 I/O Queue Elements

The RT-11 I/O queue is made up of a linked list of queue elements. A single element has the structure shown in Figure 5-1:

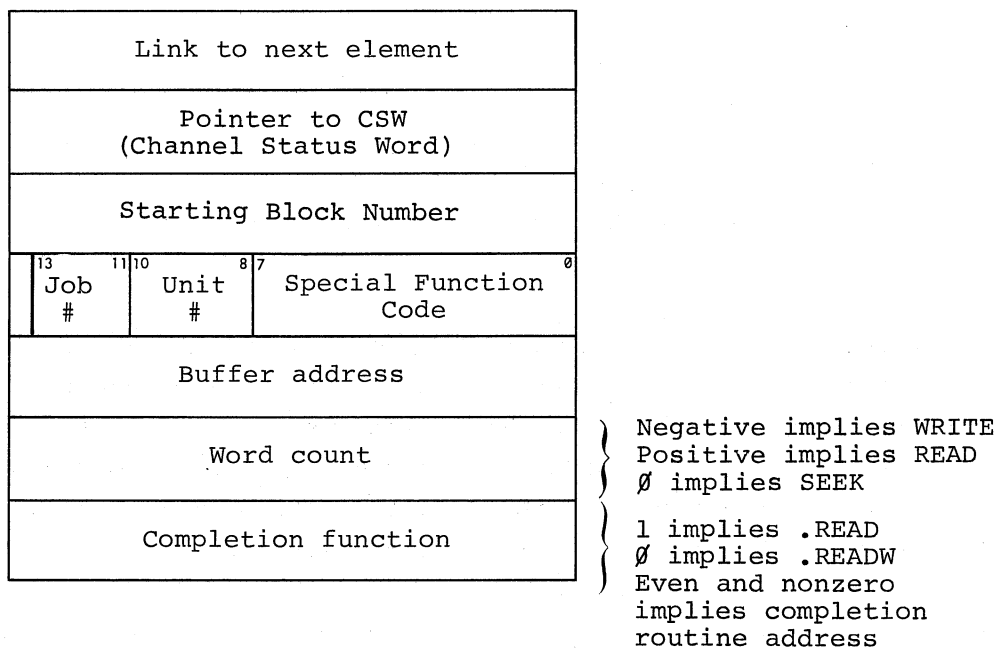
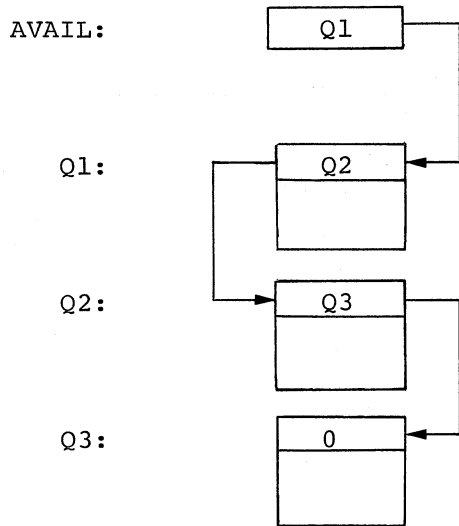


Figure 5-1  
I/O Queue Element

RT-11 maintains one queue element in the Resident Monitor. (In F/B, one element per job is maintained in the job's impure area.) This is sufficient for any program that uses wait-mode I/O (.READW/.WRITW). However, for maximum throughput, the .QSET programmed request should be used to create additional queue elements.

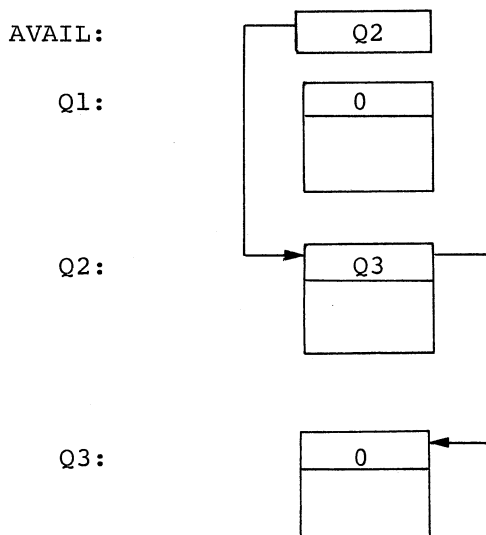
If an I/O operation is requested and a queue element is not available, RT-11 must wait until an element is free to queue the request. This obviously slows up program execution. If asynchronous I/O is desired, extra queue elements should be allocated. It is always sufficient to allocate N new queue elements, where N is the total number of pending requests that can be outstanding at one time in a particular program. This produces a total of N+1 available elements, since the Resident Monitor element is added to the list of available elements.

Diagrammatically, the I/O queue appears as follows:



AVAIL is the list header. It always contains a pointer to an available element. If AVAIL is 0, no elements are currently available.

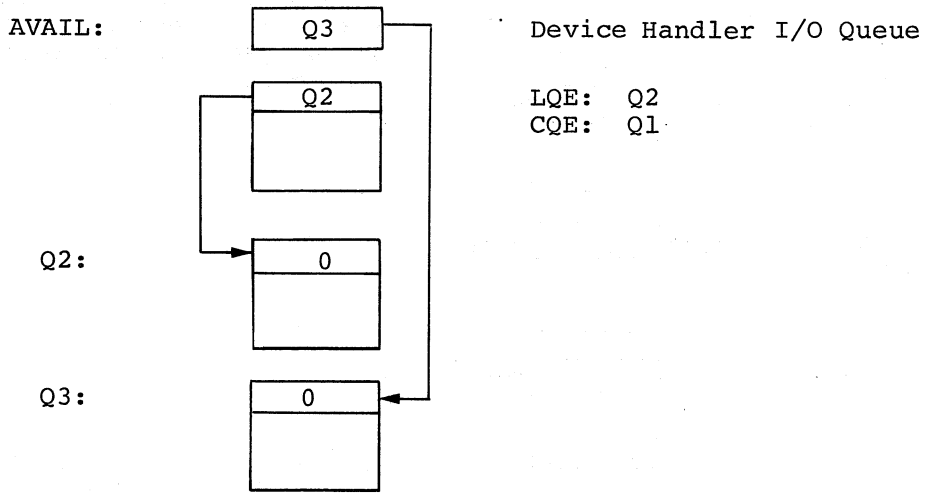
When an I/O request is initiated, an element is allocated (removed from the list of available elements) and is linked into the appropriate device handler's I/O queue. The handler's queue header consists of two pointers: the current queue element (CQE) pointer, pointing to the element at the top of the list, and the last queue element (LQE) pointer, pointing to the last element entered in the queue. The LQE pointer is used by the S/J monitor for fast insertion of new elements into the queue.



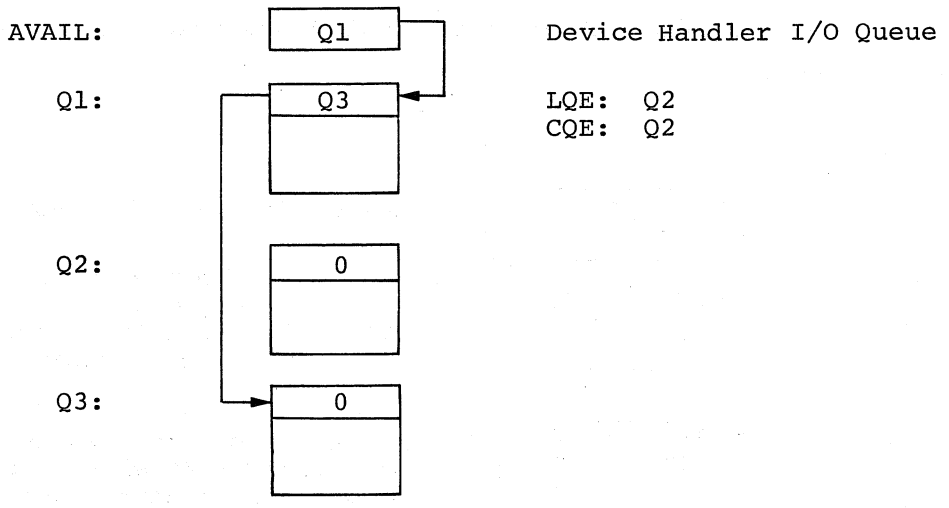
Device Handler I/O Queue

LQE: Q1 (Pointer to last queue element)  
 CQE: Q1 (Pointer to current queue element)

In this case, the device is associated with element Q1. If another request comes in for that same device before the first completes, a waiting queue is built up for that device.

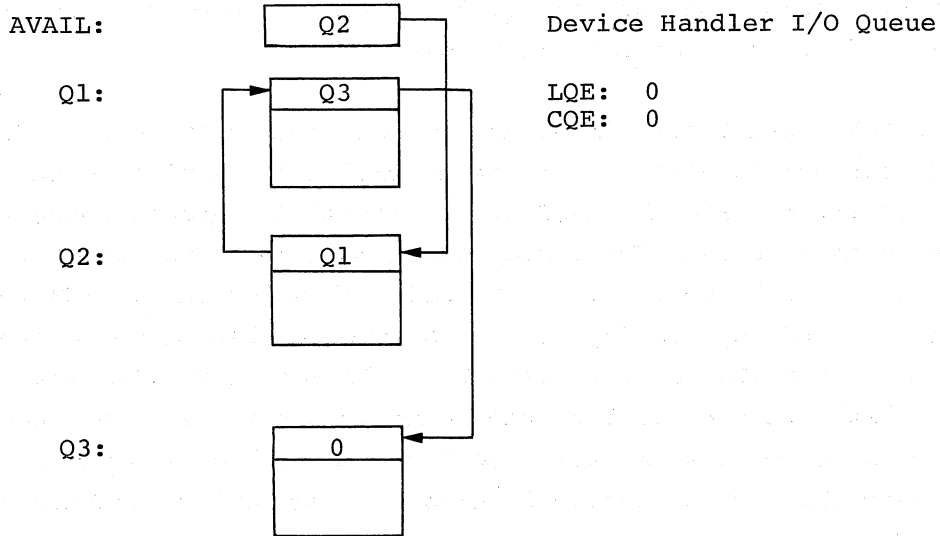


When the I/O transfer in progress completes, Q1 is returned to the list of available elements, and the transfer indicated by Q2 will be initiated:





When Q2 is completed, it too is returned to the list of available elements.



Note that the order of the queue element linkages may be altered.

A distinction between S/J and F/B operation is that F/B maintains two separate queue structures, one for each active job. The queue headers (AVAIL) are words in the user's impure area. The centralized queue manager dispatches transfers in accordance with job priority. Thus, if two requests are queued waiting for a particular device, the foreground request is honored first. At no time, however, will an I/O request already in progress be aborted in favor of a higher priority request; the operation in progress will complete before the next transfer is initiated.

Another difference between S/J and F/B operation is that the F/B scheduler will suspend a job pending the availability of a free queue element and will try to run another job.

### 5.1.2 Completion Queue Elements

The F/B Monitor maintains, in addition to the queue of I/O transfer requests, a queue of I/O completion requests. When an I/O transfer completes and a completion routine has been specified in the request (i.e., the seventh word of the I/O queue element is even and non-zero), the queue completion logic in the F/B Monitor transfers the request node (element) to the completion queue, placing the channel

status word and channel offset in the node. This has the effect of serializing completion routines, rather than nesting them. Completion routines are called by the completion queue manager on a *first-in/first-out* basis, and the completion routines are entered at priority level 0 rather than at interrupt level.

The .SYNCH request also makes use of the completion queue. When the .SYNCH request is entered, the seven-word area supplied with the request is linked into the head of the completion queue, where it appears to be a request for a completion routine. The .SYNCH request then does an interrupt exit. The code following the .SYNCH request is next called at priority level 0 by the completion queue manager. To prevent the .SYNCH block from being linked into AVAIL (the queue of available elements), the word count is set to -1. The completion queue manager checks the word count before linking a queue element back into the list of available elements, and skips elements with the -1 word count.

Figures 5-2 and 5-3 show the format of the completion queue and .SYNCH elements.

OFFSET	
0	QUEUE LINK
2	
4	
6	
10	CHANNEL STATUS WORD
12	CHANNEL OFFSET
14	COMPLETION ROUTINE ADDRESS

Figure 5-2  
Completion Queue Element

OFFSET	
0	QUEUE LINK
2	JOB NUMBER
4	
6	
10	SYNCH I.D.
12	-1
14	SYNCH RETURN ADDRESS

Figure 5-3  
.SYNCH Element

### 5.1.3 Timer Queue Elements

Another queue maintained by the F/B Monitor is the timer queue. This queue is used to implement the .MRKT request, which schedules a completion routine to be entered after a specified period of time. The first two words of the element are the high- and low-order time and the seventh word is the completion routine address. An optional sequence number can be added to the request to distinguish this timer request from others issued by the same job.

The F/B Monitor uses the timer queue internally to implement the .TWAIT request. The .TWAIT request causes the issuing job to be suspended and a timer request is placed in the queue with the .RSUM logic as the completion routine. Refer to Figure 5-4 for the format of the timer queue element.

OFFSET		
0	HIGH-ORDER TIME	C.HOT
2	LOW-ORDER TIME	C.LOT
4	LINK TO NEXT ELEMENT	C.LINK
6	JOB # OF OWNER	C.JNUM
10	OWNER'S SEQUENCE #	C.SEQ
12	0	
14	COMPLETION ADDRESS	C.COMP

Figure 5-4  
Timer Queue Element

## 5.2 DEVICE HANDLERS

This section contains the information necessary to write an RT-11 device handler. It is illustrated with an example, a driver for the RS64 fixed-head disk (with RC11 controller). A source listing is included in Appendix A, Section A.1; portions of this listing are referenced throughout the remainder of this section and in future sections.

The user should refer to the PDP-11 Peripherals Handbook for details regarding the operation of any particular peripheral.

### NOTE

All RT-11 handlers must be written in position independent code (PIC). Consult the PDP-11 Processor Handbook for information on writing PIC.

#### 5.2.1 Device Handler Format

The first five words of any device handler are header words. The format is:

<u>Word #</u>	<u>Contents</u>
1	Address of first word of device's interrupt vector.
2	Offset from current PC to interrupt handler.
3	Processor status word to be used when interrupt occurs. Must be 340 (priority 7).
4,5	Zero. These are the queue pointers.

See area C in the example handler (Section A.1).

A word must be provided at the end of the handler. When the handler is .FETCHed, the monitor places a pointer to the monitor common interrupt entry code in the last word of the handler. This requires that the handler size in the monitor's \$HSIZE table be exact or the handler will malfunction. See area M in the example in Section A.1.

The word preceding the interrupt handler entry point must be an unconditional branch to the handler's abort code. The abort code is used by the F/B Monitor to stop I/O for the device. The abort entry point is shown at area G in the example and the abort code is at area K. (See the RT-11 System Reference Manual, Section H.2, for further information.)

### 5.2.2 Entry Conditions

The device handler is entered directly from the monitor I/O queue manager, at which time it initiates the data transfer. The fifth word of the header contains a pointer into the queue element to be processed. This word (called CQE, for Current Queue Element) points to the third word of the queue element, which is the block number to be read or written. Referring to the example, location RCCQE contains the address of the third word of the queue element to be processed. It is generally advisable to put the pointer into a register, as that greatly facilitates picking up arguments to initiate the transfer. In the example, the entry point is at the location marked by E. Notice that registers need not be saved.

### 5.2.3 Data Transfer

Most handlers use the interrupt mechanism when transferring data. The handler initiates the transfer and then returns immediately to the monitor with an RTS PC, shown at area F. When the transfer is completed, the device interrupts. When the interrupt routine determines that I/O is complete or that an error has occurred, it jumps to the monitor completion routine in the manner shown at area J in the listing.

If the interrupt mechanism is not used, the data transfer must be completed before returning to the monitor. The handler must loop on a device flag with the interrupt disabled. When I/O is complete, the driver returns to the monitor with a jump to the monitor completion code, similar to that shown at area J in the example.

### 5.2.4 Interrupt Handler

Once the transfer has been initiated and control has passed back to the monitor, data interrupts will occur.

Information in the header of the handler causes the interrupt to be vectored to the interrupt handling code within the handler. The code at the interrupt location should keep the transfer going, determine when the transfer is complete, and detect errors.

When the transfer is done, control must be passed to the monitor's I/O queue manager, which performs a cleanup operation on the I/O queue.

Restrictions that apply to the interrupt code are:

1. The common interrupt entry into the monitor must be taken. Interrupt routines linked into a program use the .INTEN request described in Chapter 9 of the RT-11 System Reference Manual. Handlers made part of the system have a more efficient method of entry. The last word of the handler is set to point to the monitor common interrupt entry code when the handler is fetched. Upon reception of an interrupt, the handler must execute this code by performing a JSR R5, @\$INPTR, where \$INPTR is the tag commonly used by RT-11 handlers for the pointer word. See areas I and N in the example. The JSR instruction must be followed by the complement of the priority at which the handler will operate. See area I for an easy method to make the assembler compute the complement. On return from the monitor's interrupt entry code, R4 and R5 have been saved and may be used by the handler. Other registers must be saved and restored if they are to be used.
2. A check must be made to determine if the transfer is complete. However, with nonfile-structured devices, such as paper tape, line printer, etc., an interrupt occurs whenever a character has been processed. For these devices, the byte count, which is in the queue element, is used as a character count.

Nonfile-structured input devices should be able to detect an end of file condition, and pass that on to the monitor.

#### NOTE

The queue element contains a word count, not a byte count. The initial entry to the handler should change the word count to a byte count if the device interrupts at each character. The transfer is complete when the byte count decrements to 0.

Before the conversion to bytes is made, the sign of the word count must be determined since it specifies whether this transfer is a Read or Write. A negative word count implies a Write and should be complemented before converting to bytes.

3. Check for occurrence of an error. If a hardware error occurred, the hard error bit in the channel status word (CSW) should be set, the transfer should be aborted, and the monitor completion code executed. The address of the channel status word is in word 2 of the queue element. The error bit is bit 0 of the CSW. Generally, it is advisable to retry a certain number of times if an error occurs. RT-11 currently retries up to eight times before deciding an error has occurred. (Note that this is true for file-structured devices only.) It is

desirable, in case an error occurs, to do a drive or control reset, where appropriate, to clear the error condition before a retry is initiated. See the area between I and H in the example.

4. If the transfer is not complete and no error has occurred, registers used should be restored, and an RTS PC executed.

To pass an EOF (End of File) to the monitor, the 2000 bit in the CSW should be set. Refer to the sample handler in Appendix A for an example of setting the EOF bit. When EOF is detected on non-file structured devices, the remainder of the input buffer must be zeroed.

5. When the transfer is complete, whether an error occurred or not, the monitor I/O completion code must be entered to terminate activity and/or enter a completion routine. When return is made to the monitor, R4 must point to the fifth word of the handler (RCCQE in the example). See area J in the example for the method of returning to the monitor completion routine.

Handlers should check for special error conditions that can be detected on the initial entry to the handler. For example, trying to write on a read-only device should produce a hard error. It must be emphasized that the user handlers should interface to the system in substantially the same way as the handler in Section A.1. This handler is included as a guide and an example.

### 5.3 ADDING A HANDLER TO THE SYSTEM

When the handler has been written and debugged, it may be installed in the system by following the procedures in this section. The process consists of inserting information about the handler into the monitor tables listed below.

<u>Table to be Changed</u>	<u>Contents</u>
\$HSIZE	Size of handler (in bytes).
\$DVSIZ	Size of device in 256-word blocks. If nonfile device, entry = 0.
\$PNAME	Permanent name of the device (should be two alphanumeric characters entered in .RAD50 notation, left-justified).
\$STAT	Device status table. Refer to Section 2.5.2.2 for the format of \$STAT table.
LOWMAP	Low memory protection map; refer to Section 2.5.4.

There is no restriction on handler names; any 2-letter combination not currently in use may be chosen for the new handler and the name may be inserted in any unused slot in the \$PNAME table, or in a slot occupied by a nonexistent device (i.e., a device not installed on the user's system). Note that the name must be entered in .RAD50. Since PATCH does not have a .RAD50 interpretation switch, the name must be entered to PATCH in its numerical form. Appendix C of the RT-11 System Reference Manual contains a .RAD50 conversion table; ODT can also be used to perform .RAD50 conversions.

As an example, assume again the handler for the RC11/RS64 disk (the sample handler in Section A.1) is to be inserted in the system. First, the values of the table entries for this device are determined (the addresses used in the example are for illustrative purposes only; consult Table 2 of RT-11 System Release Notes (V02C) for the correct table addresses for the version in use):

\$HSIZE:	316	After assembly, the handler was found to take up 316 bytes. See area 0 in the example listing.
\$DVSIZ:	2000	The disk has 1024 (decimal) 256-word blocks for storage.
\$PNAME:	.RAD50 /RC/ or 70370	The name assigned is RC. The .RAD50 value of RC is 70370.
\$STAT:	100023	The device is file-structured, is a read/write device, and uses the standard RT-11 file structure. The identifier (selected by the user) is 23. Refer to Section 2.5.2.2 for the format of the \$STAT table.
LOWMAP:	14	Protect RC vector 210,212 at byte 336 of LOWMAP (refer to Section 2.5.4.).

Once these values have been decided, the steps for inserting the device handler are:

1. Assemble the handler, using either MACRO or ASEMBL.
2. Link the handler at 1000. The name of the handler should be whatever the \$PNAME entry is, with the .SYS extension appended:

```
.R LINK
  RC.SYS=RC   where RC.OBJ is the handler object
  UNDEF GLBLS module. The default link address is
              1000.
```



NOTE

If the handler being linked is one that could also be a system device handler, the user can expect one undefined global, \$INTEN.

3. Run PATCH to modify the tables and protect the interrupt vectors.

For this example, assume that the table addresses are found to be:

<u>Table</u>	<u>S/J Address</u>	<u>F/B Address</u>
\$HSIZE	13624	14556
\$DVSIZ	13660	14612
\$PNAME	16470	17630
\$STAT	16524	17664

NOTE

The addresses above are for illustration only. Consult Table 2 of RT-11 System Release Notes (V02C) for current table addresses and for the address of the monitor base location, BASE.

The tables have room for fourteen (decimal) device entries; all are already assigned by the monitor. Assuming that a given configuration never has all supported devices, however, at least one slot should be available to be overlaid. For example, assume the twelfth slot is occupied by a device not installed on the system, and therefore available for change. The octal offset is 26, which, added to the table addresses above, gives the address of the empty slot:

S/J Monitor:

.R PATCH<CR>

PATCH Version number

FILE NAME--

```
*MONITR.SYS/M<CR>          [/M is necessary;
*BASE;ØR<CR>              Monitor base;
*Ø,13652/ 4ØØØ 316<CR>    $HSIZE table;
*Ø,137Ø6/  Ø  2ØØØ<CR>    $DVSIZ table;
*Ø,16516/ 625Ø 7037Ø<CR>  $PNAME table;
*Ø,16552/  4  1ØØØ23<CR>  $STAT table;
*Ø,16336\  77  <CR>      Check that vectors in
*E                          permanent map are protected;
:                             Exit to monitor]
```

## F/B Monitor

.R PATCH

PATCH Version number

FILE NAME--

\*MONITR.SYS/M<CR>

[/M is necessary;

\*BASE;ØR<CR>

Monitor base;

\*Ø,14556/ 4ØØØ 316<CR>

\$HSIZE table;

\*Ø,14612/ Ø 2ØØØ<CR>

\$DVSIZ table;

\*Ø,17630/ 625Ø 7Ø37Ø<CR>

\$PNAME table;

\*Ø,17664/ 4 1ØØØ23<CR>

\$STAT table;

\*Ø,17336 77 <CR>

Check that vectors in

\*E

permanent map are protected;

.

Exit to monitor]

At this point, the system should be re-bootstrapped to make the modified monitor resident. The device RC will then be available for use.

### 5.4 WRITING A SYSTEM DEVICE HANDLER

This section describes the procedures for writing a new system device handler. A system device is the device on which the monitor and handlers are resident. RT-11 currently supports the RK, RF, DP, DS, and DX disks, and DEctape as system devices. The procedures for writing the handler and creating a new monitor are explained, illustrated by the example in Section A.1, the RC11/RS64 handler.

The basic requirements for a system device are random access and read/write capability. These requirements are met by the RC11 disk, which is a multiple platter, fixed-head disk. When writing the driver, the procedures in Section 5.2 should be followed. Because the system handler is linked with the monitor, the additional tagging and global conventions described here must also be followed.

#### 5.4.1 The Device Handler

The following conditions must be observed when writing a system handler. Refer to the example listing in Section A.1.

1. The handler entry point must be tagged xxSYS, where xx is the 2-letter device name. For the RC disk, this is RCSYS. See area D in the listing.  
Important: Note that the tag is placed after the third word of the header block.
2. The entry points of all current system devices must be referenced in a global statement. These

currently include RKSYS, RFSYS, DSSYS, DXSYS, DPSYS, DTSYS and RCSYS. See Area A.

3. The entry point tags of all other system devices must be equated to zero. See area B in the listing.
4. A .CSECT SYSHND must be included at the top of the handler code. It is located above area C in the example.
5. The last word of the handler is used for the common interrupt entry address. This should have the tag \$INPTR and should be set to the value \$INTEN. See areas M and N in the example listing. These tags should be global. See area A.
6. The interrupt entry point should have the tag xxINT, or RCINT for this example, and this must be a global. See areas A and H.
7. The handler size must be global, with the symbolic name xxSIZE, or RCSIZE. See area A. This step is not necessary if the monitor sources are available and are being reassembled, since the global will be generated by the HSIZE macro. See Step 3 in Section 5.4.3.

#### 5.4.2 The Bootstrap

This section describes the procedure for modifying the system bootstrap to operate with a new system device. Either the bootstrap source must be acquired, or the listing in Section A.2 may be used. Again, the RC11/RS64 disk is used for an example. The references in this section, however, are to the bootstrap listing found in Section A.2 of Appendix A.

The following changes must be made to the bootstrap to support a new system device:

1. Add a new conditional, \$xxSYS, to the list at point AA. Here xx is the 2-letter device name, and in this case the conditional is \$RCSYS.
2. Add a simple device driver for the device inside a \$xxSYS conditional. This is shown at area CC. Because the RC11 is similar to the other disks, it is possible to share code with the other device drivers, reducing the implementation effort. To do this, the \$RCSYS conditional is added at area BB and the device specific code is at area FF. This code merges with the common code at area GG.

3. The device driver has these characteristics:
  - a. The SYSDEV macro must be invoked for the device. The macro arguments are the 2-letter device name and the interrupt vector address. For this example, the arguments are "RC" and "210", shown at area DD on the listing.
  - b. The device driver entry point must have the tag READ. See area EE.
  - c. When the driver is entered:
    - R0 = Physical Block Number
    - R1 = Word Count
    - R2 = Buffer Address
    - R3,R4,R5 = are available for use by the driver routine
  - d. The driver must branch to BIOERR if a fatal I/O error occurs.

#### 5.4.3 Building the New System

This section describes the procedure for building a new monitor using the system device handler and bootstrap just developed. Again, the example used is the RC11/RS64 disk, and the appropriate listings are those in Sections A.1 and A.2.

The procedure is:

1. Assemble the handler, producing an object module with the name xx.OBJ, where xx is the 2-letter device name. In this example, the name is RC.

```
.R MACRO
*RC.OBJ=RC.MAC
```

2. Assemble the bootstrap, defining the conditional \$xxSYS (where xx is again the device name; e.g., \$RCSYS). Define the conditional BF if an F/B bootstrap is desired. Let BF be undefined for an S/J bootstrap. For the S/J bootstrap:

```
.R MACRO <CR>
*RCBTSJ=TT:,DK:BSTRAP <CR>
^$RCSYS=1 <CR>
^Z^$RCSYS=1 <CR>
^ZERRORS DETECTED: 0
FREE CORE: 15608. WORDS
```

For the F/B bootstrap:

```
.R MACRO<CR>
*RCBTFB=TT:,DK:BSTRAP<CR>
^$RCSYS=1<CR>
BF=1<CR>
^Z^$RCSYS=1<CR>
BF=1<CR>
^ZERRORS DETECTED: Ø
FREE CORE: 1558Ø. WORDS
```

3. If the monitor sources are available, the DEVICE macro described in Section 2.5.2.9 can be invoked for the new device by editing the macro call into RMONFB.MAC and RMONSJ.MAC and reassembling the monitor. For the RC device, the macro would be:

```
DEVICE RC 2000 100020 RCSYS
```

The HSIZE macro, described in the same section, must also be invoked. For the RC device, the macro would be:

```
HSIZE RC,316,SYS
```

Monitor assembly instructions are in Chapter 5 of the RT-11 System Generation Manual. If this approach is used, the table patching procedure in step 5 is not necessary.

4. Link the monitor with the new bootstrap and device handler.

For S/J:

```
.R LINK
*RCMNSJ.SYS,MAP=RCBTSJ,RT11SJ,RC
```

For F/B:

```
.R LINK
*RCMNFBSYS,MAP=RCBTFB,RT11FB,RC
```

5. If step 3 was not done and step 4 used the current monitor object modules, then the monitor tables must be patched to enter the device information. The monitor device tables are located using the procedure in Section 5.3. An additional table, the \$ENTRY entry point table, must also be patched. For this example, assume the table addresses are:

<u>Table</u>	<u>S/J Address</u>	<u>F/B Address</u>
\$HSIZE	13674	14602
\$DVSIZ	13730	14636
\$PNAME	16516	17640
\$STAT	16552	17674
\$ENTRY	16612	17612

NOTE

These table addresses are for illustration only. Consult Table 2 of RT-11 System Release Notes (V02C) for the table addresses of the current monitor release and for the address of BASE.

A link map was made during the linking sequence in Step 4. Locate the value of the system handler entry point, xxSYS. For this example, the tag is RCSYS and its value is found to be 56266 for F/B. This value is put in the \$ENTRY table. The other values were determined in Section 5.3:

```
$HSIZE = 316
$DVSIZ = 2000
$PNAME = 70370
$STAT  = 100023
$ENTRY = 56266 (F/B) 45056 (S/J)
```

The patch procedure for the S/J monitor, using the twelfth slot, would then be:

```
.R PATCH<CR>
_
PATCH Version number

FILE NAME--
*RCMNSJ.SYS/M<CR>           [The /M is necessary;
*BASE;0R<CR>                Monitor base;
*0,13674/____4000____316<CR> $HSIZE table;
*0,13730/____0____2000<CR>  $DVSIZ table;
*0,16516/____6250____70370<CR> $PNAME table;
*0,16552/____4____100023<CR> $STAT table;
*0,16612/____0____45056<CR> $ENTRY table;
*E                           Exit to monitor]
_
:
```

For the F/B monitor:

```
.R PATCH<CR>
_
PATCH Version number

FILE NAME--
*RCMNFBSYS/M<CR>
*BASE;0R<CR>
*0,14602/____4000____316<CR>  [$HSIZE table;
*0,14636/____0____2000<CR>  $DVSIZ table;
*0,17640/____6250____70370<CR> $PNAME table;
*0,17674/____4____100023<CR> $STAT table;
*0,17564/____0____56266<CR> $ENTRY table;
*E                           Exit to monitor]
_
:
```

The new monitor is now complete and may be used by transferring it to an RC disk and renaming it to MONITR.SYS.

## 5.5 DEVICES WITH SPECIAL DIRECTORIES

The RT-11 monitor can interface to devices having nonstandard (that is, non RT-11) directories. This section discusses the interface to this type of device.

### 5.5.1 Special Devices

Special devices are file-structured devices that do not use an RT-11 directory format. Examples are magtape and cassette as supported under RT-11. They are identified by setting bit 12 in the device status word. The USR processes directory operations for RT-11 directory-structured devices; for special devices, the handler must process directory operations (LOOKUP, ENTER, CLOSE, DELETE), as well as data transfers.

5.5.1.1 Interfacing to Special Device Handlers - There are three types of processes that a special device handler must perform:

1. Directory operations (.LOOKUP, .ENTER, etc.)
2. Data transfer operations (.READ, .WRITE)
3. Special operations (rewind, backspace, etc.)

The particular process required is passed to the handler in the form of a function code, located in the even byte of the fourth word of the I/O queue element (see Section 5.1.1). The function code may be positive or negative. Positive codes are used for processes of types 1 and 2 above; negative codes indicate device-dependent special functions.

The positive function codes are standard for all devices and include:

<u>Code</u>	<u>Function</u>
Ø	Read/Write
1	Close
2	Delete
3	Lookup
4	Enter

These functions correspond to the programmed requests .READ/.WRITE, .CLOSE, .DELETE, .LOOKUP, and .ENTER, described in Chapter 9 of the RT-11 System Reference Manual. The .RENAME request is not supported for special devices.

A queue element for a special handler will look identical to an element for a standard RT-11 handler when the function is a .READ/.WRITE (negative word count implies a .WRITE). For the remaining positive functions, word 5 of the queue element (the buffer address word discussed in Section 5.1.1) will contain a pointer to the file descriptor block, containing the device name, file name, and file extension in .RAD5Ø format.

Negative function codes are used for device-dependent special functions. Examples of these are backspace and rewind for magtape. Because these functions are characteristic of each device type, no standard definition of negative codes is made; they are defined uniquely for each device.

Software errors (for example, file not found or directory full) occurring in special device handlers during directory operations are returned to the monitor through the procedure described next. A unique error code is chosen for each type of error. This error code is directly returned by placing it in SPUSR (special device USR error), located at a fixed offset (272) into RMON. (Section 2.5.1 discusses monitor fixed offsets.) Hardware errors are returned in the usual manner by setting bit Ø in the channel status word pointed to by the second word of the queue element.

5.5.1.2 Programmed Requests to Special Devices - Programmed requests for directory operations and data transfers to special devices are handled by the standard programmed requests. When a .LOOKUP is done, for example, the monitor checks the device status word for the special device bit. If the device has a special directory structure, the proper function code is inserted into the queue element and the element is directly queued to the handler, by-passing any processing by the RT-11 USR. Device independence is maintained, since .READ, .WRITE, .LOOKUP, .ENTER, .CLOSE, and .DELETE operations are transparent to the user.

Requests for device-dependent special functions having negative function codes, must be issued by using the .SPFUN special function programmed request, described in Chapter 9 of the RT-11 System Reference Manual. Devices which need to use the .SPFUN requests must have a bit set in the device status table (see Section 2.5.2.2).



## 5.6 ADDING A SET OPTION

The Keyboard Monitor SET command permits certain device handler parameters to be changed from the keyboard. For example, the width of the line printer on a system can be SET with a command such as:

```
SET LP WIDTH=80
```

This is an example of a SET command that requires a numeric argument. Another type of SET command is used to indicate the presence or absence of a particular function. An example of this is a SET command to specify whether an initial form feed should be generated by the LP handler:

```
SET LP FORM           (generate initial form feed)
SET LP NOFORM        (suppress initial form feed)
```

In this case, the FORM option may be negated by appending the NO prefix.

The SET command is entirely driven by tables contained in the device handler itself. Making additions to the list of SET options for a device is easy, requiring changes only to the handler, and not to the monitor. This section describes the method of creating or extending the list of SET options for a handler. The example handler used is the LP/LS11 line printer handler, listed in Appendix A in Section A.3. The SET command is described in Chapter 2 of the RT-11 System Reference Manual.

Device handlers have a file name in the form xx.SYS, where xx is the 2-letter device name; e.g., LP.SYS. Handler files are linked in save image format at a base address of 1000, in which a portion of block 0 of the file is used for system parameters. The rest of the block is unused, and block 0 is never FETCHed into memory. The SET command uses the area in block 0 of a handler from 400 to 776 (octal) as the SET command parameter table. The first argument of a SET command must always be the device name; e.g., LP in the previous example command lines. SET looks for a file named xx.SYS (in this case LP.SYS) and reads the first two blocks into the USR buffer area. The first block contains the SET parameter table, and the second block contains handler code to be modified. When the modification is made, the two blocks are written out to the handler file, effectively changing the handler.

The SET parameter table consists of a sequence of 4-word entries. The table is terminated with a zero word; if there are no options available, location 400 must be zero. Each table entry has the form:

.WORD	value	
.RAD50	/option/	[2 words of RAD50]
.BYTE	<routine-400>/2	
.BYTE	mode	

where:

value	is a parameter passed to the routine in register 3.
option	is the name of the SET option; e.g., WIDTH or FORM.
routine	is the name of a routine following the SET table that does the actual handler modification.
mode	indicates the type of SET parameter: a. Numeric argument - byte value of 100 b. NO prefix valid - byte value of 200

The SET command scans the table until it finds an option name matching the input argument (stripped of any NO prefix). For the first example command string, the WIDTH entry would be found (area 2 in the listing in Section A.3). The information in this table entry tells the SET processor that O.WIDTH is the routine to call, that the prefix NO is illegal and that a numeric argument is required. Routine O.WIDTH is located at area 4 on the listing. It uses the numeric argument passed to it to modify the column count constant in the handler. The value passed to it in R3 from the table is the minimum width and is used for error checking.

The following conventions should be observed when adding SET options to a handler:

1. The SET parameter tables must be located in block 0 of the handler file and should start at location 400. This is done by using an .ASECT 400 (area 1 on the listing).
2. Each table entry is four words long, as described previously. The option name may be up to six .RAD50 characters long, and must be left-justified and filled with spaces if necessary. The table terminates with a zero (area 3 on the listing).

3. The routine that does the modification must follow the SET table in block 0 (area 4 on the listing). It is called as a subroutine and terminates with an RTS PC instruction. If the NO prefix was present and valid, the routine is entered at entry point +4. An error is returned by setting the C bit before exit. If a numeric argument is required, it is converted from decimal to octal and passed in R0. The first word of the option table entry is passed in R3.
4. The code in the handler that is modified must be in block 1 of the handler file, i.e., in the first 256 words of the handler. See areas 6 and 7 on the listing for code modified by the WIDTH option.
5. Since an .ASECT 400 was used to start the SET table, the handler must start with an .ASECT 1000. See area 5 on the listing.
6. The SET option should not be used with system device handlers, since the .ASECT will destroy the bootstrap and cause the system to malfunction.

#### 5.7 CONVERTING USER-WRITTEN HANDLERS

User-written device handlers must, in all cases, conform to the standard practices for Version 2 (2B and 2C). General programming information is discussed in Appendix H of the RT-11 System Reference Manual. Points to consider when converting user-written device handlers (written under Version 1 of the RT-11 system) follow; the details of these procedures have already been discussed.

1. The last word of a device handler is used by the monitor, thus the user must be sure to include one extra word at the end of his program when indicating the handler size.
2. The third header word of the handler should be 340, indicating that the interrupt should be taken at level 7.
3. It is not necessary to save/restore registers when the handler is first entered, although to do so is not harmful.
4. When an interrupt occurs, the handler must execute an .INTEN request or its equivalent. On return from .INTEN, R4 and R5 may be used as scratch registers. Device handlers may not do EMT requests without executing a .SYNCH request.
5. The handler must return from an interrupt via an RTS PC.
6. When the transfer is complete, the handler must exit to the monitor to terminate the transfer or enter a completion routine. When return is made to the monitor, R4 should point to the fifth word of the handler.

7. The handler should contain an abort entry point (located at INTERRUPT SERVICE -2) to which control is transferred on forced exit. The abort entry point should contain a BR instruction to code that will perform the necessary operations (stop device action and exit to monitor completion code).

## CHAPTER 6

### F/B MONITOR DESCRIPTION

The RT-11 Foreground/Background Monitor permits two jobs to simultaneously share memory and other system resources. The foreground job has priority and executes until it is blocked (i.e., execution is suspended pending satisfaction of some condition, such as I/O completion). When the foreground job is blocked, the background job is activated and executes until it finishes or until the foreground blocking condition is removed.

#### 6.1 INTERRUPT MECHANISM AND .INTEN ACTION

All interrupt handlers must be entered at priority level 7 and must execute a .INTEN request on entry. The handler will then be called (as a co-routine of the monitor in system state) at its normal priority level. This is essential to the operation of RT11 for two reasons:

1. As a co-routine of the monitor, the interrupt handler exits to the monitor, which then does job scheduling.
2. Because of the above condition, there is a danger that interrupt processing may be postponed due to a context switch. For example, if a disk interrupts a lower priority device handler and goes to I/O completion, the monitor may switch to the foreground job and delay the lower priority interrupt until the foreground job is again blocked. By requiring the .INTEN request of all interrupt handlers, the monitor can assure that all interrupts are processed before the context switch is made.

The .INTEN request is implemented as a JSR R5 to the first fixed-offset location of RMON, which contains a jump to the interrupt entry code. This code saves R4 (R5 was saved by the JSR) and increments the system state counter. If the interrupt occurred on a job stack, the stack pointer is switched to use the system stack. The priority is lowered

to the handler's requested priority and control returns to the handler via another JSR instruction.

The handler interrupt code now executes in system state, with several results: any further interrupts are handled on the system stack, preventing their loss by a context switch to another job's stack; a context switch or completion routine cannot occur until all pending interrupts are processed; any error occurring in the handler occurs in system state, causing a fatal halt. When the handler exits via an RTS PC instruction, control returns to the monitor, which can now enter the scheduling loop if all interrupts have been processed.

## 6.2 CONTEXT SWITCH

When passing control from one job to another, the F/B Monitor does a complete context switch, changing the machine environment to that of the new job. The current context is saved on the stack of the current job and is replaced by the context of the new job.

The information saved on the stack includes:

1. The general registers (R0-R5)
2. The system communication area (memory locations 34-52)
3. The FPP registers, if used
4. The list of special locations supplied by the job (via .CNTXSW), if any

In addition, the stack pointer (R6) is saved in the job's impure area at offset I.SP (=50). The switch requires a minimum of  $23_{10}$  words of stack, not including the special swap list.

The following are the minimum calculated times to context switch between jobs. The assumptions are that the F/G job is waiting for I/O completion, the handler completes an I/O request, and there are no user I/O completion routines.

Processor	$\frac{11}{20}$	$\frac{11}{40}$	$\frac{11}{45}$
(core memory)	.66 ms	.36 ms	.28 ms

### 6.3 BLOCKING A JOB

The F/B Monitor gives priority to the foreground job, which runs until it is blocked by some condition. In this case, the background, if runnable (i.e., not blocked itself), is scheduled. The conditions which may block a job are flagged in the I.JSTA word, which is located in the job's impure area:

<u>Tag</u>	<u>Bit in I.JSTA Word</u>	<u>Condition</u>
TTIWT\$	14	Waiting for terminal input
TTOWT\$	13	Waiting for room in output buffer
CHNWT\$	11	Waiting for channel to complete
SPND\$	10	Suspended
NORUN\$	9	Not loaded
EXIT\$	8	Waiting for all I/O to stop
KSPND\$	6	Suspended from KMON
USRWT\$	4	Waiting for the USR

### 6.4 JOB SCHEDULING AND USE OF .SYNCH REQUEST

The F/B Monitor uses a scheduling algorithm to share system facilities between two jobs. The goal of the scheduler is to maximize system utilization, with priority given to the foreground job. The scheduler is generalized to use job numbers for scheduling, the higher job number having the higher priority. The background job is assigned job number 0 and the foreground job number 2. Job numbers must be even.

The foreground job runs until it is blocked by some condition (see Section 6.3), at which point the scheduler is initiated. The job list is scanned top down (from highest to lowest priority) for the highest priority job that is runnable. A job is runnable if it is not blocked, or if it is only blocked pending completion and is not suspended. If no jobs are currently runnable, the idle loop is entered.

If the new job is runnable, a context switch is made. The context switch routine tests for the completion pending condition (i.e., I/O is finished and a user completion routine was queued). In this case, a pseudo-interrupt is placed on the job's stack to call the completion queue manager when the scheduler exits to the job.

The scheduler is event driven and is entered from the common interrupt exit path whenever an event has occurred which requires action by the scheduler. The set of such events include:

1. An .EXIT or .CHAIN request
2. A job abort from the console, or an error abort
3. I/O transfer completed
4. Expiration of timed wait
5. A blocking condition encountered:
  - a. .TWAIT request or SUSPEND command
  - b. .TTYIN or .CSI waiting for end of line
  - c. .TTYOUT or .PRINT waiting for room in output buffer
  - d. Attempt to use busy channel
6. A blocking condition removed
7. No queue elements available
8. .SYNCH request (see below).

The .SYNCH request is used in interrupt routines to permit the issuing of other programmed requests. The .SYNCH macro is expanded as a JSR R5 to the .SYNCH code in the F/B resident monitor. The .SYNCH routine uses the associated 7-word block as a queue element for the completion queue.

If the .SYNCH block is not in use, register R5 is incremented to the successful return address and placed in the block as the completion address. The word count is set to -1 to prevent the block from being linked into the AVAIL queue. The block is placed in the completion queue, at its head, and the job associated with the .SYNCH request is flagged to have a completion routine pending. A request for a job switch is entered before the .SYNCH logic exits with an interrupt return.

On exit from the interrupt with a job switch pending, the scheduler is entered and the completion queue manager is called. When control finally returns to the code following the .SYNCH request, it is executing as a completion routine at priority level 0. It can now issue programmed requests without fear of being interrupted. If another interrupt comes in, and it requests a completion routine, the completion routine will be queued pending return of the current



completion routine, since the .SYNCH block is freed before calling the completion routine. Further interrupts will be rejected by the .SYNCH code, unless provision is made for supplying extra .SYNCH blocks.

#### 6.5 USR CONTENTION

The directory operations handled by the USR are not re-entrant, particularly since the directory segment is buffered within the USR. Therefore, to use the USR in F/B, a job must have ownership of the USR. To facilitate this, the F/B monitor maintains a USR queuing mechanism.

Before issuing a USR request, a job must request ownership of the USR. If the USR is in use by another job, even of lower priority, the requesting job is blocked and must wait for the USR. The USRWT\$ flag is set in the I.JSTA word (see Section 6.3) and the job cannot continue until the USR is released and the blocking bit cleared. When the USR is released, the job list is scanned for jobs waiting for the USR, starting with the job having highest priority.

Because of the impact this may have on system performance, CSI requests are handled differently in the F/B system than in the S/J Monitor. If the command string is to come from the console keyboard, the prompting asterisk is printed and then the USR is released, pending completion of command line input. This prevents a job doing a CSI request from locking up the USR and blocking another, perhaps higher priority, job from executing. A job can determine if the USR is available by doing a .TLOCK request (see Chapter 9 of the RT-11 System Reference Manual).

#### 6.6 I/O TERMINATION

Because of the multi-job capabilities of RT-11 F/B, termination of I/O on job exit or abort must be handled differently than in the S/J Monitor. The use of the RESET instruction is unacceptable, and a form of I/O rundown must be used. This is done by the IORSET routine, called when doing an abort or hard exit.

The IORSET routine searches the queue of every resident handler for elements belonging to the aborted job. If a handler is found to be resident and active (i.e., there are elements on its queue), the IORSET routine "holds" the handler from initiating a new transfer by setting bit 15 of the LQE word (entry point) in the handler. The

current transfer may complete, but the hold bit will prevent the queue manager from initiating a new transfer.

While it is held, the handler's queue is examined for the current request. If it belongs to the aborted job, the handler's abort entry point is called to stop the transfer. The queue of pending I/O requests is then examined and any elements belonging to the aborted job are discarded. The hold flag is cleared and a test is made to see if the current transfer completed while the handler was held. If it did, the completion queue manager, COMPLT, is again called to return the completed element and initiate the next transfer. At this point, any elements belonging to the aborted job will have been removed from the queue.

After the device handlers are purged, the internal message handler is examined for waiting messages that were originated by the aborted job. All such messages are discarded. Finally, all mark time requests belonging to the aborted job are cancelled.

## CHAPTER 7

### RT-11 BATCH

The RT-11 BATCH system is composed of a BATCH compiler and a run-time handler. The BATCH compiler converts BATCH Job Control language into a format comprehensible to the BATCH run-time handler. The compiler creates a control (CTL) file (from the BATCH language statements) which is then scanned by the handler; the CTL format is a versatile programming language in its own right. The result is a BATCH system that is simple to use, and yet easily customized to handle different situations.

#### 7.1 CTL FORMAT

The BATCH run-time handler uses a unique language format that includes many programming features, such as labels, variables, and conditional branches. The directives are explained in detail in Chapter 12 of the RT-11 System Reference Manual.

Each directive consists of a backslash character followed by one or more other characters. For example, to run PIP and generate a listing, the CTL directives \E (execute) and \D (data line) are used:

```
\ER PIP
\DLP:=/L
```

Messages are sent to the console device by using the \@ directive:

```
\@ PLEASE MOUNT DT2
```

Labels and unconditional branches are implemented with the \L (label) and \J (jump) directives:

```
\JEND 1
.
.
\LEND
```

Each BATCH command is sent to the log as it is executed, using the \C (comment) directive:

```
\C
$JOB
```

In this case, every character up to the next backslash is sent to the log.

## 7.2 BATCH RUN-TIME HANDLER

The BATCH run-time handler (BA.SYS) is constructed as a standard RT-11 device handler. To use the handler, it must be made permanently resident via the monitor LOAD command. The handler links itself into the monitor, intercepting certain EMTs described later.

The linking occurs the first time the BATCH compiler is run after the BA handler is loaded. The compiler does a .READW to the BA handler, which then links itself to the monitor and returns a table of addresses to the BATCH compiler. The linking is achieved by replacing the addresses of monitor EMT routines with corresponding addresses in the BATCH handler. Those EMTs that are diverted include:

<u>EMT</u>	<u>BATCH Handler Routine</u>
.TTYIN	B\$TIN
.TTYOUT	B\$TOT
.EXIT	B\$EXT
.PRINT	B\$PRN

Once the link is established, the BATCH handler cannot be unloaded. The links must first be undone by again running the BATCH compiler and specifying the /U switch. The compiler removes the links and prints a prompting message, after which the UNL BA command can be issued.

With the BA handler linked to the monitor, all console terminal communication is diverted to BA, along with program exits. The BA handler then dispatches the program request to the monitor routine or diverts it to a routine in BA, depending on the values of switches in BATSW1. The switches are:

<u>TAG</u>	<u>BIT</u>	<u>DESCRIPTION</u>
HELP	0	0 = Do not log terminal input (.TTYIN) 1 = Log terminal input
DESTON	1	0 = EMT is going directly to monitor 1 = BA intercepts the EMT
SOURCE	2	0 = Character input by monitor from console terminal 1 = Character input comes from BATCH stream
COMWAT	3	0 = No command 1 = Command is waiting
ACTIVE	4	0 = Console terminal inactive 1 = Console terminal is active; i.e., BA is waiting for input from console terminal
DATA	5	0 = Characters are going to KMON; i.e., KMON is active in B/G 1 = Characters are going to B/G programs
BDESTN	6	0 = Output characters are going to console terminal 1 = Output characters are going to LOG
BGET	7	0 = Normal mode 1 = Get mode (\G); input comes from console terminal until <CR><LF> is encountered
NOTTY	8	0 = Log terminal output 1 = Do not log terminal output (.TTYOUT, .PRINT)
	9-13	Reserved
BSOURC	14	0 = BA directives come from console terminal 1 = BA directives come from CTL file
BEXIT	15	1 = A program has done an .EXIT while DATA switch was set

The BATSW1 word, located six bytes past the handler entry point, determines the state of the system at any given moment. If the word is zero, RT-11 operates normally. When the DESTON bit is set, EMTs are diverted to routines in BA for action, but the specific action taken by those routines is determined by the other switch bits.

For example, if the BDESTN bit is set, output from .TTYOUT and .PRINT is diverted from the console terminal to the log device. If SOURCE is set, the characters for the .TTYIN request are taken from the BATCH stream rather than from the console terminal via the monitor ring buffer. Directives for the BA handler itself may come from either the CTL file or the console terminal, depending on the state of the BSOURC bit.

The state of the background is reflected in the DATA bit. Either the KMON is active (DATA=0) or a program is active (DATA=1). If a program issues an .EXIT request while in DATA mode, the BEXIT state is entered until the BA handler encounters the next KMON directive (\E) in the BATCH stream, causing any unused \D lines to be ignored. A program can be aborted by diverting any of the .TTYIN, .TTYOUT or .PRINT requests to the .EXIT code in the monitor.

### 7.3 BATCH COMPILER

The obvious function of the BATCH compiler is to convert BATCH Standard Commands into the BA handler directives mentioned in Section 7.1, creating a control (CTL) file. BATCH jobs entered from a card reader or a file-structured device are compiled into a CTL file stored on a file-structured device for execution by the BA handler. However, the BATCH Compiler has other important functions; these are described in this section along with details on the initiation and termination of BATCH jobs.

#### 7.3.1 BATCH Job Initiation

The following sequence of actions is performed by the BATCH Compiler when setting up a job for execution:

1. A check is made to ensure that LOG and BA device handlers are loaded and assigned properly. The LOG handler must be assigned the logical name LOG;; the BATCH Compiler may be run several times during the course of a job to do special tasks for the BA handler, and it will reference LOG:.
2. A nonfile-structured .LOOKUP is done on BA and a .READW is issued. If this is the first time BATCH has been run since BA was loaded, the handler links itself to the monitor (see Section 7.2). BA returns a list of

eleven pointers to important parameters within BA.  
These include:

- BA state word (BATSW1)
- CTL file savestatus area (INDATA)
- LOG file savestatus area (ODATA)
- Output (LOG) buffer (OUTBUF)
- Output buffer pointer (BATOPT)
- Output character counter (BATOCT)
- Input character counter (BATICT)
- Monitor EMT dispatch address save areas

3. A command string is collected from the console terminal and is processed by .CSISPC. An input file must be specified.
4. If the input file is a .BAT file to be compiled, a .CTL file is entered. If the LOG: device is file-structured, a fixed-size enter is done and then the file is initialized by writing zeroes in all blocks.
5. A .LOOKUP is done on all input files.
6. The .LOG file is .CLOSED so that a .LOOKUP and .SAVE-STATUS may be done. The savestatus data is placed in the ODATA area in BA.
7. If the input file is a .BAT file, it is now compiled, with output going into the .CTL file.
8. The .CTL file is closed, again so that a .LOOKUP and .SAVSTATUS may be done. The .SAVSTATUS data is transferred to the INDATA area in BA. Buffer pointers and counters in BA are initialized.
9. The BA handler is activated by setting the SOURCE, DESTON, BSOURC and BDESTN bits in the BATSW1 state word in BA. Control passes to BA when the compiler does an .EXIT, assuming an abort is not requested.
10. If an abort is requested (an error occurred during compilation or the /N switch was used), the .LOG file is .REOPENed and all \$ command lines are logged out with any error diagnostics. The BATSW1 word is then cleared before exiting, preventing the execution of the job.

The following switches are used by the BATCH system during job initiation and continuation, and should not be typed by the user:

- /B BATCH continuation of jobs in input stream
- /D Print the physical device name assigned a logical device name in a \$DISMOUNT command
- /M Make a temporary source file
- /R Return from \$CALL
- /S \$CALL subroutine

### 7.3.2 BATCH Job Termination

Every BATCH job must be terminated with an \$EOJ statement. The \$EOJ statement causes the compiler to insert the CTL directives:

```
\R BATCH
\D/R
```

The /R switch for the BATCH compiler, which is legal only when entered from a BATCH stream, is used to terminate a BATCH job. This switch causes the compiler to pop the BATCH stack up a level. If the stack was empty, the stream is finished and the compiler cleans up, clears the BATSW1 word in BA, and exits. If the stack is not empty, the /R switch implies a return from a \$CALL. The stack contents are used to restore parameters in the BA handler so that control will return to the calling BATCH stream at the next statement after the \$CALL.

### 7.3.3 BATCH Compiler Construction

The BATCH Compiler is constructed in two pieces: a data area and a program area. The data area is located in low memory, in a .CSECT named UNPURE. The contents are described in the accompanying table (Table 7-1). The program section, located in the .CSECT named PROGRM, starts at the symbol START. The general register R4 always points to UNPURE and all references to the data base are made as indexed references relative to R4.

Locations in the data base are created with the ENTRLO macro. For example,

```
ENTRLO BOTLCT,0
```

allocates one word in the data base and initializes it to zero. The symbol BOTLCT is an offset into the data base, so that references to BOTLCT are made in the form BOTLCT(R4).



Table 7-1  
 BATCH Compiler Data Base Description

Tag	Byte Offset	Description
BATSWT	0	BATCH Control Switches  ABORT = 10000      ABORT after compile DATDOL = 40000    DATA or DOLLARS set NO        = 20000     "NO" prefix on switch CTYOUB = 10000     Output to CTY (\ @) LOGOUB = 4000      Output to LOG (\ C) DATOUB = 2000      Output to user prog (\ D) COMOUB = 1000     Output to monitor (\ E) JOB        = 400      \$JOB encountered MAKEB    = 200      /B switch on command COMMA    = 100      Comma terminates command BFORLI    = 40      Next link requires FORTRAN library  UNIQUE = 20      UNIQUE command option set BANNER = 10      Print BANNER on \$JOB, \$EOJ RT11    = 4      RT11 default on NO '\$' in Column 1  TIME     = 2      Print time of day MAKE    = 1      Create a source file
BATSW2	2	More BATCH Control Switches  ABORT =100000      Second time through ABORT FIRST = 10000      First card processed SBIT    = 4000      /S switch on command SEQ     = 2000      \$SEQ card processed LSTBIT = 1000      Request temporary listing file  COMSWB = 400      Command switches MAKEB    = 200      Same as BATSWT STARFD = 100      Asterisk in FD field STAROK = 40      Wild card option is valid BNOEOJ = 20      \$JOB or \$SEQ before \$EOJ LSTDAT = 10      List DATA sections BEOF    = 4      EOF encountered on .BAT file  XSWT    = 2      /X switch set EOJ     = 1      \$EOJ encountered
TMPSWT	4	Temporary command switches
COMSWT	6	Current command switches
LINSIZ	10	Input line buffer size
BINLCT	12	Last buffer character count
INSTAT	14	Input buffer status (see OTSTAT)
ICHRPT	16	Input character pointer
BINCTR	20	Input buffer counter

(continued on next page)

Table 7-1 (Cont.)  
 BATCH Compiler Data Base Description

Tag	Byte Offset	Description
BINARG	22	Input file EMT argument list
BATIBK	24	Input file block number
BATIBP	26	Input buffer address
	30	Input buffer size
	32	Wait I/O
BOTLCT	34	Last output buffer character count
OTSTAT	36	Output buffer status  BFREE = 1      0 → Buffer is free BWAIT = 2      In I/O wait BEOF = 4       End of file
OCHRPT	40	Output character pointer
BOTCTR	42	Output character count
BOTARG	44	Output file EMT argument list
BATOBK	46	Output file block number
BATOBP	50	Output buffer address
	52	Output buffer size
	54	Wait I/O
STACK	56	Compiler stack pointer save area  These are the arguments passed between BATCH and BA:
BATSW1	60	Pointer to BATSW1 in BA.SYS
INDATA	62	Pointer to INDATA
ODATA	64	Pointer to ODATA
OUTBUF	66	Pointer to BATCH handler output buffer
BATOPT	70	Pointer to output character pointer
BATOCT	72	Pointer to output character counter
BATICT	74	Pointer to input character counter

(continued on next page)

Table 7-1 (Cont.)  
 BATCH Compiler Data Base Description

Tag	Byte Offset	Description
		Pointers to EMT intercept pointers:
O\$EXT	76	.EXIT
O\$TIN	100	.TTYIN
O\$TOT	102	.TTYOUT
O\$PRN	104	.PRINT
		CSI Buffer:
SPC0	106	Channel 0
SPC1	120	1
SPC2	132	2
SPC3	144	3
SPC4	154	4
SPC5	164	5
SPC6	174	6
SPC7	204	7
SPC8	214	10
LINIMP	224	Pointer to command line buffer (LINIMM)
LINIMM	226	Command line input buffer
LINIMS	350	Command line buffer save area
LIBLST	470	ASCII name of FORTRAN default library plus a line buffer
BATIBF	610	BATCH Compiler input buffers (INBSIZ * 2)
BATOBF	2610	BATCH Compiler output buffers (OTBSIZ * 2)
QSET	4610	Seven I/O queue elements for double/buffering
SOUTMP	4700	Source temporary file descriptor
OBJTMP	4714	Object temporary file descriptor
LOGTYP	4730	LOG device status word (word 0 of .DSTATUS)
ARGARG	4732	EMT argument list for BA handler initialization

(concluded on next page)

Table 7-1 (Cont.)  
 BATCH Compiler Data Base Description

Tag	Byte Offset	Description
STKBLK	4744	EMT argument list for READ/WRITE of BATCH stack
DEFCHN	4756	Default channel numbers
DEVSPC	4770	Pointer to device handler space
WDBLK2	4772	Two-word EMT argument block
WDBLK5	5000	Five-word EMT argument block
FTLPC	5012	Contents of PC on BATCH fatal error
AREA0	5014	Pointer to impure area
LSTTMP	5016	Listing temporary file descriptor
SWTMSK	5026	Switch mask for this BATCH directive
FD0	5030	File descriptor 0 for BATCH directive
FD1	5034	1
FD2	5040	2
FD3	5044	3
FD4	5050	4
FD5	5054	5

#### 7.4 BATCH EXAMPLE

The following example demonstrates how the compiler converts BATCH Standard Commands into RT-11 BATCH handler directives. The example consists of a main BATCH stream, EXAMPL.BAT, and a BATCH subroutine file, EDITIT.BAT. EXAMPL creates a program, assembles and runs it. The program, called FILE.MAC, prints a message that is diverted to the log. The listing file from the assembly is printed and then deleted. The BATCH variable S is then tested and, if it is zero, the BATCH subroutine EDITIT is called. The EDITIT stream uses EDIT to edit the file FILE.MAC, changing the message to be printed. After return from EDITIT, the stream branches unconditionally to label L1, repeating the assembly and execution of FILE.MAC. EDITIT increments the variable S before returning, so that the BATCH stream, on encountering the IF statement again, now branches to label L2, skipping the call to EDITIT. \$DIRECTORY and \$DELETE operations are performed before finally exiting from BATCH.

Note the following about the .CTL files created:

1. The \$JOB command produces a comment for the log (the \C directive, but no action directives). Its function is to initialize the BATCH compiler.
2. The \$CREATE command produces directives that run the BATCH compiler, using the file name to be created with a /M switch. This is a special function of the BATCH compiler used to create data files. The compiler will enter the data that follows in the CTL file into the newly created file, until an EOF (CTRL/Z) is encountered. The data is fed to the compiler by the BATCH handler through the .TTYIN programmed request. After the EOF character is encountered, the BATCH compiler closes the new file and exits, returning control to the BATCH handler through the .EXIT request. In this example, the file created is called FILE.MAC.
3. The \$MACRO command has the /RUN switch appended, which forces the compiler to generate a series of assembly, link and execute instructions. A temporary execution file, 000000.SAV, is created from the assembled object module, FILE.OBJ. After execution with the monitor R command, the temporary execution file is deleted with PIP.
4. PIP is used to implement \$PRINT, \$DELETE, \$COPY, and \$DIRECTORY. The compiler translates these commands into the appropriate PIP command strings.

5. The variable S is defined to be zero with the LET statement. This translates into the BATCH handler directive,

```
\KS1<null>
```

which instructs the BATCH handler to set variable S to the value in the byte following the character 1.

6. Labels are implemented by inserting a \L directive followed by the 6-character label name into the CTL stream where the label was declared. The label is also logged out with the \C directive so that the labels will appear in the log.
7. The unconditional branch, or GOTO command, is implemented with the \J directive immediately followed by the label. Note that the BATCH programmer must indicate whether the branch is forward or reverse. In this case, the branch is a backward reference and a minus sign is prefixed to the label:

```
GOTO -L1
```

There is no error checking done by the compiler. If an error is made (e.g., the minus sign is left off the L1), the BATCH handler searches forward in the CTL stream until it finds the label. Since an error was made, the label will not be found. The search (and consequently the BATCH job) terminates when the label stopper (\L\$\$\$\$\$\$) is encountered at the end of the CTL file.

8. The IF conditional branch is implemented with the \I directive. The \I directive is followed by the name of the variable to be tested, the value to be tested against, and three label fields. Each label field consists of the 6-character label name with a reference character appended. The character 1 indicates the label is a forward reference, a 0 indicates a backward reference. The test value is subtracted from the current value of the variable and the appropriate branch is taken. If no label is specified for a field, it is filled with spaces and causes the BATCH stream to fall through to the next command if that branch is elected.
9. The \$CALL command is very useful and permits a BATCH stream to call another BATCH file as a subroutine, with control returning to the command following the \$CALL. The \$CALL is implemented by simply running the BATCH compiler, passing it the name of the \$CALLED routine with a /S switch appended. Another BATCH compile/execute sequence will follow, but the /S switch will cause the compiler to save certain locations in the BATCH handler in an internal stack in the BA.SYS file. In this example, the \$CALL EDITIT statement causes the file EDITIT.BAT to be compiled and executed.

10. BATCH variables may be used to enter ASCII values into a job stream. In the file EDITIT, the variable A is set equal to the value of the ESC (or ALT MODE) character. The variable A is inserted into a string of EDIT commands in place of the ALT MODE character.
11. The \$EOJ must terminate every BATCH job. The \$EOJ command generates the stopper label, \L\$\$\$\$\$, and then produces directives to run the BATCH compiler again, this time with a /R switch. The compiler, when given a /R switch, checks the BATCH stack. If it is empty, the compiler exits. Otherwise, the stack is popped to restore conditions in the BATCH handler prior to the \$CALL causing the push, and the BATCH stream continues. The \$EOJ finally generates a \E to bring in the KMON and a \F<CR> to terminate the BATCH stream.

EXAMPL.BAT

```
$JOB
$MESSAGE EXAMPLE BATCH STREAM
$CREATE FILE,MAC
      ,MCALL .REGDEF,.PRINT,.EXIT
      .REGDEF

START: .PRINT #MSG
      .EXIT
      .NLIST REY
MSG:   .ASCIZ /THIS MESSAGE COMES FROM THE BATCH STREAM/
      .EVEN
      .LIST REY
      .END START

SEOD
SRT11
      LET S=0

L1:
$MACRO/RUN FILE,LST/LIST FILE,MAC/INPUT FILE/OBJECT
$PRINT FILE,LST
$DELETE FILE,LST
SRT11
      IF(S=0) ,L2
$CALL EDITIT !CALL EDITIT TO EDIT FILE,MAC
SRT11
      GOTO -L1

L2:
$DIRECTORY FILE.*
$DELETE FILE.*
SEOD
```

EDITIT.BAT

```
$JOB/RT11
$! JOB TO EDIT FILE,MAC
      %S
      LET A=33
      !INCREMENT S TO PREVENT RECURSION
      !A IS ALT MODE

,R EDIT
*EBFILE,MAC'A'R'A'A'
*GMSG:'A'KI .ASCIZ /MODIFIED BY EDITOR RUN BY BATCH/
*'A'EX'A'A'
SEOD
```



EXAMPL.CTL

```
\C
$JOB

$MESSAGE EXAMPLE BATCH STREAM
\E\@ EXAMPLE BATCH STREAM
\C
$CREATE FILE.MAC
\ER BATCH
\DFILE.MAC/M=
    .MCALL .RFGDEF,.PRINT,.EXIT
    .REGDEF
START: .PRINT #MSG
    .EXIT
    .NLIST BEY
MSG: .ASCIZ /THIS MESSAGE COMES FROM THE BATCH STREAM/
    .EVEN
    .LIST BEY
    .END START

\C
$EOD

$RT11
    LET S=0
\KS1 \LL1 \CL1:

$MACRO/RUN FILE.LST/LIST FILE.MAC/INPUT FILE/OBJECT
\ER MACRO
\DFILE,FILE.LST=FILE.MAC
\F\D\ER LINK
\D000000=FILE
\ER \D000000
\ER PIP
\D000000.SAV/D
\C
$PRINT FILE.LST
\ER PIP
\DLST:*,*/X=FILE.LST
\F\D\C
$DELETE FILE.LST
\ER PIP
\DFILE.LST/D
\C
$RT11
    IF(S=0) ,L2
\IS 1 1L2 1\L \C
$CALL EDITIT !CALL EDITIT TO EDIT FILE.MAC
\F\ER BATCH
\DEDITIT/S
\C
$RT11
    GOTO -L1
\JL1 0\LL2 \CL2:
```

EXAMPL.CTL (Cont)

```
SDIRECTORY FILE.*
\ER PIP
\DFILE.* /L
\F\D\C
SDELETE FILE.*
\ER PIP
\DFILE.* /D
\C
SEQJ
\LS$$$$$F\ER BATCH
\D/R
\EF
```

EDITIT.CTL

```
\C
$JOB/RT11

$! JOB TO EDIT FILE.MAC
      %S          !INCREMENT S TO PREVENT RECURSION
\KS0\C LET A=33  !A IS ALT MODE
\KA1 \ER EDIT
\DEBFILE.MAC\KA2R\KA2\KA2
\DGMSG!\KA2KI .ASCIZ /MODIFIED BY EDIT RUN BY BATCH/
\DKA2EX\KA2\KA2
\C
SEQJ
\LS$$$$$F\ER BATCH
\D/R
\EF
```

EXAMPL.LOG

\$JOB

\$MESSAGE EXAMPLE BATCH STREAM

\$CREATE FILE.MAC

\$EOD

\$RT11

LET S=0

L1 L1:

\$MACRO/RUN FILE.LST/LIST FILE.MAC/INPUT FILE/OBJECT

\*ERRORS DETECTED: 0  
FREE CORE: 15100. WORDS

\*

EXAMPL.LOG (Cont.)

.MAIN. RT-11 MACRO VM02-10 10-APR-75 10:33:45 PAGE 1

```
1          .MCALL .REGDEF,.PRINT,.EXIT
2 000000   .REGDEF
3 000000   START: .PRINT #MSG
4 000006   .EXIT
5          .NLIST  BEX
6 000010   124 MSG: .ASCIZ /THIS MESSAGE COMES FROM THE BATCH STREAM/
7          .EVEN
8          .LIST   BEX
9          000000' .END  START
```

EXAMPL.LOG (Cont.)

.MAIN. RT-11 MACRO VM02-10 10-APR-75 10:33:45 PAGE 1+  
SYMBOL TABLE

MSG	000010R	PC	=X000007	R0	=X000000
R1	=X000001	R2	=X000002	R3	=X000003
R4	=X000004	R5	=X000005	SP	=X000006

START 000000R  
. ABS. 000000 000  
000062 001

ERRORS DETECTED: 0  
FREE CORE: 15100. WORDS

FILE,FILE.LST=FILE,MAC

THIS MESSAGE COMES FROM THE BATCH STREAM

SPRINT FILE.LST

SDELETE FILE.LST

SRT11  
IF(S=0) ,,L2

SCALL EDITIT !CALL EDITIT TO EDIT FILE,MAC

SJOB/RT11

S! JOB TO EDIT FILE,MAC  
XS !INCREMENT S TO PREVENT RECURSION  
LET A=33 !A IS ALT MODE

\*EBFILE,MACSRSS

\*  
GMSG:SKI .ASCIZ /MODIFIED BY EDITOR RUN BY BATCH/  
SEXSS

SE0J  
SSSSSS

SRT11  
GOTO -L1

L1:

SMACRO/RUN FILE.LST/LIST FILE,MAC/INPUT FILE/OBJECT

\*ERRORS DETECTED: 0  
FREE CORE: 15136. WORDS

\*

EXAMPL.LOG (Cont.)

.MAIN. RT=11 MACRO VM02-10 10-APR-75 10:34:08 PAGE 1

```
1
2 000000          .MCALL  .REGDEF,.PRINT,.EXIT
3 000000          .REGDEF
4 000006          .PRINT  #MSG
5                .EXIT
6 000010          .NLIST  BEX
7                .ASCIZ  /MODIFIED BY EDITOR RUN BY BATCH/
8                .EVEN
9 000000          .LIST   BEX
                .END    START
```

EXAMPL.LOG (Cont.)

.MAIN. RT-11 MACRO VM02-10 10-APR-75 10:34:08 PAGE 1+  
SYMBOL TABLE

MSG	000010R	PC	=X000007	R0	=X000000
R1	=X000001	R2	=X000002	R3	=X000003
R4	=X000004	R5	=X000005	SP	=X000006

START 000000R  
. ABS. 000000 000  
000050 001

ERRORS DETECTED: 0  
FREE CORE: 15136. WORDS

FILE,FILE,LST=FILE,MAC

MODIFIED BY EDITOR RUN BY BATCH

\$PRINT FILE.LST

\$DELETE FILE.LST

\$RT11  
IF(S=0) ,,12  
L2:

\$DIRECTORY FILE.\*

10-APR-75  
FILE .BAK 1 10-APR-75  
FILE .MAC 1 10-APR-75  
FILE .OBJ 1 10-APR-75  
3 FILES, 3 BLOCKS  
417 FREE BLOCKS

\$DELETE FILE.\*

\$EOJ

## 7.5 CTT TEMPORARY FILES

In certain cases the BATCH compiler will produce temporary files with the extension CTT and the file name of the BAT file being compiled. These files occur when a multiple input file command string is issued, or when an unexpected \$JOB or \$SEQ statement occurs in a BATCH stream, or when multiple jobs are run from the card reader or a .BAT file.

The CTT file is actually a CTL file used to link together execution of several BATCH jobs. Each CTT file contains the BA directives:

```
\ER BATCH  
\D/B
```

which execute the BATCH compiler, passing it the /B switch.

The CTT file also contains the following information:

1. Current input channel number (range is 3-10<sub>8</sub>)
2. Current input file block number
3. The CTL file descriptor block (device, file name and file size)
4. The LOG file descriptor block (device, file name, and file size)
5. The set of input (BAT) file descriptor blocks (device and file name)

When the CTT file is executed, the compiler restores the input channel number and block number and the entire set of file descriptor blocks from the CTT file. If, for example, the input channel number is 4, the second of a string of .BAT files is compiled and executed.



APPENDIX A

SAMPLE HANDLER LISTINGS

A.1 RC11/RS64 DEVICE HANDLER

RC11 V01-01 (FIXED HEAD DISK) RT-11 MACRO VM02-09 8-APR-75 12104126 PAGE 1

1 .TITLE RC11 V01-01 (FIXED HEAD DISK)  
2 RT-11 RC11/RS64 DEVICE HANDLER

3  
4 JDFC-11-XXXX-A

5  
6 JJFG

7  
8 JOCTOBER 1974

9  
10 JCOPYRIGHT (C) 1975

11  
12 DIGITAL EQUIPMENT CORPORATION  
13 MAYNARD, MASSACHUSETTS 01754

14  
15 THIS SOFTWARE IS FURNISHED UNDER A LICENSE FOR USE ONLY  
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24 TO CHANGE WITHOUT NOTICE AND SHOULD NOT  
25 BE CONSTRUED AS A COMMITMENT BY DIGITAL  
26 EQUIPMENT CORPORATION. DIGITAL ASSUMES NO  
27 RESPONSIBILITY FOR ANY ERRORS THAT MAY APPEAR  
28 IN THIS DOCUMENT.

29  
30 DIGITAL ASSUMES NO RESPONSIBILITY FOR THE  
31 USE OR RELIABILITY OF ITS SOFTWARE ON  
32 EQUIPMENT WHICH IS NOT SUPPLIED BY  
33 DIGITAL.  
34  
35

```

1 .SMTL HANDLER DEFINITIONS
2 .MCALL .REGDEF, .V2..
3
4
5
6
7
8

```

```

9
10
11
12
13
14
15
16
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24
25
26
27
28
29
30
31
32
33
34
35
36
37
38
39
40

```

```

000000
000001
000002
000003
000004
000005
000006
000007

```

```

IRPGISTER DEFINITION
LIST ME
REGDEF

MONITOR REFINED CONSTANTS
MONLOW = 44
IPINITER TO G MANAGER
COMPLETION ENTRY
HERR = 1
RTRY = 8.

```

```

PRIORITY CONSTANTS
PR7 = 300
PR5 = 200

```

```

RC-11 COMMUNICATION CONSTANTS
WR = 103
RD = 104
TNHCA = 1000
ABORT = 400
RTRYFR = 060000
RBYT 14 = 1 ==> DATA ERROR
RBYT 13 = 1 ==> ADDRESS ERROR

```

```

CONTROL REGISTERS
RCLA = 177440
RCPA = 177442
RCPR = 177444
RCCS = 177446
RCWC = 177450
RCCA = 177452
RCMN = 177454
RCDB = 177456

```

```

41 000210                                RCVEC = 210      INTERRUPT VECTOR ADDRESS
42
43
44
45
46
47
48
49
50

```

IRC SYSTEM DEFINITIONS

```

A  'GLOBAL RCYSYS, RKSYS, RFSYS, DPSYS, DSSYS, DTSYS, DXSYS
   'GLOBAL SINTPR, SINTEN, RCINT          ISIZE IS REQUIRED BY BSTRAP
   .GLOBAL RCISIZE

```

B

```

RKSYS = 0

```

RC11 VM01-01 (FIXED HEAD DISK) RT-11 MACRO VM02-09 A-APR-75 12104126 PAGE 2+

```

50 000000
51 000000
52 000000
53 000000
54 000000

```

```

RFSYS = 0
DPSYS = 0
DSSYS = 0
DTSYS = 0
DXSYS = 0

```

RC11 VM01-01 (FIXED HEAD DISK) RT-11 MACRO VM02-09 A-APR-75 12104126 PAGE 3

```

1  ;SMTL RC11 DEVICE HANDLER
2  ;(MAXIMUM SUPPORT 1 CONTROLLER AND 4 R864 DT8KS)
3  ;(1924 BLOCKS OF 256 WORDS)
4  ;CSECT SYSHND
5
6 000000
7
8
9
10 000000 000210
11 000002 000060
12 000004 000340
13 000006 000000
14 000006 000000
15 000010 000000
16
17
18 000012 012727 000010
19 000016 000000
20 000020 004067 000226
21 000024 000000
22 000026 062704 000010
23 000032 012514
24 000034 012544
25 000036 012705 000103

```

RC11 DEVICE HANDLER

C

```

'WORD RCVEC
'WORD RCINT=
'WORD PRY

```

D

```

RCYSYS
RCLOGI
RCOGEI

```

E

```

#ENTRY POINT
MOV #RCTRY, (RC)+
#WORD 0
RCOETRI JSR R0, RCOO01
ZEROI
ADD #10, R4
MOV (R5)+, 0R4
MOV (R5)+, -(R4)
MOV #WR, R5

```

RC11 DEVICE HANDLER

LOAD POINT

ADDRESS OF INTERRUPT VECTOR  
OFFSET TO INTERRUPT ROUTINE  
PRIORITY 7

POINTER TO LAST 0 ENTRY  
POINTER TO CURRENT 0 ENTRY

SPECIFY THE REPLY COUNT  
REPLY COUNTER  
SET UP THROUGH COMMON ROUTINE  
END ZERO FILL INITIATION  
POINT TO RCCA  
SET BUFFER ADDR (RCA)  
SET WORD COUNT ADDR (RCWC)  
ASSUME WRITE FUNCTION

```

26 000042 000714
27 000044 001427
28 000046 100402
29 000050 129525
30 000052 005414
31 000054 010544
32 000056 000207

```

```

TST 0R4
REQ RCHOME
RMT 18
CMPB (RS)+,(RS)+
NEG 0R4
MOV RS,=(R4)
RTS PC

```

```

PCHECK WORD COUNT
INC I/O
WRITE SPECIFIED
ISPT READ FUNCTION CODE
MAKE WORD COUNT (-) (RCWC)
ISPT PROPER FUNCTION IN RCCS
GO-AWAY FOR I/O

```



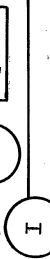
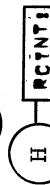
RC11 V01-01 (FIXED HEAD DISK) RT-11 MACRO VM02-09 A-APR-75 12:04:26 PAGE 4  
 RC-11 INTERRUPT ROUTINE

.SRTTL RC-11 INTERRUPT ROUTINE

```

1
2
3
4 000060 000430
5
6
7 000062 004577 000226
8 000066 000100
9
10
11
12 000070 016705 177714
13 000074 019704 177446
14 000100 005714
15
16
17 000102 100023
18 000104 032714 060000
19 000110 001403
20 000112 005367 177700
21 000116 003300
22 000120 052755 000001
23
24
25 000124 010704
26 000126 062704 177662
27 000132 013705 000054
28 000136 000175 000270
29
30 000142 019737 000400
31 000150 000765
32
33
34 000152 004067 000100
35 000156 001000 001000
36 000160 014505 000002

```



```

JSR R5,0SINPTR
.WORD 2C4RS>&PR7
PCHECK FOR ERROR ON RC & RTRY IF APPLICABLE

```

ABORT ROUTINE FOR F/0  
 MONITOR

NOTIFY MONITOR & RET  
 PRIORITY TO LEVEL 5  
 IRC RUNS AT PRIORITY 5  
 IF APPLICABLE

```

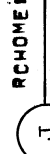
IGET CURRENT QUEUE ADR
IPPOINT TO DSK CNTRL & STATUS REG
IPDATA ADDRESS OR WRITE
PCHECK ERROR OR WRITE=LOCK
INON-EXISTENT DISK?
INC.
IF ONLY DATA & ADR ERRORS MERIT A RETRY
INOPF - GO BACK WITH ERROR
IPRETRY FATAL ERROR 8 TIMES
IF OK DO IT AGAIN.
IF CAN'T GET BY HARD ERROR
ISPT ERROR FLAG +
IGO BACK TO MONITOR

```

```

RCHOME: MOV PC,R4
ADD #RCCOE,.,R4
MOV #MONLOW,R4
JMP #OFFSET(R5)

```



```

RCAVRT: MOV #ABORT,0RCCS
BR RCHOME

```



```

IF FILL REMAINDER OF DISK BLACK WRITTEN TOO WITH ZEROS.
PCFILL: JSR R0,RCCOMN
.WORD INHCA
MOV P(R5),R5

```

IPYC ADDR OF CURRENT QUEUE ENTRY  
 IF EXIT TO MONITOR QUEUE COMPLETION  
 ABORT CURRENT OPERATION  
 IF EXIT TO MONITOR

```

37 000164 100357
38 000166 104705
39 000170 001755
40 000172 005405
41 000174 010546
42
43
44
45
46 000176 013746
47 000202 00A205
48 000204 10A066
49 000210 105316
50 000212 001373
51 000214 005726
52 000216 001401
53 000220 005205
54 000222
55
56 000222 060524
57 000224 012605

```

```

;CURRENT QUE ELEMENT (RCQOE)
;NFILL FOR READS
;EVEN # BLOCKS WRITTEN
;YFS - FILL NOT NECESSARY.
;WRITE WORD COUNTS ARE NEGATIVE IN
;THE QUE.
; CALCULATE THE # OF SECTORS IN
; THE CURRENT OPERATION
; (32 WORDS = ONE SECTOR)

```

```

MOV #5,=(SP)
;PUSH REPEAT COUNT ONTO STACK
;DTRDOP THE # WORDS
;CHECK FOR SECTOR OVERFLOW
;INCREMENT REPEAT COUNT
;SECTOR OVERFLOW ?
;NO
;INCLUDE NEXT SECTOR

```

```

; CALCULATE CURRENT DISK ADR. (RCDA)

```

```

RC11 V01-01 (FIXED HEAD DISK) RT-11 MACRO VM02-09 8-APR-75 12104126 PAGE 4+
RC-11 INTERRUPT ROUTINE

```

```

58 000226 052705 177400
59
60
61 000232 005724
62 000234 012724
63 000240 010524
64 000242 010714
65 000244 062714 177560
66 000250 000207

```

```

;WRITE MUST BE LESS THAN
;A BLOCK (RCWC TAKES
;2'S COMPLEMENT NEG. VALUE.)
;POINT TO RCDS
;SET WRITE FUNCTION
;SET WORD COUNT (RCWC)
;POINT MEMORY ADDRESS TO A ZERO.
;PTC (INTO RCCA)
;EXIT

```



```

1  .SRTL COMMON SUBROUTINE
2  IRCOMN
3  ICOMMON SUBROUTINE USED BY INTERRUPT
4  IAND ENTRY ROUTINES
5
6  000252 017704 177446      #RCC6,R4
7  000256 031014          #R0,R4
8  000260 001402          IS
9  000262 012600          (SP)+,R0
10 000264 000717          RCHOME
11 000266 014705 177516  (R4)+, -(SP)
12 000272 012546          #SP
13 000274 006316          #SP
14 000276 004316          #SP
15 000300 004316          (R0)+,R4
16 000302 052014          -(R4),-(R4)
17 000304 024444          (SP)+,R4
18 000306 012614          YST
19 000310 005725          RTS
20 000312 000200
21
22
23
24 000314 000000G
25
26 000316 000001
27
28 000001
29
30
31
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```

```

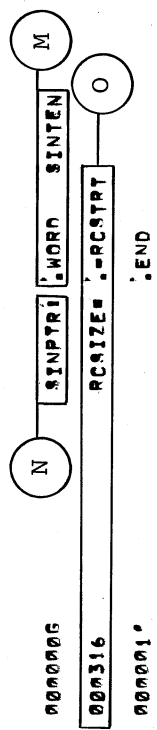
IPT TO DSK CNTRL & STATUS REG
IPTLL IN PROGRESS
INQ
IPOP R0
IPTNTS FILL OF BLK WITH 0'S
IPTR TO CURRENT QUEUE ENTRY
IGET BLOCK NUMBER
ICALCULATE DISK ADDRESS FOR RCDA
UNIT, TRACK# + SECTOR ADDRESS)
I142*8+2561
IINHIB CURR. ADR INC (IF NEEDED)
IPOINT TO RCDA
ISPT DISK ADR FOR TRANSFER
ISGNORE UNIT #
IPRTURN TO CALLER

```

```

IMONTOR ENTRY ADDR.
ISIZE OF HANDLER

```



SYMBOL TABLE

ABORT = 000400	DPYS = 000000 G	DSSYS = 000000 G	DTSYS = 000000 G	DXSYS = 000000 G
ADERR = 000001	INUCA = 001000	MONLOW = 000054	OFFSET = 000270	PC = X0000007
PR5 = 000240	PRY = 000340	RCABRT 000142R	RCCA = 177452	RCCOMN 000256R
RCCOM1 = 000252R	RCCOE = 000010R	RCCS = 177446	RCCA = 177442	RCCD = 177456
RCDR = 177444	RCFILL = 000152R	RCHOME 000124R	RCNT = 000068R	RCLA = 177440
RCLQE = 000006R	RCMN = 177454	RCRETR 000020R	RCSTRT = 000000R	RCSTRT = 000000R
RCSYS = 000006R	RCNRY = 000010	RCVEC = 000010	RCWC = 177450	RD = 000105
RETRY = 000016R	RFSYS = 000000 G	RKSYS = 000000 G	RYER = 060000	RD = X0000000
R1 = X0000001	R2 = X0000002	R3 = X0000003	R4 = X0000004	R5 = X0000005
SP = X0000006	WR =	ZERO = 000024R	SINPTR 000314R	SINTEN = ***** G
P..V2 = 000001				
P..ABS. = 000000				
SYSHND 000316				
ERRORS DETECTED: 0				
FREE CORE: 15627, WORDS				

RC.LP3/NITM/CIRC



BOOT V02R-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 A-APR-75 11149104  
 TABLE OF CONTENTS

2-	2	MACROS, GLOBALS
3-	1	ASPECT
7-	29	BOOTSTRAP I/O DRIVER - RC11
10-	1	BOOTSTRAP CORE DETERMINATION
11-	1	READ MONITOR, LOOKUP HANDLERS
12-	1	RELOCATION LIST

BOOT V02R-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 A-APR-75 11149104 PAGE 1

```

1  SRCSYS=1
2  ;TITLE  BOOT V02R-01    RT-11 BOOTSTRAP
3  ; RT-11 BOOTSTRAP
4  ;
5  ; DEC-11-ORSTA-D
6  ; COPYRIGHT (C) 1975
7  ;
8  ; DIGITAL EQUIPMENT CORPORATION
9  ; MAYNARD, MASSACHUSETTS 01754
10 ;
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24 ; OR RELIABILITY OF ITS SOFTWARE ON EQUIPMENT
25 ; WHICH IS NOT SUPPLIED BY DIGITAL.
26 ;
27 ;
    
```

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.SRTTL MACROS, GLOBALS
.MCALL .:V1..
.V1..
.MCALL .:EXIT, .:LOOKUP, .:PRINT, .:SAVESTATUS

!*****
! CONDITIONAL ASSEMBLY OF ROOT FOR SINGLE USER OR RF SYSTEM
! DEFAULT TO SINGLE USER
!IF NDF BF
BF=0

! GLOBAL REFERENCES TO MONITOR
!GLOBAL SOVREC, SENTRY, SINPTR, SKMLOC, SMONBL, SPNAME, SBL0T
!GLOBAL SSWPRL, SUSRLC, SPNAMO
!GLOBAL RSTRING, CORPTR, DKASSG, FILLER, HWFPS, HWDSPS, KMLOC
!GLOBAL KMON, KMONSZ, KWILLS, MAPOFF, GCOMP, RT1SZ
!GLOBAL RTLEN, RTSIZE, SWAPSZ, SYENDO, SYNCH, SYASSG
!GLOBAL SYSLOW, TTIBUF, YTRBUF, USRL0C, USRSZ, MAXSYH

!GLOBAL RELLST

! FOLLOWING ARE GLOBALS FOR EITHER BF OR SU SYSTEM, BUT NOT BOTH
!IF NE BF
!GLOBAL RCNTXT, RKGN1, RKGN2, RKGN3, CNTXT, FUDGE1, FUDGE2
!GLOBAL MSGENT, RMONSP, SWPTR, SWPTR, TTIUSR, TTOUSR, .SCRN
!IF
!GLOBAL AVAIL, I.CSW, PPRADD, PPRIGN, MONLOC, TRAPLC, TRAPER
.ENDC

PERM = 2000
EMBLK = 4000
JSW = 44
SR = 177570

! REGISTER DEFINITIONS!
R0=X0
R1=X1
R2=X2
R3=X3
R4=X4
R5=X5
R6=X6
R7=X7

! MONITOR OFFSET CONSTANTS

```

```

HARDWARE CONFIGURATION WORD
SYSTEM UNIT #
CLOCK STATUS REGISTER
GT40 LOCATION
KEYBOARD STATUS
" " BUFFER
PRINTER STATUS
" " BUFFER

```

```

CONFIG = 300
SYUNIT = 274
LKCS = 177546
GT40 = 172000
TKR = 177560
TPR = 177562
TPR = 177564
TPR = 177566

```

```

000300
000274
177546
172000
177560
177562
177564
177566

```

ROOT V028-01 RT-11 ROOTSTRAP RT-11 MACRO VM02-09 A-APR-75 11:49:04 PAGE 3

```

:STARTL ASECT
:IF NDF SRESYS
:IF NDF SDXSYS
:IF NDF SDPSYS
:IF NDF SDSYS
:IF NDF SRESYS
:IF NDF SDXSYS
SRKSYS=0
PENDC
PENDC
PENDC
PENDC
PENDC

```



IT MUST BE AN RK SYSTEM

```

:IF NDF SRESYS
:IF NDF SDXSYS
:IF NDF SDPSYS
:IF NDF SDSYS
:IF NDF SRESYS
:IF NDF SDXSYS
SRKSYS=0
PENDC
PENDC
PENDC
PENDC
PENDC

```

:STARTL ASECT

1

2

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ASECT

TURN ON SRKSYS IF ALL OTHERS ARE OFF

```

:ASECT
= 0
240
RR

```

IBOOT VALIDATION PATTERN

IBRANCH TO REAL BOOT

ROOT1

```

:IF NDF SDXSYS
:JMP

```

ROOT

ROOT11

ROOT11

ROOT11

ROOT11

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ROOT11

ROOT11

IPUT THE JUMP BOOT IN TRAP VECTOR

ISTART THE BOOTSTRAP

ISTART FUNCTION

ILEMPTY BUFFER

ILEAD SECTOR

ILEAD SECTOR

ILEAD SECTOR

ILEAD SECTOR

ILEAD SECTOR

ILEAD SECTOR

ILEAD SECTOR

ILEAD SECTOR

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```
INITIALIZE BPT AND IOT VECTORS  
ION BPT INTERRUPT TO READS ROUTINE  
IPS SET TO 0  
ION IOT INTERRUPT TO WAIT ROUTINE  
IREAD FROM UNIT 0, SETS WEIRD BUT OK PS  
IWRITE FROM UNIT 1  
I34-52 USEABLE  
UNTRD(R0),RDCMD;SPT READ FUNCTION FOR CORRECT UNIT  
INIT SP WITH NEXT INSTRUCTION  
IAREA TO READ IN NEXT PART OF BOOT  
ISET TRACK NUMBER  
IOUT OF ROOM HERE, GO TO CONTINUATION  
IPAPER TAPE VECTORS  
ISET TO BIG WORD COUNT  
ISET TO ABSOLUTE TRACK 1  
IBRANCH TO CONTINUATION  
IPROGRAMMABLE CLOCK  
IABSOLUTE SECTOR 3 FOR NEXT PART  
ICALL READS SUBROUTINE  
IBRANCH TO CONTINUATION  
ILOTS OF UNUSED VECTORS, (DR-11B?)
```

A-12

ROOT V02R=01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 8-APR-75 11:49:04 PAGE 3+

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```
READS: MOV #RXCS,R4  
MOV R4,R5  
MOV (PC)+,(R5)+  
RDCMD: TOT  
MOV R3,R5  
TOT  
MOV R0,R5  
TOT  
MOV #CSGO+CSEBUIF,R4;LOAD EMPTY BUFFER FUNCTION INTO RXCS  
TOT  
TSTB R4  
RPL RTTRET  
MOV B  
DEC R1  
RGT 48  
CLR R2  
RR 48  
TST R4  
REQ  
RMI  
RTTRET: RTI
```

```
IR4 -> RX STATUS REGISTER  
IR5 WILL POINT TO RX DATA BUFFER  
IINITIATE READ FUNCTION  
IGETS FILLED WITH READ COMMAND  
ICALL WAIT SUBROUTINE  
ILOAD SECTOR NUMBER INTO RXDB  
ICALL WAIT SUBROUTINE  
ILOAD TRACK NUMBER INTO RXDB  
ICALL WAIT SUBROUTINE  
ICALL WAIT SUBROUTINE  
IIS TRANSFER READY UP?  
IBRANCH IF NOT, SECTOR MUST BE LOADED  
IMOVE DATA BYTE TO MEMORY  
ICHECK BYTE COUNT  
IFLOOP AS LONG AS WORD COUNT NOT UP  
IKLUDGE TO SLUFF BUFFER IF SHORT WD CNT  
IFLOOP  
IIS TR, ERR, DONE UP? INT ENR CAN'T BE  
IFLOOP TILL SOMETHING  
ISTART AGAIN IF ERROR  
IRETURN
```

```

81 . = 200
82 CMPB (R3)+,(R3)+
83 RPT
84 CMPB (R3)+,(R3)+
85 RPT
86 RPT #CSUNIT,RDCMD
87 RNE IS
88 CLR R0
89 MOV R0,(PC)+
90 .WORD 0
91 .TRWAIT,#20
92 JMP ROOT
93 .ENDC
94

```

ROOT V02B-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 8-APR-75 11229104 PAGE 4

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```

; FOLLOWING ARE THE BOOTSTRAP I/O DRIVERS FOR EACH VALID
; SYSTEM DEVICE.
; CALLING SEQUENCE:
; R0 = PHYSICAL BLOCK TO READ/WRITE
; R1 = WORD COUNT
; R2 = BUFFER ADDRESS
; R3,R4,R5 ARE AVAILABLE AND MAY BE DESTROYED BY THE DRIVER
; THE DRIVER MUST GO TO RIGERR IF A FATAL I/O ERROR OCCURS.
; IT MUST ALSO INVOKE THE MACRO SYSDEV
; MACRO SYSDEV NAME,VECTOR
; GLOBL NAME,INT, NAME,SIZE
; SYMNAME = 0
; IRPC X,<NAME>
; SYMNAME = <SYMNAME>+X-1000+50
; ENDR
; SYVEC = VECTOR
; = SYVEC
; = SYMNAME,NAME,INT,340
; = SYSIZE
; =WORD NAME,SIZE
; = 402
; SYMTO = VECTOR / 20
; SYMITS = *B11000000
; RPT <VECTOR & 17> / 4
; SYMITS = SYMITS / 4
; ENDR
; ENDM
; SYSDEV

```

```

;SECTOR 2 OF RX BOOT
;BUMP TO SECTOR 5
;CALL READS SUBROUTINE
;BUMP TO SECTOR 7
;CALL READS SUBROUTINE
;CHECK UNIT ID
;BRANCH IF BOOTING UNIT 1, R0=1
;SET TO UNIT 0
;SAVE UNIT BOOTED FROM FOR LATER
;SAVE THE UNIT HERE
;LETS HANDLE ERRORS DIFFERENTLY
;NOW WE ARE READY TO DO THE REAL BOOT

```

```

1  .IF DF SDSYS IRS SYSTEM
2  .SRTTL BOOTSTRAP T/O DRIVER = RS11
3
4  / RS11 DISK HANDLER
5
6  .IF DF MBUSSC
7  SYSDEV DS,150
8  RSCS2 = 176310
9  .ENDC
10 .IF NDP MBUSSC
11 SYSDEV DS,204
12 RSCS2=172050
13 .ENDC
14
15 READ1
16 MOV R0,R4
17 MOV #RSCS2,R5
18 MOV (R5),-(SP)
19 MOV #40,(R5)
20 RIT #2,16(R5)
21 RNE 15
22 ASL R4
23 RLC #57,(SP)
24 MOV (SP)+,(R5)
25 MOV R4,=(R5)
26 MOV R2,=(R5)
27 MOV R1,=(R5)
28 NEG @R5
29 MOV #71,=(R5)
30 RIT #100200,@R5
31 BEQ 28
32 RMI R1DERR
33 RTS PC
34
35 .ENDC

```

ICOPY BLOCK NUMBER  
 IPOINT TO REGISTERS  
 ISAVE UNIT #  
 ICONTROLLER CLEAR  
 IWHAT IS IT?  
 IIT'S AN R504  
 IIT'S AN R503  
 ICONVERT TO TRACK/SECTOR  
 ISTRIP TO UNIT BITS  
 ISET UNIT  
 ISET BLOCK  
 IGO, READ, NO INTERRUPT  
 IWAIT FOR DONE OR ERROR  
 IBOOT ERROR

ASECT

```

1  / IF DF SDSYS          /CONDITIONAL FOR RP11 DISK
2  .SRTTL BOOTSTRAP I/O DRIVER - RP11
3
4
5  / RP11 DISK DRIVER
6
7  SYSDEV DP,254
8  RPS= 176714
9  RPS= 176710
10 RPA= 176724
11 CS,GO= 000001
12 CS,RO= 000004
13 CS,DRV= 003400
14 DS,ATT= 000377
15
16 READ:  MOV R0,R3
17        JSR R2,DIV
18        .WORD 10,
19        MOV R4,=(SP)
20        MOV R5,R3
21        JSR R2,DIV
22        .WORD 20,
23        SWAB R4
24        RIS (SP)+,R4
25        MOV #RPA,R3
26        MOV R5,=(R3)
27        MOV R2,=(R3)
28        MOV R1,=(R3)
29        NEG @R3
30        RIC #CS,DRV)+,(R3) /CLEAR ALL BUT UNIT #
31        RIS #CS,RO+CS,GO,@R3 / AND START READ
32        TSTB @R3
33        RPL 18
34        TST @R3
35        RMT @DERR
36        MOVB #DS,ATT,@RPS
37        CLRB @RPS
38        RTS PC
39
40 / DIVIDE ROUTINE FOR RP HANDLER.
41 / R5 = R3 / @R2, REMAINDER IN R4
42
43 DIVE  CLR R5
44        CLR R4
45        TST R3
46        REG 48
47        COM R5
48        ROL R3
49
50 /DEVICE IS RP. IT VECTORS TO 254
51 /RP11 DEVICE CONTROL REG
52 /RPO3 STATUS REG
53 /RPO3 DISK ADRS REGISTER
54 /GO BIT IN CONTROL & STATUS
55 /READ FUNCTION CODE
56 /UNIT SELECT BITS
57 /UNIT ATTN BITS
58
59 IR3 = BLOCK #
60 IGET SECTOR NUMBER
61 IBY DIVIDING BY 10
62 ISAVE SECTOR
63 ISET NEW DIVIDEND
64 IAND COMPUTE CYL & TRACK
65 IBY DIVIDING BY 20
66 IPOSITION TRACK IN HIGH BYTE
67 IAND INSTALL SECTOR
68 IR3 -> DISK ADRS REG
69 ISET TRACK & SECTOR
70 IAND CYLINDER
71 IAND BUS ADDRESS
72 IAND WORD COUNT
73 IMAKE NEGATIVE
74 I#C<CS,DRV)+,(R3) /CLEAR ALL BUT UNIT #
75 IWAIT UNTIL TRANSFER COMPLETE
76
77 IANY ERRORS?
78 IYES
79 I#CLEAR UNIT ATTN FOR BOTH
80 IOLD & ECO'D CONTROLLERS
81 IELSE JUST RETURN
82
83 IQUOT. = 0
84 IREM. = 0
85 IIS DIVIDEND 0?
86 IYES = JUST RETURN
87 IQUOT. = -1 & SET CARRY
88 INORMALIZE

```

```

49      RCC          18
50      ROL          R4
51      CMP          R4,0R2
52      RLO          38
53      SUB          0R2,R4
54      ROL          R5
55      RSL          R3
56      RNE          28
57      COM          R5

      /SHIFT & SUBTRACT

      /FIX QUOTIENT

```

```

BOOT V02R-01      RT-11 BOOTSTRAP RT-11 MACRO VM02-09      8-APR-75 11:49:04 PAGE 6+
ASECT

```

```

58      TST          (R2)+
59      RTS          R2
60
61      .ENDC

```

```

BOOT V02B-01      RT-11 BOOTSTRAP RT-11 MACRO VM02-09      8-APR-75 11:49:04 PAGE 7
ASECT

```

```

1      .IP DF      SRKSYS:SRFSYS:SRCSYS      BB
2
3      /IFDF      SRFSYS
4      .SRTTL    BOOTSTRAP I/O DRIVER - RF11
5
6      / RF11    DISK HANDLER
7
8      SYSDEV    RF,204
9      RFCS      = 177460
10     RFMC      = 177462
11     RFMA      = 177464
12     RFDA      = 177466
13     RFDE      = 177470
14     RFDB      = 177472
15
16     READI     #RFDA,R3
17     MOV       R0,R5
18     SWAB     R5
19     MOV       R5,R4
20     CLR      R5
21     MOV       R5,(R3)+
22     BIC      #17740,R4
23     MOV       R4,(R3)
24     TST      -(R3)
25
        /DEVICE IS RF. IT VECTORS TO 204.
        /CONTROL & STATUS REGISTER
        /WORD COUNT
        /MEMORY ADDRESS
        /DISK ADDRESS
        /DISK ADDRESS EXTENSION
        /DATA BUFFER
        /POINT TO DISK ADDRESS
        /COPY BLOCK NUMBER
        /MULTIPLY BY 256 TO GET WORD # ON DISK
        /SAVE HIGH ORDER DISK ADDRESS
        /MAKE DA AN EVEN BLOCK NUMBER
        /PUT LOW ORDER ADDRESS IN CONTROLLER
        /ISOLATE HIGH ORDER ADDRESS
        /PUT IT IN CONTROLLER
        /RESET POINTER

```



CC — .ENDC  
 .IFDF SRC3VS  
 #CONDITIONAL FOR RC (R964) DISK

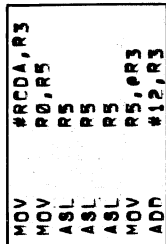
.SRTTL BOOTSTRAP I/O DRIVER - RC11

DD — / RC11 DISK HANDLER  
 SYSDEV RC,210  
 / DEVICE IS RC. IT VECTORS TO 210

EE — / RC11 CONTROL REGISTERS

```

177440 RCLA= 177440 / LOOK AHEAD REGISTER
177442 RCDAA= 177442 / DISK ADDRESS REGISTER
177444 RCER= 177444 / DISK ERROR STATUS REGISTER
177446 RCCS= 177446 / DISK CONTROL & STATUS REGISTER
177450 RCWC= 177450 / WORD COUNT REGISTER
177452 RCCA= 177452 / CURRENT ADDRESS REGISTER
177454 RCMN= 177454 / MAINTENANCE REGISTER
177456 RCDB= 177456 / DATA BUFFER REGISTER
  
```



EE — READ1

```

45 000402 012703 177442 / PT TO DISK ADR REGISTER
46 000406 010005 / GET BLOCK NUMBER
47 000410 006305 / CALCULATE DISK ADR FOR RCDAA
48 000412 006305 / UNIT, TRACK # & SECTOR ADR
49 000414 006305 / 132 * 8=256J
50 000416 010513 / FIND PROPER DISK ADR
51 000420 062703 000012 / PT TO CURRENT ADR REG + 12
  
```

EF — .ENDC  
 .IFDF  
 .SRTTL BOOTSTRAP I/O DRIVER - RK05  
 / RK05 DISK HANDLER

BOOT V02B-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 8-APR-75 1149104 PAGE 7 +  
 BOOTSTRAP I/O DRIVER - RC11

```

58 SYSDEV RK,220 / DEVICE IS RK. IT VECTORS TO 220
59 RKDA = 177412 / RK DISK ADDRESS
60 READ1 MOV #14,R3 / PHYSICAL BLOCK TO RK DISK ADD.
61 RR 28 / ENTER BLOCK # COMPUTATION
62 ADD #20,R3 / CONVERT DISK ADDRESS
63 SUB #14,R0 / 
64 RPL 15 / 
65 ADD R3,R0 / RM HAS DISK ADDRESS
66 MOV #RCDAA,R3 / POINT TO HARDWARE DISK ADDR REGISTER
67 PIC #1777,R3 / LEAVE THE UNIT NUMBER
68 BIS R0,(R3) / PUT DISK ADDRESS INTO CONTROLLER
  
```

```

69 
70 
  
```

```

71
72
73
74 000424 010243
75 000426 010143
76 000430 005413
77 000432 012743 000005
78 000436 105713
79 000440 105376
80 000442 005713
81 000444 100401
82 000446 000207
83
84

```

GG

```

.ENDC
; THIS CODE IS COMMON TO RK05,RC11 AND RF11 HANDLERS
;BUFFER ADD.
;WORD COUNT
;(NEGATIVE)
;START DISK READ
;WAIT UNTIL COMPLETE

SS:
RPL
TST
RMI
RTS

;ANY ERRORS?
;HARD HALT ON ERROR
PC

```

BOOT V02B-01 RT-11 BOOTSTRAP I/O DRIVER - RC11  
 8-APR-75 11149104 PAGE 8

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```

.IF DF SDTSYS
.SRCTL BOOTSTRAP I/O DRIVER = DECTAPE
; DECTAPE BOOTSTRAP HANDLER
SYSDEV DT,214
TCCH = 177342
TCDT = 177350
TCST = 177340

READ: MOV #TCCH,R4
MOV #TCDT,R3
MOV R0,R5
SUB #2,R5
MOV #4003,R4
RIT #100200,R4
REG 28
RMI DTERR
CMP R5,R3
BLT DTSRCH
MOV #3,R4
RIT #100200,R4
REG 48
RMI DTERR
CMP R0,R3
RGT DTFWRD
BLT DTSRCH
MOV R2,-(R3)
NEG R1
MOV R1,-(R3)

```

```

;DEVICE IS DT. IT VECTORS TO 214.
;COMMAND REGISTER
;DATA REGISTER
;STATUS REGISTER

IR4 => COMMAND REG
IR3 => DATA REG
;COPY BLOCK NUMBER
;SEARCH FOR 2 EARLIER
;REVERSE,RNUM
;WAIT TILL BLOCK FOUND

;IS IT THE DESIRED BLOCK
;NO,CONTINUE SEARCHING
;SEARCH FORWARD (RNUM)
;WAIT

;DESIRED BLOCK
;NO-SEARCH FORWARD
;NO-SEARCH REVERSE
;BUFFER ADDRESS
;WORD COUNT

```

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```

D74I      MOV      #5,R4
          BIT      #100200,R4
          BEQ      DT4
          ANI      #0ERR
          CLR      #R4
          RTS      PC
          DTERR:  TST      #TCS1
          RPL      #0ERR
          BIT      #4000,R4
          ANE      DTFWRD
          BR       DTSRCH
          IREAD
          IWAIT FOR COMPLETION
          IREAD ERROR
          ISTOP DT
          IWHAT KIND OF ERROR ?
          INOT END ZONE
          IREVERSE?
          ITHEN GO SEARCH FORWARD
          ELSE SEARCH REVERSE
    
```

.ENDC

BOOT V020-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 8-APR-75 11149104 PAGE 9  
 BOOTSTRAP I/O DRIVER - RC11

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```

          IF DF SDXSYS
          :SRTTL BOOTSTRAP T/O DRIVER = FLOPPY
          SYSDEV DX,264
          READI:  R0
          ASL     R0
          ASL     R0
          ASL     R1
          MOV     R0,=(SP)
          MOV     R0,R3
          MOV     R0,R4
          CLR     R0
          BR      39
          SUR     #23,,R3
          INC     R0
          SUB     #26,,R4
          RPL     28
          CMP     #-14,,R4
          ROL     R3
          SUR     #26,,R3
          BPL     48
          ADD     #27,,R3
          BPT     (SP)+,R0
          INC     R0
          TST     R1
          RGT     18
          RTS     PC
          TRWAIT: TST     #R4
          RPL     RTIRET
          :FLOPPY VECTORS THROUGH 264
          IFCONVERT BLOCK TO LOGICAL SECTOR
          ILSN=BLOCK*4
          IMAKE WORD COUNT BYTE COUNT
          ISAVE LSN FOR LATER
          IWE NEED 2 COPIES OF LSN FOR MAPPER
          IINIT FOR TRACK QUOTIENT
          IJUMP INTO DIVIDE LOOP
          IPERFORM MAGIC TRACK DISPLACEMENT
          IBUMP QUOTIENT, STARTS AT TRACK 1
          ITRACK=INTEGER(LSN/26)
          ILOOP = R4REM(LSN/26)-26
          ISET C IF SECTOR MAPS TO 1-13
          IPERFORM 211 INTERLEAVE
          IADJUST SECTOR INTO RANGE =1.-26
          I(DIVIDE FOR REMAINDER ONLY)
          INOW PUT SECTOR INTO RANGE 1-26
          ICALL READS SUBROUTINE
          IGET THE LSN AGAIN
          ISET UP FOR NEXT LSN
          IWHATS LEFT IN THE WORD COUNT
          IBRANCH TO TRANSFER ANOTHER SECTOR
          IRETURN
          INEW WAIT SUBROUTINE, PRINTS ERRORS
          IRETURN FROM INTERRUPT
    
```

\*\*\*\*\* THIS MUST FALL INTO BIODERR \*\*\*\*\*  
.ENDC

BIODERR: JSR R0,REPORT ;SAY THAT WE GOT ERROR  
;ASCIZ '<15><12>\?R=I/O ERROR\<12>

```
077  
111  
040  
122  
012  
000450 004067 000024  
000454 015 012  
102 055 111  
000462 057 117  
000465 105 122  
000470 117 122  
000473 000
```

.EVEN

40

BOOT V02R-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 8-APR-75 11149104 PAGE 10  
BOOTSTRAP CORE DETERMINATION

```
1  
2  
3 000474 112037 177566  
4 000500 105737 177564  
5 000504 100375  
6 000506 105710  
7 000510 001371  
8 000512 000005  
9 000514 000000  
10 000516 000776  
11 000520 012706 010000  
12 000524 012700 000002  
13 000530 012701 000400  
14 000534 012702 001000  
15 000540 004767 177636  
16  
17  
18 000544 012703 000004  
19 000550 011305  
20 000552 012723 000620  
21 000556 000013  
22  
23  
24  
25  
26  
27  
28  
29  
30  
31
```

.SMTTL BOOTSTRAP CORE DETERMINATION

```
REPORT1: MOV8 (R0)+,#PTR ;PUT ANOTHER CHARACTER OUT  
REPORT1: TSTB #PTR ;WAIT FOR TYPER READY  
RPL ;  
TSTB #R0 ;ANYTHING MORE ?  
RNE REPORT1 ;YES, LOOP  
RESET ;STOP ALL DEVICES  
HALT ;KEEP HIM FROM CONTINUING  
RR ;  
MOV #1000,R3 ;SET STACK POINTER  
MOV #2,R0 ;READ IN SECOND PART OF BOOT  
MOV #<ROOTSZ-1>+*400,R1 ;EVERY BLOCK BUT THE ONE WE ARE IN  
MOV #1000,R2 ;INTO LOCATION 1000  
JSR PC,READ  
;BOOTSTRAP BLOCK 0 TOO BIG  
;POINT TO TRAP LOCATIONS  
MOV #4,R3 ;SAVE TRAP LOC  
MOV #R3,R5  
MOV #NRM,(R3)+ ;SET TRAP FOR NON EXISTENT MEMORY  
CLR #R3
```

.ITF GT .=-1000,  
MOV #4,R3  
MOV #R3,R5  
MOV #NRM,(R3)+  
CLR #R3

```
); THIS BOOTSTRAP CAN SIMULATE ANY SIZE PDR-11.  
); IF LOCATION 'FIDDLE' IS A HALT, THE CPU WILL STOP DURING THE BOOT.  
); ON CONTINUE, THE TOP 5 BYTS OF THE SWITCH REGISTER ARE USED TO  
); SET THE TOP OF AVAILABLE CORE AS A MULTIPLE OF 1K.  
); IF THE BR IS >= 160000 OR IF FIDDLE IS A BR 18,  
); THE BOOTSTRAP WILL DO A NORMAL CORE DETERMINATION.  
);  
);  
); FIDDLE! BR 18 ;CHANGE TO HALT FOR FIDDLING
```

```

32 000562 013702 177570
33 000566 042702 003777
34 000572 020227 160000
35 000576 103410
36 000600 005002
37 000602 062702 004000
38 000606 020227 160000
39 000612 001402
40 000614 005712
41 000616 000771
42 000620 012723 001476
43 000624 011367 177160
44 000630 012701 001604
45 000634 010100
46 000636 005737 177546
47 000642 052140
48 000644 005737 172000
49 000650 052110
50 000652 170000
51 000654 052110
52 000656 010523
53 000660 012723 000340
54 000664 010523
55 000666 012723 000341
56 000672 162702 000000G
57 000676 062702 000000G

```

```

#SR,R2
#377,R2
R2,#160000
NXM
R2
#4000,R2
R2,#160000
NXM
#R2
28
#BCLR, -(R3)
#R3,10
#TLIST,R1
R1,R0
#LKCS
(R1)+, -(R0)
#0740
(R1)+, #R0
(R1)+, #R0
R5,(R3)+
#340,(R3)+
R5,(R3)+
#341,(R3)+
#RTSIZE,R2
#FILLER,R2

```

```

MOV
R1C
CMP
RLO
CLR
ADD
CMP
REG
TST
BR
MOV
MOV
MOV
TST
BIS
TST
BIS
CFCC
RIS
MOV
MOV
MOV
MOV
SUB
ADD

```

```

GET SWITCH VALUE
ISOLATE TOP 5 BITS (1K INCREMENTS)
SHOULD WE DO NORMAL CHECK ?
INO, USE THE SR VALUE
LOOK FOR TOP OF CORE
MOVE TO NEXT 1K BANK
REACHED 28K YET ?
YES, DO A 28K SYSTEM
INO, SEE IF THIS LOCATION EXISTS
KEEP GOING IF WE DIDN'T TRAP
NONMEMORY TRAPS HERE
BAD INSTRUCTIONS TRAP HERE
BITS FOR CLEARING ON ERROR TRAPS
CHECK PRESENCE OF CLOCK
ADVANCE LIST
CHECK FOR DISPLAY
CHECK FOR FPU
RESTORE TRAP
ITRAP TO 4 IS PR7, CARRY OFF
RESTORE 10
ITRAP TO 10 IS PR7, CARRY ON
IR2 NOW POINTS TO WHERE WE WANT THE KMON
ABOUT IT AGAINST THE TOP OF CORE

```

ROOT V02B-01 RT-11 ROOTSTRAP RT-11 MACRO VM02-09 8-APR-75 11:49:04 PAGE 10+

```

58 000702 162702
59 000704
60 000706
61 000706 062702 000000G
62 000712 020227 010000
63 000716 103466
64 000720 010246
65 000722 012700 000001
66 000726 006300
67 000730 062700 000004
68 000734 012701 001000
69 000740 012702 001634
70 000744 004767 177432
71 000750 012701 001644
72 000754 012100
73 000756 010102

```

```

RECOVER UNUSED CORE FROM SY;
SYSTEM HANDLER SIZE PUT HERE
THIS WAY BECAUSE NO GLOBAL ARITH.
IS IT JUST TOO TINY ?
YES
PUT LOAD ADDRESS ON STACK
INOW READ FIRST DIRECTORY BLOCK
IDIRECTORY STARTS AT 6

```

```

(PC)+,R2
#MAXSYH,R2
R2,#10000
TOOSML
R2, -(SP)
#1,R0
R0
#4,R0
#1000,R1
#BUFFER,R2
PC,READ
#BUFFER+10,R1
(R1)+,R0
R1,R2

```

```

SUR
SYSIZE = .
. = .+2
ADD
CMP
RLO
MOV
MOV
ASL
ADD
MOV
MOV
JSR
MOV
MOV
MONFI

```

```

IR2 NOW POINTS TO WHERE WE WANT THE KMON
ABOUT IT AGAINST THE TOP OF CORE

```

```

74 000760 032721 002000
75 000764 001411
76 000766 162721
77 000770 051646
78 000772 162721
79 000774 035562
80 000776 162711
81 001000 075273
82 001002 001002
83 001004 054141
84 001006 001447
85 001010 032712 004000
86 001014 001010
87 001016 066200 000010
88 001022 062702 000016
89 001026 066702 000060
90 001032 010201
91 001034 000750
92 001036 016700 000574
93 001042 001331
94 001044 004067 177450
95 001050 001050 015 012
001053 102 055
001056 117 040
001061 117 116
001064 124 122
001067 123 131
001072 012 000
96
97 001074 004067 177400
98 001100 015 012
001103 102 055
001106 117 124
001111 105 116
001114 125 107
001117 040 103
001122 122 105
001125 000
99

BIT #PERM,(R1)+
REG 1S
SUB (PC)+,(R1)+
.RAD50 /MON/
SUB (PC)+,(R1)+
.RAD50 /ITR/
SUB (PC)+,(R1)
.RAD50 /SYS/
RNE 1S
RIS -(R1),=(R1)
REG MONFND
BIT #ENDBLK,(R2)
RNE 2S
ADD 10(R2),R0
ADD #16,R2
ADD BUFFER+6,R2
MOV R2,R1
BR MONF
MOV BUFFER+2,R0
RNE DFND
JSR R0,REPORT
.ASCIIZ <15><12>\?R=NO MONTR.SYS\<12>

IIS IT A PERMANENT FILE?
INO. WE ARE TRYING TO FIND THE
IFILE MONTR.SYS, AS THAT IS
ITHE CURRENT MONITOR.

IPLAST WAS NOT .SYS EXTENSION
IBOTH MUST BE 0
IFOUND THE MONITOR
IIS THIS ALL IN SEGMENT?
IYES. READ NEXT, IF ANY.
IINCREASE START BLOCK
IGET TO NEXT ENTRY
IPOINT R1 TO NEXT
ISEE IF NEXT IS AVAILABLE
IYES. CONTINUE
IHE AIN'T GOT A MONITOR

```

```

1 001126 011602
2 001130 062700
3 001134 010046
4 001136 012701
5 001142 166701
6 001146 062701
7 001152 006201
8 001154 003401
9 001156 062701
10 001162 062700
11 001166 004767
12 001172 012700
13 001176 012604
14 001200 012604
15 001202 010164
16 001206 062701
17 001212 062701
18 001216 010164
19 001222 060430
20 001226 103774
21 001230 012005
22 001234 060405
23 001242 060415
24 001244 012005
25 001246 001374
26 001250 013700
27 001254 000054
28 001258 000000
29 001262 000000
30 001266 000000
31 001270 000000
32 001274 000000
33 001278 000000
34 001282 000000
35 001286 000000
36 001290 000000
37 001294 000000
38 001298 000000
39 001302 000000
40 001306 000000
41 001310 000000
42 001314 000000
43 001318 000000
44 001322 000000
45 001326 000000
46 001330 000000
47 001334 000000
48 001338 000000

```

```

:SRITL READ MONITOR, LOOKUP HANDLERS
MONFND: MOV
ADD #BOOTSZ,R0
MOV #MAXSYH,R1
SUB SYSIZE,R1
ADD #FILLER,R1
ASR R1
NEG R1
ADD #RTLEN,R1
ADD #SWAPSZ,R0
JSR PC,READ
MOV #RELLST,R0
MOV (SP)+,R1
MOV (SP)+,R4
SUB #KMON,R4
MOV R1,$SWPBL(R4)
ADD #SWAPSZ,R1
MOV R1,$MONBL(R4)
ADD R4,#(R0)+
CMP R0,#RELLST
BLO 18
MOV (R0)+,R5
ADD R4,R5
ADD R4,#R5
MOV (R0)+,R5
RNE 28
MOV #54,R0
:IF DF SRKSYS!SDXSYS!SDPSYS!SDSSYS !THE RK,RX,RP,RJS03/4 CAN BOOT FROM ANY UNIT
:IF DF SRKSYS
MOV #RCKDA,R1
ROL R1
ROL R1
ROL R1
ROL R1
PIC #C7,R1
:ENDC
:IF DF SDSSYS
MOV #RSCS2,R1
PIC #C7,R1
:ENDC
:IF DF SDPSYS
MOV #RPCS,R1
PIC #C4C6.DRV>,R1
SWAB R1

```

```

:RECALL LOAD LOCATION
:BUMP R0 OVER BOOT RECORDS
:SAVE SWAP BLOCK POINTER
:DO GLOBAL ARITHMETIC HERE
:R1 = MAXSYH-SYSIZE (BYTES)
:ADD AMOUNT OF EXTRA STUFF
:(WORDS)
:(TO SUBTRACT)
:LENGTH TO LOAD (WORDS)
:POINT TO BLOCK WITH KMON
:READ THE MONITOR INTO PLACE
:POINT TO LIST OF THINGS TO RELOCATE
:R1 = SWAP BLOCK NUMBER
:R4 => KMON IN CORE
:SUBTRACT LOCATION KMON WAS LINKED TO
:R4 = BIAS. SET UP SWAP BLOCK #

:SET USR BLOCK #
:RELOCATE A POINTER IN THE ASECT
:DONE YET ?
:END
:GET POINTER TO THING IN MONITOR
:BIAS THE POINTER
:INOW RELOCATE THE WORD
:GET NEXT POINTER
:POINT TO MONITOR

```

```

:EXTRACT IT
:DP SRKSYS
:CODE FOR RJS03/4
:UNIT # INTO R1
:STRIP TO 3 BITS

:RPI1
:GET CONTROLLER STATUS REG INTO R1
:STRIP TO UNIT NUMBER
:UNIT # INTO BITS 2=0

```

```

49          .ENDC
50          .IF DF
51          MOV STUNIT,R1
52          .ENDC
53          ADD R1,DKAS8G(R0)
54          ADD R1,SYAS8G(R0)
55          MOV R1,SYUNIT+1(R0)
56          .IF DF
57          .ENDC
58          .IF DF
59          MOV STUNIT,R1
60          .ENDC
61          ADD R1,DKAS8G(R0)
62          ADD R1,SYAS8G(R0)
63          MOV R1,SYUNIT+1(R0)
64          .IF DF
65          .ENDC
66          .IF DF
67          .ENDC
68          .IF DF
69          .ENDC
70          .IF DF
71          .ENDC
72          .IF DF
73          .ENDC
74          .IF DF
75          .ENDC
76          .IF DF
77          .ENDC
78          .IF DF
79          .ENDC
80          .IF DF
81          .ENDC
82          .IF DF
83          .ENDC
84          .IF DF
85          .ENDC
86          .IF DF
87          .ENDC
88          .IF DF
89          .ENDC
90          .IF DF
91          .ENDC
92          .IF DF

```

```

BOOT VM02R-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 A-APR-75 11149104 PAGE 11+
READ MONITOR, LOOKUP HANDLERS

```

```

58          .ENDC
59          .ENDC
60          .ENDC
61          .ENDC
62          .ENDC
63          .ENDC
64          .ENDC
65          .ENDC
66          .ENDC
67          .ENDC
68          .ENDC
69          .ENDC
70          .ENDC
71          .ENDC
72          .ENDC
73          .ENDC
74          .ENDC
75          .ENDC
76          .ENDC
77          .ENDC
78          .ENDC
79          .ENDC
80          .ENDC
81          .ENDC
82          .ENDC
83          .ENDC
84          .ENDC
85          .ENDC
86          .ENDC
87          .ENDC
88          .ENDC
89          .ENDC
90          .ENDC
91          .ENDC
92          .ENDC

```

```

R0,R1,R2,R3,R4,R5,R6,R7,R8,R9
R10,R11,R12,R13,R14,R15,R16,R17,R18,R19
R20,R21,R22,R23,R24,R25,R26,R27,R28,R29
R30,R31,R32,R33,R34,R35,R36,R37,R38,R39
R40,R41,R42,R43,R44,R45,R46,R47,R48,R49
R50,R51,R52,R53,R54,R55,R56,R57,R58,R59
R60,R61,R62,R63,R64,R65,R66,R67,R68,R69
R70,R71,R72,R73,R74,R75,R76,R77,R78,R79
R80,R81,R82,R83,R84,R85,R86,R87,R88,R89
R90,R91,R92,R93,R94,R95,R96,R97,R98,R99

```



```

93 001424 001360 BNE 6S
94 001426 012737 100000 000044 MOV #100000,#JSW
95 001434 005000 .PRINT #BSTRNG
96 001442 012720 CLR R0
97 001444 012720 MOV (PC)+,(R0)+
98 001446 040000 BIC R0,R0
99 001450 012720 MOV (PC)+,(R0)+
100 001452 .EXIT
101 001454 032767 000000 000120 BIT #KW11LS,RCNFG
102 001462 001403 REC 10S
103 001464 012737 000100 177546 MOV #100,#LKCS
104 001472 005000 CLR R0
105 001474 .EXIT
106 001476 .DSABL LSR
107
108
109 001476 005011 ACLR: CLR 0R1
110 001500 000002 RTI

```

```

I PRINT BOOT HEADER
I AND IF HE HAS A CLOCK
I WE TURN IT
I ON
I TRAP MEANS THIS CONFIGURATION NYET
I UNTRAP

```

BOOT V028-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 8-APR-75 11149104 PAGE 12

```

RELOCATION LIST
1 001502 000004 .SRTTL RELOCATION LIST
2 001504 000010 RELST: 4
3 001506 000030 10
4 001510 000054 30
5 001512 000060 54
6 001514 000064 60
7 001516 000100 64
8 001520 000210 100
9 001522 000244 SYVEC
10 001524 000000 244
11 001526 000000 USRLOC
12 001528 000000 SUSRLC
13 001530 000000 GCOMP
14 001532 000000 SKMLC
15 001534 000000 TTIBUF
16 001536 000020 TTIBUF+2
17 001540 000060 TTIBUF+6
18 001542 000100 TTIBUF+10
19 001544 000000 TTBUFF
20 001546 000040 TTBUFF+4
21 001550 000060 TTBUFF+6
22 001552 000000 SYSLOW
23 001554 000020 CORPTR+2
24 001556 000000 SINPTR
25 001560 000000 SYNCH

```

```

ILLEGAL MEM AND INST. TRAPS
MEMT
ADDRESS OF RMON
TTY VECTORS
CLOCK VECTOR
SYSTEM DEVICE VECTOR
LOCATION OF FPU TRAP
ADDRESS OF USR NOW
QUEUE COMPLETION
ADDRESS OF KMON
TTY RING BUFFER--INPUT
TTY RING BUFFER--OUTPUT
LOWEST USED LOCATION
FREE CORE LIST
POINTER TO SINTEN IN RESIDENT HANDLER
SYNCHRONIZATION ADDRESS

```

```

26 .IF NE DF
27 MSGENT
28 TTUSER
29 TTUSER
30 FUDGE1
31 FUDGE2
32 BKGND1
33 BKGND2
34 BKGND3
35 CNTXT
36 PCNTXT
37 RMONSP
38 SWIPTR
39 SHOPTR
40 .SCRN
41
42 .IFF
43 TRAPLC
44 TRAPER
45 PPPADD
46 PPIGN
47 MONLOC
48 Y.CSW
49 AVAIL
50 .ENDC
51
52 BCFGI
53 TSLIST1
54
55
56
57
58
59
60

```

RELOCATE OR STUFF HERE  
 FLOGS FOR TRAPS TO 4/10  
 IFPP SERVICE FOR MONITOR  
 IWHERE USER WILL SIT  
 ISINGLE USER STUFF HERE  
 IMONITOR FREE Q POINTER  
 IEND OF LIST  
 IBOOT CONFIGURATION WORD-DO NOT MOVE  
 KMI1LS,HWDSPS,HWPUS IBITS IN CONFIG WORD  
 I BLOCK IS THE ARGUMENT AREA FOR AN RT-11 LOOKUP.  
 I BLOCK1 .RAD50 /SY /

ROOT VM02R-01 RT-11 BOOTSTRAP RT-11 MACRO VM02-09 8-APR-75 11149104 PAGE 12+  
 RELOCATION LIST

ADDRESS	WORD	RAD50	BLKW	FILENAME	STATUS
58	001614	000000	000000	FILENAME	GOES HERE
59	001620	075273			
60	001622		5	STATUS	GOES HERE

1 001634 . . . . .  
2 000002 . . . . . + 777 / 1000  
3 002000 . . . . . \* ROOTSZ \* 1000  
4 000001 . . . . . :END

AVAIL = ***** G	BCLR = 001476	RCNFG = 001602	RF = 000000	BI0ERR = 000450
BLOCK = 001612	BOOT = 000520	ROOTSZ = 000002	ROOT1 = 000034	BSTRNG = ***** G
BUFFR = 001634	CBLOK = 001622	CONFIG = 000300	CORPTR = ***** G	DFND = 000726
DKASSG = ***** G	ENDBLK = 000000	FIDDLE = 000500	FILLER = ***** G	FNAM = 001614
FPPADD = ***** G	FPPIGN = ***** G	GT40 = 172000	HWDSPS = ***** G	HMPPUS = ***** G
I.CSW = ***** G	JSW = 000044	KMLOC = ***** G	KMON = ***** G	KMONSZ = ***** G
KHILLS = ***** G	LKCS = 177546	MAPOFF = ***** G	MAXSYH = ***** G	MONF = 000756
MONFND = 001126	MONLOC = ***** G	NXM = 000620	PC = *X000007	PERM = 002000
QCOMP = ***** G	RCCA = 177452	RCCS = 177446	RCDA = 177442	RCDB = 177456
RCER = 177444	RCINT = ***** G	RCLA = 177440	RCMN = 177454	RCSIZE = ***** G
RCUC = 177450	READ = 000402	RELLST = 001502	RELIST = 001524	REPORT = 000500
REFOR1 = 000474	RLEN = ***** G	RTSIZE = ***** G	RT11SZ = ***** G	R0 = *X000000
R1 = *X000001	R2 = *X000002	R3 = *X000003	R4 = *X000004	R5 = *X000005
SP = *X000006	SR = 177570	SWAPSZ = ***** G	SYASSG = ***** G	SYBITO = 000010
SYBITS = 000014	SYENTO = ***** G	SYINDO = ***** G	SYNAME = 070370	SYNCH = ***** G
SYSIZE = 000704	SYSLOW = ***** G	SYUNIT = 000274	SYVEC = 000210	TKR = 177562
TKS = 177560	TO08ML = 001074	TPR = 177566	TPS = 177564	TRAPER = ***** G
TRAPLC = ***** G	TSLIST = 001604	TTIBUF = ***** G	TTIOBUF = ***** G	USRLOC = ***** G
USRZ = ***** G	SDVREC = ***** G	SENTRY = ***** G	SINPTR = ***** G	USKLOC = ***** G
SMONBL = ***** G	SPNAME = ***** G	SPNAMO = ***** G	SRCSYS = 000001	SSLOT = ***** G
SWHBL = ***** G	SUSRLC = ***** G	...V1 = 000001		
. ABS. 002000 000				
000000 001				
ERRORS DETECTED: 0				
FREE CORE: 14985. WORDS				

RCRTSJ,LPI/NITTM/C=RCISYS,BSTRAP

A.3 LP/LS11 DEVICE HANDLER

LP V02-03 25-JUN-74 RT-11 MACRO VM02-10 14-APR-75 10:05:11 PAGE 1

```
1 .TITLE LP V02-03 25-JUN-74
2
3
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25
26
27
28
29
30
31
```

RT-11 LINE PRINTER (LP/LS11) HANDLER

DEC-11-ORTLA-D

PGR/FP/ARC/EF

MARCH 1973/FEBRUARY 1974

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```

1 000000
2 000001
3 000002
4 000003
5 000004
6 000005
7 000006
8 000007
9
10
11
12 177514
13 177516
14 000200
15
16
17 000001
18 000054
19 000270
20 000340
21 000200
22
23
24
25 000015
26 000012
27 000014
28 000011
29
30 000204
31
32

```

```

R0=X0
R1=X1
R2=X2
R3=X3
R4=X4
R5=X5
SP=X6
PC=X7

; LINE PRINTER CONTROL REGISTERS
LPS = 177514 ;LINE PRINTER CONTROL REGISTER
LPR = 177516 ;LINE PRINTER DATA BUFFER
LPVEC = 200 ;LINE PRINTER VECTOR ADDR

;CONSTANTS FOR MONITOR COMMUNICATION
HDFRR = 1 ;HARD ERROR BIT
MONLOW = 54 ;BASE ADDR OF MONITOR
OFFSFT = 270 ;PRINTER TO G MANAGER COMP ENTRY
PR7 = 340
PR4 = 200

; ASCII CONSTANTS
CR = 15
LF = 12
FF = 14
HT = 11
COLSTZ = 132. ;132 COLS
;GLOBAL COLCNT

```

THE FOLLOWING ARE THE PARAMETERS FOR INTERFACE TO THE MONITOR 'SET' COMMAND

1	000000	000000	1	.ASECT	
2	000400	000400	2	.WORD	30. MINIMUM WIDTH
3	000400	000400		.RAD50	/ WIDTH /
4	000400	000400		.WORD	<0.WIDTH-400>/2+40000 JNO 'NO' OPTION, NUMBER REQUIRED
5	000400	000400		NOP	JNO CR => NOP CROPT
6	000400	000400		.WORD	<0.CR-400>/2+100000 ALLOW 'NO'
7	000400	000400		NOP	JNO FORM0 => NOP FF0PT
8	000400	000400		.RAD50	/FORM0 /
9	000400	000400		.WORD	<0.FORM0-400>/2+100000 ALLOW 'NO'
10	000400	000400		RMT	LPERR=ERR0PT+. JNO HANG RY GOING TO ERROR
11	000410	000240		.RAD50	/HANG /
12	000412	000000		.WORD	<0.HANG-400>/2+100000 ALLOW 'NO'
13	000415	100033		NOP	JFOR NO LC, CONVERT LC TO UC
14	000420	000240		.WORD	40
15	000420	000240		.RAD50	/LC /
16	000422	000000		.WORD	<0.LC-400>/2+100000
17	000426	100045		0	END OF LIST
18	000430	100500		N.WIDTH:MOV	R0, COLCNT
19	000432	031066		MOV	R0, RSTC+2
20	000436	100045		CMP	R0, R3
21	000440	000000		RTS	PC
22	000442	045570		0.CR:	MOV (PC)+,R3
23	000446	100052		MOV	BFO RSTC-CROPT+
24	000450	000000		MOV	BFO R3,CROPT
25	000452	000000		RTS	PC
26	000452	010047		0.FORM0:MOV	(PC)+,R3
27	000456	010067		MOV	ALK0-FF0PT+
28	000462	020003		CMP	R3,FF0PT
29	000464	000207		RTS	PC
30	000466	012703		0.HANG:	MOV (PC)+,R3
31	000470	001403		MOV	BFO R3,FF0PT
32	000472	010367		MOV	BFO R3,FF0PT
33	000476	000207		RTS	PC
34	000480	012703		0.LC:	CLR R3
35	000482	001403		MOV	BFO R3
36	000484	010367		MOV	BFO R3,LCOPT
37	000486	000207		RTS	PC
38	000500	012703		0.WIDTH:MOV	R0, COLCNT
39	000502	001403		MOV	R0, RSTC+2
40	000504	010367		CMP	R0, R3
41	000510	000207		RTS	PC
42	000512	012703		0.HANG:	MOV (PC)+,R3
43	000514	001403		MOV	BFO R3,FF0PT
44	000516	010367		MOV	BFO R3,FF0PT
45	000522	000207		RTS	PC
46	000524	005003		0.LC:	CLR R3
47	000526	000240		MOV	BFO R3
48	000530	010367		MOV	BFO R3,LCOPT
49	000534	000207		RTS	PC

1	001000			5		
2	000200				7 LOAD POINT	
3	000034				LOADPT: .WORD	
4	0001002				.WORD	
5	0001004				.WORD	
6	0001006				.WORD	
7	0001010				.WORD	
8						
9					7 ENTRY POINT	
10						
11	0010102	177772			LP: MOV	LPCOF,R4
12	000364	000006			ASL	A(R4)
13	001022	103115			RCC	LPFRR
14	001024	052737	177514		RIS	#100, #LPS
15	001032	000207			RTS	PC
16						
17					7 INTERRUPT SERVICE	
18	001034	000512			RR	LPNONE
19						
20					RR	LPNONE
21	001036	000577	000242		LPNT: JSR	RS, @INTEN
22	001042	000140			MOV	=C@PR4>RPR7
23	001044	01A704	177740		TST	LPFGE,R4
24	001050	000737	177514		ERRPT: RMT	#LPS
25	001054	100441			TST	(R4)+
26	001056	000724			RLX0	
27	001060	001471			FFOPT: REC	(R4)+
28	001062	000724			TST	#LPS
29	001064	100737	177514		LPNEXT: TSTB	RPL
30	001070	100033			ASL8	(PC)+
31	001072	104327			TARFLG: .WORD	
32	001074	000000			RNF	
33	001076	001057			IGNORE: MOV8	@(R4)+,R5
34	001100	114405	177600		RIC	#177600,R5
35	001102	042705			TST	(R4)
36	001106	005714			REC	LPNONE
37	001110	001464			TNC	(R4)
38	001112	005214			TNC	(R4)
39	001114	000244			CMPB	RS, #40
40	001116	120527	000040		RLC	CHRST
41	001122	103417			CMPB	#100,R5
42	001124	122705	000140		RHTS	PCHAR
43	001130	103002			SUR	(PC)+,R5
44	001132	162705			LCOPT: 40	
45	001134	000040			PCHAR: DEC	(PC)+
46	001136	005327			COLCNT: .WORD	COLSIZ
47	001140	000204			RLT	IGNORE
48	001142	002756				

5

6

49 001144 10A327  
 50 001146 000001  
 51 001150 001423  
 52 001152 11M537 177516  
 53 001156 00M742  
 54 001160 000207  
 55 001162 12M527 000011  
 56 001166 001420  
 57 001170 12M527 000012

TARCNT: .WORN 1  
 PC1: MOV#LPB  
 RET: RTS  
 CHRTST: CMPB  
 TARSET: CMPB  
 CMPB

ASLB (PC)+  
 RSTTAB  
 R5,#LPB  
 LPNEXT  
 PC  
 R5,#HT  
 TARSET  
 R5,#LF

UPDATE TAR COUNT  
 RSTTAB  
 PRINT THE CHAR  
 TRY FOR NEXT CHAR  
 IS CHAR A TAB?  
 RSTTAB  
 IS IT LF?

LP V02-03 25-JUN-74 RT-11 MACRO VM02-10 14-APR-75 10:05:11 PAGE 4+

58 001174 001406  
 59 001176 12M527 000015  
 60 001202 000240  
 61 001204 12M527 000014  
 62 001210 001333  
 63 001212 01P767 000204  
 64 001220 01P767 000001  
 65 001226 000751  
 66 001230 01A767 177712  
 67 001236 01P705 000040  
 68 001242 000735  
 69  
 70 001244 005244  
 71 001246 022424  
 72 001250 01P705 000014  
 73 001254 000756  
 74  
 75 001256 05P754 000001  
 76  
 77  
 78  
 79 001262 005037 177514  
 80 001266 010704  
 81 001270 06P704 177520  
 82 001274 01P705 000054  
 83 001300 000175 000270  
 84  
 85 001304 000000  
 86  
 87 000306  
 88  
 89 000001  
 .END

REG RSTC  
 CMPB R5,#CR  
 NOP R5,#FF  
 CMPB IGNORE  
 MOV #COLSIZ,COLCNT  
 RSTTAB: MOV #1,TARCNT  
 RR PC1  
 MOV TARCNT,TABFLG  
 MOV #40,R5  
 RR PCHAR

RLK0: TNC  
 CMP (R4)+,(R4)+  
 MOV #FF,R5  
 RR RSTC

LPERR: RIS #DERR,0-(R4)  
 I OPERATION COMPLETE

LPERR: CLR #LPS  
 MOV PC,R4  
 ADD #LPCOE-,R4  
 MOV #MONLOW,R5  
 JMP @OFFSET(R5)

INTEN: 0  
 LP\$IZE = .-LOADPT  
 .END

RYES-RESTORE COLUMN COUNT  
 IS IT CR?  
 IGNORE UNLESS MODIFIED  
 IS IT A FF?  
 AND-CHAR IS NON-PRINTING  
 PRE-INITIALIZE COLUMN COUNTER  
 RSTTAB: MOV #1,TARCNT  
 PRINT THE CHAR  
 SET UP TAR  
 PRINT SPACES

MAKE SURE WE ONLY COME HERE ONCE  
 PRINT TO ADDR OF NEXT CHAR  
 PRINT INITIAL FF  
 SET HARD ERROR BIT

TURN OFF INTERRUPT  
 ADDR OF COE IN R3  
 JUMP TO Q MANAGER



LP V02-03  
SYMBOL TABLE

25-JUN-74

RT-11 MACRO VM02-10

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BLK0 001244  
CROPT 001202  
HT = 000011  
LOADPT 001000  
LPERR 001236  
LPSIZE= 000306  
O.FORM 000500  
PCHAR 001136  
RSTC 001212  
R3 =X000003  
TARCNT 001146  
ABS. 001306 000  
000000 001  
ERRORS DETECTED: 0  
FREE CORE: 18070. WORDS

CHRST 001162  
FRROPT 001054  
IGNORE 001100  
LP 001012  
LPTNT 001036  
LPVEC = 000200  
O.HANG 000512  
PC1 001152  
RSTTAB 001220  
R4 =X000004  
TARFLG 001074

COLCNT 001140 G  
FF 000014  
TNTEN 001304  
LPA = 177516  
I.PLOF 001006  
MONLOW = 000054  
O.LC 000524  
PR4 = 000200  
R0 =X000000  
R5 =X000005  
TARSFT 001230

COLSZ= 000204  
FFOPT 001060  
LCOPT 001134  
LPCOF 001010  
LPNEXT 001064  
OFFSFT= 000270  
O.WINT 000452  
PR7 = 000340  
R1 =X000001  
SP =X000006

CR = 000015  
HOFRR = 000001  
LF = 000012  
LPDONE = 001262  
LPS = 177514  
O.CR 000466  
PC =X000007  
RET 001160  
R2 =X000002  
TAR 001236

.LPI/NITM/C=LP

A.4 CR11 DEVICE HANDLER

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```

1 .TITLE CR.SYS
2 RT-11 CARD READER (CR11) HANDLER
3
4
5 DEC-11-OCRMA-D
6
7 FCP, ARC, PRR
8 MARCH 1975
9
10 COPYRIGHT (C) 1974, 1975
11
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13 MAYNARD, MASSACHUSETTS 01754
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```

CR.SYS RT-11 MACRO VM02-10 28-APR-75 16:00:18 PAGE 2

```

1 MISCFLANEOUS EQUATES
2
3
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25
26
27
28
29
30

```

13	000230		
14	177140		: INTERRUPT VECTOR
15	177162		: CARD READER STATUS REGISTER
16	177162		: DATA RUFFER 1
17	177164		: DATA RUFFER 2
18			
19	000001		: CONSTANTS FOR MONITOR COMMUNICATIONS
20	000054		: HARD ERROR
21	000270		: BASE ADDRESS OF MONITOR
22	177776		: OFFSET TO HANDLER RETURN
23	000340		: PROGRAM STATUS WORD
24	000300		: PRIORITY 7
25			: PRIORITY 4
26			
27	000015		: ASCII CONSTANTS
28	000012		: CARRIAGE RETURN
29	000040		: LINE FEED
30	000041		: SPACE
31			: END OF FILE
32			
33	000001		: CARD READER CONTROL AND STATUS BITS
34	000002		: READ
35	000100		: DETECT CARD
36	000200		: INTERRUPT ENARLF
37	000400		: COLUMN DONE
38	001000		: READY
39	002000		: BUSY
40	004000		: ONLINE
41	010000		: DATA LATE
42	020000		: MOTION CHECK(CM11 ONLY)
43	040000		: HOPPER CHECK(CM11 ONLY)
44	100000		: CARD NONE
45			: ERROR

```

1  .SRTTL CONFIGURATION SECTION
2
3  ; THE FOLLOWING CODE IS EXECUTED WHEN A "SFT CR" CONSOLE COMMAND IS
4  ; GIVEN.
5
6  .ASETT
7  .=4000
8  .CONFIGURATION AREA
9
10 ; SET CR (NO) CRIF
11 ; APPEND/DO NOT APPEND CARRIAGE RETURN/LINE FEED TO EACH CARD IMAGE
12
13 RR .+IXCRLF=XCRLF
14 .RAD50 /CRLF/ 022600
15 .WORD <CRLF-400>/2+100000
16
17 ; SET CR (NO) TRTM
18 ; TRTM/DO NOT TRTM TRAILING BLANKS FROM CARD IMAGES
19
20 RR .+IXTRTM=XTRTM
21 .RAD50 /TRIM/ 050500
22 .WORD <TRIM-400>/2+100000
23
24 ; SET CR (NO) HANG
25 ; HANG/RETURN HARD ERROR IF READER NOT READY AT START OF TRANSFER
26
27 RNF .+IERROR=XHANG
28 .RAD50 /HANG/ 025700
29 .WORD <HANG-400>/2+100000
30
31 ; SET CR CODE (S) 024 (029)
32 ; SET TRANSLATION TO 026 (029) MODE
33
34 .WORD 024.
35 .RAD50 /CODE/ 017500
36 .WORD <CODE-400>/2+400000
37
38 ; SET CR (NO) IMAGE
39 ; TRANSMIT EACH COLUMN AS 12 BITS (ONE WORD/COLUMN)
40
41 .WORD NOTMAG=IMRASF
42 .RAD50 /IMAGE/ 026210
43 .WORD <IMAGE-400>/2+100000
44 .WORD 0
45 ; END-OF-OPTIONS FLAG

```

```

1  .SRTTL  CONFIGURATION SUBROUTINES
2
3  CRIF:  MOV 012703 (PC)+,R3 ;NOP IF POSITIVE
4  NOP 000240
5  MOV 010367 000246 ;ENTRY POINT FOR NO
6  PTS 000207
7
8  TRTM:  MOV 012703 (PC)+,R3 ;NOP IF POSITIVE
9  NOP 000240
10 MOV 010367 000236 ;ENTRY POINT FOR NO
11 PTS 000207
12
13 HANG:  MOV 012703 (PC)+,R3 ;NOP IF POSITIVE
14 NOP 000240
15 MOV 010367 000036 ;ENTRY POINT FOR NO
16 PTS 000207
17
18 CODE:  MOV 010701 PC,R1 ;RI -> CONVERSION TABLE FOR 024
19 ADD 062701 #SFT026=.,R1 ;026 REQUESTED?
20 SUB 160300 ;NOPF - ERROR (NOTE THAT C IS SET)
21 RMT 100407 ;YES, IT'S 026 - GO DO IT
22 MOV 001407 SETCON ;YES, IT'S 026 - GO DO IT
23 000050 #SFT029=SET026,R1 ;RI -> CONVERSION TABLE FOR 029
24 022700 #3,R0 ;WAS IT 029?
25 001402 SETCON ;YES
26 000261 ;ELSE AN ERROR - INDICATE SUCH
27 000207 ;AND RETURN TO KMON
28 010703 PC,R3 ;RI -> CHARACTER TABLE
29 062703 177716 ;PICK UP NEXT OFFSET TO BE MODIFIED
30 005000 ; FROM APPROPRIATE TABLE
31 000552 152100 ;ALL DONE (NOTE: C IS CLEAR)
32 001771 ;PRINT TO BYTE TO MODIFY
33 000554 060300 ;AND PLUS IN NEW VALUE
34 000556 112110 ;CONTINUE
35 000562 000772
36
37 TMAG:  ADD 062703 000010
38 000570 060703
39 000572 00176 ;
40 000576 012367 000176 ;
41 000602 012367 000154 ;
42 000606 012367 000156 ;
43 000612 000207 000164 ;
44
45 NOTMAG: MOV 001402 ;NEXTCHR-XTM1
46 000616 114537 000306 ;CTRL-XTM2(R5),*(PC)+
47 000622 000001 ;WORD
48 000624 000402 ;
49 000626 012337 177162 ;
50 000632 000002 ;WORD

```

SRITL 026, 029 CONVERSION TARLFS

026 CONVERSION TABLE  
 MODIFIERS CHARACTER TABLE TO ACCEPT 026 KEYPUNCH CODES

1	000634	012	137	.BYTE	012,137	BACK ARROW	(8-2)
2	000636	013	075	.BYTE	013,075	EQUAL	(8-3)
3	000640	015	136	.BYTE	015,136	UP ARROW	(8-5)
4	000642	016	047	.BYTE	016,047	POSTROPH	(8-6)
5	000644	017	134	.BYTE	017,134	BACKSLASH	(8-7)
6	000646	052	073	.BYTE	052,073	SFMTCOLN	(0-8-2)
7	000650	054	050	.BYTE	054,050	LEFT PAREN	(0-8-4)
8	000652	055	042	.BYTE	055,042	QUOTES	(0-8-5)
9	000654	056	043	.BYTE	056,043	LR STGN	(0-8-6)
10	000656	057	045	.BYTE	057,045	PERCENT	(0-8-7)
11	000640	112	072	.BYTE	112,072	COLN	(11-A-2)
12	000642	115	133	.BYTE	115,133	L BRACKET	(11-A-5)
13	000644	116	076	.BYTE	116,076	GREATER THAN	(11-A-6)
14	000646	117	046	.BYTE	117,046	AMPERSAND	(11-A-7)
15	000652	200	053	.BYTE	200,053	PI US	(12)
16	000654	212	077	.BYTE	212,077	QUESTION	(12-A-2)
17	000656	214	051	.BYTE	214,051	RIGHT PAREN	(12-A-4)
18	000674	215	135	.BYTE	215,135	R BRACKET	(12-A-5)
19	000700	216	074	.BYTE	216,074	LESS THAN	(12-A-6)
20	000702	000000		.WORD	0	END OF TABLE **	

029 CONVERSION TABLE  
 MODIFIERS CHARACTER TABLE TO ACCEPT 029 KEYPUNCH CODES

1	000704	012	072	.BYTE	012,072	COLN	(8-2)
2	000706	013	043	.BYTE	013,043	LR STGN	(8-3)
3	000710	015	047	.BYTE	015,047	POSTROPH	(8-5)
4	000712	016	075	.BYTE	016,075	EQUAL	(8-6)
5	000714	017	042	.BYTE	017,042	QUOTES	(8-7)
6	000716	052	134	.BYTE	052,134	BACKSLASH	(0-8-2)
7	000720	054	045	.BYTE	054,045	PERCENT	(0-8-4)
8	000722	055	137	.BYTE	055,137	BACK ARROW	(0-8-5)
9	000724	056	076	.BYTE	056,076	GREATER THAN	(0-8-6)
10	000726	057	077	.BYTE	057,077	QUESTION	(0-8-7)
11	000730	112	135	.BYTE	112,135	R BRACKET	(11-A-2)
12	000732	115	051	.BYTE	115,051	RIGHT PAREN	(11-A-5)
13	000734	116	073	.BYTE	116,073	SFMTCOLN	(11-A-6)
14	000736	117	136	.BYTE	117,136	UP ARROW	(11-A-7)
15	000740	200	046	.BYTE	200,046	AMPERSAND	(12)
16	000742	212	133	.BYTE	212,133	L BRACKET	(12-A-2)
17	000744	214	074	.BYTE	214,074	LESS THAN	(12-A-4)
18	000746	215	050	.BYTE	215,050	LEFT PAREN	(12-A-5)
19	000750	216	053	.BYTE	216,053	PI US	(12-A-6)
20	000752	000000		.WORD	0	END OF TABLE **	

.ITF GE <CHRTBI-LOADPT+216>=1000, .ERROR TABLE NOT IN BLOCK 1

```

1 00000000
2
3
4
5
6 000000 000230
7 000002 000050
8 000004 000300
9 000006 000000
10 000010 000000
11
12
13
14 000012 014705 177772
15 000016 014704 000176
16 000022 004365 000006
17 000026 101555
18 000030 032737 001400 177160
19 000036 000240
20 000040 005725
21 000042 001525
22 000044 005725
23 000046 000514
24 000050 000547
25
26
27
28
29 000052 004577 001246
30 000056 000040
31 000060 005367 000264
32 000064 013705 177160
33 000070 100541
34 000072 105705
35 000074 100041
36 000076 013705
37 000102 001423
38 000104 013704
39 000110 026767 000044 000072
40 000116 001002
41 000120 011647 000100
42 000124 042716 177003
43 000130 011646
44 000132 005416
45 000134 042626
46 000136 001402
47 000140 012705 000377
48 000144 014727 000010

```

```

.SRTTL HANDLER PROPR
.CSECT CR11

? LOAD POINT
LOADPT: .WORD CRVECT
        .WORD CRINT-.
        .WORD 300
CRIGF: .WORD 0
CRDGF: .WORD 0

? ENTRY POINT
CRAND: MOV     CRDF,P5
        MOV     CRPTR,R4
        ASI     4(R5)
        PLDS   LERRR
        #READY+RUSY,#CRST JTS READER READY?
        NOP
        YHANG: NOP
        TST
        RFD
        TST
        RR
        ABORT

? INTERRUPT ENTRY POINT
CRINT: TSP     P5,@SINPTR
        NEG     C<PR6>&PR7
        MOV     @CRST,R6
        RMT     FRDP
        TSTR   P5
        RPI    CAPD
        MOV     @CR2,R5
        RFD    TNCOLT
        RMP     @CR1,=(SP)
        RNF    CRPTR,R1F0P
        TSTPIN
        MOV     @SP,CHAR12
        TSTPIN: RIC #17703,@SP
        MOV     @SP,=(SP)
        NEG
        RIC     (SP)+,(SP)+
        YIM1:  RER  NATCHR
        MOV     #377,R5
        NATCHR: MOV     CRPTR,(PC)+

```

```

? INTERRUPT VECTOR
? OFFSET TO INTERRUPT SERVICE
? PS
? LAST RUFUF ENTRY
? CURRENT QUEUE ENTRY

? POINT TO 0 ELEMENT
? POINT INTO CARD IMARE
? CONVERT WORD COUNT TO RYTE COUNT
? NULL REQUEST OR WRITE IS LOGIC FRD
? * PATCH HERE TO ISSUE HARD ERROR
? BLOCK 0 REQUEST ?
? YES, GO INITIATE REQUEST
? NO, POINT TO BUFFER PTRS
? GO SEE IF ANY STUFF IS LEFT IN OLD CARD
? ABORT

? ENTER SYSTEM STATE
? COUNT DOWN INTERRUPTS THIS CARD
? GET STATUS
? WHOOPS -- ERROR
? CHECK FOR COLUMN NONE
? BRANCH IF NOT COLUMN NONE
? IT'S BLANK
? GET COMPRESSED CHAR
? FIRST COLUMN?
? NONE
? ELSE SAVE FOR EOF CHECK
? CHECK ONLY ROWS 1-7
? CHECK FOR INVALID PUNCHES
? BY CHECKING FOR 2 OR MORE PUNCHES
? IT'S OKAY
? ELSE FORCE TRANSLATION INTO 134
? MEMBER POSITION OF NON-BLANK

```

```

49 000150 000000
50 000152 060705
51 000154 114537 000306
52 000160 000000
53 000162 063767 000001 177770
54 000170 052737 000100 177140
55 000176 000207
56
57

```

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```

58
59
60 000200 032705 040000
61 000204 001400
62 000206 012704
63 000210 000000
64 000212 014705 177572
65 000216 025525
66 000220 022727 007417
67 000224 000000
68 000226 001471
69 000230 010106
70 000232 014701 177722
71 000236 000200
72 000240 014701 177704
73 000244 005201
74 000246
75 000246 000240
76 000250 112721 000015
77 000254 112721 000012
78 000260 010147 177644
79 000264 012601
80 000266 000400
81 000270 112435
82 000272 005315
83 000274 001457
84 000276 005205
85 000300 024704 177644
86 000304 101371
87 000306 032737 001400 177140
88 000314 001325
89 000316 010704
90 000320 062704 000542

```

? CARD DONE OR ERROR

```

CARD: RIT #CARDN,RS
REC FRP1
MOV (PC)+,R4
RUPTR: .WORD
CMP (RS)+,(RS)+
CMP #7417,(PC)+
CHAR12: .WORD
REC FNAFTL
MOV R1,(SP)
MOV CHARTR,R1
YTOIM: NOP
MOV FNDPTR,R1
TNC R1
LXTRIM: NOP
YCOLF: #CR,(R1)+
#LF,(R1)+
R1,ENDPTR
IXCRIF: MOV (SP)+,R1
RR CONT
FITBIUF: MOV (R4)+,@(RS)+
#RS
#ETMON
REC #R1
TNC #R1
CMP FNDPTR,R4
RHT FIBUF
RIT #READY+R1USV,#CRST
RNF INTRET
MOV PC,R4
ADD #CHRAUF-.,P4

```



```

91 000324 010467 177630
92 000330 010467 177614
93 000334 010467 177650
94 000340 005047 177660
95 000344 012727 000120
96 000350 000000
97 000352 012727 000101
98 000360 000207
99
100
101
102 000362 014705 177422
103 000366 052755 000001
104 000372 000416
105
106 000374 032705 004000
107 000400 001370
108 000402 005767 177742
109 000406 100337
110
111 000410 000673
112
113
114

```

VARIOUS ERRORS

```

ERROR: MOV CRCPF,R5
RIS PHERR,0-(PS)
RR ABORT

```

```

ERROR: RIT #DATLAT,R5
RNF LERRR
TST COICNT
RPI FRP1
RR CAPD

```

END OF FILE CARD FOUND

```

!START FILLING FROM NEW CARD
!WHICH IS AS YFT EMPTY
!ESTABLISH BUFFER POINTER
!CLEAR EOF FLAG
!SFT COLUMN COUNT
!COUNT OF COLUMNS REMAINING IN CARD
!FIRST START A CARD GAIN
!BYE

```

```

!POINT TO BUFFER ELEMENT
!YOU CAN'T WRITE ON A READER

```

```

!DATA RATE IS ONLY NOT CURABLE
!ISSUE HARD ERROR IF SO
!DONE WITH DATA COLUMNS?
!NOPE -- MUST REPTCK CHECK, ETC.
! START A NEW READ TO CORRECT CONDITION.
!ELSE ASSUME CARD DONE

```

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HANDLER PROPER

```

115 000412 012504
116 000414 105024
117 000416 005315
118 000420 001375
119 000422 052775 020000 177770
120 000430 005067 177514
121
122
123
124 000434 010467 177520
125 000440 005037 177140
126 000444 010704
127 000446 062704 177342
128 000452 013705 000054
129 000456 000175 000270

```

```

FNFTL: MOV (R5)+,R4
CLRBUF: CLR (R4)+
RNF #R5
RIS #20000,0-10(PS)
ABART: CLR FNOPTR

```

```

! RETURN TO MONITOR (REQUEST DONE, EOF, OR ERROR)

```

```

REYMON: MOV P4,CHRPTP
CLP #FRST !SAVE POSITION IN CARD
MOV PC,R4 !NO INTERRUPTS
ADD #CRCOE--,R4 !THE USUAL MONITOR RETURN
MOV #MONLOW,R5
IMP #OFFSET(PS)

```

```

!POINT INTO HIS BUFFER
!CLEAR IT ALL

```

```

!SFT EOF BIT IN CHANNEL
!FORCE A READ NEXT TIME

```

```

!SAVE POSITION IN CARD
!NO INTERRUPTS
!THE USUAL MONITOR RETURN

```

.SRTL CHARACTER TABLE

! THE FOLLOWING MACRO TAKES AS ARGUMENTS THE ASCII TRANSLATION  
 ! DESIRED AND THE LIST OF PUNCH COMBINATIONS FOR THAT CHARACTER.

```

.MACRO C $LIST
  TEM
  X,<LIST>
  .IF NE X,
  .IF LE X,7
  TETV,
  .IFF
  U=10
  .RPT X,=A,
  U=U+1
  .ENDR
  TET+1
  .ENDC
  .IFF
  TET+40
  .ENDC
  .ENDR
  .=CTRL+T
  .BYTE $CHAR
  $CHAR = $CHAR + 1
  .ENDM
  C
  
```

! THE FOLLOWING TABLE TRANSLATES 029 KEYPUNCH CODES TO ASCII.

```

CTRL: 254
.RPT .BYTE 130 ;DEC STANDARD FROM CHARACTER
.ENDR
.=CTRL
  
```

1  
2  
3  
4  
5  
6  
7  
8  
9  
10  
11  
12  
13  
14  
15  
16  
17  
18  
19  
20  
21  
22  
23  
24  
25  
26  
27  
28  
29  
30  
31  
32  
33  
34  
35

000462  
000000  
000462

START AT OCTAL 000

34	000000	SCHAR E C
37	000462	<12,0,9,A,1>
38	000754	<12,0,1>
40	000704	<12,0,2>
41	000705	<12,0,3>
42	000706	<9,7>
43	000512	<0,9,A,5>
44	000540	<0,9,A,6>
45	000541	<0,9,A,7>
46	000542	<11,9,A>
47	000611	<12,0,5>
48	000710	<12,0,A,3>
49	000540	<12,0,A,0>
50	000716	<12,0,A,5>
51	000717	<12,0,A,6>
52	000720	<12,0,A,7>
53	000721	<12,11,9,8,1>
54	000722	<11,0,1>
55	001014	<11,0,2>
56	000604	<11,0,3>
57	000605	

FNULI  
 CTRL-A  
 CTRL-B  
 CTRL-C  
 CTRL-D  
 CTRL-E  
 CTRL-F  
 CTRL-G  
 CTRL-H  
 CTRL-I  
 CTRL-K  
 CTRL-L  
 CTRL-M  
 CTRL-N  
 CTRL-O  
 CTRL-P  
 CTRL-Q  
 CTRL-R  
 CTRL-S

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CHARACTER TARLF

58	000606	C	CTRL-T
59	000517	C	<9,8,4>
60	000520	C	<9,8,5>
61	000505	C	<9,2>
62	000551	C	<0,9,6>
63	000613	C	<11,9,A>
64	000614	C	<11,0,A,1>
65	000522	C	<9,8,7>
66	000572	C	<0,9,7>
67	000617	C	<11,9,A,4>
68	000620	C	<11,0,A,5>
69	000621	C	<11,0,A,6>
70	000622	C	<11,0,A,7>
71	000463	C	<>
72	000702	C	<12,A,7>
73	000502	C	<8,7>
74	000476	C	<8,3>
75	000576	C	<11,A,3>
76	000537	C	<0,8,4>
77	000663	C	<12,7>
78	000500	C	<8,5>
79	000700	C	<12,A,5>
80	000600	C	<11,A,5>
		C	<11,A,4>

CTRL-U  
 CTRL-V  
 CTRL-W  
 CTRL-X  
 CTRL-Y  
 CTRL-Z  
 ALTMODE (FSCAPE)  
 CTRL-  
 SPACE  
 !  
 #  
 \$  
 %  
 &  
 ' "  
 ( )  
 \*



123 000526  
 124 000527  
 125 000530  
 126 000531  
 127 000532  
 128 000533  
 129 000543  
 130 000675  
 131 000535  
 132 000575  
 133 000602  
 134 000540  
 135 000474  
 136 000724  
 137 000725  
 138 000726  
 139 000727  
 140 000730  
 141 000731  
 142 000732  
 143 000733  
 144 000743  
 145 000764  
 146 000765  
 147 000766  
 148 000767  
 149 000770  
 150 000771  
 151 000772  
 152 000773  
 153 001043  
 154 000625  
 155 000626  
 156 000627  
 157 000630  
 158 000631  
 159 000632  
 160 000633  
 161 000643  
 162 000723  
 163 000763  
 164 000623  
 165 000624  
 166  
 167  
 168  
 169 001062  
 170 001324  
 171

C <0,4>  
 C <0,5>  
 C <0,6>  
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FU  
 FV  
 FW  
 FX  
 FY  
 FZ  
 FL  
 FA  
 FE  
 FF  
 FACCENT GRAVE  
 FFC A  
 FFC R  
 FFC C  
 FFC D  
 FFC F  
 FFC F  
 FFC G  
 FFC H  
 FFC T  
 FFC J  
 FFC K  
 FFC L  
 FFC M  
 FFC N  
 FFC O  
 FFC P  
 FFC Q  
 FFC R  
 FFC S  
 FFC T  
 FFC U  
 FFC V  
 FFC W  
 FFC X  
 FFC Y  
 FFC Z  
 FOPEN BRACE  
 FVERTICAL BAR  
 FCLOSE BAR  
 FITTLE  
 FDFL

\* = CHPTPL + 256.  
 CHBUIF: BIKW B1.  
 \$INPTR: .WORD 0

001062  
 000000

IPLURGED TO POINT TO COMMON ENTRY

CR.SYS RT-11 MACRO VM02-10 28-APR-75 16:00:38 PAGE 7+

CHARACTER TABLE

172 001326  
 173  
 174 000001'

CRSIZE=-LOADPT

.END

CR.SYS RT-11 MACRO VM02-10 28-APR-75 16:00:38 PAGE 7+

SYMBOL TABLE

ABORT	000430P	002	RUFPTR	000210P	002	RUSY	= 001000	CARD	000200P	002	CRDN	= 040000
CHAR12	000224R	002	CHRBUF	001042P	002	CHPTR	000140P	CHRPL	000442R	002	CLRBUF	000414R
CODE	000510	002	CONEXT	000540	002	COLCNT	000350P	COLD	= 000200	002	CONT	000300P
CR	= 000015	002	CR1	= 177142	002	CR2	= 177144	CRGF	000010P	002	CRHAND	000012P
CRINT	000052P	002	CRIF	000452	002	CRLOF	000006P	CRSIZ	= 001326	002	CRST	= 177160
CRVCT	000230	002	DATLAT	= 004000	002	FJECT	= 000002	ENDFL	000412R	002	ENDPTR	000150P
FOF	= 000041	002	FRR	= 100000	002	FROF	000374P	FRR1	000306P	002	FILBUF	000270P
HANG	000476	002	HDFRR	= 000001	002	HDPCK	= 020000	TMAGF	000564	002	IMRARE	000572
INCOLT	000152P	002	INTER	= 000100	002	IXTRFT	000170P	IERRR	000342R	002	IF	= 000012
LOADPT	000000P	002	IXTRLF	000240P	002	IXTRTM	000246P	MONLW	= 000054	002	MOTIN	= 010000
NOTMAG	000614	002	NXTCR	000144R	002	NFFSFT	000270	ONI IN	= 002000	002	PC	= X000007
PR4	= 000300	002	PR7	= 000340	002	PS	= 177776	PEAD	= 000001	002	READP	000316P
READY	= 000400	002	PETHON	000434R	002	R0	= X000000	P1	= X000001	002	R2	= X000002
R3	= X000003	002	R4	= X000004	002	R5	= X000005	SCNDF	000550	002	SETCND	000542
SET026	000634	002	SET029	000704	002	SP	= X000006	SPACE	= 000040	002	T	= 000227
TRTM	000464	002	TSTPIN	000124P	002	U	= 000020	YCRLF	000246P	002	YHANC	000036R
XIMI	000116R	002	YIM2	000154P	002	YIM3	= 000142P	YTRIM	000236P	002	YATMAG	000624
XCHAR	= 000200	002	XINPTR	001324R	002							

.ABS. 000754 000  
 . 000000 001  
 CR11 001326 002  
 ERRORS DETECTED: 0  
 FRFE CORR: 17698. WORDS

.LP:/N:ITM/CBCP

NT VM2-07 12-APR-74 RT-11 MACRO VM02-1M 28-APR-75 16:04:24 PAGE 1

```

1 .TITLE NT VM2-07 12-APR-74
2
3 ? RT-11 DECTAPE (TC11) HANDLER
4
5 ? DEC-11-OPTDA-D
6
7 ? PB/EF
8
9 ? AUGUST, 1974
10
11 ? COPYRIGHT (C) 1974,1975
12
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31

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.CRECT SYSMND
.ENARL ISR
R0=X0
R1=X1
R2=X2
R3=X3
R4=X4
R5=X5
R6=X6
PC=X7
PS=177774

; DECTAPE CONTROL REGISTERS
TCRT = 177350
TCST = 177340
TCFM = 177342
TCWC = 177344
TCRA = 177346
TCVEC = 210

; CONSTANTS FOR MONITOR COMMUNICATION
HDFRR = 1
MONLOW = 54
OFFSET = 270
;GLOBL DTSYS, PKSYS, PFSYS,
;GLOBL SIMPTR, SINTEN, OTINT

;DATA REGISTER
;CONTROL AND STATUS REGISTER
;COMMAND REGISTER
;WORD COUNT REGISTER
;BUS ADDRESS REGISTER
;TC11 INTERRUPT VECTOR

;HARD FRDP BIT
;MONITOR BASE POINTER
;POINTER TO Q MANAGER COMP ENTRY

;PK TS NOT RESIDENT
;NFIFTH IS RF

;ENTERS AT LVL 7

```



```

1 000000
2 000000
3 000000
4 000000
5 000002
6 000004
7 000006
8 000006
9 000010
10 000010
11 000012
12 000016
13 000020
14 000024
15 000030
16 000030
17 000030
18 000032
19 000034
20 000040
21 000042
22 000044
23 000050
24 000054
25 000056
26 000060
27 000064
28 000070
29 000072
30 000074
31 000100
32 000110
33 000112
34 000114
35 000120
36 000124
37 000126
38 000132
39 000136
40 000142
41 000144
42 000144
43 000144
44 000144
45 000144
46 000144
47 000144
48 000144

```

: LOAD POINT  
 :WORD TCVER  
 :WORD INTINT-  
 :WORD PR7  
 DTSYS:  
 DTIOF: :WORD 0  
 DTIOF: :WORD 0  
 I ENTRY POINT  
 MOV #8.,(PC)+  
 DTTRY: :WORD 0  
 MOV #PS, -(SP)  
 JSP PC,DTINT  
 RTS PC  
 I INTERRUPT SERVICE  
 RR DTSTOP  
 JSP PS, @SIMPTR  
 :C<300>&PR7  
 :WORD  
 MOV RTIOF,R0  
 MOV #TCCM,R4  
 MOV (R0)+, -(SP)  
 MOV (R0)+,PS  
 RLC #C<3400>,PS  
 RIT #100100,(R4)  
 RMT RTERR  
 RER RETRY  
 RT #2,(R4)  
 RER RTDONE  
 CMP #TCCT,RWANT  
 RER RLKEND  
 RER FORWARD  
 REVERSF:RIS #4000,PS  
 SUR #2,(SP)  
 RER FORWARD  
 RIS #10000,RS  
 FORWARD:RIS #103,RS  
 MOV (SP)+,RWANT  
 RETRNI: MOV PS,(R4)  
 MOV (SP)+,R0

:ADDRESS OF INTERRUPT VECTOR  
 :OFFSET TO INTERRUPT SERVICE  
 :PRIORITY 7  
 :POINTER TO LAST 0 ENTRY  
 :POINTER TO CURRENT 0 ENTRY  
 :INIT THE RETRY COUNT  
 :RETRY COUNTER  
 :FAKE AN INTERRUPT  
 :BACK TO MONITOR  
 :ABORT CALL FROM RF SYSTEM  
 :NOW JSR TO COMMON CODE  
 :RIN AT LEVEL 4  
 :RD POINTS TO 0 PLFMENT  
 :RD POINTS TO CONTROL REGISTER  
 :DESTROYED BLOCK # ONTO STACK  
 :UNIT # INTO RS  
 :ISOLATE UNIT NUMBER  
 :ERRRR RIT ON?  
 :YES  
 :IF INTERRUPT IS OFF, WE ARE INITIATING  
 :A REQUEST  
 :SEARCHING  
 :IN-A READ OR WRITE JUST COMPLETED  
 :COMPARE ACTUAL BLOCK TO DESTROYED BLOCK  
 :FOUND IT  
 :SEARCH IN THE FORWARD DIRECTION  
 :SFT REVERSE RIT  
 :SEARCH FOR TWO BLOCKS BEFORE ONE  
 :ACTUALLY DESIRED (TO ALLOW  
 :SPACE FOR THE TURN-AROUND  
 :DON'T SFT DELAY INHIBIT  
 :START IS ALREADY MOVING FORWARD  
 :SO INHIBIT HARDWARE DELAY  
 :INTERRUPT ENARL,RNIM,AND GO  
 :REMEMBER THE BLOCK WE ARE LOOKING FOR  
 :TELI CONTROLLER TO GO  
 :RESTORE R0

```

49 000146 000207
50
51 000150 032714 004000 004000 (R4)
52 000154 001757 REVERSE TAPE
53 000156 032714 004000 004000 (R4)
54 000162 001363 FORWARD
55
56
57

```

BACK INTO MONITOR

IFRF WE IN REVERSE?  
 IN-REVERSE TAPE  
 IFRF WE GOING FORWARD?  
 IN-WE HAVE TO TURN AROUND

INITIATE READ/WRITE REQUEST

DT 002-07 12-APR-74 RT-11 MACRO VM02-10 28-APR-75 16:04:24 PAGE 14

```

58 000164 052705 010115 RIS #10115,R5
59 000170 012037 177346 MOV (R0)+,#TCRA
60 000174 011016 MOV (R0),(SP)
61 000176 100404 RMT 1$
62 000200 001423 REC DTDONE
63 000202 005416 NEG (SP)
64 000204 042705 RIC #10,R5
65 000210 012637 MOV (SP)+,#TCWC
66 000214 000752 RR PTRNI
67
68
69 000216 032737 104000 177340 DTERR: RIT #104000,#TCST
70 000224 100003 RPI NOTE7
71 000226 032714 000002 RIT #2,(R4)
72 000232 001346 RNF FNDZR
73 000234 005367 177556 DEF DTRY
74 000240 003017 RGT RETRY
75 000242 052770 000001 177772 RIS #DNERR,0-6(R0)
76
77
78
79 000250 005726 DTDONE: TST (SP)+
80 000252 012600 MOV (SP)+,R0
81 000254 112737 000011 177342 DTSTOP: MOV #11,#TCRM
82 000262 010704 MOV PC,R4
83 000264 062704 ADD #DTCRE-,R0
84 000270 013705 MOV #MONLOW,R5
85 000274 000175 JMP #OFFSET(R5)
86
87
88
89 000300 105737 177340 RETRY: TSTB #*TCST
90 000304 100710 RMT FORW1
91 000306 062767 000004 000002 ADD #4,BWANT
92 000314 022716 CMP (PC)+,(SP)

```

ASSUME WRTF  
 CORF ADDRESS  
 WORD COUNT (OVER BLOCK #)  
 WRITE WAS A GOOD GUFFS  
 IF ZERO, SEEK  
 READ-NEGATE WORD COUNT  
 SET READ FUNCTION  
 SET WORD COUNT

END7 FROOP?  
 INT END7  
 WRF WE SPARCHING?  
 YES-REVERSE TAPE  
 MORE TRTER LEFT?  
 YES  
 IN-SET HARD ERROR BIT

OPERATION FINISHED

POP BLOCK  
 RSTORE R0  
 STOP SELECTED DRIVE

ANDR OF ONE IN R4

RETRY CODE

TAPE UP TO SPEED?  
 YES-AVOID STOPPING TAPE  
 IN-TT TAKES 4 BLOCKS TO START AND STOP  
 MAKE AN ATTEMPT TO START IN THE

RIGHT DIRECTION BASED ON LAST BLOCK  
 :DFSTRFD

RWANT: 0 DIRECT  
 RR DIRECT  
 \$INPTR: .WORD \$INTFN  
 DT\$IZE = .-REC :SIZE OF DT HANDLER  
 .END

93 000316 000000  
 94 000320 000674  
 95 000322 000000  
 96 000324 000001  
 100

DT V02-07 12-APR-74 RT-11 MACRO VM02-10 28-APR-75 16:00:24 PAGE 3+

SYMBOL TABLE

REF	000000	002	RLKEND	000156R	002	RWANT	000316R	002	DIRECT	000112R	002	DPAYS	= 000000 G
DSSYS	= 000000 G	002	DTCOF	000010R	002	DTDONE	000250R	002	DTFRP	000216R	002	DTINT	000034RG
DTLOF	000006R	002	DT\$IZE	= 000324R	002	DTSTOP	000254R	002	DTSYS	000006RG	002	DTTRY	000016R
DTSYS	= 000000 G	002	FNDZP	000150R	002	FORWAR	000172R	002	FORW1	000126R	002	HDFFR	= 000001
MONLOW	= 000054	002	NOTEZ	000234R	002	OFFSFT	= 000270	002	PC	=X000007	002	PR7	= 000340
PS	= 17776	002	PETRN1	000142R	002	PETRY	000300R	002	P2	=X000002	002	RFYS	= 000000 G
RKSYS	= 000000 G	002	R0	=X000000	002	R1	=X000001	002	PEVERS	000114R	002	R3	=X000003
R4	=X000004	002	R5	=X000005	002	SP	=X000006	002	TCRA	= 177346	002	TCM	= 177342
TCOT	= 177350	002	TCST	= 177340	002	TCVEC	= 000214	002	TCWC	= 177344	002	\$INPTR	000322RG
\$INTFN	= ***** G	000			000			000			000		
.ABS.	000000	001			001			001			001		
SYSHND	000324	002			002			002			002		
ERRORS DETECTED:	0												
FRPE CORR:	10076.	WORDS											

.LP:/N:TTM/C=DT



## APPENDIX B

### FOREGROUND TERMINAL HANDLER

The following listing is a terminal handler for the foreground. The user can write his own handler using this code as an example, or use the copy provided in the software kit. Instructions for its use are found on the second and third pages of the listing.

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```

.TITLE KB.MAC V01-01  
 / RT-11 V2 DEVICE INDEPENDENT TERMINAL HANDLER, KB.  
 / DEC-11-ORKRA-D  
 / COPYRIGHT (C) 1975  
 / DIGITAL EQUIPMENT CORPORATION  
 / MAYNARD, MASSACHUSETTS 01754  
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 / MARCH 1975  
 / RGB

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/ RT-11 V2 DEVICE INDEPENDENT TERMINAL HANDLER, KB. KB  
 / CAN BE USED BY EITHER THE FOREGROUND OR BACKGROUND (BUT NOT  
 / BOTH SIMULTANEOUSLY) TO READ AND WRITE TO ANY DL-11A OR KL-11A  
 / CONTROLLED TERMINAL.  
 / THIS HANDLER HAS THE FOLLOWING CHARACTERISTICS:  
 / 1) CARTRIDGE RETURN CAUSES THE REMAINDER  
 / OF THE INPUT BUFFER FOR THE CALLING READ REQUEST TO BE  
 / ZERO-FILLED, AND THE READ IS COMPLETED. THUS, THE HANDLER

```

11 TRANSFERS ONE LINE AT A TIME, NO MATTER HOW LONG THE
12 INPUT BUFFER IS FOR THE READ REQUEST, THE UNUSED PORTION
13 OF THE BUFFER IS ZERO-FILLED, CARRIAGE RETURN ECHOES
14 CARRIAGE RETURN, /LINE-FEED, AND INSERTS CR AND LF CHARACTERS
15 IN THE BUFFER IF THERE IS ROOM, ELSE ONLY CR IS PLACED IN THE BUFFER.
16
17 2)FORM FEED ECHOES 7 LINE FEEDS, AND INSERTS A FF CHARACTER IN
18 THE BUFFER.
19
20 3)RUBOUT ECHOES "^" AND DELETES THE LAST CHARACTER IN THE BUFFER.
21 IF THERE ARE NO CHARACTERS IN THE BUFFER, RUBOUT DOES NOT ECHO
22 AND IS IGNORED.
23
24 4)TAB ECHOPS ENOUGH SPACES TO POSITION THE PRINT HEAD AT THE
25 NEXT TAB STOP, AND INSERTS A TAB CHARACTER IN THE BUFFER.
26
27 5)CTRL U ECHOES "u" AND ERASES THE CURRENT LINE.
28
29 6)CTRL Z ECHOES "z" AND CAUSES THE HANDLER TO REPORT END-OF-FILE.
30
31 7)THE CTRL Z CHARACTER IS NOT INSERTED IN THE BUFFER,
32 THE LOW-SPEED READER WILL RUN IF IT IS TURNED ON WHILE A READ
33 REQUEST IS PENDING TO THE HANDLER. IF THE TAPE BEING READ HAS
34 MANY TABS, HOWEVER, THE TIME NECESSARY TO ECHO THE TABS WILL
35 CAUSE CHARACTERS FOLLOWING THE TABS TO BE LOST. TO DISABLE THE
36 ECHOING OF TABS, THE "SET" COMMAND CAN BE USED AS FOLLOWS:
37
38 "SET KR LSR" WILL DISABLE TAB ECHOING, ALLOWING A TAPE
39 TO BE READ WITHOUT CHARACTER LOSS.
40
41 "SET KB NOLSR" WILL ENABLE TAB ECHOING, FOR NORMAL KEYBOARD
42 INPUT. THIS IS THE DEFAULT.
43
44 8)WHEN THE HANDLER RECEIVES A READ REQUEST, A "*" CHARACTER IS
45 PRINTED IN THE LEFT MARGIN OF THE TERMINAL TO SIGNIFY THAT THE
46 HANDLER IS READY FOR INPUT. THIS CHARACTER CAN BE CHANGED, OR THE
47 PROMPT FEATURE CAN BE REMOVED, BY REASSIGNING THE SYMBOL
48 "PROMPT" TO THE ASCII VALUE OF THE
49 DESIRED CHARACTER. SETTING PROMPT TO "0" WILL CAUSE NO CHARACTER
50 TO BE PRINTED.
51
52 9)IF NO READ REQUEST IS ACTIVE, THE HANDLER WILL NOT ACCEPT INPUT,
53 AND THE KEYBOARD WILL NOT ECHO. IF IT DOES ECHO, THE HANDLER IS
54 ACCEPTING INPUT.
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1 )THE HANDLER IS INSTALLED VIA THE FOLLOWING PROCEDURES:
2 )
3 ) 1) ASSEMBLE IT AS FOLLOWS:
4 ) DEFINE FILLER CONDITIONALS IF NECESSARY
5 ) .R MACRO
6 ) *KR=KB
7 )
8 ) 2) LINK IT AS FOLLOWS:
9 ) .R LINK
10 ) *KB.SYS=KB
11 )
12 ) 3) INSTALL IT AS DEVICE "KB", AS DESCRIBED IN SECTION XXXXX
13 ) OF THE RT-11 V2 SOFTWARE SUPPORT MANUAL. REMEMBER THAT
14 ) THE VECTORS FOR THE TERMINAL, MUST BE PROTECTED IN THE BIT MAP
15 ) AS DESCRIBED IN THAT SECTION.
16 ) THE VALUES FOR THE VARIOUS TABLE ENTRIES SHOULD BE
17 ) HSIZ=VALUP OF SYMBOL "KBSIZE" ON LAST LINE OF LISTING
18 ) DVSIZ=0 (NON-FILE STRUCTURED DEVICE)
19 ) PNAME=4242M (RAD50 FOR "KB ")
20 ) STATE HIGH ORDER BYTE=0, LOW ORDER BYTE=ANY DEVICE NUMBER
21 ) AVAILABLE. NOTE THAT IT CANNOT BE 4, A VALUE >50
22 ) IS RECOMMENDED.
23 )
24 ) 4) ONCE INSTALLED, KB: WILL BE AVAILABLE WHEN THE SYSTEM IS REBOOTED.
25 )
26 ) THE HANDLER ITSELF IS ACTIVATED WITH READ AND WRITE REQUESTS, AS ARE ALL
27 ) RT-11 DEVICE HANDLERS. WHEN USING SYSTEM PROGRAMS WHICH OPERATE ON
28 ) LARGE BUFFERS, SEVERAL LINES MAY ACCUMULATE IN THE BUFFER BEFORE
29 ) THEY APPEAR ON THE TERMINAL, AND THEN ALL AT ONCE. TO AVOID THIS PROBLEM,
30 ) EACH OUTPUT BUFFER CAN BE ZERO-FILLED AND SENT TO THE TERMINAL TO PRINT
31 ) EACH LINE-THE HANDLER WILL IGNORE NULLS ON OUTPUT.
32 ) IN FORTRAN, EACH LINE CAN BE FORCED IN OR OUT BY USING A REWIND
33 ) FOLLOWING EACH READ OR WRITE TO THE DEVICE. FOR EXAMPLE:
34 ) LOGICAL*1 INPUTL(80)
35 ) CALL ASSIGN (7,'KB1/C')
36 ) WRITE (7,1)
37 ) REWIND 7
38 ) WRITE (7,2)
39 ) REWIND 7
40 ) READ (7,3) INPUTL
41 ) REWIND 7
42 )
43 )
44 )
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;MCALL :RPGDEF,..V2.....INTEN
;RPGDEF
;..V2...

;VECTOR AND DEVICE REGISTER ADDRESSES-EDIT THESE TWO TO RECONFIGURE
;KEYBOARD VECTOR
;KEYBOARD CONTROL REGISTER

000300
174500
;OTHER DEVICE ADDRESSES
;TRVEC=KBVEC+4
;KBUP=KBCCR+2
;TPCSR=KBCCR+4
;TPBUP=KBCCR+6

;CONSTANTS
;OFFSET=270
;OFFSET TO MONITOR COMPLETION ENTRY

;EOF BIT IN CSM
;PR=340
;PR=200
;IPSN VALUE FOR PRIORITY 7
;IPSN VALUE FOR PRIORITY 4

;TAB
;LTNE FEED
;FORN FEED
;CARRIAGE RETURN=15
;CTRL/U
;CTRL/2
;SPACE
;RUBOUT

;LENGTH OF ECHO BUFFER
;PROMPT='>'

;SFT LSR CODE
;THE FOLLOWING IN THE HANDLER INTERFACE TO THE MONITOR SET_COMMAND.
;FOR DETAILS OF INTERFACING TO THE SFT COMMAND,SEE THE RT-11 V2 SOFTWARE
;SUPPORT, MANUAL
;ASECT
;=400
;HT
;RAD50 /LSR /
;WORD
;=40PLSR-400>/2+100000
;
;OPLSR:
;MOV (PC)+,R3
;MOV LSR,SFT LSR0PT TO 377
;MOV R3,LSR0PT
;RTS PC
;CSECT

000270
020000
000340
000200
000011
000012
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000015
000025
000032
000040
000177
000024
000076
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000400
000011
000400
047012
100005
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012703
000377
010367
000640
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1 000000 000304
2 000002 000116
3 000004 000340
4 000006 000000
5 000010 000000
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000012 010700
000014 062700
000020 010037
012737 000300 +000302
000024 012737
000032 005067
000036 014705
000042 022525
000044 011567
000050 012567
000054 006315
000056 011567
000062 001466
000064 103420
000066 011504
000070 014505
000072 105025
000074 004304
000076 001375
000100 105267
000104 004167
000110 0076
000112 000167
000246
000300
000340 +000302
000702
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000670
000654

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010700
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010037
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000300 +000302
005067
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THIS IS THE HANDLER HEADER AREA, USED BY FETCH AND THE QUEUE MANAGER TO STORE VARIABLES CRITICAL TO HANDLER OPERATION.

KBSTRT: FWORD TPVEC  
 FWORD TPINT.  
 FWORD PR7  
 FWORD 0  
 FWORD 0

ILAST QUEUE ENTRY  
 ICURRENT QUEUE ENTRY

FOLLOWING IS THE TRANSFER INITIATION CODE:

IF THE FIRST WORD OF THIS ROUTINE IS THE ENTRY POINT FOR ALL TRANSFER REQUESTS, THE KEYBOARD VECTOR IS SET UP (FETCH ONLY SETS UP THE PRINTER VECTOR), AND THE PARAMETERS FOR THE TRANSFER ARE ESTABLISHED.

IF THE REQUEST IS A WRITE, CONTROL TRANSFERS TO THE PRINTER ROUTINE TO OUTPUT THE FIRST CHARACTER FROM THE USER BUFFER. IF IT IS A READ, THE ENTIRE USER BUFFER IS ZEROED, A FLAG (READFL) IS SET TO SHOW THAT A READ IS IN PROGRESS, AND A PROMPT CHARACTER IS ECHOED FOR THE TERMINAL BEFORE THE KEYBOARD INTERRUPT IS ENABLED.

PC,R0  
 #KBINT-.,R0  
 R0,#KBVEC  
 #PR7,#KBKRVFC+2  
 READFL  
 #R5)+,(R5)+  
 (R5),UBPTR  
 (R5)+,UBPTR1  
 (R5),BYCNT  
 DONE  
 TPOUT2  
 (R5),R4  
 -(R5),R5  
 (R5)+  
 R4  
 DEC  
 RNF  
 INCB  
 JSR  
 .BYTE  
 JMP

ICALCULATE ABSOLUTE ADDRESS OF KEYBOARD INTERRUPT SERVICE  
 ISET UP KEYBOARD VECTOR  
 IINIT READ FLAG AND TAB COUNT  
 IPPOINT TO CURRENT ELEMENT  
 IADD 4 TO R5  
 ISET UP POINTER TO USER BUFFER FOR LATER  
 IAND SAVE ORIGINAL POINTER INTO BYTE COUNT  
 IMAKE WORD COUNT INTO BYTE COUNT  
 IAND SAVE IT  
 IWORD COUNT OF 0 IS SEEK, WHICH IS NOP IN THIS HANDLER  
 IF NEGATIVE, WRITE TO PRINTER  
 IBYTE COUNT TO R4  
 IUSER BUFFER POINTER IN R5  
 IZERO USER BUFFER BEFORE STARTING TRANSFER

IF R5 IF NOT DONE  
 ISET "READ IN PROGRESS" FLAG  
 IPROMPT INPUT WITH "B"  
 IENABLE KEYBOARD INTERRUPT AND RETURN

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000120 004577 000622
000124 000140
000126 03P737 000200 176504 TPOUT21 REG
000134 001436

000136 010705
000140 062705 000550
000144 111504
000146 001410
000150 012746 000024
000154 114525 000001
000160 005316
000162 003374
000164 005726
000166 000412

000170 104767 000544
000174 001016
000176 117704
000202 002267 000540
000206 005267 000534
000212 003012 000524
000214 104704

THIS IS THE ABORT ENTRY POINT-THE HANDLER IS ENTERED AT THIS ADDRESS
IF THE MONITOR RECEIVES A REQUEST TO ABORT ANY TO TRANSFER IN PROGRESS
RR
ABORT

THIS IS THE TERMINAL OUTPUT INTERRUPT SERVICE. AFTER ENTERING SYSTEM STATE,
IT DETERMINES IF THERE ARE ANY CHARACTERS IN THE ECHO BUFFER TO BE
PRINTED. IF NOT, IT THEN DETERMINES WHETHER A WRITE REQUEST IS IN PROGRESS
FOR NOT. IF SO, THE NEXT CHARACTER IN THE USER BUFFER IS PRINTED.
IF NOT, THE INTERRUPT IS DISMISSED.
IF THERE ARE CHARACTERS IN THE ECHO BUFFER, THE FIRST CHARACTER IN THE
LIST IS FETCHED INTO R4, THE LIST IN THE ECHO BUFFER IS "SLID UP"
BY ONE CHARACTER, AND THE CHARACTER IN R4 IS THEN PRINTED.
IF THE FILLER CONDITIONAL CODE IS INCLUDED AT ASSEMBLY TIME,
IF THE CHARACTER IN R4 IS COMPARED AGAINST THE CHARACTER TO BE FILLED AFTER.
IF THE SAME, A COUNT OF NECESSARY FILLS IS STOPPED IN "FILCN1" AND THE
CHARACTER IS PRINTED. THE INTERRUPT SERVICE THEN CHECKS THE NUMBER
OF FILLS NEEDED AS THE FIRST ITEM, AND PRINTS NULLS IF ANY ARE LEFT
ENTER SYSTEM STATE
R5,0SINPTR
WC<PR0>6PR7
#200,0#TPCSR
RTSPC

IJS THE PRINTER READY
IYES-THEN WAIT FOR INTERRUPT TO PRINT ANYTHING

CONDITIONAL CODE FOR FILLER
IANY FILLS NEED TO BE OUTPUT?
IBRANCH IF NOT
IYES-DECREASE NUMBER BY ONE
INULL IS FILLER
IGO PRINT IT

ICALL ABSOLUTE ADDRESS
IOP ECHO BUFFER
IGET CHAR TO ECHO FROM ECHO BUFFER
IBRANCH IF BUFFER EMPTY
INUMBER OF CHARS IN ECHO BUFFER ON STACK
ISLIDE ECHO LIST UP
IDCREASE COUNT OF CHARS TO SLIDE
IBRANCH IF NOT FINISHED
IDONE-CLEAN UP STACK
IAND PRINT CHAR

IARE WE READING OR WRITING?
IBRANCH IF READING
IGET CHAR FROM USER BUFFER INTO R4
IBUMP BUFFER POINTER
IAND DECREASE TRANSFER COUNT
IBRANCH IF TRANSFER COMPLETE
IDON'T PRINT NULLS
R4

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48 000216 001743
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54 000220 012737 000100 176504 TPUT31
55 000226 110437 176506
56 000232 000207
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61 000234 005037 174504
62 000240 005037 176500
63 000244 010704
64 000246 062704 177542
65 000252 013705 000054
66 000256 000175 000270
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IRREQUEST TERMINATION AND ABORT CODE
ITTHIS ROUTINE IS ENTERED WHEN THE I/O TRANSFER IS
ICOMPLETED OR ABORTED. THE DEVICE INTERRUPTS ARE DISABLED, AND
ISTANDARD MONITOR COMPLETION ENTRY CODE IS EXECUTED.
IABORT: CLR 00TPCSR
IDONE: CLR 00KBCSR
IMOV PC,R4
IADD #RAC0E-.R4
IMOV 0054,R5
IJMP 00FFSET(05)

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11 000262 004577 000460
12 000266 000140
13 000270 113704 176502
14 000274 042704 177600
15 000300 120427 000177
16 000304 001020
17 000306 026767 000430 000430

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IKEYBOARD INTERRUPT SERVICE
ITTHIS IS THE KEYBOARD INTERRUPT SERVICE ROUTINE. AFTER ENTERING
ISYSTEM STATE, IT GETS THE TYPED CHARACTER INTO R4, THEN
IPROCEEDS DOWN A CHAIN OF CHECKS FOR THE SPECIAL CASE CHARACTERS
I(RUBOUT, CTRL U, CTRL Z, CR, FF). IF IT IS ONE OF THE SPECIAL
ICHARACTERS, THE ROUTINE "ECHO" IS CALLED TO ECHO APPROPRIATE
ICHARACTERS ON THE TERMINAL, THEN APPROPRIATE ACTION FOR THE SPECIAL
IIS TAKEN. IF A NORMAL CHARACTER IS TYPED, IT IS ECHOED AND PLACED
IIN THE USER BUFFER BEFORE THE INTERRUPT IS DISMISSED.
IKBINT: JSR R5,0SINPTR
IWORD 0C<PR0>8PRY
IMOV 00KBRUP,R4
IBIC #177600,R4
ICMP R4,#DELETE
IANE 115
ICMP UBPTR,UBPTR1

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18	000314	001520	000420	KBTN	INO=IGNORE RUBOUT
19	000316	000367	000240	UBPTR	IBACK UP POINTER INTO USER BUFFER
20	000322	000167	000240	R1,ECHO	
21	000326	134	377	%,377	
22	000330	100267	000405	TARCNT	IBUMP TAB COUNTER FOR "A"
23	000334	100077	000402	UBPTR	IZERO RUBBED OUT CHAR
24	000340	000267	000372	RYCNT	IAND INCREASE TRANSFER COUNT TO REFLECT LOST CHAR
25	000344	000504		KBTN	IRE=ENABLE INTERRUPTS AND EXIT
26	000346	120427	000014	R4,#FF	IIS THIS CHAR A FORM FEED?
27	000352	001006		AS	IBRANCH IF NOT
28	000354	000167	000206	R1,ECHO	IYES=ECHO Y LINE FEEDS
29	000360	012	012	LF,L,F,L,F,L,F,L,F,L,F,377	
	000363	012	012		
	000366	012	377		
30	000370	120427	000015	R4,#CR	IIS THIS CHAR A CR?
31	000374	001017		78	IBRANCH IF NOT
32	000376	000167	000164	R1,ECHO	IYES=ECHO CR,LP
33	000402	015	012	CR,L,F,0,377	
	000405	377	000		
34	000406	110477	000330	R4,#UBPTR	IPUT CR IN USER BUFFER
35	000412	000267	000324	UBPTR	IBUMP USER BUFFER POINTER
36	000416	000367	000314	RYCNT	IROOM IN BUFFER FOR LF TOO?
37	000422	001706		DONE	IF DON'T INSERT IT IF NOT
38	000424	110777	000012	#LF,#UBPTR	IF ELSE ADD LF TO BUFFER
39	000432	000702	000310	DONE	

1	000434	120427	000025	R4,#CTRLU	IIS CHAR CTRL U?
2	000440	001007		88	IBRANCH IF NOT
3	000442	000167	000120	R1,ECHO	IECHO "SU"
4	000446	136	015	%,0,CR,L,F,0,377	
5	000451	012	000		
6	000454	000167	17352	RETRY	IAND RESTART READ
7	000460	120427	000032	R4,#CTRLZ	IIS CHAR CTRL Z?
8	000464	001013		98	IBRANCH IF NOT
9	000466	000167	000074	R1,ECHO	IECHO "SZ"
10	000472	136	015	%,0,CR,L,F,0,377	
	000475	012	000		
11	000500	010705	17304	KBCGE,R5	IPOINT R5 TO 0 ELEMENT
12	000504	052775	020000	#EOF,0-2(R5)	IAND SET EOF FLAG IN CSW
13	000512	000052	177776	DONE	ISTOP TRANSFER
14					
15	000514	120427	000040	R4,#40	IIS THIS A PRINTING CHAR?
16	000520	002002		218	IBRANCH IF NOT
17	000522	104267	000213	TARCNT	IYES-INCREASE TAB POSITION

```

18 000526 110467 000004
19 000532 004167 000030
20 000536 000 377
21 000540 110477 000176
22 000544 005267 000172
23 000550 005367 000162
24 000554 001631
25 000556 012737 000101 174500 KBTNI
26 000564 000207 RTS

```

18 SET UP TO ECHO CHAR

```

MOV B R4,203
JSR R1,ECHO
LD B,0,377
MOV B,R4,0UBPTR
TNC UBPTR
DEC ANYCNT
NEG NONE
MOV #101,00KRCAR
RTS PC

```

19 PUT CHAR IN USER BUFFER  
20 BUMP BUFFER POINTER  
21 ANY MORE TO TRANSFER  
22 BRANCH IF NOT  
23 DECREMENT KEYBOARD INTERRUPT  
24 RETURN TO MONITOR

```

1
2
3
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25
26
27
28
29
30
31
32
33
34
35

```

```

SUBROUTINE ECHO
THIS SUBROUTINE SERVES TO PLACE THE SPECIFIED CHARACTERS IN THE
ECHO BUFFER AND START THE PRINTER IN CASE IT IS IDLE.
THE CALLING SEQUENCE IS
JSR R1,ECHO
LD B,CHAR1,CHAR2,CHAR3,...,CHARN,377
NOTE THAT THERE MUST BE AN EVEN NUMBER OF BYTES IN THE ARGUMENT LIST
AND THEREFORE THE NUMBER OF CHARACTERS EXCLUDING THE 377
MUST BE ODD.
WHEN ENTERED, ECHO SCANS THE ECHO BUFFER TO FIND THE END OF THE
ECHO LIST, WHICH IS MARKED BY A NULL BYTE. WHEN THE END OF THE LIST
IS FOUND, IT IS DETERMINED IF THERE ARE AT LEAST 8 FREE SLOTS IN THE LIST
TO ACCOMMODATE A POSSIBLE LINE FEED OR FORM FEED. IF NOT, THE
CHARACTER JUST TYPED IS IGNORED. IF SO, THE CHARACTERS FROM THE
ARGUMENT LIST FOLLOWING THE CALL ARE INSERTED IN THE BUFFER.
IF THE PRINTER IS STOPPED IF IT IS IDLE, AND THE ROUTINE RETURNS.
NOTE THAT TAB IS A SPECIAL CASE; IF R4 CONTAINS A TAB CHARACTER
WHEN THIS ROUTINE IS ENTERED, THE ARGUMENT LIST IS NOT USED. RATHER,
AN APPROPRIATE NUMBER OF SPACES TO MOVE THE PRINT HEAD TO THE
NEXT TAB STOP ARE PLACED IN THE ECHO BUFFER, AND THE ROUTINE RETURNS
ENARL LSR
MOV PC,R5
ADD #RSTRT-.1,R5
MOV R5,TEMP
ADD #RLENGTH-1,TEMP
TSTB (R5)+
RNE 43
DEC R5
SUR R5,TEMP
CMP TEMP,#8.
RGT 38
MOV (SP)+,R1
RR KBTN

```

PC,R5 ICALC ABSOLUTE ADDRESS

10P ECHO BUFFER  
15AVE ADDRESS OF ECHO BUFFER  
TEMP POINTS TO END OF ECHO BUFFER  
IS THIS END OF ECHO LIST?

BRANCH IF NOT  
YES-R5 POINTS TO FIRST FREE SLOT IN ECHO LIST  
IF END NUMBER OF FREE SLOTS IN ECHO LIST  
IS THERE ENOUGH ROOM TO ECHO TAB OR FF?  
BRANCH IF YES  
NO-IGNORE THIS CHAR THEN  
DTSMISS INTERRUPT

```

36 000634 010446
37 000636 120427
38 000640 000011
39 000642 001013
40 000644 110725
41 000650 100267
42 000654 130767
43 000662 001370
44 000664 000721
45 000666 100015
46 000670 000403
47 000672 110125
48 000674 100376
49 000676 100045
50 000700 000767
51 000704 010004
52 000706 000201
53
54

```

R4,=(SP)  
R4,(PC)+  
IS  
#SPACE,(PS)+  
TARCNT  
#7,TABCNT  
58  
(RI)+  
(RM)  
AS  
(RI)+,(RM)+  
IS  
-(R5)  
PC,TPOUT2  
(SP)+,R4  
R1  
DSARL LSR

I SAVE CHAR  
I IS THIS CHAR A TAB?
I THIS COMPARE OPERAND CAN BE CHANGED BY SET LSR
I BRANCH IF NOT
I ECHO A SPACE
I BUMP POSITION COUNTER
I AT TAB STOP YET?
I BRANCH IF NOT
I YES-ARTIFICIALLY BUMP RETURN
I END ECHO LIST
I AND START ECHO
I MOVE CHAR INTO ECHO LIST
I BRANCH IF END-OF-LIST NOT SEEN
I ELSE USE 0 TO MARK END OF ECHO LIST
I PRINT A CHAR TO START PRINTER
I RESTORE CHAR
I RETURN

KB.MAC V01=01 RT=11 MACRO VM02=09 8-APR-75 12333151 PAGE 10

```

1
2
3
4 000710 000
5
6
7
8
9
10
11
12
13 000734 000000
14 000736 000000
15 000740 000
16 000741 000
17 000742 000000
18 000744 000000
19
20
21 000746 000000
22
23 000750
24 000001

```

I DATA AREA
I ECHO RING BUFFER=ELENGTH CHARACTERS LONG
R0STR1 ,BYTE 0
,BLKR ELENGTH=1
I VARIABLE AREA
I POF FILCHR
I BYTE FILCHR
I BYTE 0
I ENDC 0
I WORD 0
I WORD 0
I WORD 0
I BYTE 0
I BYTE 0
I WORD 0
I WORD 0
I MONITOR SYSTEM STATE ENTRY LINK
SINPTR1 ,WORD 0
KB SIZE=,KBSTR
END

I FILLER CONDITIONAL
I CHARACTER TO BE FILLED AFTER
I NUMBER OF FILLS REMAINING
I TEMPORARY
I USER TRANSFER COUNT
I FLAG FOR "READ IN PROGRESS"
I TAB POSITION COUNTER
I POINTER INTO USER BUFFER
I POINTER TO START OF USER BUFFER

KB,MAC V01-01 RT-11 MACRO VM02-09 8-APR-75 12133151 PAGE 10+

SYMBOL TABLE

ABORT 000234R  
 DELET = 000177  
 FF = 000014  
 KBIN = 000556R  
 KBEVC = 000300  
 PC =X000007  
 READFL 000740R  
 R2 =X000002  
 SPACE = 000040  
 TPINT 000120R  
 TPVEC = 000304  
 BYTCNT 000736R  
 DONE 000240R  
 HT = 000011  
 KBINT = 000262R  
 LF = 000012  
 PROMPT= 000076  
 RETRY 000032R  
 R3 =X000003  
 TARCNT 000741R  
 TPOUT 000176R  
 UBPTR 000742R  
 CR = 000015  
 EBLNG= 000024  
 KBRUF = 176502  
 KBLDE 000006R  
 LSRORT 000640R  
 PR4 = 000200  
 RTSPC =X000004  
 R4 =X000004  
 TEMP 000734R  
 TPOUT1 000214R  
 UBPTM1 000744R  
 CTRLU = 000025  
 ECHO 000566R  
 KBCE 000010R  
 KBSIZE= 000750  
 OFFSET= 000270  
 PR7 = 000340  
 R0 =X000000  
 R5 =X000005  
 TPAUF = 176506  
 TPOUT2 000126R  
 SINPTR 000746R  
 CTRLZ = 000032  
 POP = 020000  
 KBCSR = 176500  
 KBCSTR 000000R  
 OPLSR 000412  
 RBSTR 000710R  
 R1 =X000001  
 SP =X000006  
 TPCSR = 176504  
 TPOUT3 000220R  
 ...V2 = 000001

. ABS. 000424 000  
 000750 001  
 ERRORS DETECTED: 0  
 FREE CORE: 15460. WORDS

KB,LPI/NITM/CHK



APPENDIX C

VERSION 1 EMT SUMMARY

Although Version 1 programmed requests are supported by Versions 2, 2B, and 2C of RT-11, it is strongly recommended that the Version 1 formats not be used. For purposes of compatibility, however, this section provides a brief review of the V1 format. The V2/V2B/V2C format is covered in detail in Chapter 9 of the RT-11 System Reference Manual.

In brief, the major distinctions between V1 and V2/V2B formats are:

1. V1 format has arguments pushed on the stack and in R0. V2/V2B/V2C requests generally accept a set of arguments, or an argument in R0.
2. V1 channel numbers are restricted to  $16_{10}$ . Also, the channel number in V1 is not a legal assembler argument; it is merely an integer in the range 0 to  $15_{10}$ .
3. V1 requests are non-reentrant because the channel number and function code are embedded within the EMT instruction.

Table C-1 lists all the Version 1 macro calls. Those in the left column have the same format as the corresponding Version 2/2B/2C request; those in the right column have a different format, shown after the table. The operations performed by the requests are the same in both versions.

Table C-1  
V1 Programmed Requests

V1 - Format Same as V2/V2B	V1 - Format Different from V2/V2B/V2C
.CSIGEN	.CLOSE
.CSISPC	.DELETE
.DATE	.ENTER
.DSTAT	.LOOKUP
.EXIT	.READ
.FETCH	.READC

(continued on next page)



```
.WAIT .chan
```

```
.WRITE }  
.WRITC } .chan,.buff,.wcnt,.crtn,.blk [ .crtn is required ]  
.WRITW } [ only for .WRITC ]
```

The system macro library (SYSMAC.SML) can be used with Versions 2 and 2B to generate Version 1 programmed requests.

Under Version 2, the `..V2..` macro is capable of handling V1 expansions. `..V2..` normally expands as:

```
.MCALL ...CM1,...CM2,...CM3,...CM4  
...V2=1
```

This causes Version 2 expansions in all cases. To allow expansion of all V1 requests in their V1 format (and all new Version 2 requests in V2 format) the `..V2..` macro should not be called, but the utility macros must still be defined:

```
.MCALL ...CM1,...CM2,...CM3,...CM4
```

Omitting both `..V2..` and the utility macros causes all old V1 requests to be expanded in V1 format; no V2 requests can be used.

Under Version 2B, the `..V1..` macro call enables expansion of all macros in Version 1 format. `..V1..` expands as:

```
...V1=1
```

To enable expansion of all Version 1 macros in V1 format and all new Version 2 macros in V2 format, these statements must be included:

```
.MCALL ..V1...CM1,...CM2,...CM3,...CM4  
..V1..
```

A listing of SYSMAC.SML is provided in the RT-11 System Reference Manual.



APPENDIX D  
FOREGROUND SPOOLER EXAMPLE

The following program is an example of a line printer spooler for the foreground. Instructions for its use follow.

1. Create the program using the Editor and store it on the system device under the name LSPOOL.MAC.
2. Next assemble it under MACRO and then link it to create the REL format output file:

```
.R MACRO  
*LSPOOL=LSPOOL  
  
.R LINK  
*LSPOOL=LSPOOL/R
```

3. Load the necessary handlers (in this case, LP and RF) and run the program. All files on device RF with the extension .LST are listed on the line printer and then deleted from RF:

```
.LOA LP,RF<CR>  
.FRU LSPOOL<CR>  
  
F>  
DEVICE TO SPOOL?  
  
B>  
.  
  
F>  
RF:*.LST<CR>
```

[Control must be redirected  
to the foreground via ^F.]

This program assumes device DK: and extension .LPT unless otherwise indicated.

```

1  .TITLE LSPPOOL - LINE PRINTER SPOOLER
2  .SYTL A USEFUL FOREGROUND PROGRAM
3
4
5  ; THIS PROGRAM FOR THE FOREGROUND IS A LINE PRINTER SPOOLER.
6  ; IT SEARCHES A SPECIFIED DEVICE FOR FILPS WITH A PARTICULAR
7  ; EXTENSION (THE DEFAULT IS LPT) AND PRINTS THEM, DELETING
8  ; AFTER PRINTING. IF NONE ARE FOUND, IT WILL GO TO SLEEP FOR
9  ; HALF A MNIUTE, PERMITTING THE BACKGROUND TO RUN.
10 ;
11 ; TO RUN LSPPOOL, FIRST LOAD LP HANDLER AND INPUT DEVICE HANDLER
12 ; IF IT IS NOT THE SYSTEM DEVICE TYPE.
13 ;
14 ; F.G.,
15 ;
16 ; .LOA LP,RF
17 ; .FRU LSPPOOL
18 ;
19 ; LSPPOOL WILL TYPE: "DEVICF TO SPOOL?"
20 ; TYPE INPUT DEVICE AND FILE DESCRIPTION, F.G.:
21 ;
22 ; RF:*.LST
23
24 .MCALE .V2...REGDEF
25 .MCALE .READW,.WRTTW,.LOOKUP,.DELTFE,.CSISPC,.TTYIN
26 .MCALE .PRINT,.TTYOUT,.SRESET,.PCTRLD,.CLOSE,.EXIT
27 .MCALE .DSTATUS,.TWAIT
28
29 .V2...
30 .REGDEF
31
32 USPSWP = 46 ;USR SWAP LOCATION POINTER
33 FRPBYT = 52 ;ERROR CODE
34 CR = 15 ;CARRIAGE RETURN
35 LF = 12 ;LTFE FEED
36
37 000000 012737 001260 000046 START: MOV #BUFF,#USPSWP ;MAKE USP SWAP OVER RUFF
38 000006 016627 000000 MOV SP,(PC)+ ;SAVE STACK POINTER FOR RESET
39 000010 000000 .WORD 0
40 000012 103403 .DSTATUS #IOR,#LP ;MUST RE IN MEMORY.
41 000024 005767 000750 PCS 15 ;ILLEGAL DEVICE
42 000026 001011 TST TOR+4 ;TEST ENTRY POINT
43 000032 000034 RNF BEGIN ;RP TO BEGIN IF LOADED
44 000034 .PRINT #MSG0 ;LP NOT IN MEMORY!

```

```

45 000042
46
47
48 000044
49 000052 01A706 177732
50
51 000056
52 000070
53 000072
54 000100 012702 001142
55 000104 010201
56 000106
57 000112 022700 000015

.EXIT
; COME HERE ON BAD COMMAND STRING
RADCOM: .PRINT #MSG2
MOV STKSAV,SP

REGIN: .CLOSE #0
.RCTRLN
.PRINT #MSG1
MOV #CSIRLK,R2
MOV R2,R1
.TTYTN
CMP #CR,R0

;BACK TO USER FOR A LOAD LP
;PRINT ERROR MESSAGE
;PUSH STACK, FALL THRU TO BEGIN
;WE WILL USE CH 0, SO CLEAN IT UP
;RSET CTRL/O FLAG SO
;PROMPTING MSG WILL PRINT.
;POINT TO COMMAND STRING BUFFER
;COPY THE POINTER AND INPUT COMMAND
;A CHARACTER AT A TIME.
;CARRIAGE RETURN?

```

LSPOOL - LINE PRINTER SPOOLER RT-11 MACRO VM02-11 26-NOV-75 00:07:21 PAGE 1+

```

58 000116 001773
59 000120 110022
60 000122 122700 000012
61 000126 001367
62 000130 105042
63 000132
64 000144 012600
65 000146 001336
66 000150 103735
67 000152
68 000204 014700 000776
69 000210 001000
70 000212 012700
71 000214 046624
72 000216 010067 000244
73 000222 014700 000752
74 000226 001002
75 000230 012700
76 000232 015326
77 000234 010011
78 000236 005061 000002
79
80 000242
81 000252 103403
82 000254 005767 000522
83 000260 001004
84 000262
85 000270 000672

REG
R0,(R2)+
#LF,R0
;
-(R2)
.CSISPC R1,#DEFEXT,R1
MOV (SP)+,R0
RNF RADCOM
RCS RADCOM
.LOOKUP #IOB,#1,#LP
MOV CSTRLK+36,,R0
RNF #3
MOV (PC)+,R0
.RAD50 /LPT/
MOV R0,LPT...
MOV CSTBLK+30,,R0
RNF #4
MOV (PC)+,R0
.RAD50 /DK0/
MOV R0,R01
CLR P(R1)

;STATUS #TOR,R1
RCS #5
TST TOR+4
RNF #6
.PRINT #MSG3
RR

;YES, IGNORE IT.
;MOVE IT INTO BUFFER AND
;TEST FOR END OF LINE.
;NO, GET ANOTHER CHARACTER.
;YES, CLEAR OUT THE LINE FEED
;PROCESS THE COMMAND
;TEST # OF SWITCHES;C BIT UNCHANGED.
;NO SWITCHES ALLOWED
;SYNTAX ERROR

;GET FILE EXTENSION TO PRINT.
;BRANCH IF USER SPECIFIED,
;ELSE USE LPT EXTENSION

;SAVE THE EXTENSION FOR LATER.
;GET THE INPUT DEVICE NAME
;BRANCH IF USER SPECIFIED,
;ELSE USE THE
;DEFAULT DEVICE.
;SET DEVICE NAME IN FILE
;DESCRIPTOR BLOCK, CLEARING
;OUT ANY FILE NAME.
;INPUT DEVICE HANDLER MUST BE RESIDENT
;ILLEGAL DEVICE
;TEST ENTRY POINT
;BRANCH IF O.K.,
;ELSE PRINT MESSAGE

```





```

127 000630      .WRITM #108,#1,R2,R4,R5  AND WRITE IT TO CHANNEL 1
128 000672      TST (R5)+
129 000674      RR 15
130 000676      TSTB #ERRRBYT
131 000702      REO #5
132 000704      .CLOSE #2
133 000716      .PRINT #ERRIN
134 000724      JMP FINDLP
135 000730      .CLOSE #2
136 000742      .DELETE #108,#3,R1
137 000772      RR
138
139 000774      LP: .RAD50 /LP0/
140 000776      TOR: .BLKW 5
141
142 001010      RIN
143 001016      .ASCTZ /NO LP/
144 001037      .ASCTZ /DEVICF TO SPOOL?/
145 001051      .ASCTZ /TRY AGAIN/
146 001064      .ASCTZ /DEVICE?/
147 001111      .ASCTZ /ERROR READING DIRECTORY/
148
149
150 001126      .LIST RIN
151 001132      MSG0: 000000
152 001134      MSG1: 04A624
153 001142      MSG2: 000000
154 001260      MSG3: 000000
155

```

LSPOLL LINE PRINTER SPOOLER RT-11 MACRO VM02-11 26-NOV-75 00:07:21 PAGE 1+

```

SYMBOL TABLE
RADCOM 000044R
CR      000015
FINDLP 000322R
MSG0    001010R
PC      =X000007
R4      =X000004
TIMBLK 001126R
. ABS.  000000
        011260
ERRORS DETECTED: 0
FREE CORE: 15035. WORDS
.LP:/N:TTM/C=LSP001.MAC

```



APPENDIX E  
S/J AND F/B MONITOR FLOWCHARTS

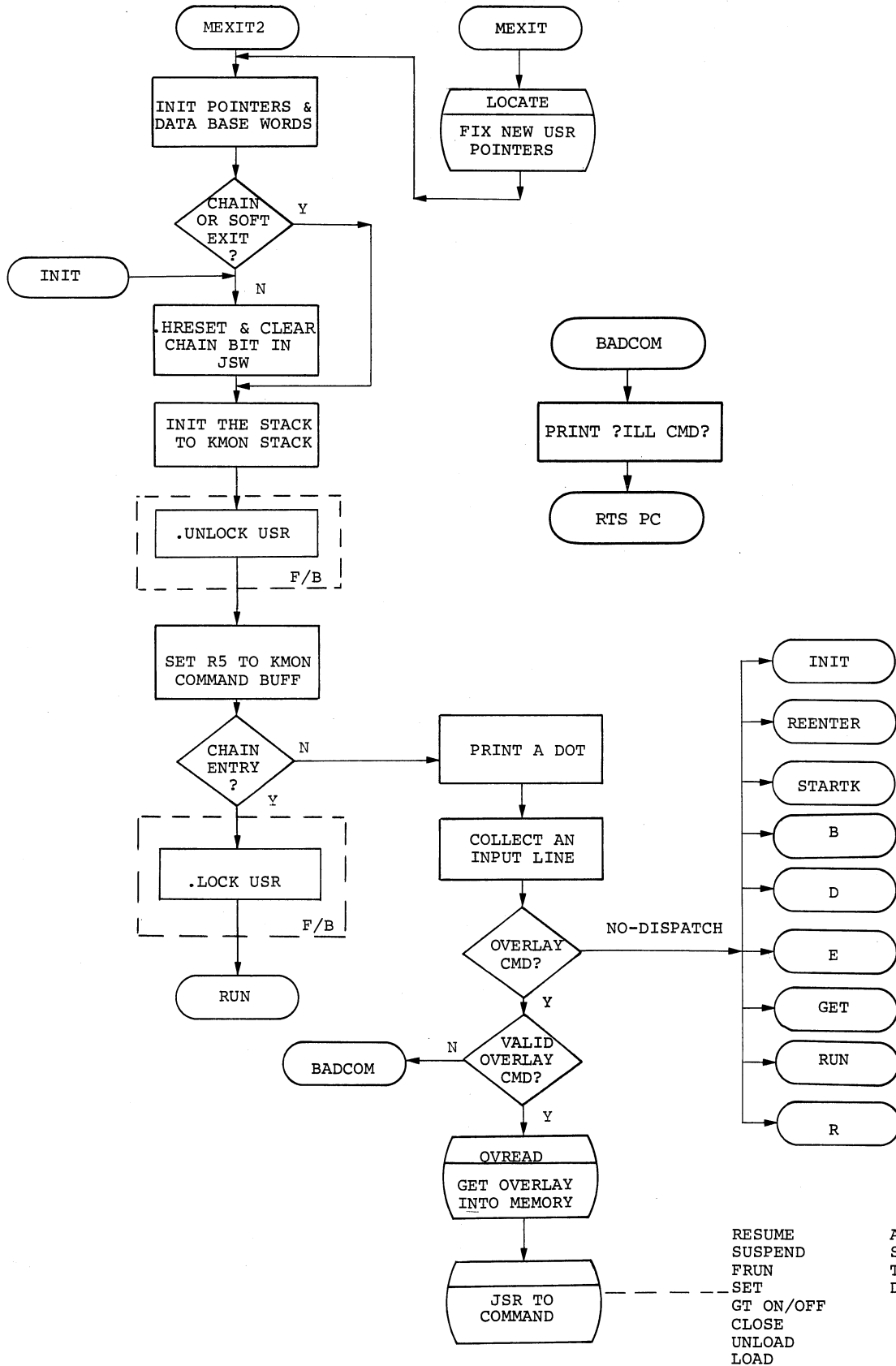
The following flowcharts are of the Single-Job and Foreground/  
Background Monitors. It is recommended that the reader have source  
listings available for reference. Steps inside    are per-  
formed only in the F/B or S/J Monitor, as noted.

An index of all entry points appears at the end of the appendix.

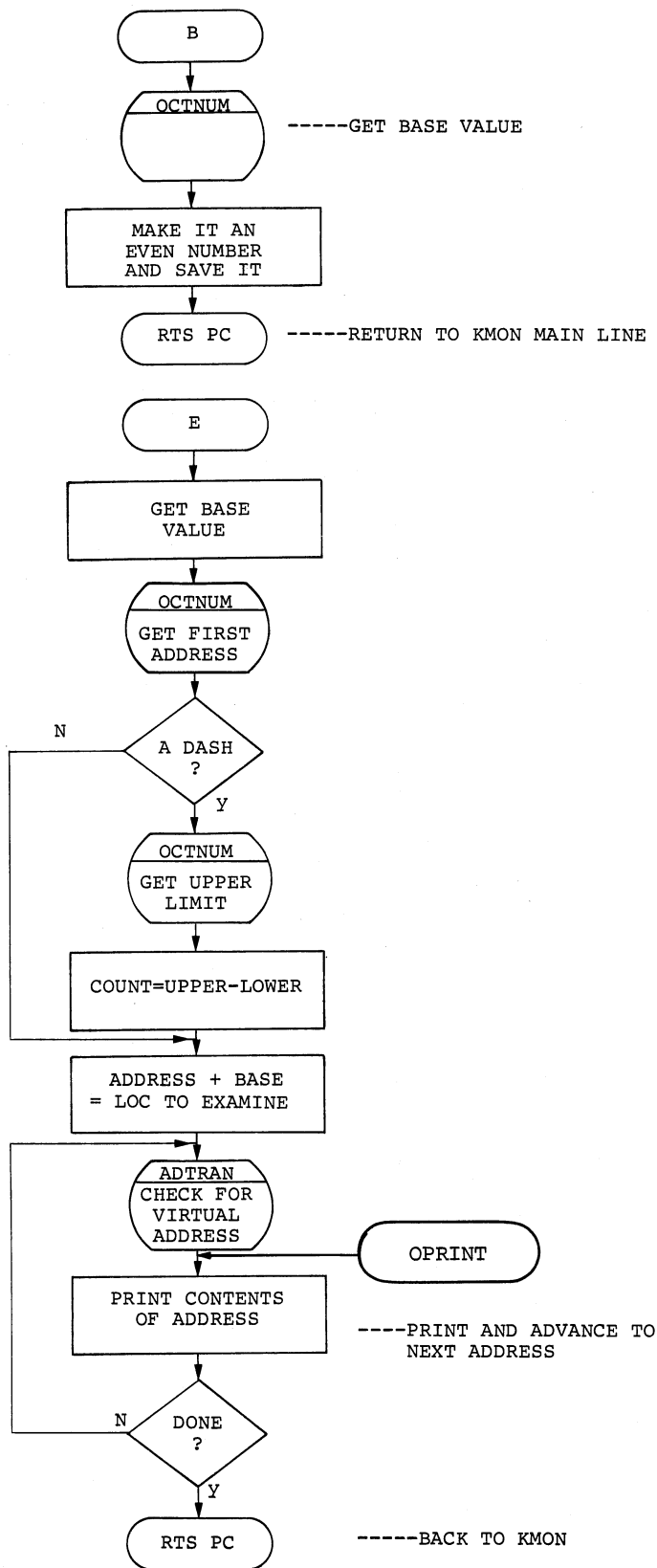


E.1 KMON (KEYBOARD MONITOR) FLOWCHARTS

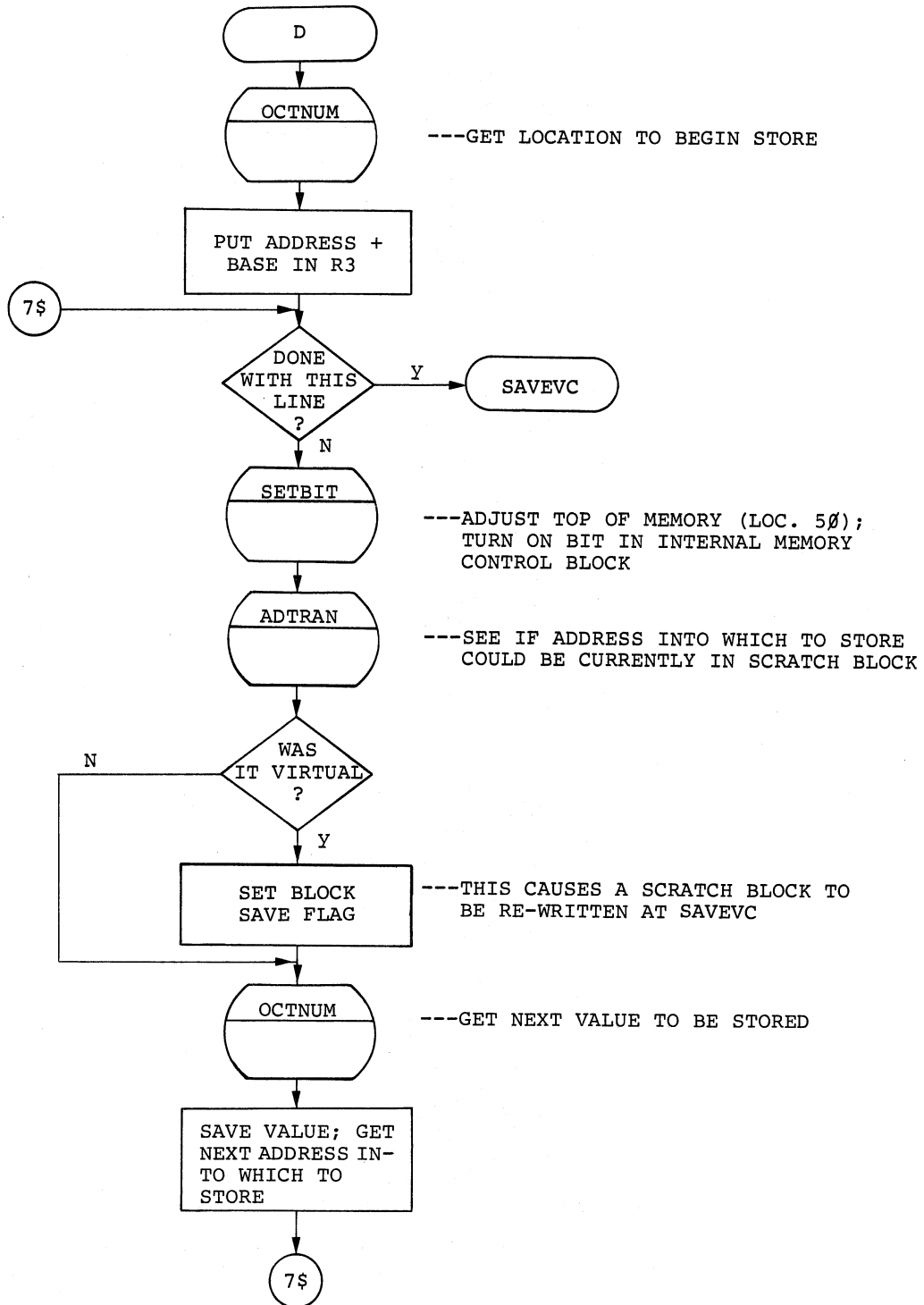
KMON



BASE/EXAMINE

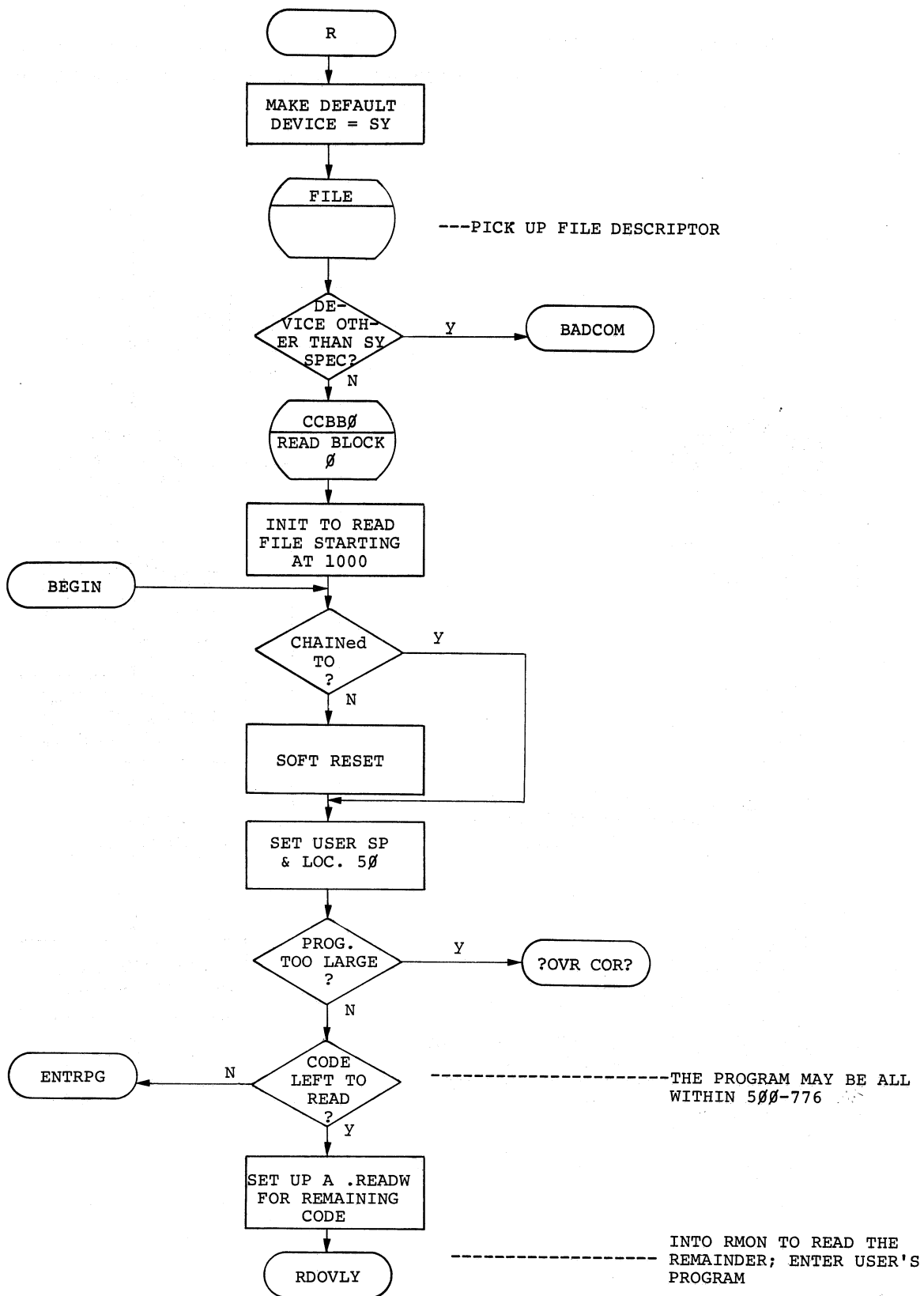


DEPOSIT



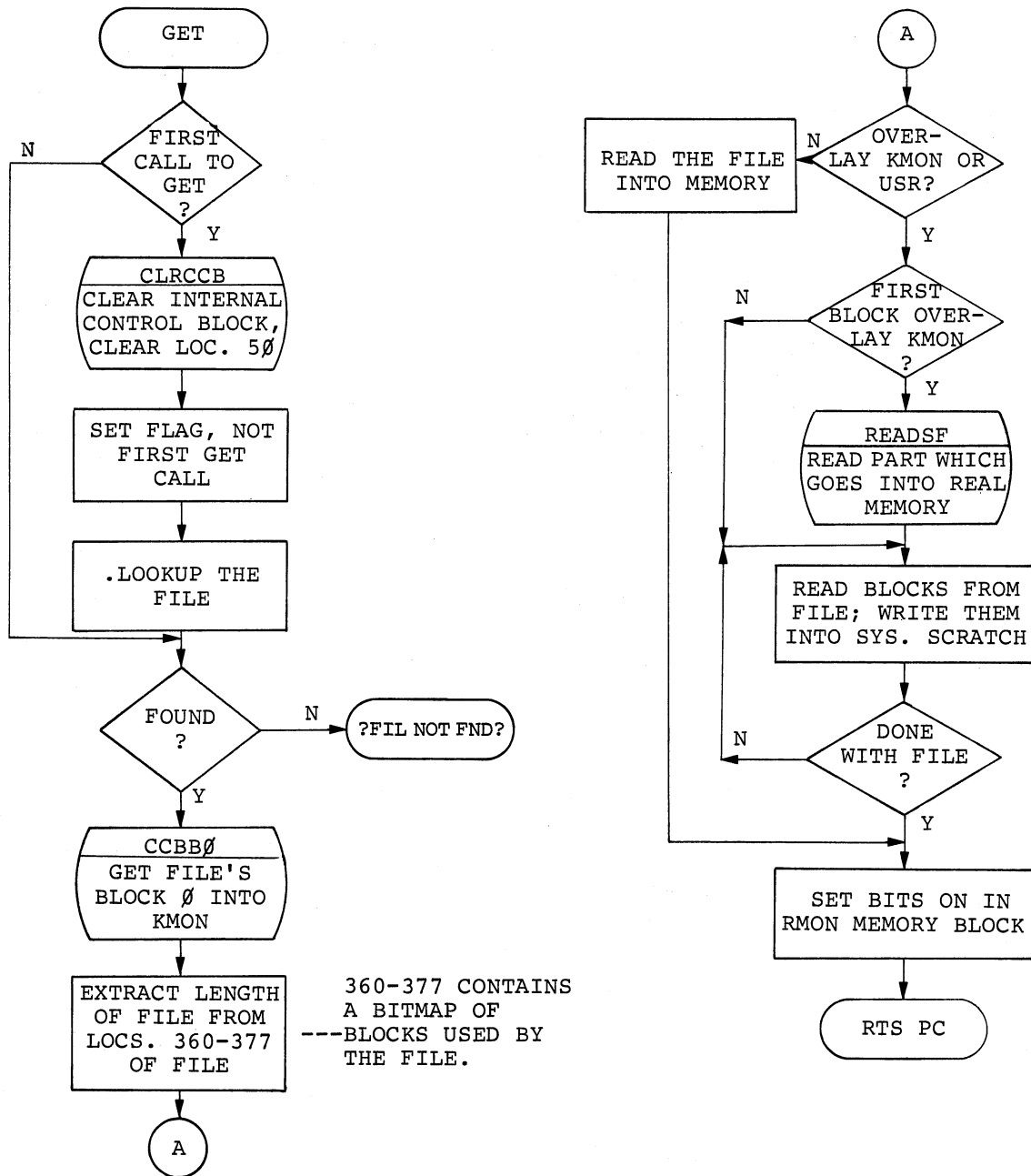
SAVEVC - Entered to rewrite the current virtual block back into the system scratch area. It also acts as the exit point for Deposit; The RTS PC will return control to KMON.



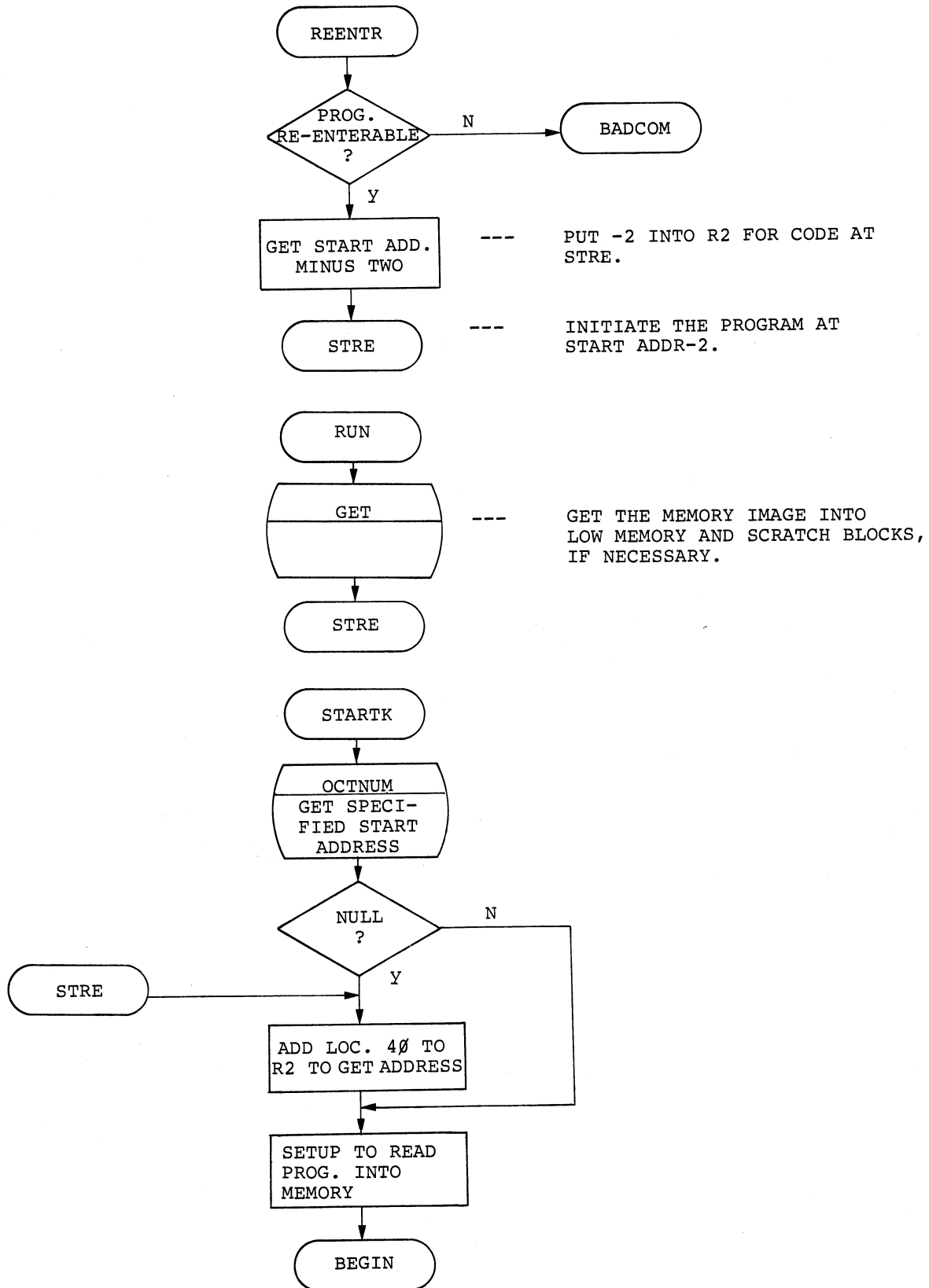


GET

GET - Used to load a .SAV image into memory. If parts of the file overlay KMON/USR, those parts are placed into system scratch blocks.



REENTER/RUN/START

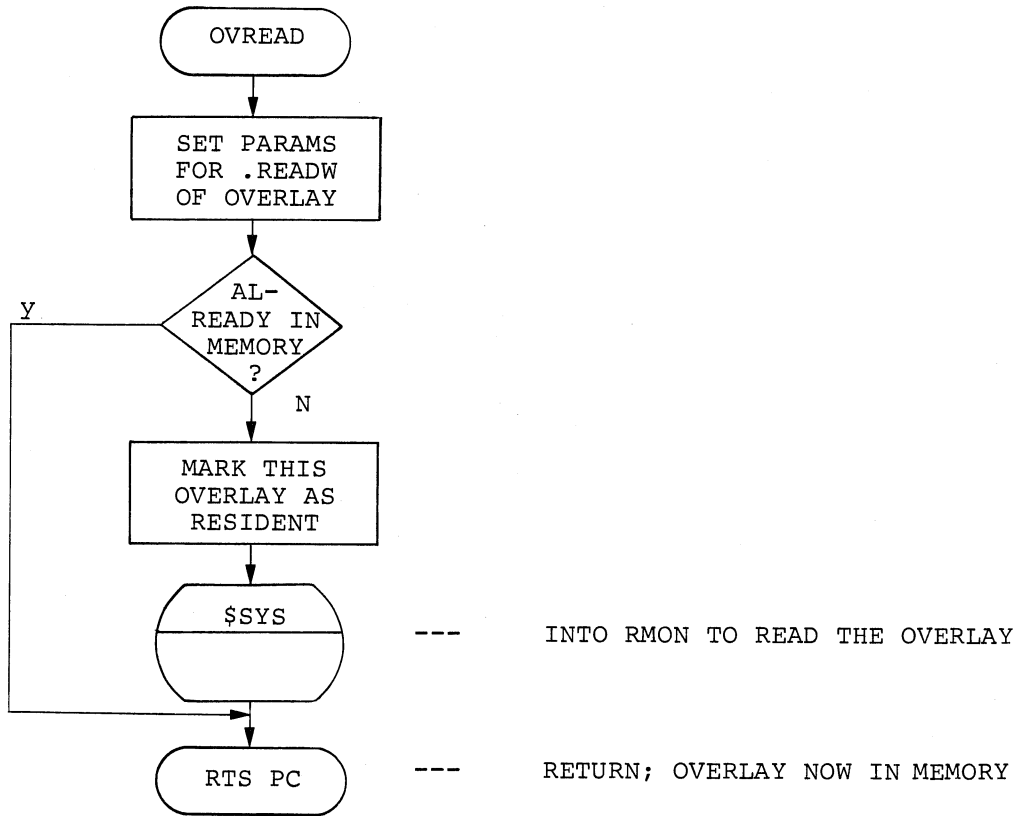




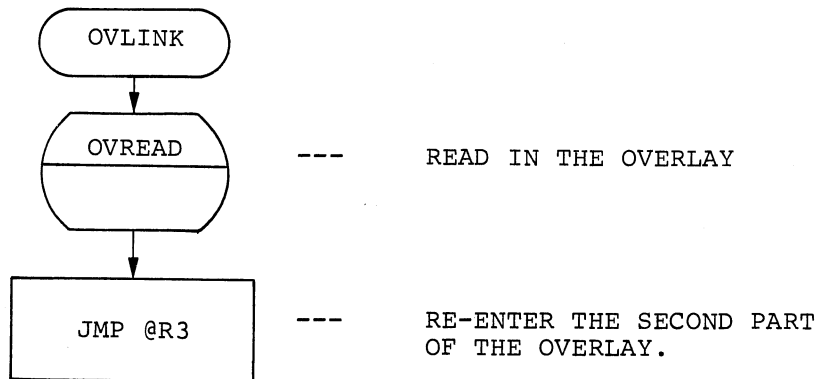
E.1.1 KMON Subroutines

OVREAD/OVLINK

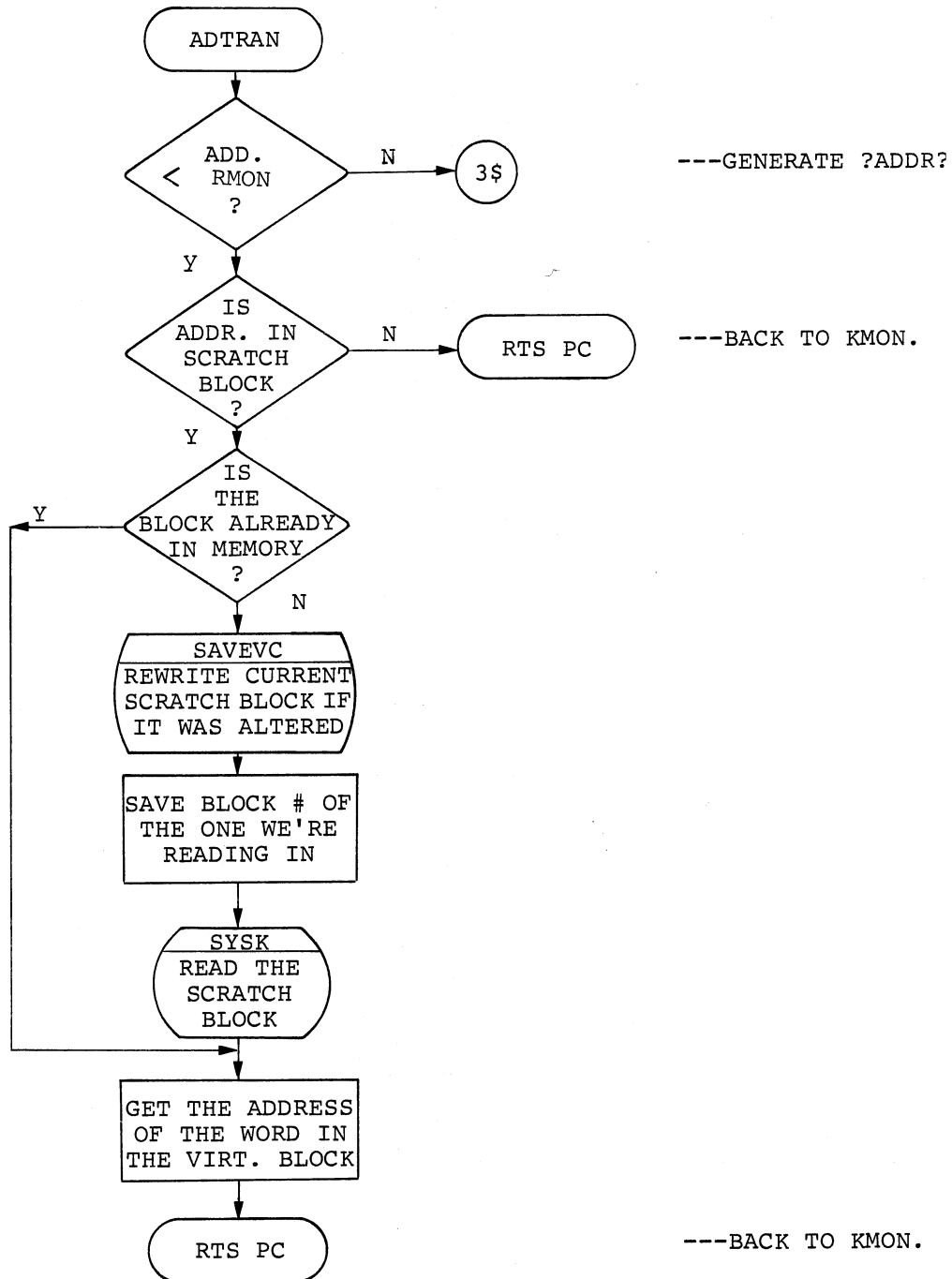
OVREAD - Used to read overlay command processors into memory.



OVLINK - Called from overlay processors to allow linking from one overlay to the other.

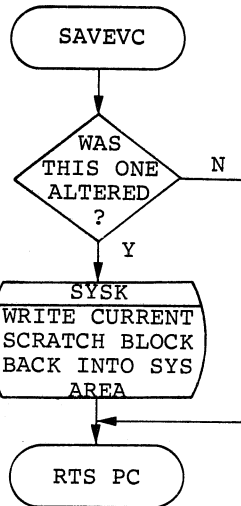


ADTRAN - Used to determine if a user-typed address is a) legal (i.e., address of RMON), b) in scratch blocks on system device.

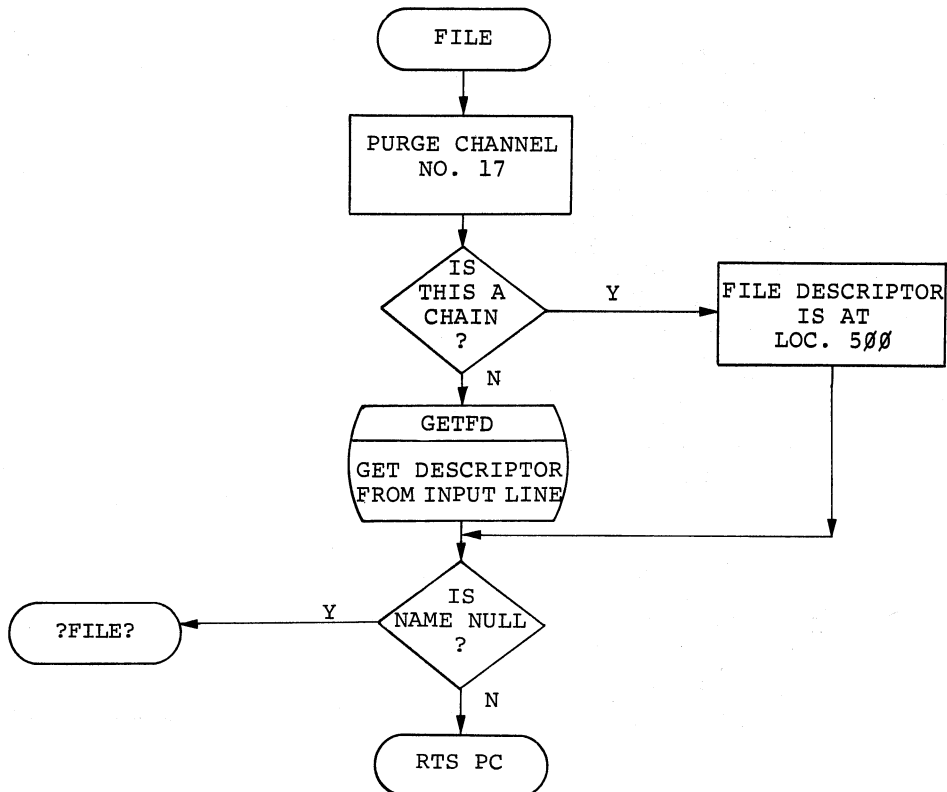


SAVEVC/FILE

SAVEVC - Rewrites a block of memory back to the system scratch area if the block's contents were altered with a Deposit.

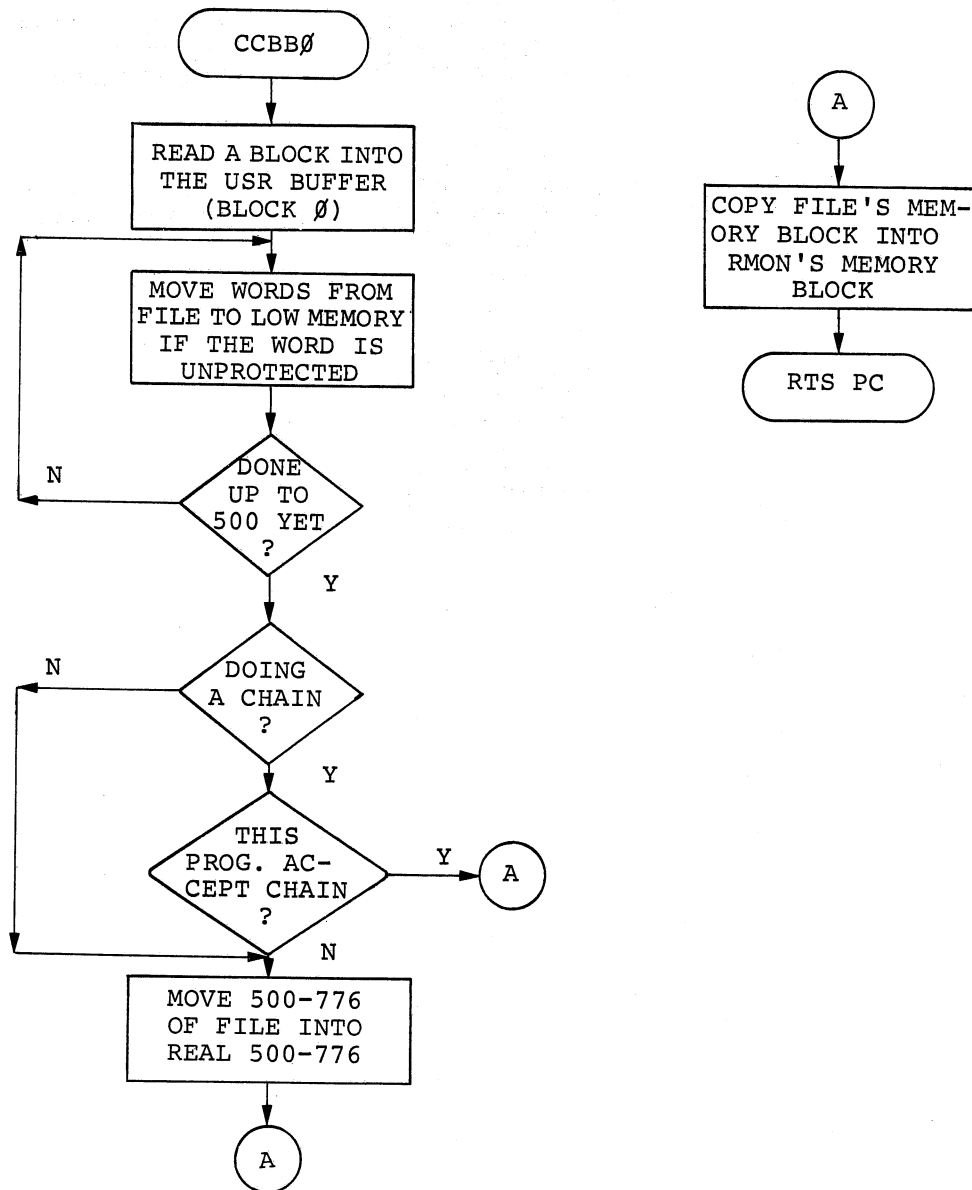


FILE - Called to pick up the .RAD50 representation of DEV:FILE.EXT. It will assume a default extension of .SAV.



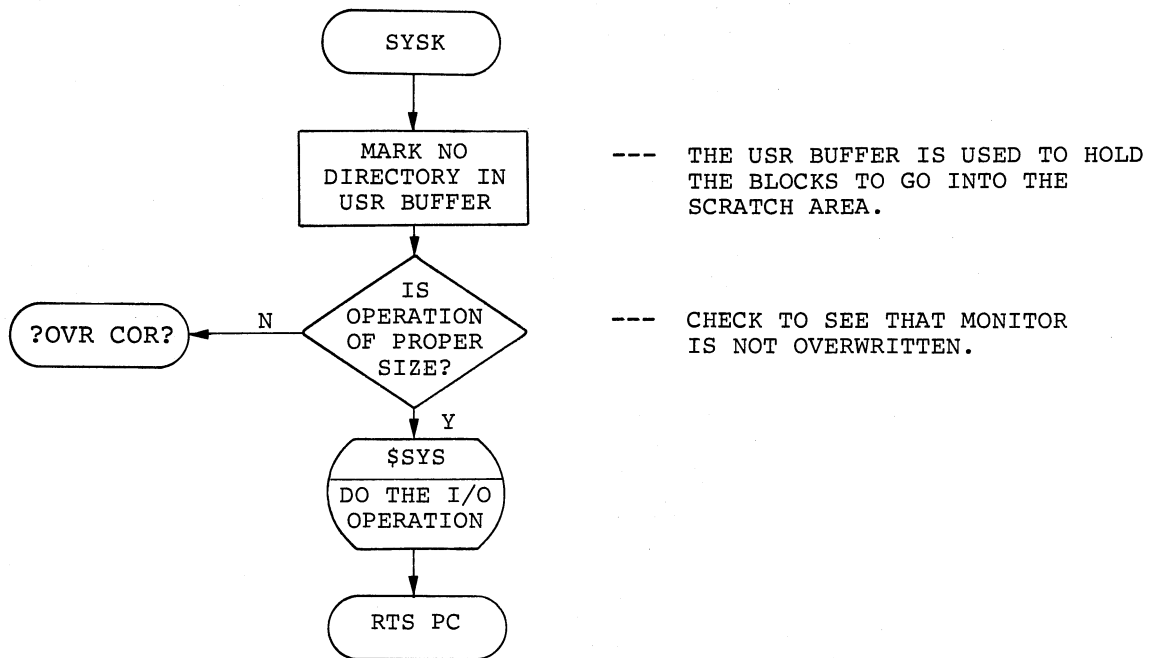


CCBBØ - The CCBBØ routine reads the first block of a .SAV file into the USR buffer, then moves selected locations from that block into the corresponding physical memory locations. The words moved are those marked with Ø's in the RMON bitmap. This procedure protects the system from having its vectors overlaid. If a chain is being done to a program which does not accept a CHAIN, 5ØØ-776 will be loaded with the contents of the file.



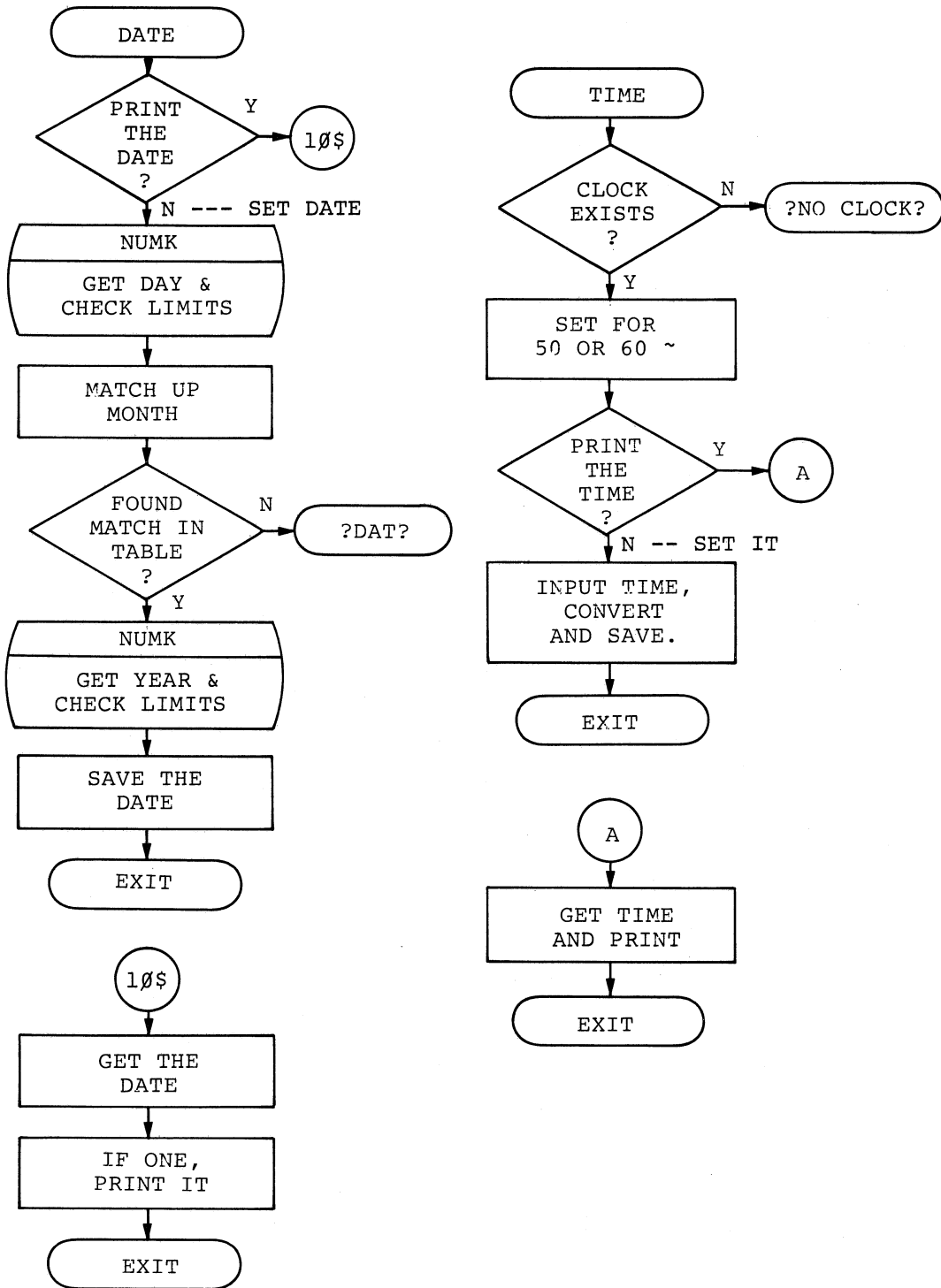
SYSK

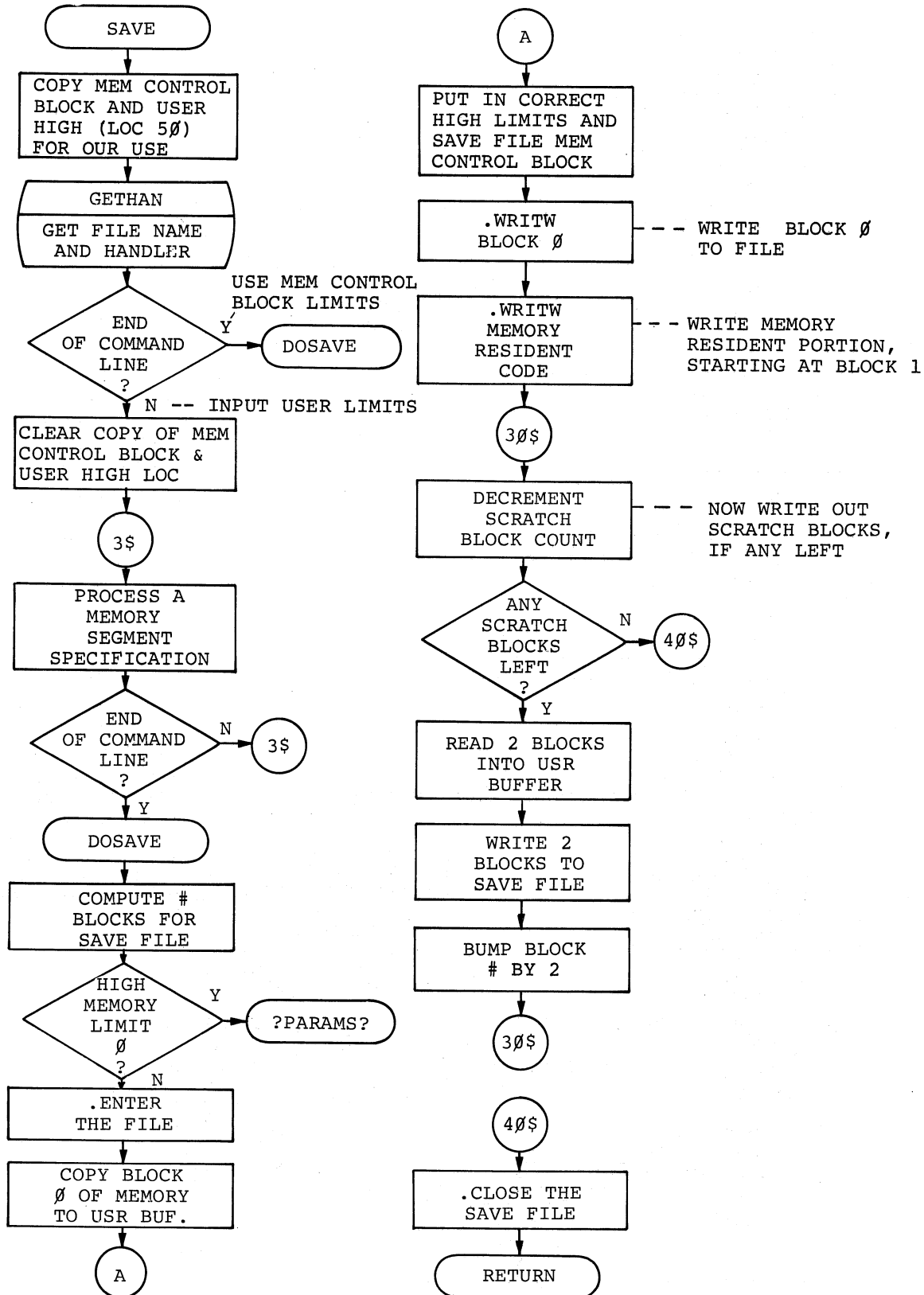
SYSK - Used to read/write blocks into and out of the system scratch area.



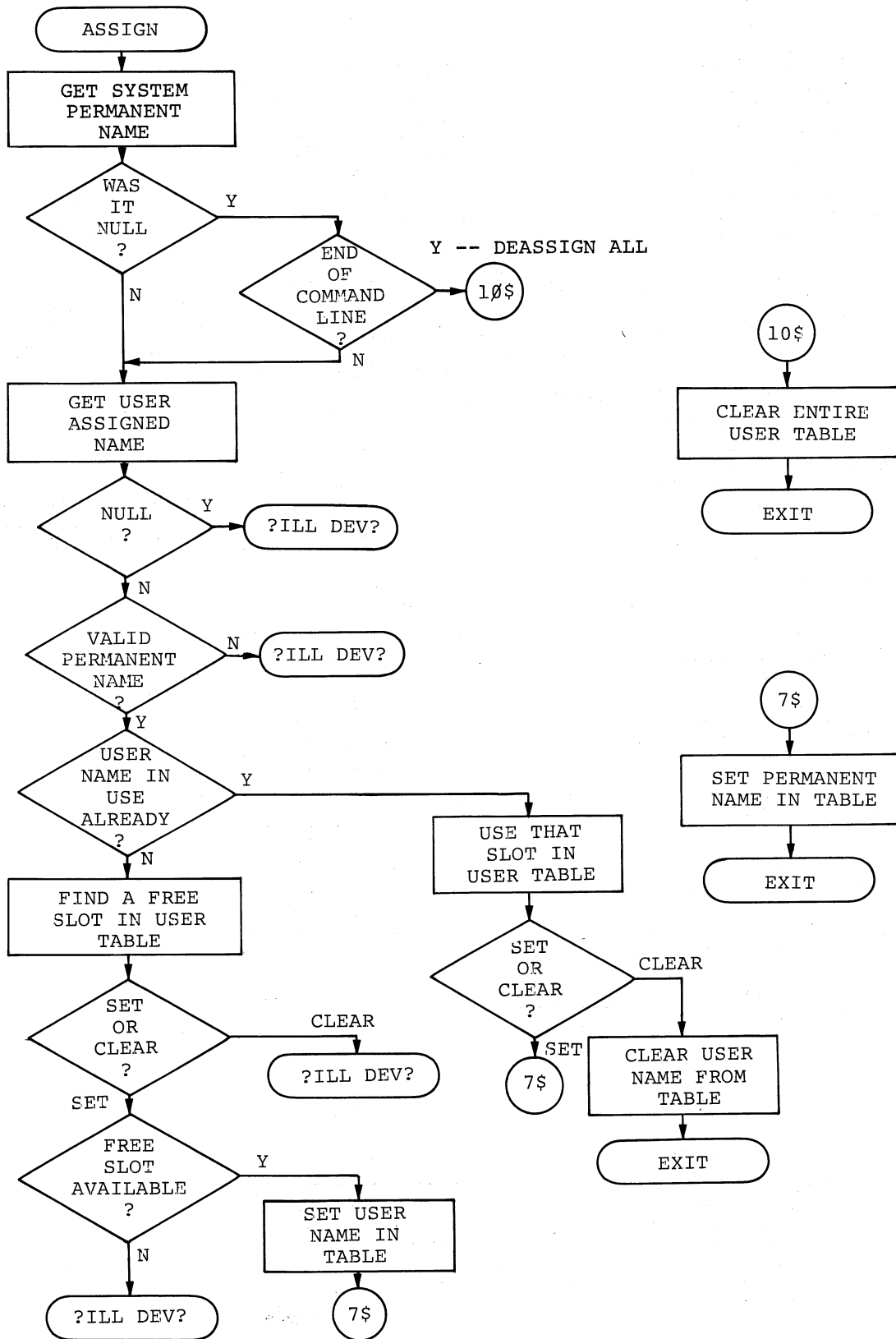
E.1.2 KMON Overlays

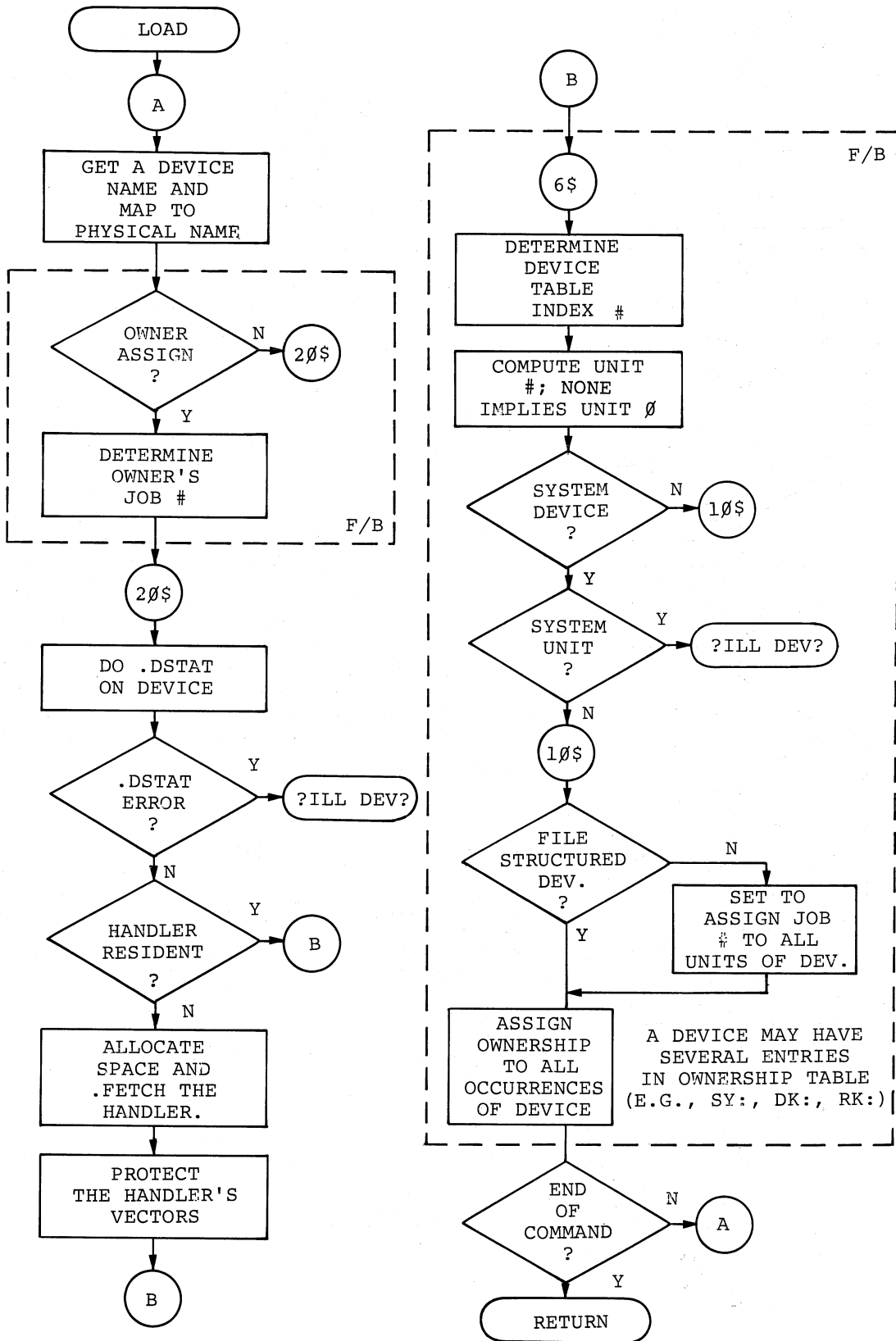
DATE/TIME



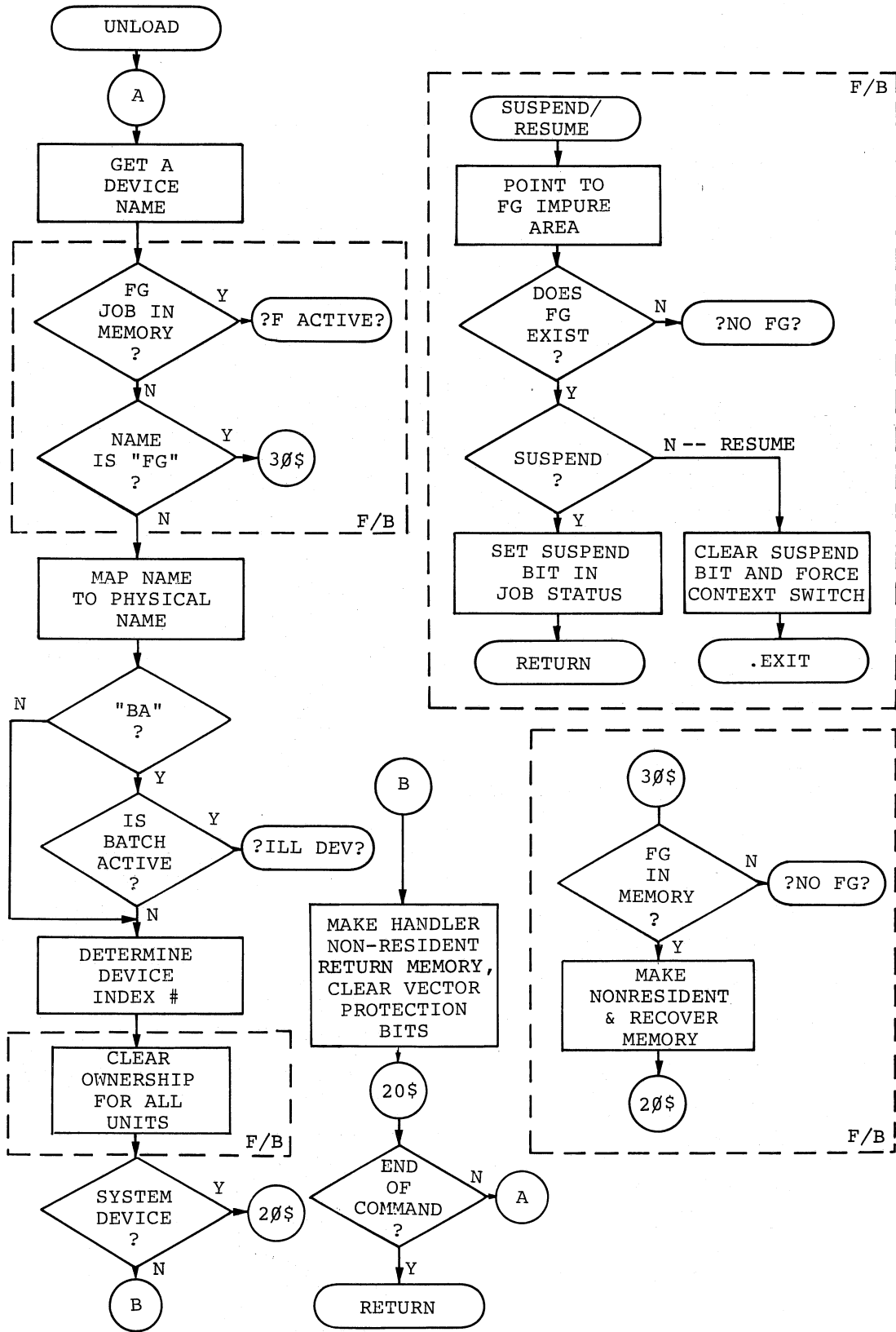


ASSIGN



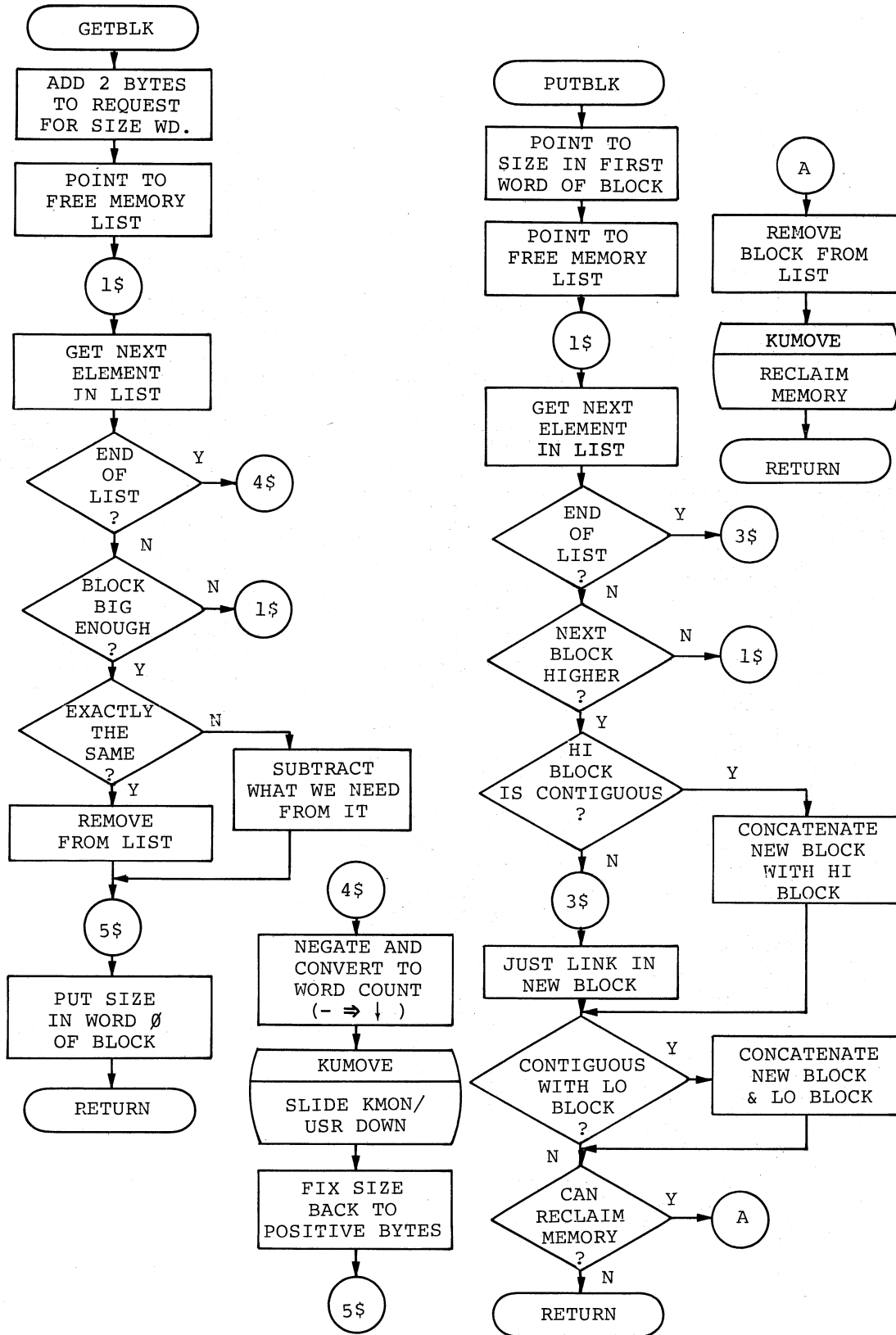


UNLOAD/SUSPEND/RESUME

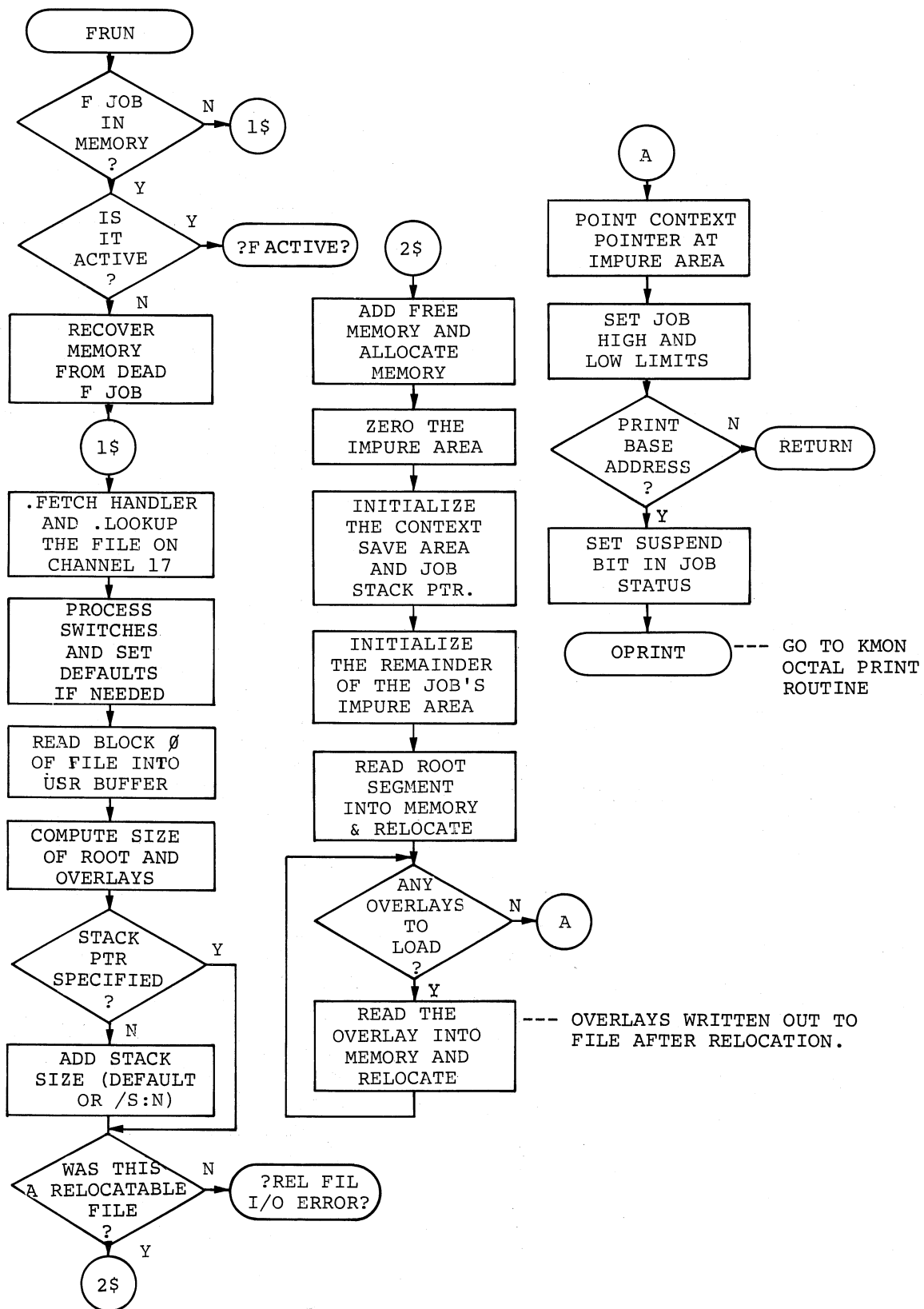




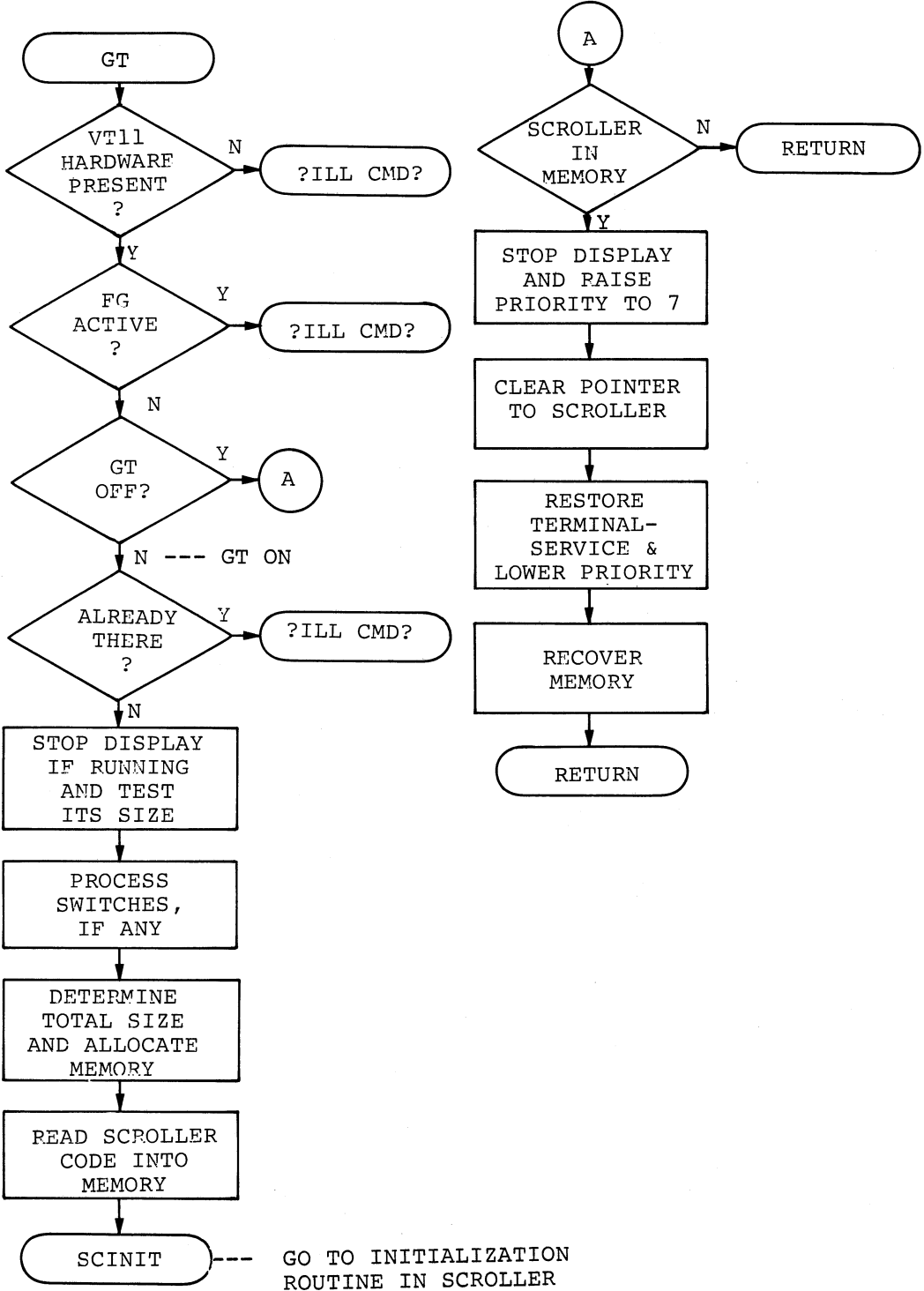
GET/PUT A BLOCK OF MEMORY



FRUN



GT ON/OFF



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E.2 USR (USER SERVICE ROUTINES) FLOWCHARTS

## USRBUF/FATAL/CDFN

The first 2 blocks of the USR are used by the USR for directory operations. They are also used by the KMON at various points for a 2-block general purpose buffer. There is, however, executable code in the buffer that can be executed every time a fresh copy of the USR is read from the system device. The functions included in the buffer are:

1. USR Relocation

This code is executed whenever the USR is newly read into memory. It serves to make certain pointers into RMON absolute.

2. Fatal error processor and fatal error messages (S/J only)

3. CDFN (channel define) EMT (S/J only)

The CDFN EMT call forces a new copy of the USR into memory to guarantee the presence of the EMT processor.

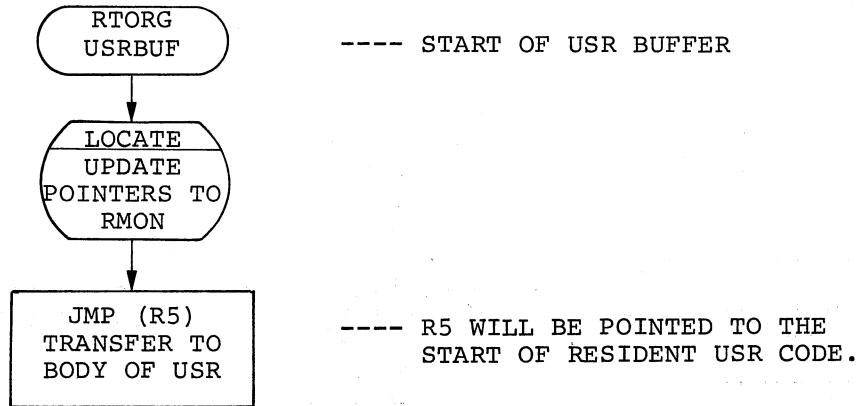
The flows for these functions follow.

### NOTE

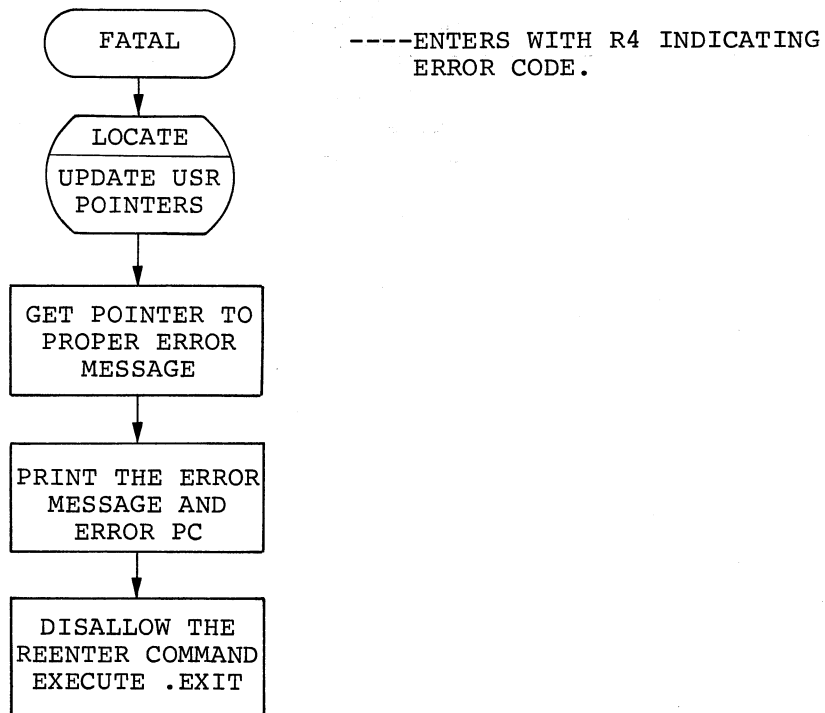
Fatal error handler and CDFN processor are RMON functions in the F/B Monitor. The only code in the buffer in the F/B system will be the USR relocation code.

USRBUF/FATAL/CDFN (CONT.)

USRBUF is the initial entry point for USR calls when the USR has just been read into memory. LOCATE sets up pointers into RMON.

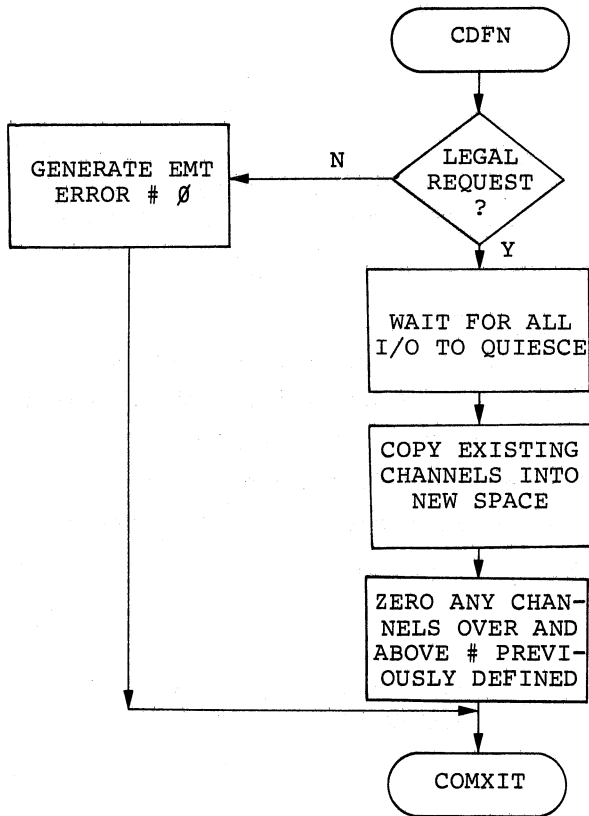


The LOCATE routine is called to update the list of pointers at RELIST. The list is initially a list of address differences (i.e., VALUE-\$RMON where VALUE is the desired location and \$RMON is the address of the start of RMON). LOCATE then makes all the differences into absolute addresses. Any errors which would generate a ?M-error use the FATAL error processor code to generate the message in the S/J system. This is a resident function in F/B.



USRBUF/FATAL/CDFN (CONT.)

CDFN - A resident function in the F/B system.



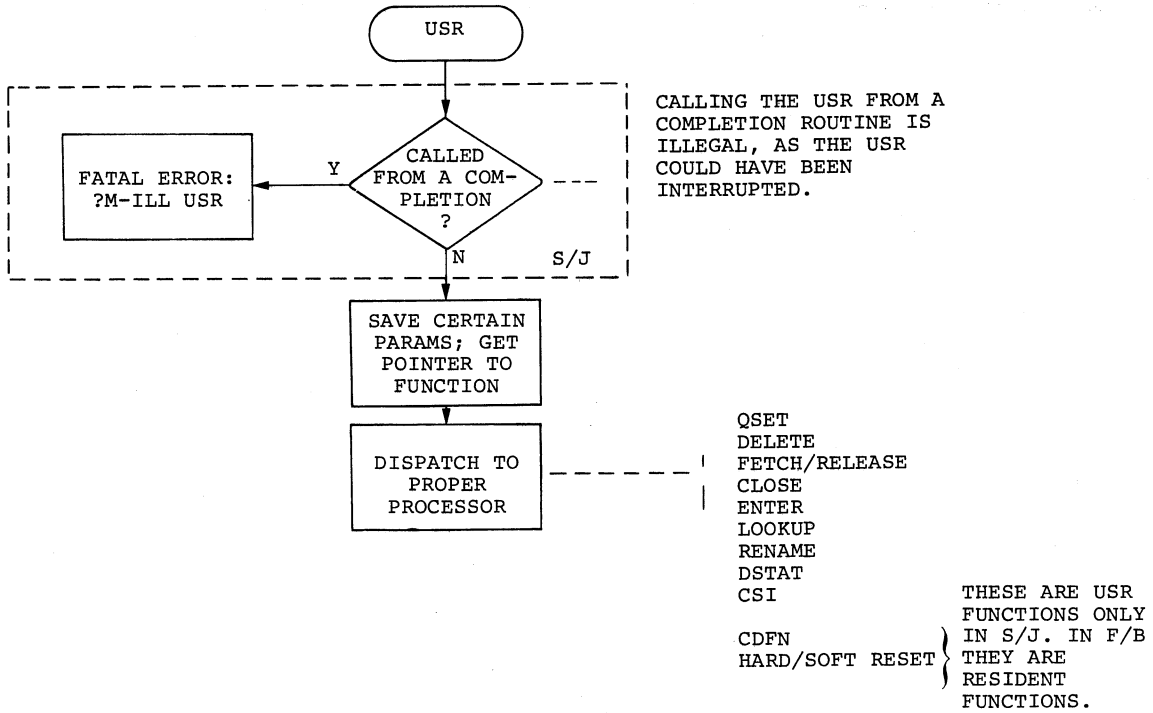
-----IF REQUEST IS FOR FEWER THAN  
THAN ALREADY EXIST, IT IS AN  
ERROR.

-----TAKE COMMON USR EXIT.

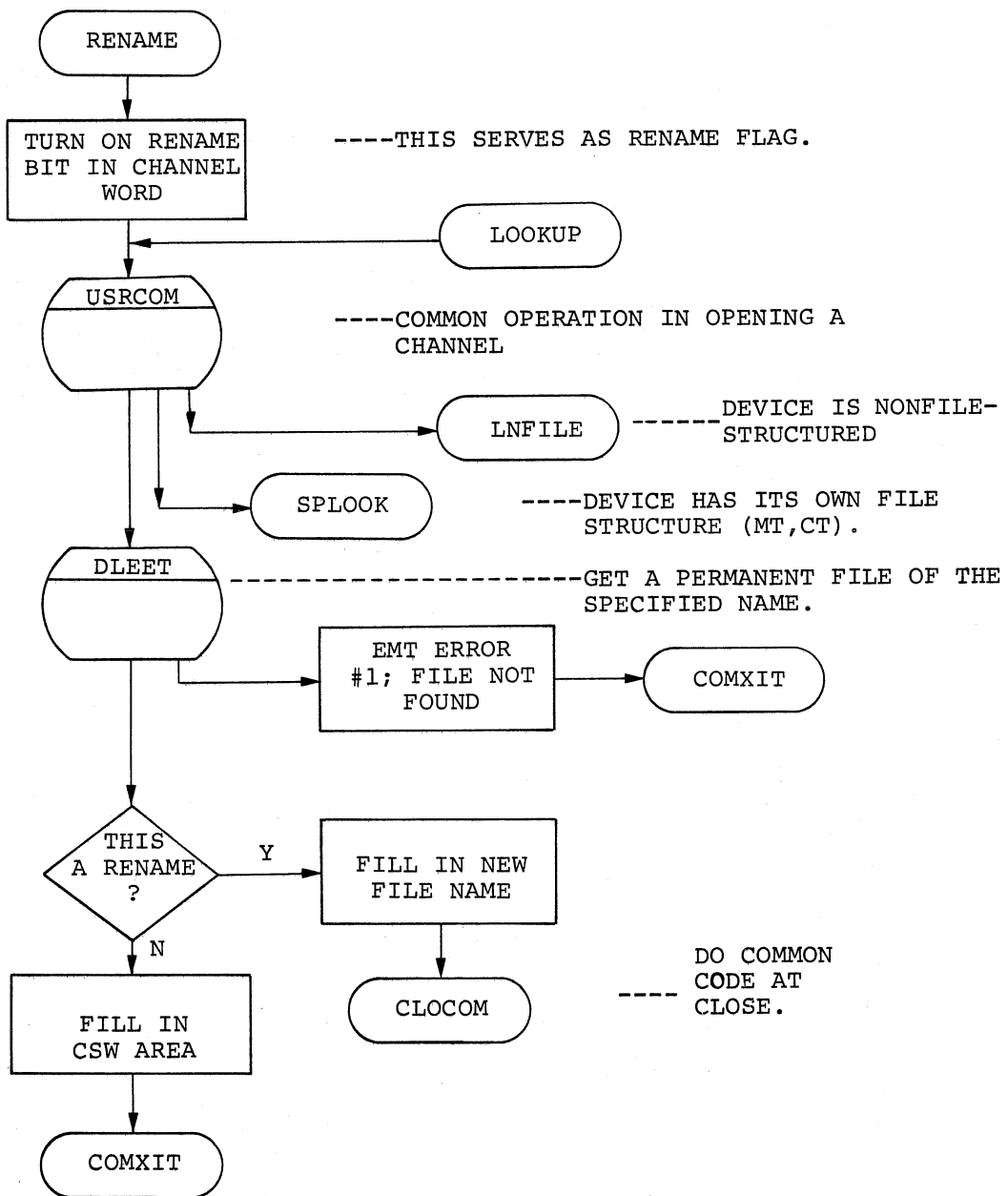


The following flowcharts detail the code contained in the main body of the USR. On entry to the USR, R2 contains an index representing the function to be performed. This is used to dispatch control to the proper processor.

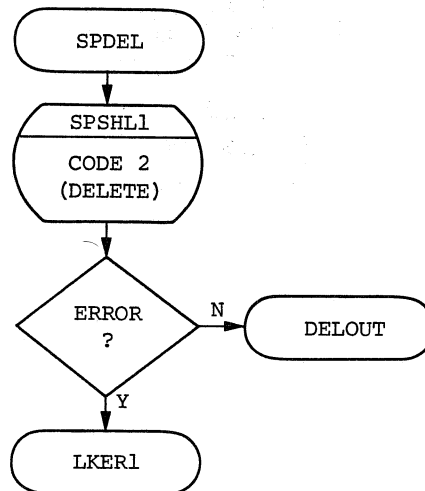
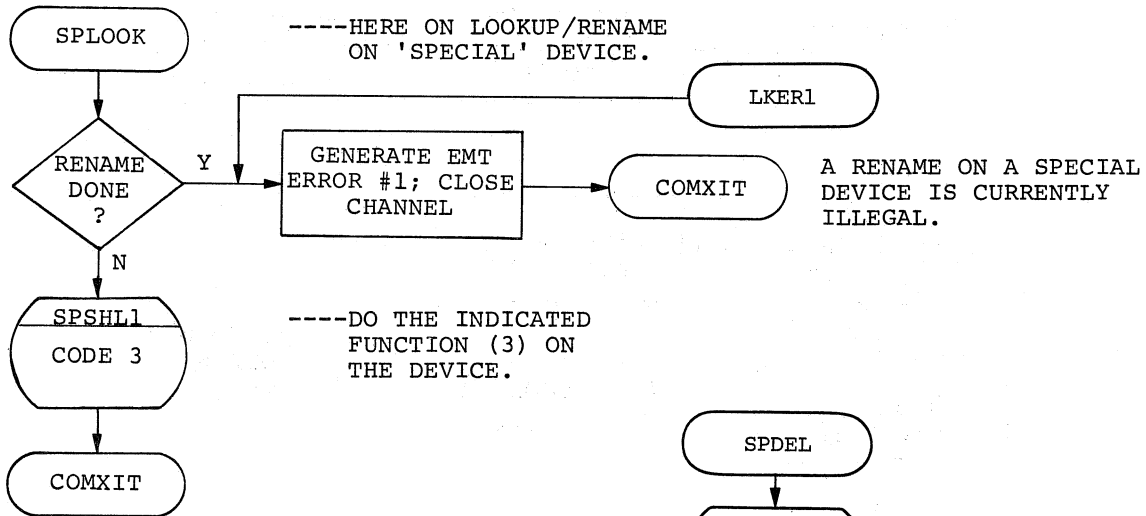
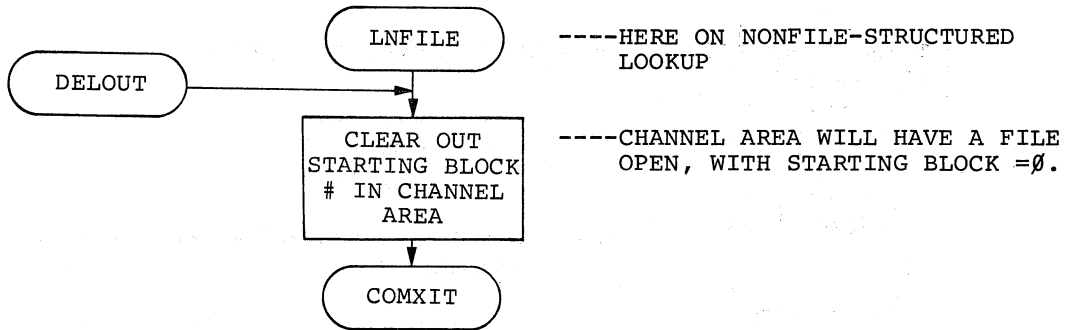
USR CODE



LOOKUP/RENAME

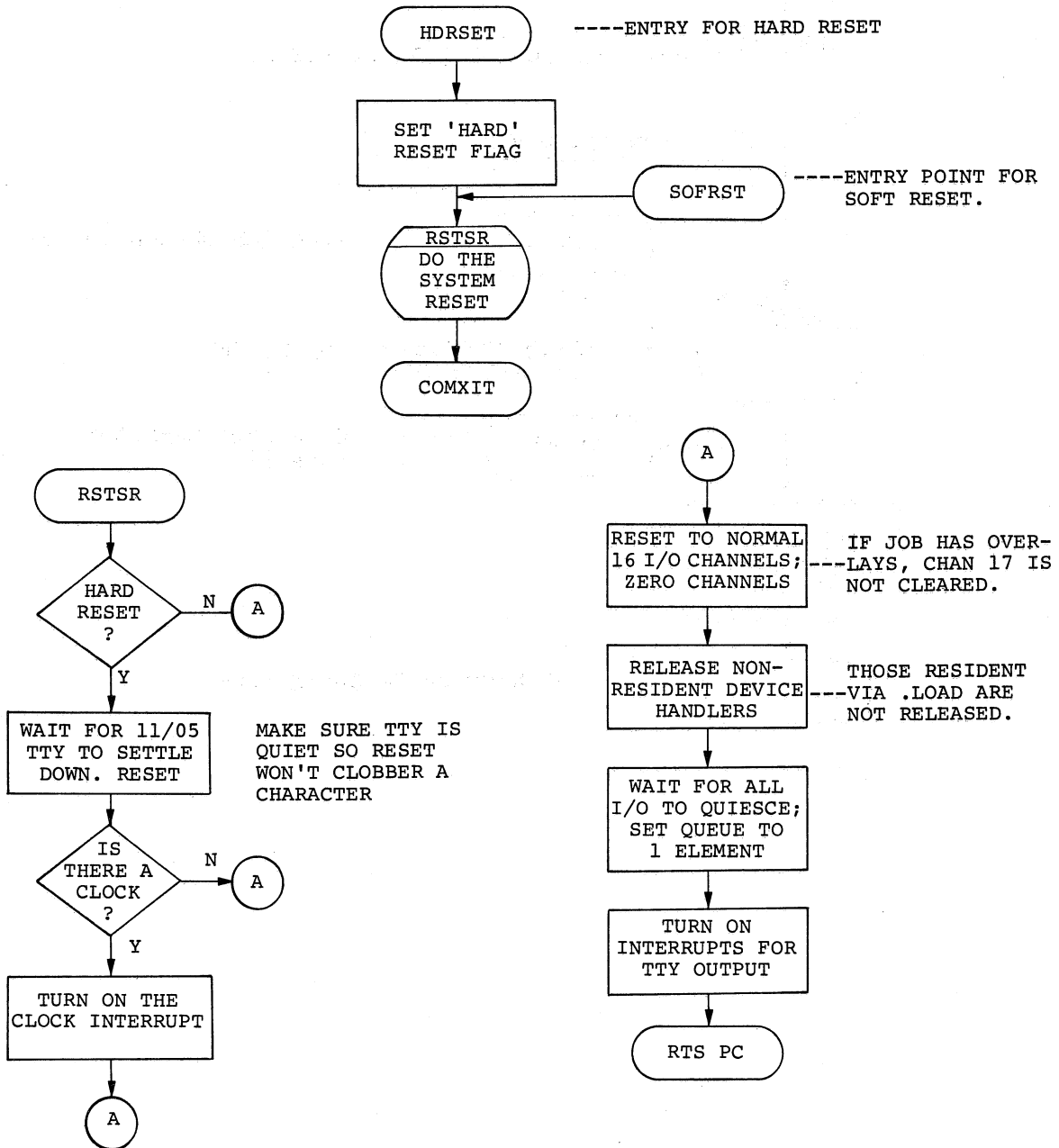


LOOKUP/RENAME (CONT.)

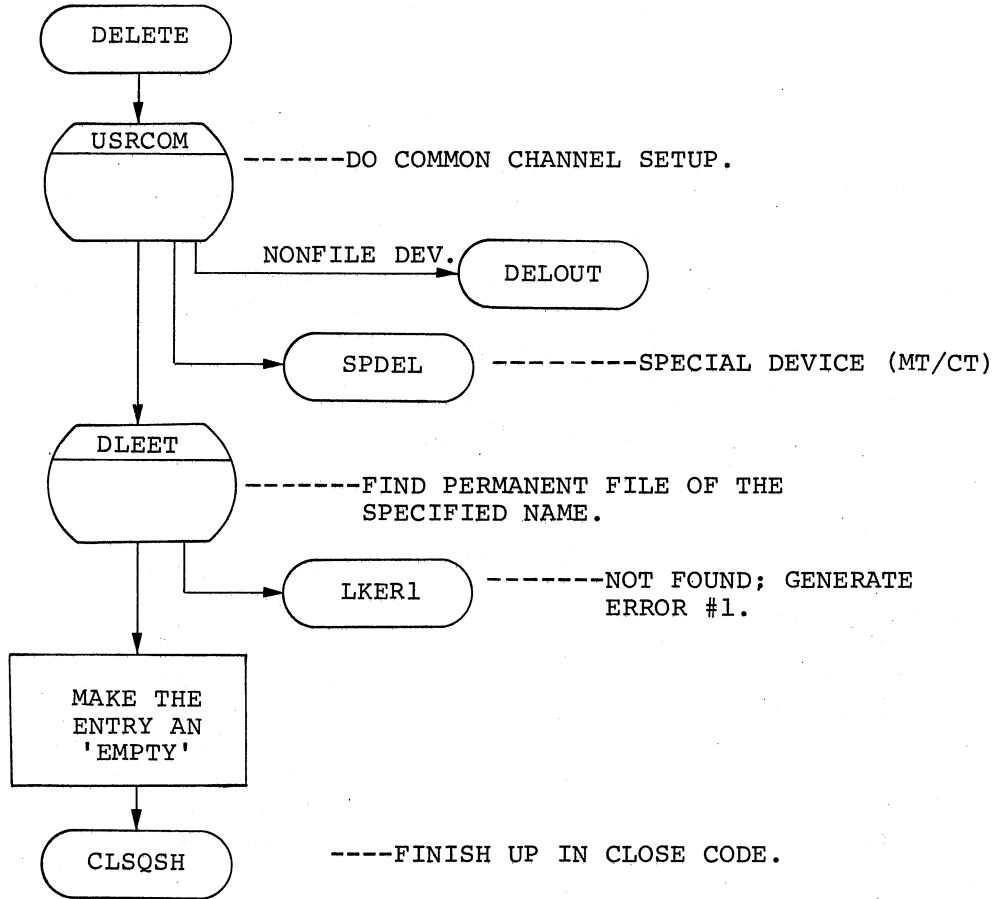


# HARD/SOFT RESET

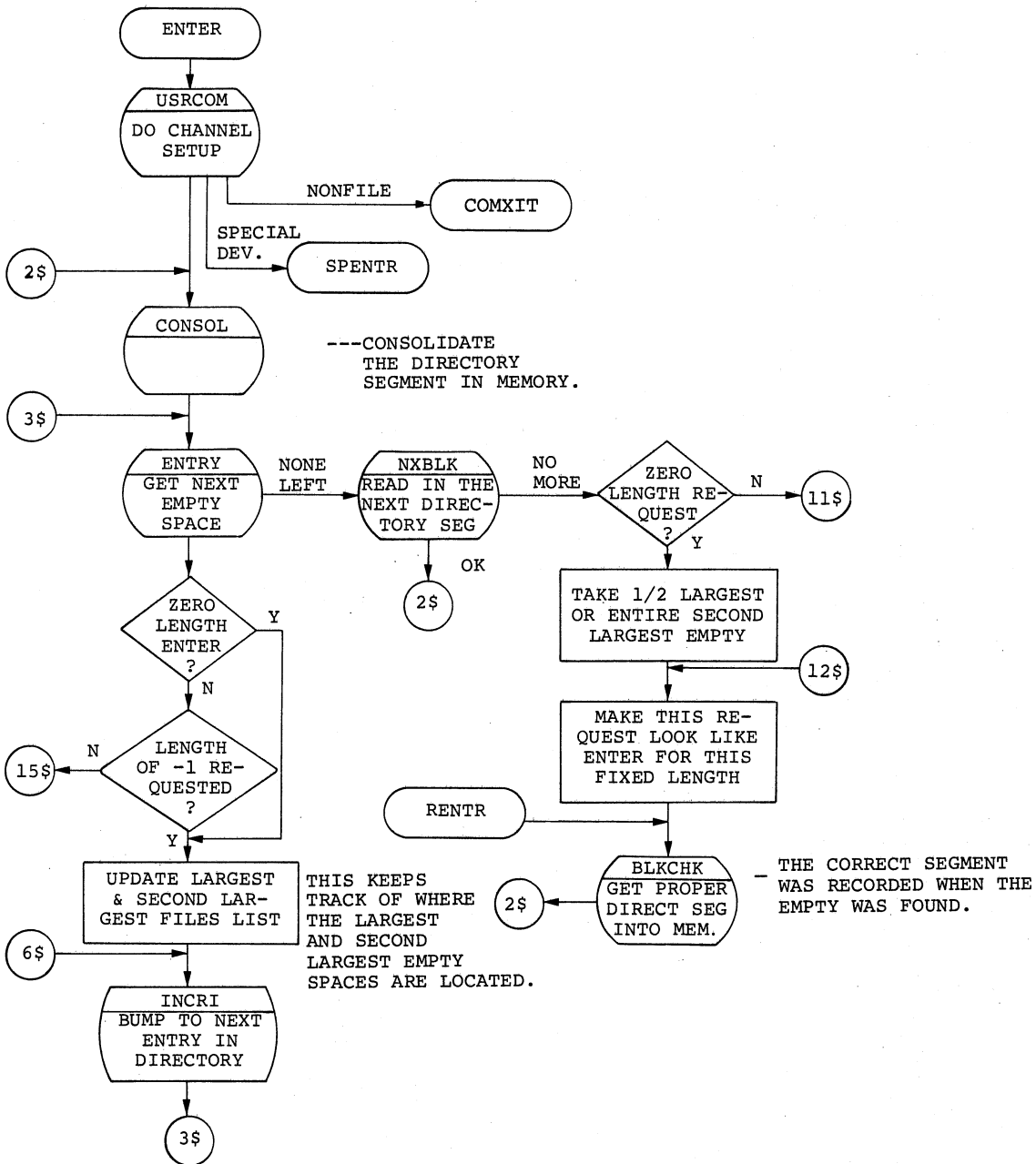
These are resident functions in F/B; USR functions in S/J.



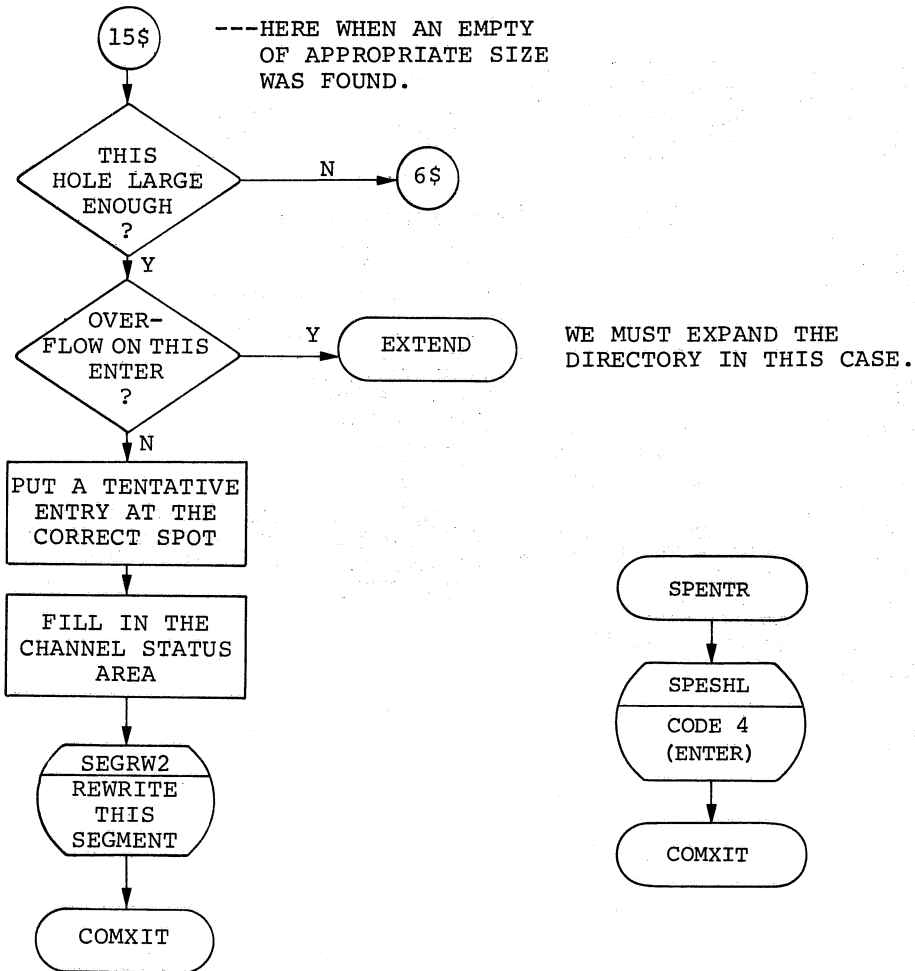
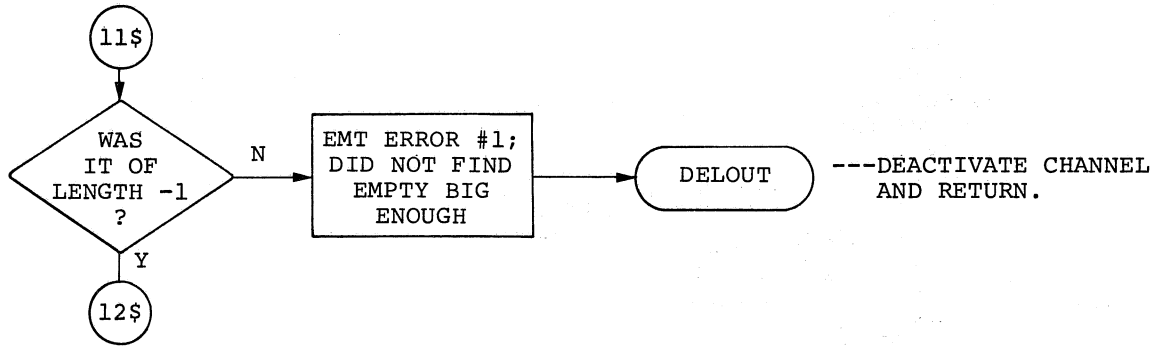
DELETE



ENTER

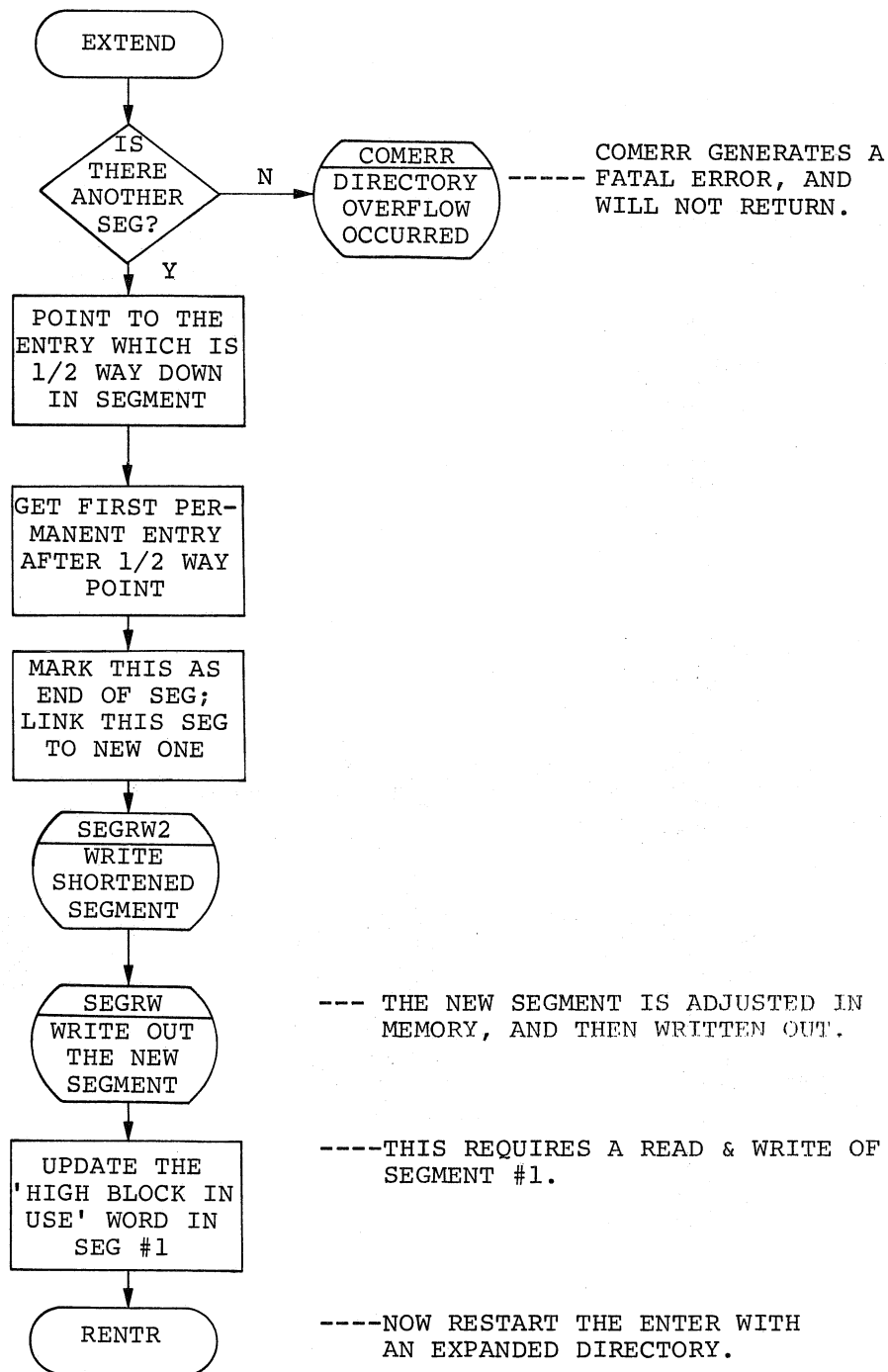


ENTER (CONT.)



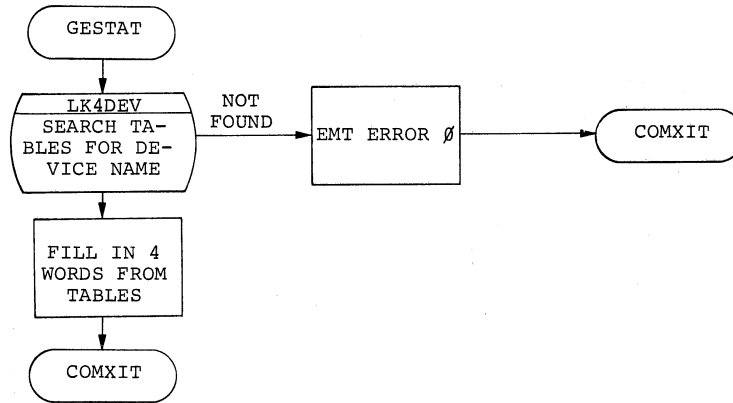


EXTEND

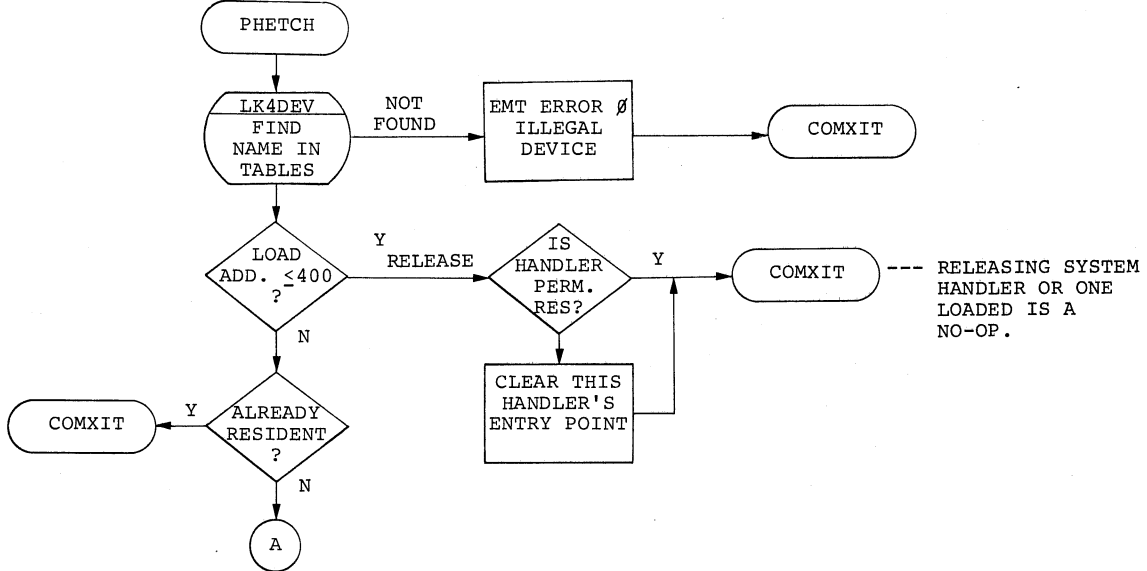


DSTAT/FETCH/RELEASE

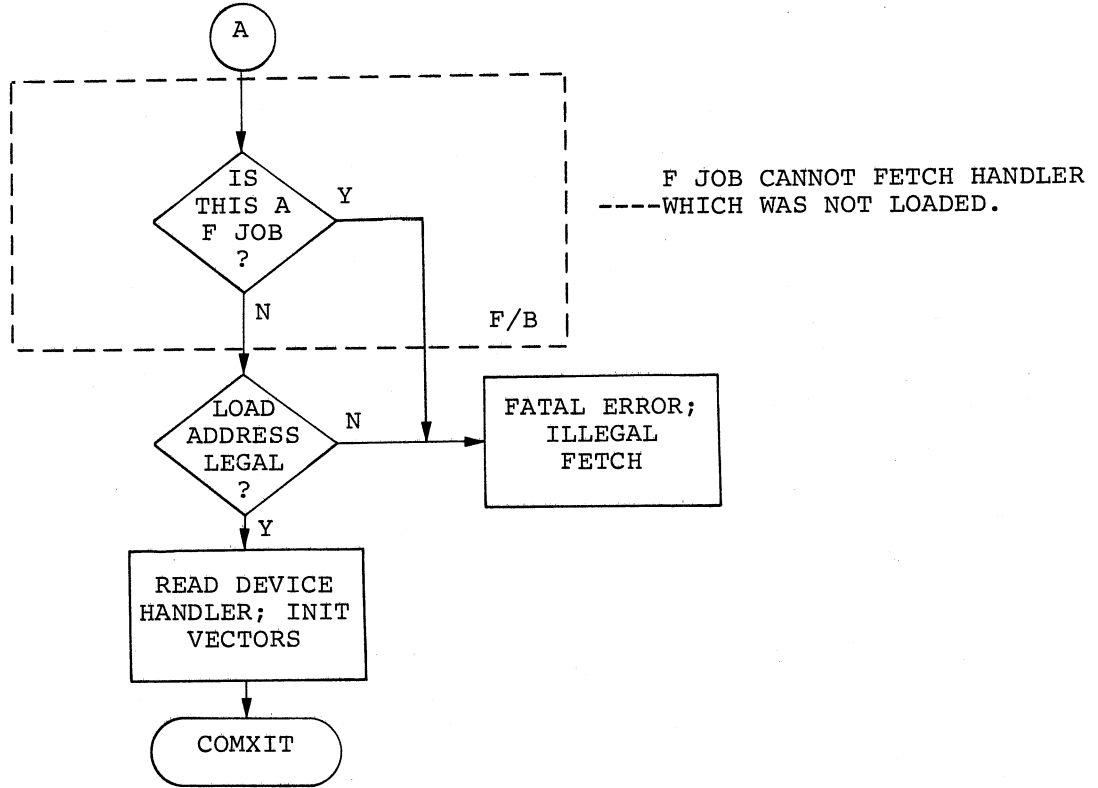
DSTAT- GET DEVICE STATUS



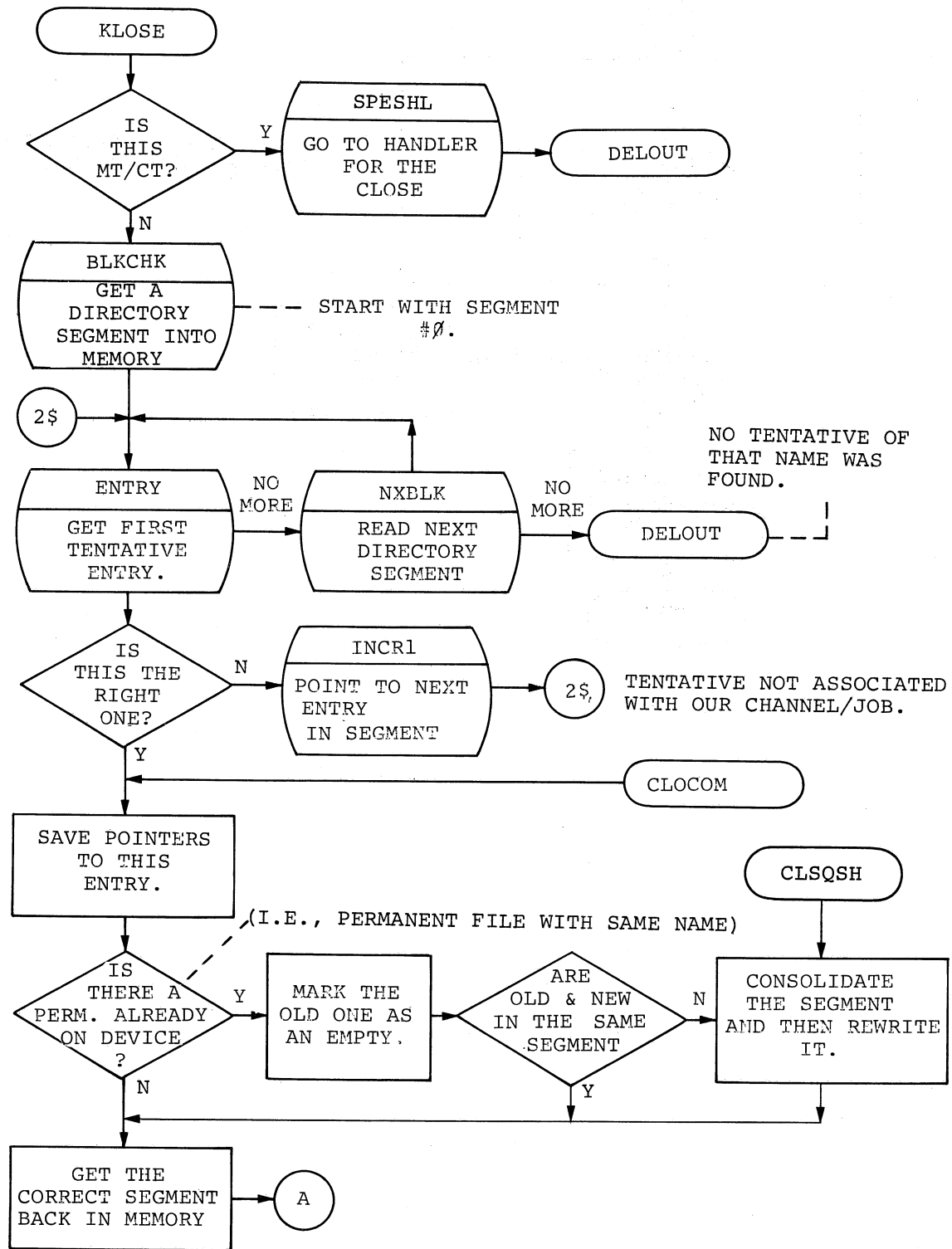
FETCH/RELEASE



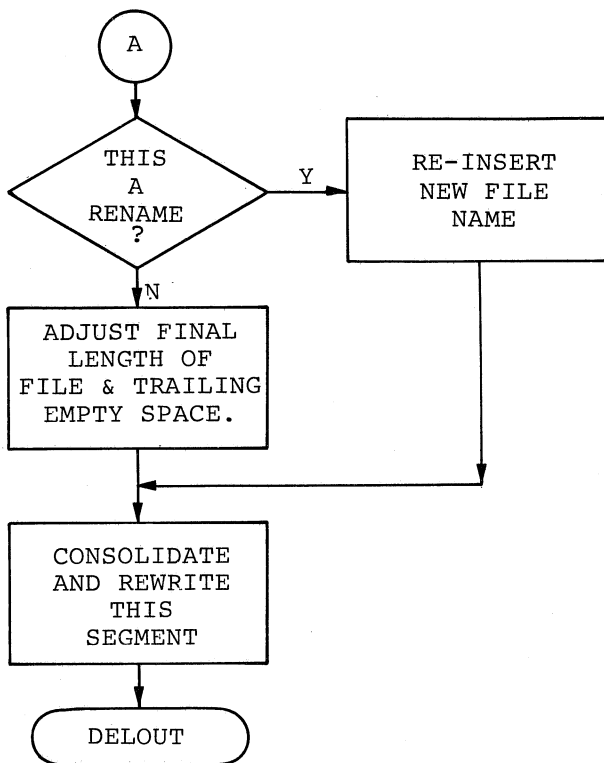
DSTAT/FETCH/RELEASE (CONT.)



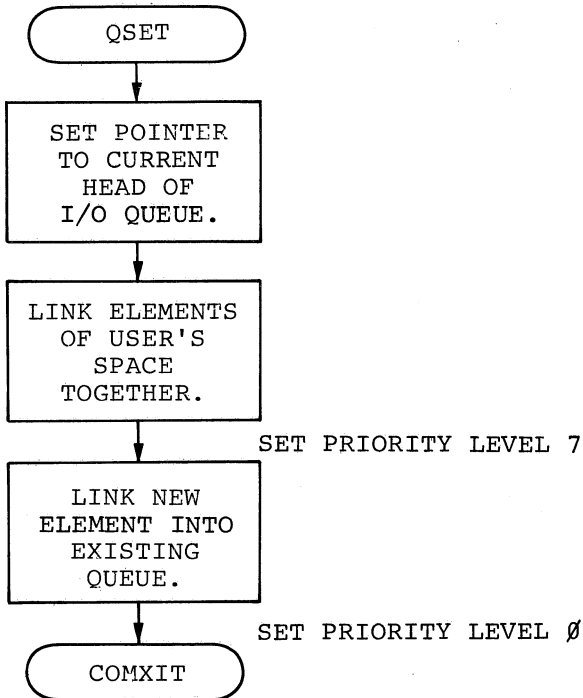
CLOSE



CLOSE (CONT.)/QUEUE EXTEND



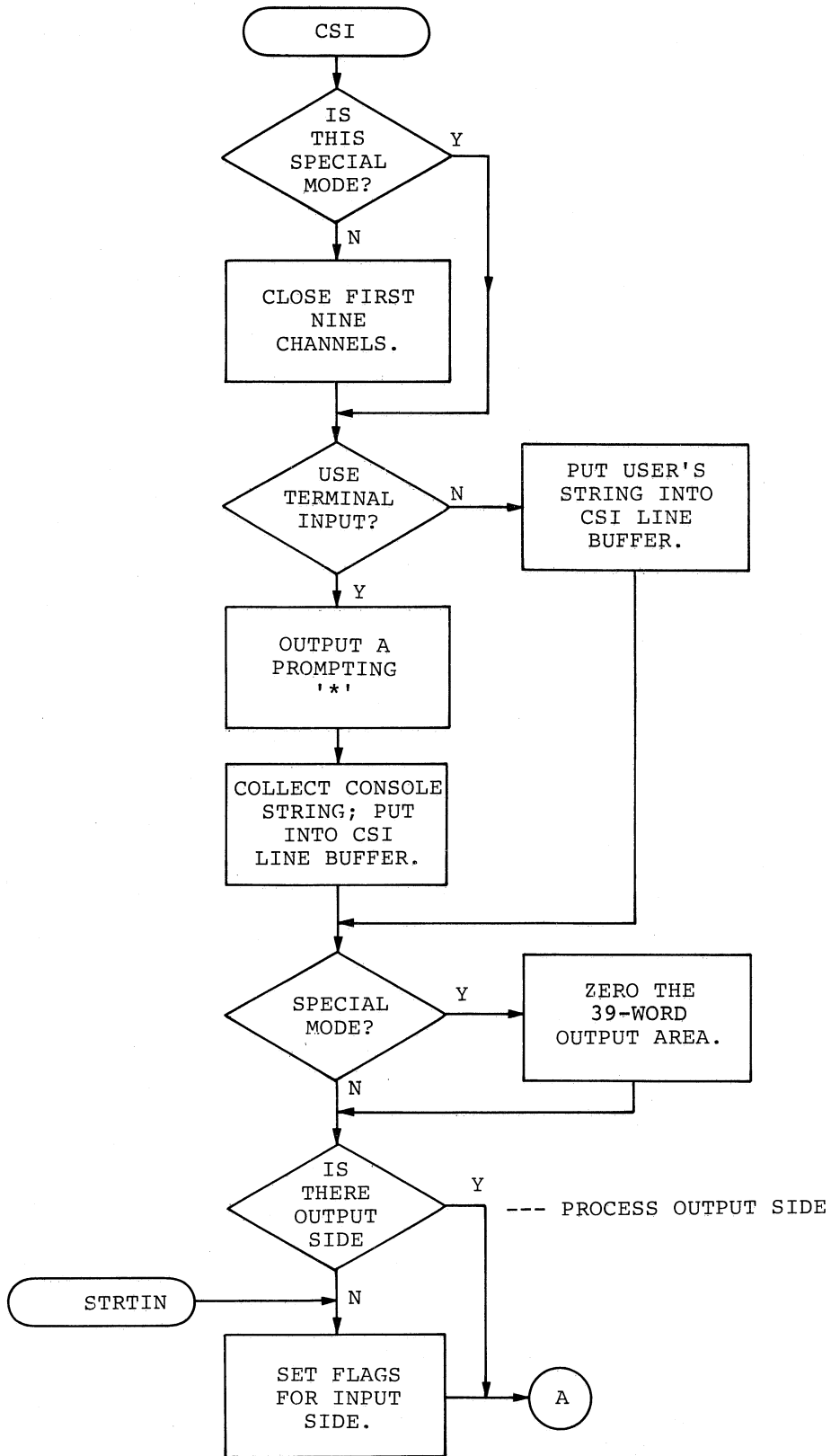
QUEUE EXTEND (QSET)



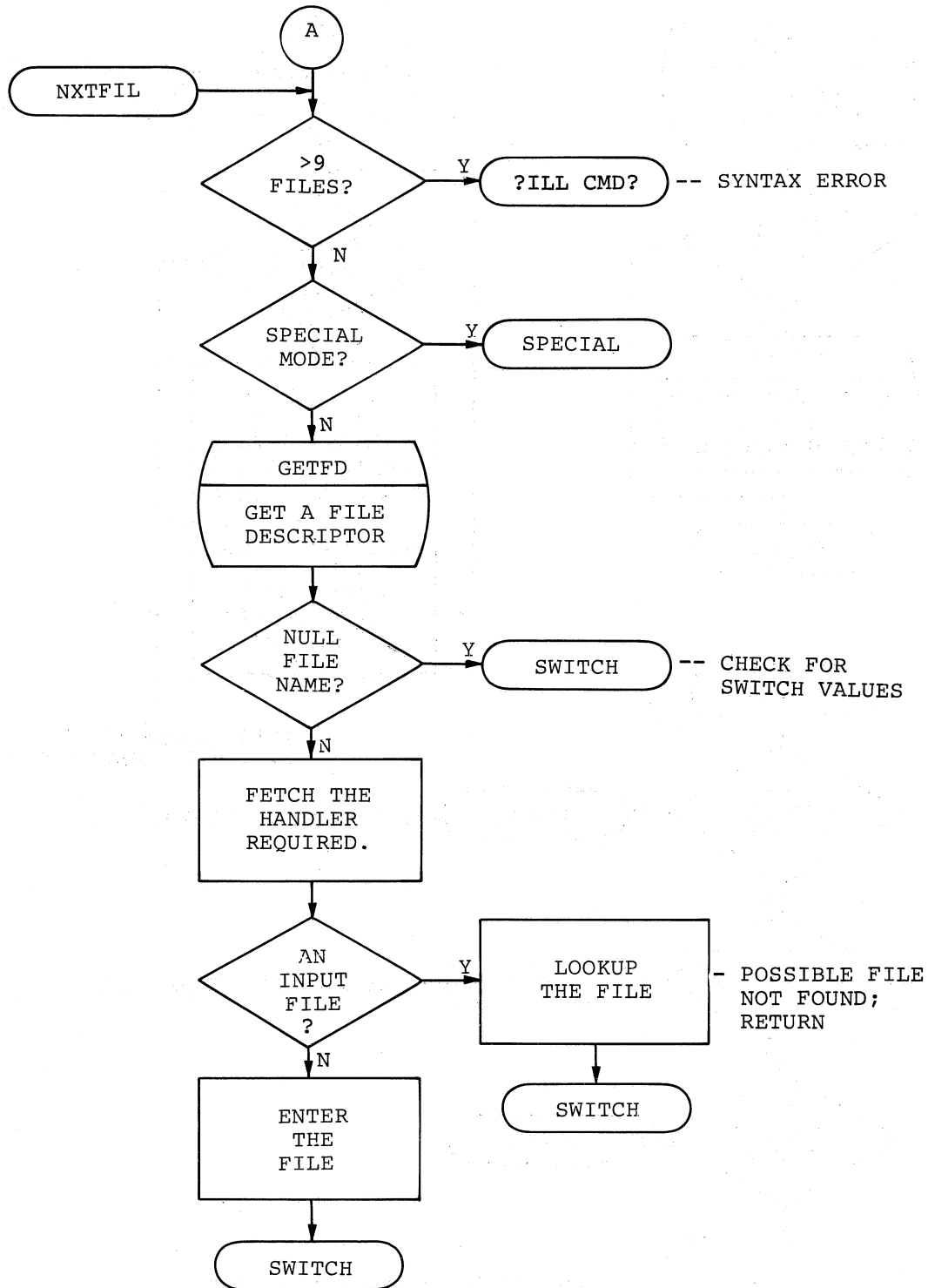


E.3 CSI (COMMAND STRING INTERPRETER) FLOWCHARTS

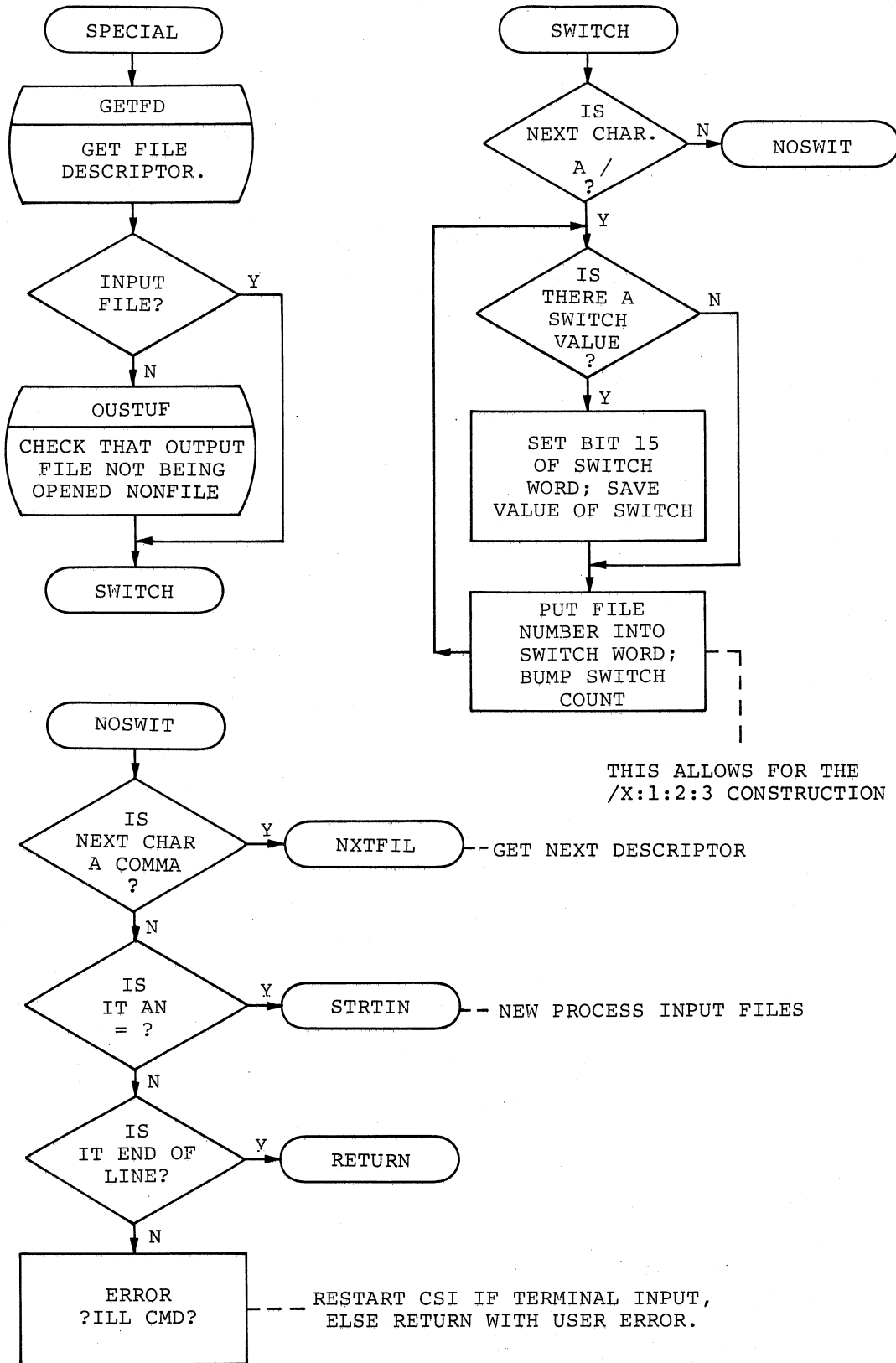
CSI CODE



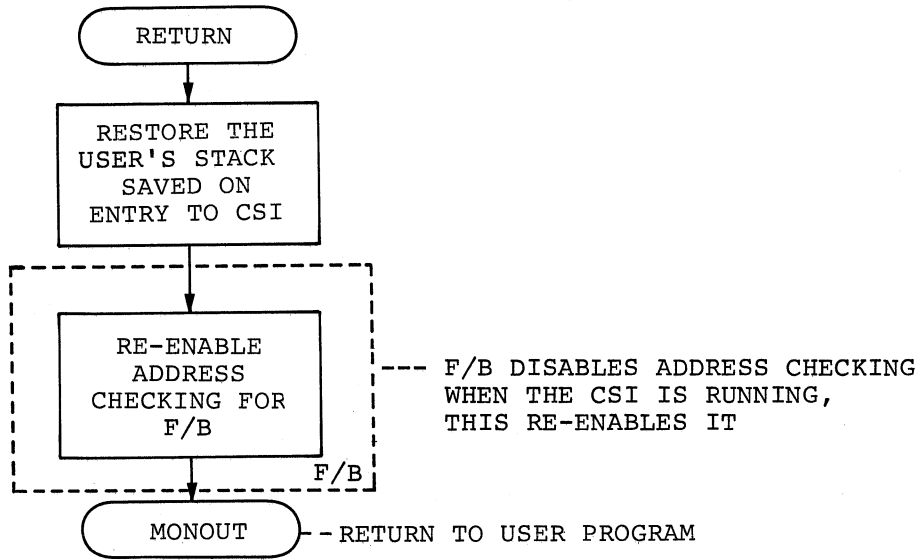




CSI CODE (CONT.)



CSI CODE (CONT.)



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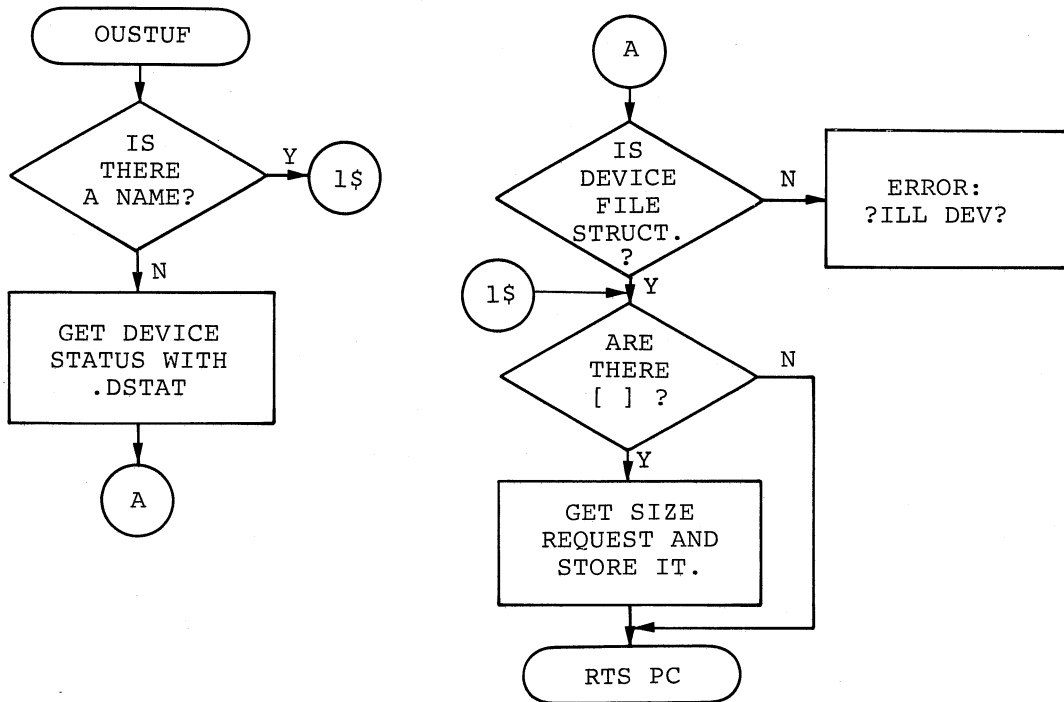
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### E.3.1 CSI Subroutines

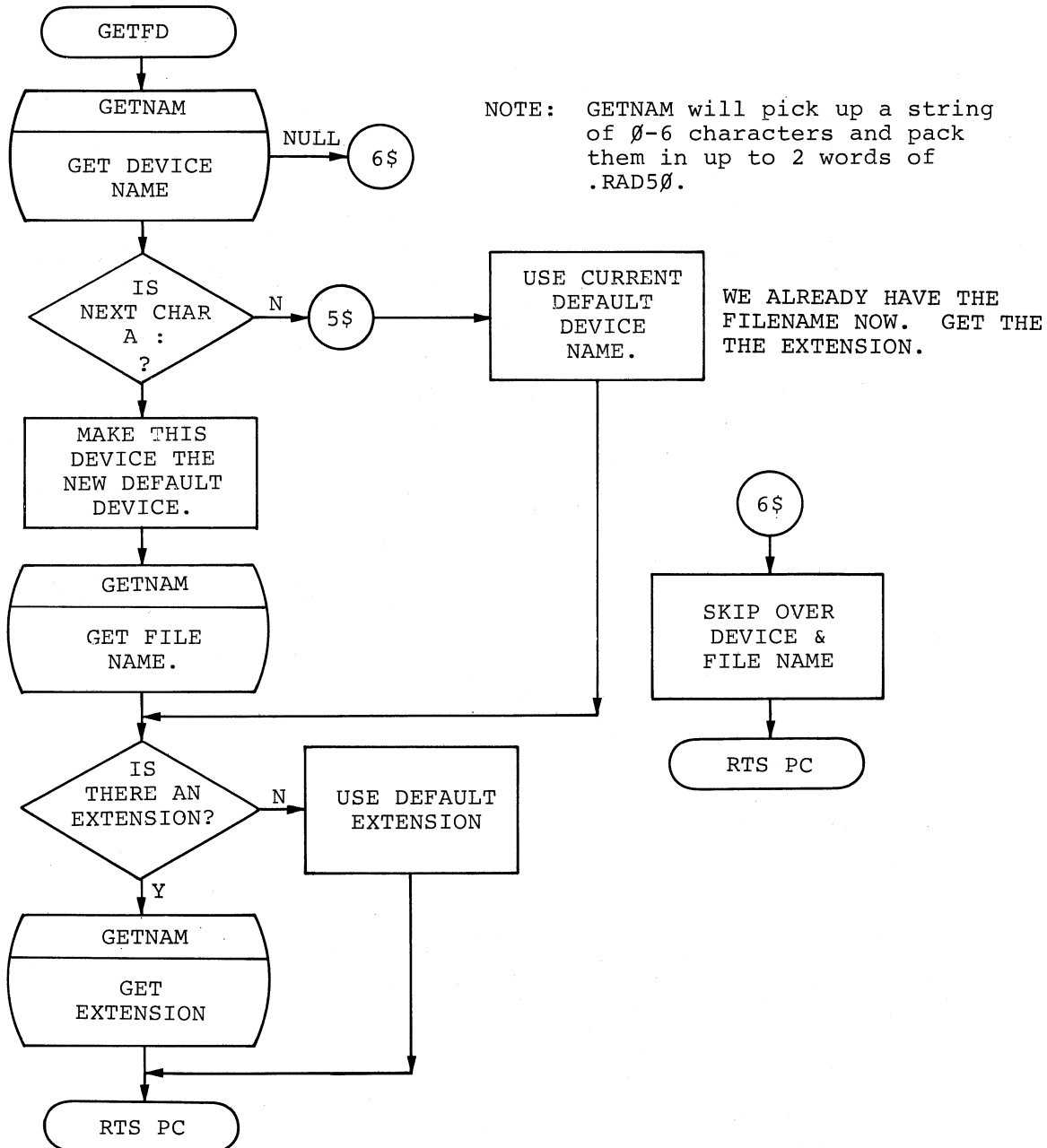
These subroutines are used by the CSI, and, in certain cases by the KMON.

OUSTUF

OUSTUF - This routine verifies that an output descriptor has a file name. If not, a syntax error is generated. It also will scan off the size in [ ] if it was specified.



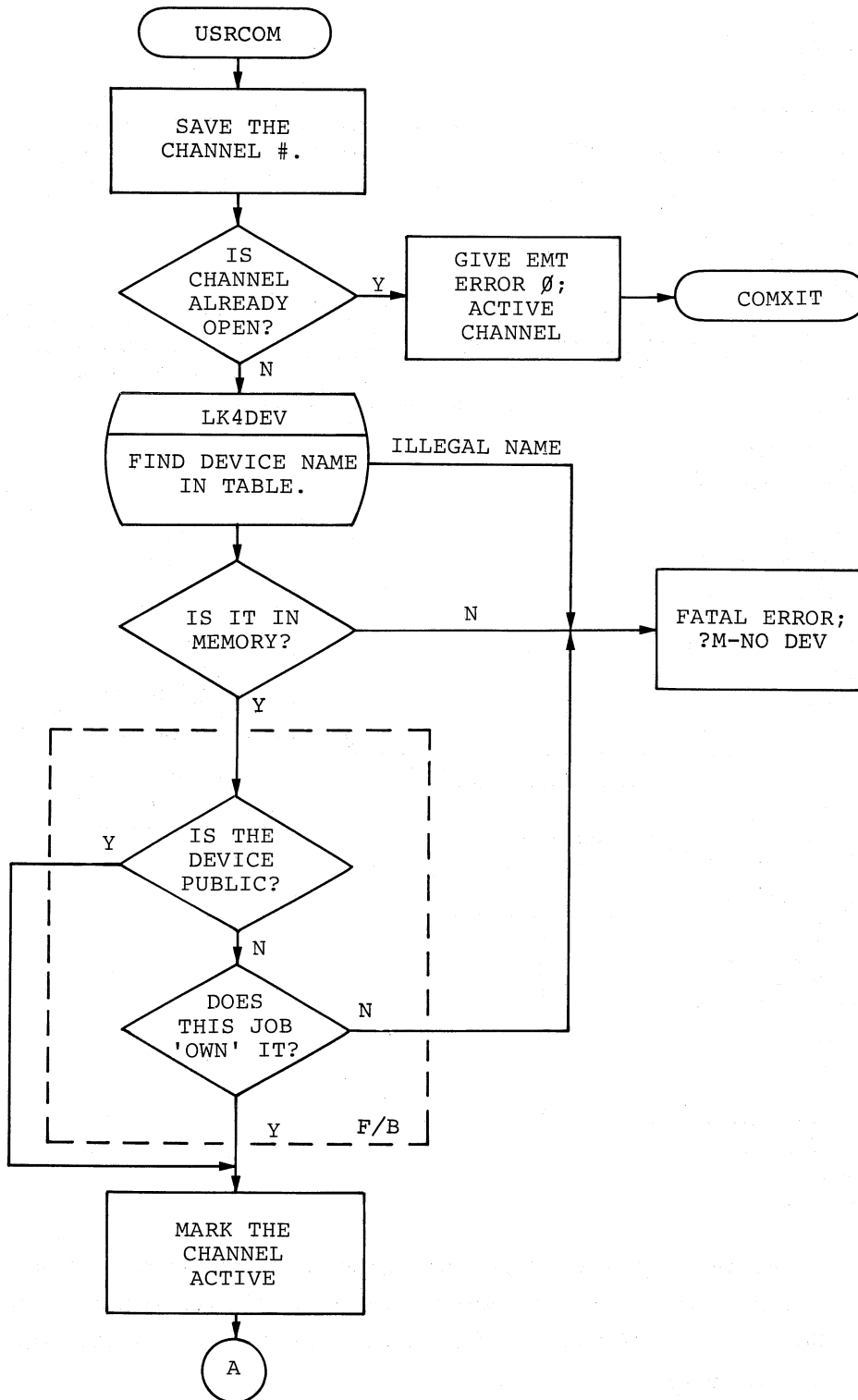
GETFD - Picks up a file descriptor (DEV:FILE.EXT) from an input string and packs it in 4 words of .RAD5Ø.



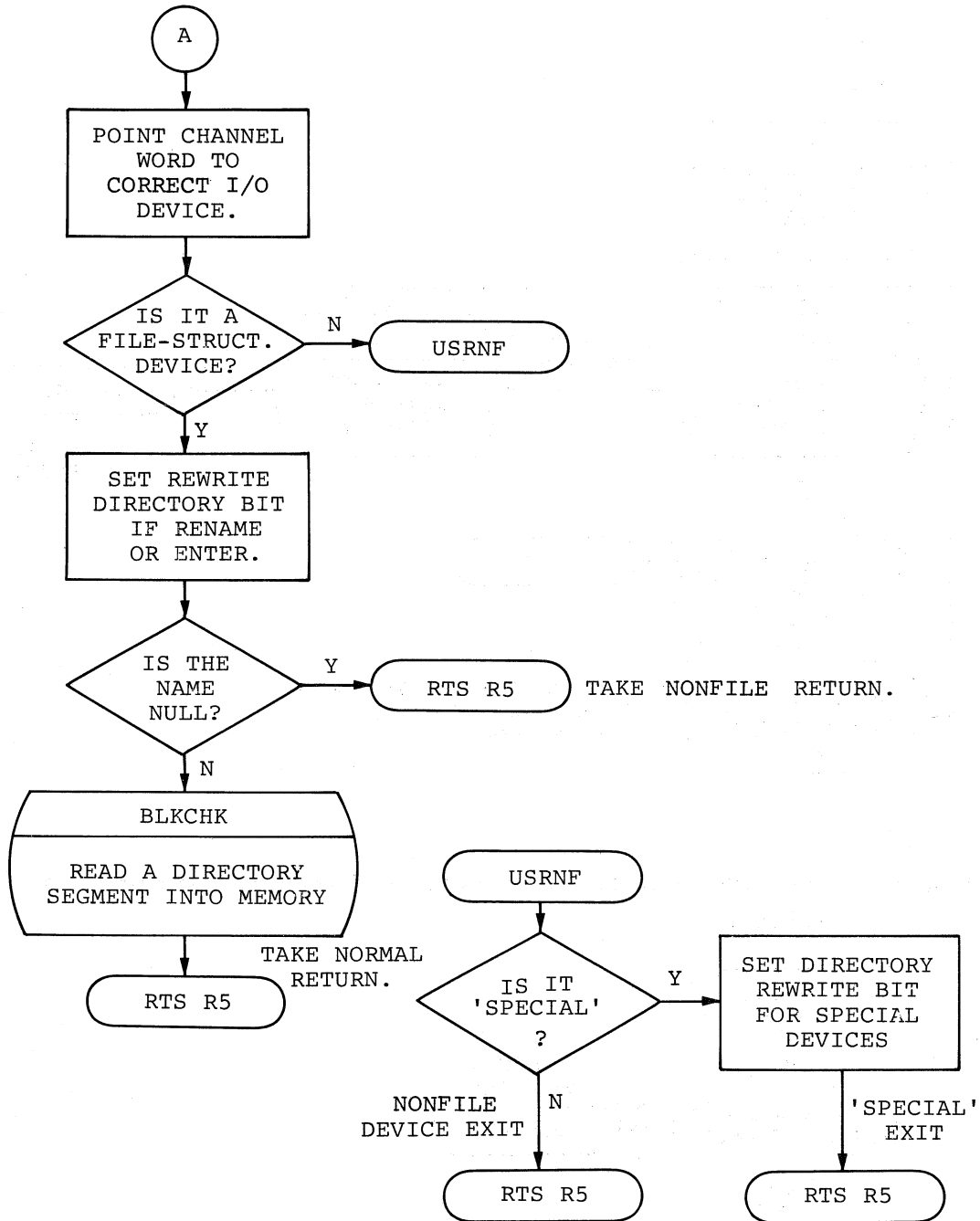
GETNAM - Converts a string of Ø-6 alphanumeric characters to a 2-word RAD5Ø group. The two words are zero filled when necessary. See code at GETNAM in the source listing if greater detail is necessary.

USRCOM

USRCOM - This routine is used to prepare a channel for I/O operations.

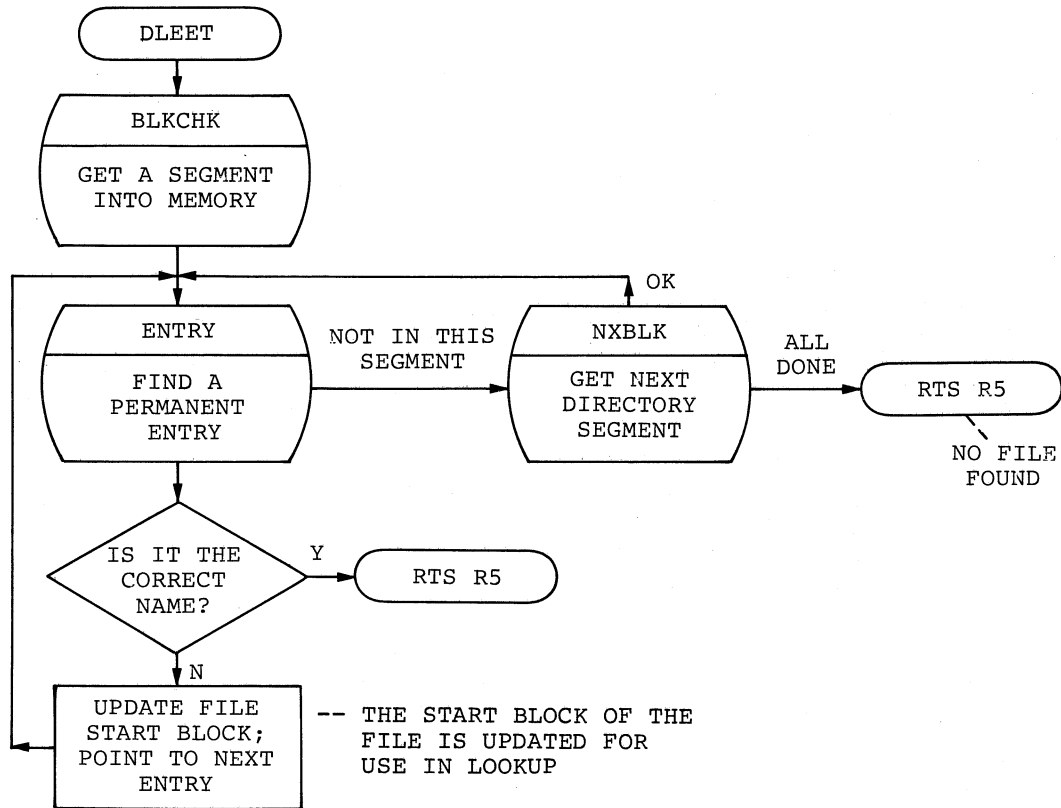




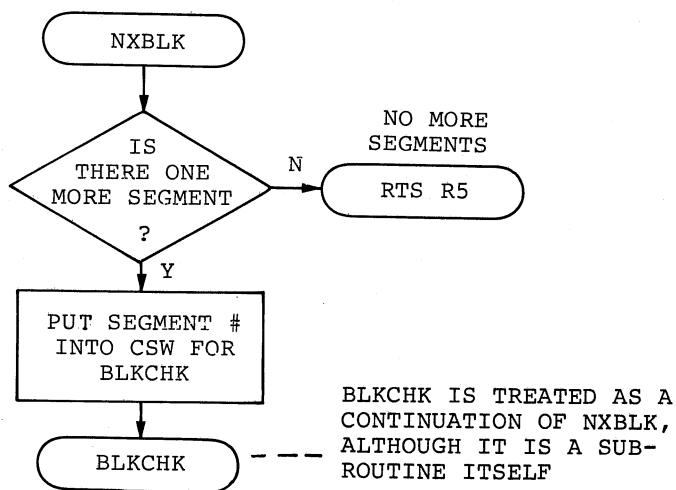


DLEET/NXBLK

DLEET - This routine scans a device directory to find a file of a specified name.

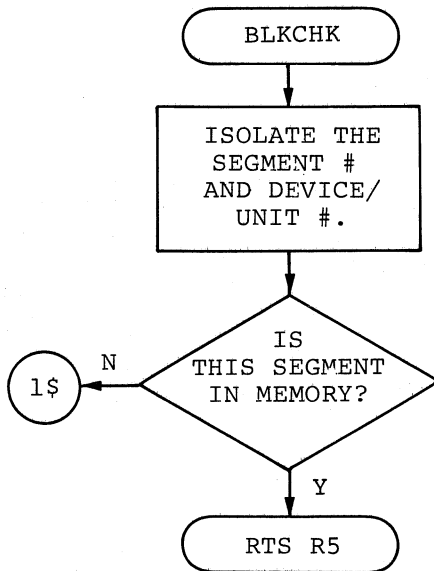


NXBLK - Gets the next in the series of directory segments, if one exists.

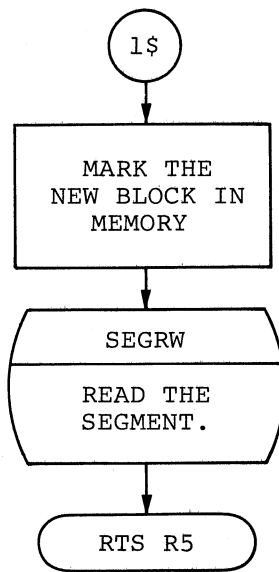


BLKCHK

BLKCHK - This routine isolates the segment number contained in bits 8-12 of the CSW, and checks to see if that segment is in memory at the current time. If not, it is read in.



Note that not only must the segment numbers agree, but also the device and unit numbers must be the same.



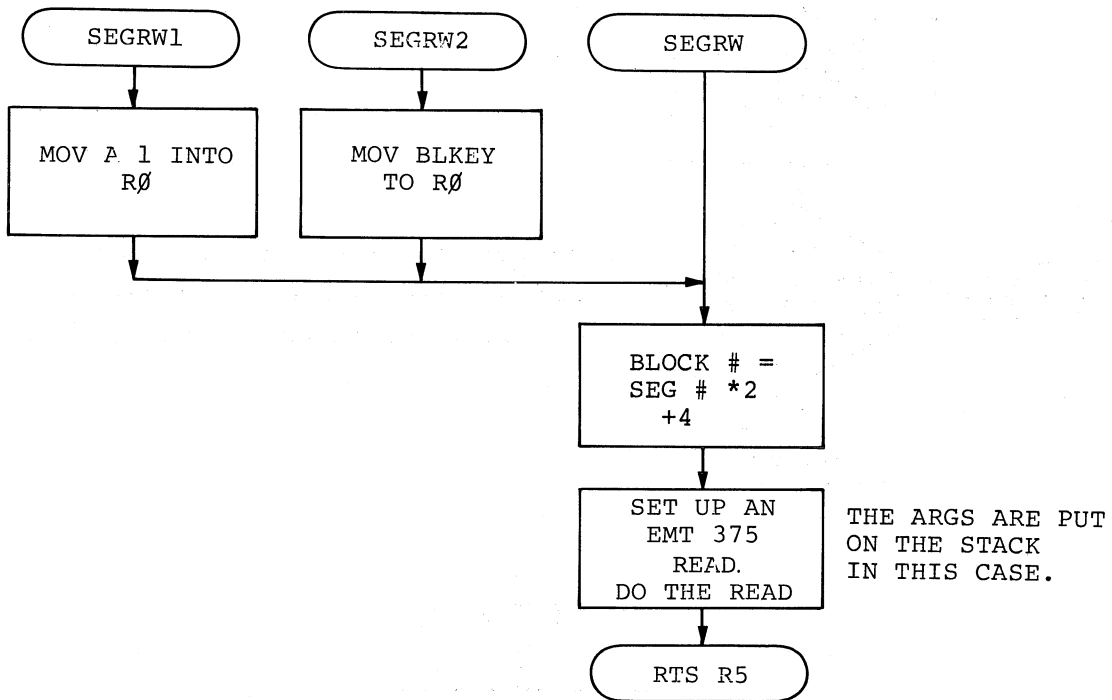
## SEGRW

SEGRW - Segment Read/Write. This routine read/writes selected directory segments. There are three entry points:

SEGRW1: Use segment #1

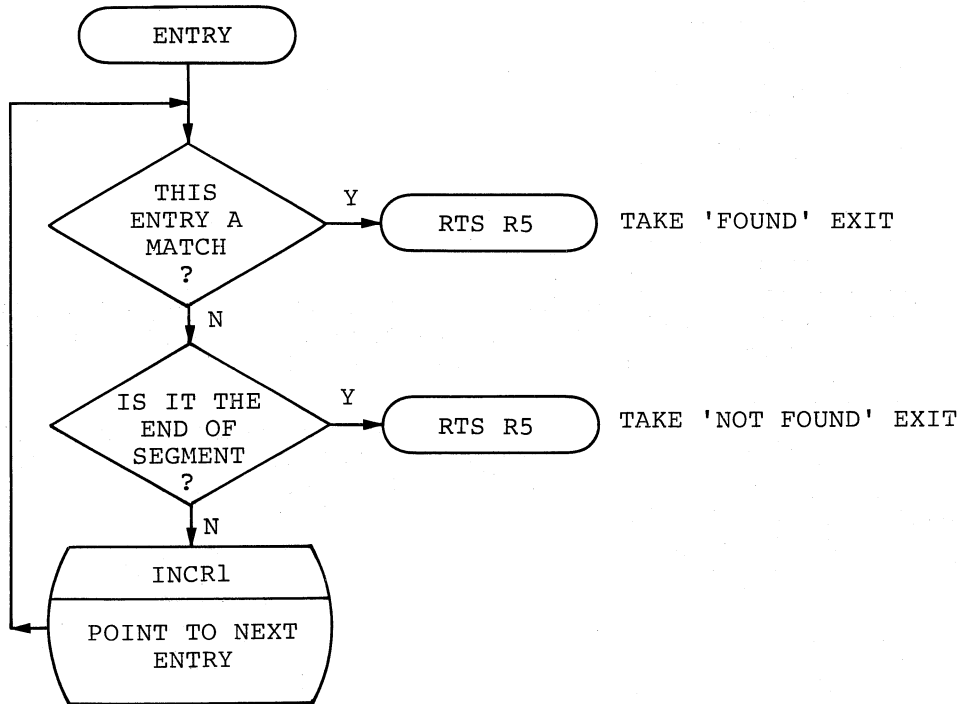
SEGRW2: Use the segment currently in memory (BLKEY)

SEGRW: Use the number in R0 as the segment #.



ENTRY/INCRL/COMERR/SPESHL

ENTRY - This routine uses R1 as a pointer into a directory segment to find a specified file type (Permanent, Tentative, Empty) or the end of segment mark.



INCRL - This routine bumps R1 to the next entry in a directory segment.

COMERR - This routine generates a fatal error from the USR. The call is:

```

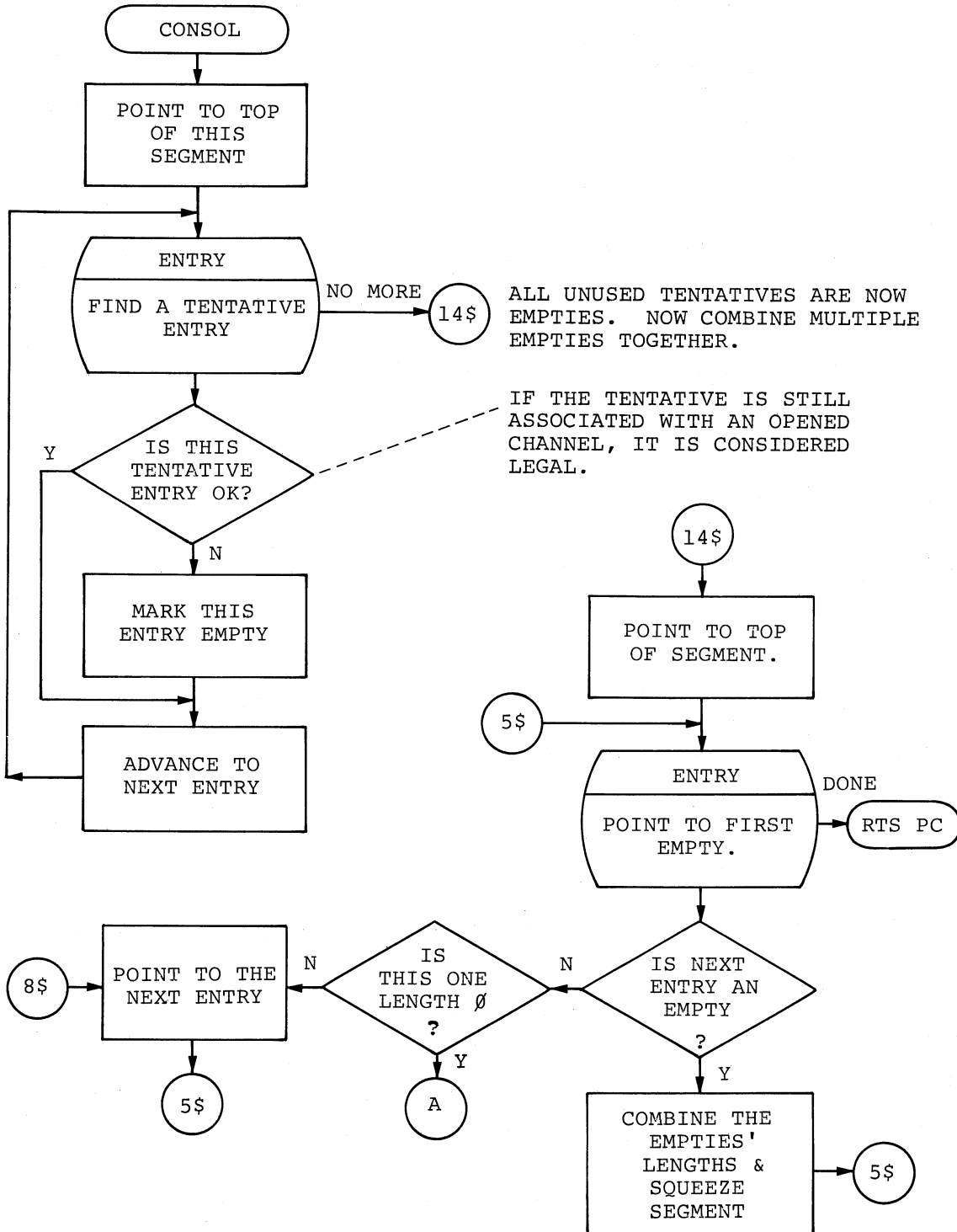
JSR R5,COMERR
code
  
```

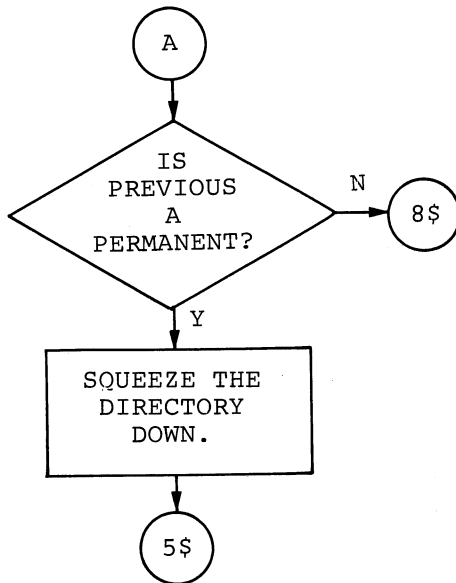
Code is used to indicate which error is to be generated. If .SERR is in effect, control passes to COMXIT, which returns to RMON.

SPESHL - This routine is used to effect file operations on MT/CT. This is done by passing a READ request to the Q manager. The even byte of the completion function will contain a 377. The queue manager detects this, and modifies the I/O queue element to indicate that the handler should perform a USR function.

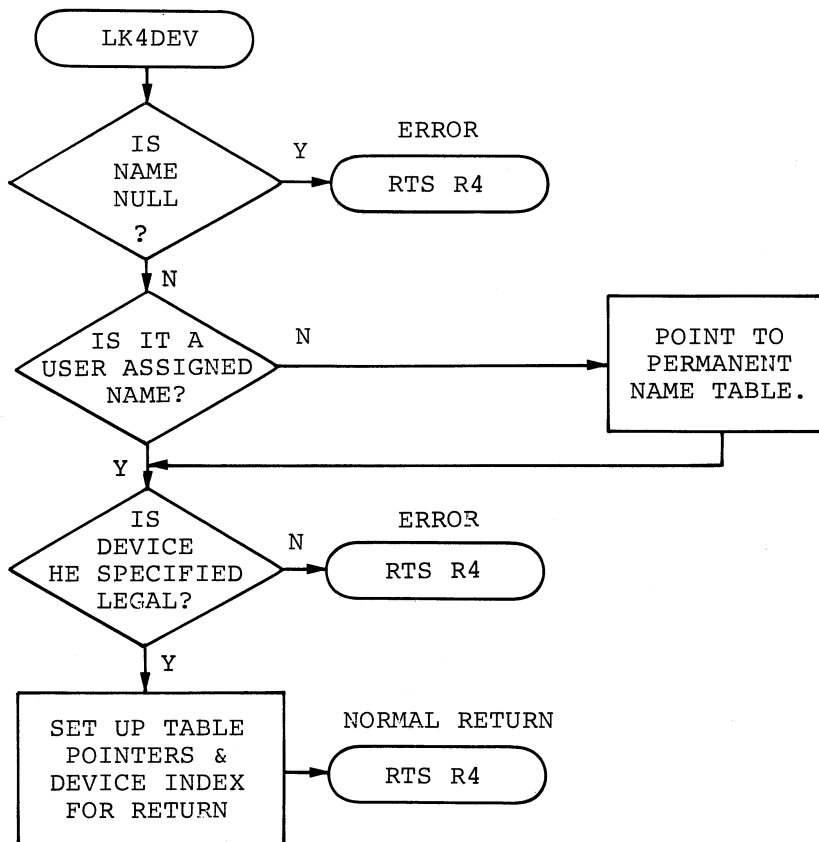
CONSOL

CONSOL - This routine is used to compact a directory segment. It combines consecutive empties into one, and makes empties out of tentative files which are not associated with an active channel.





LK4DEV - This routine looks up a specified device name in the system tables. It first attempts to find the name in the user assigned name table; failing that, the permanent name table is searched.



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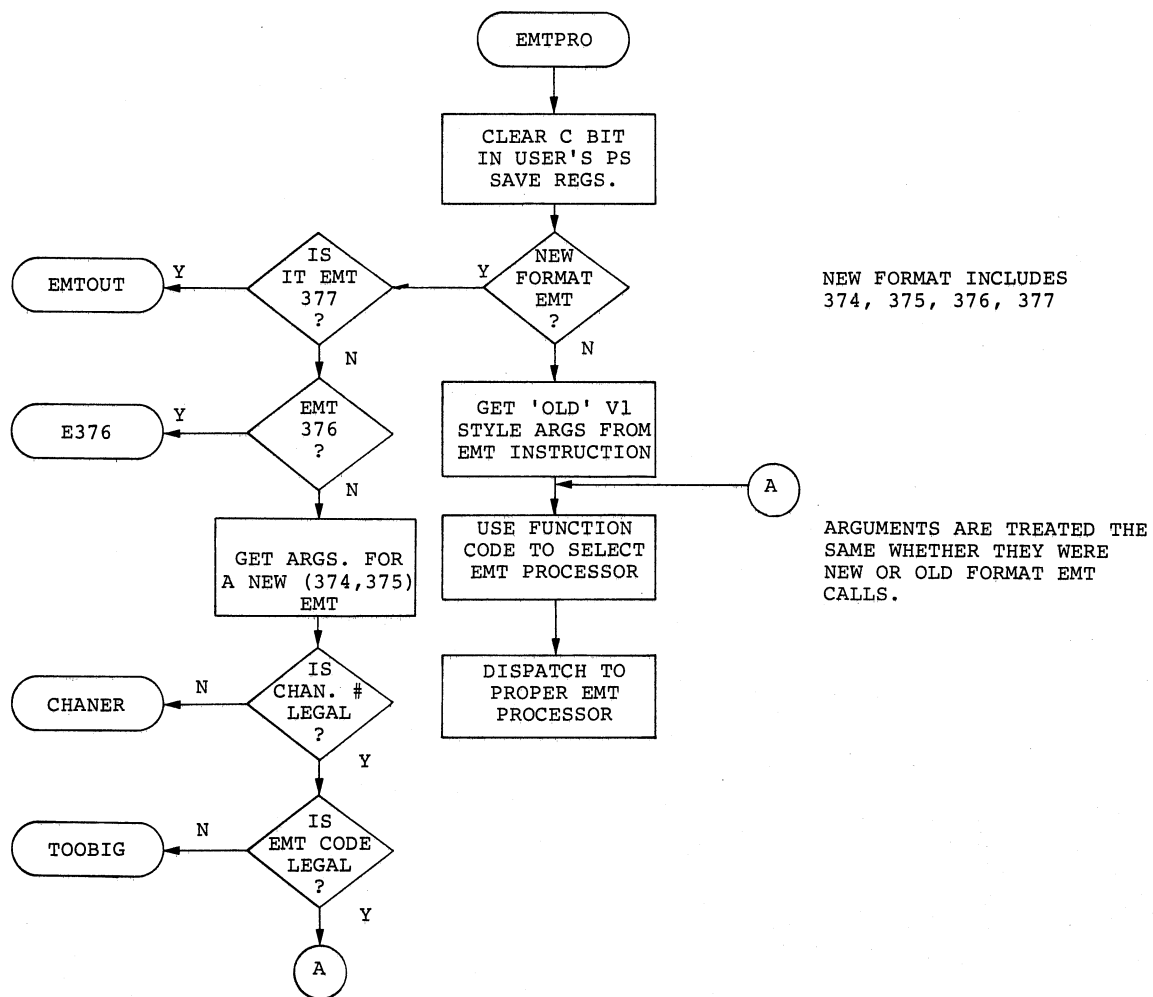
(



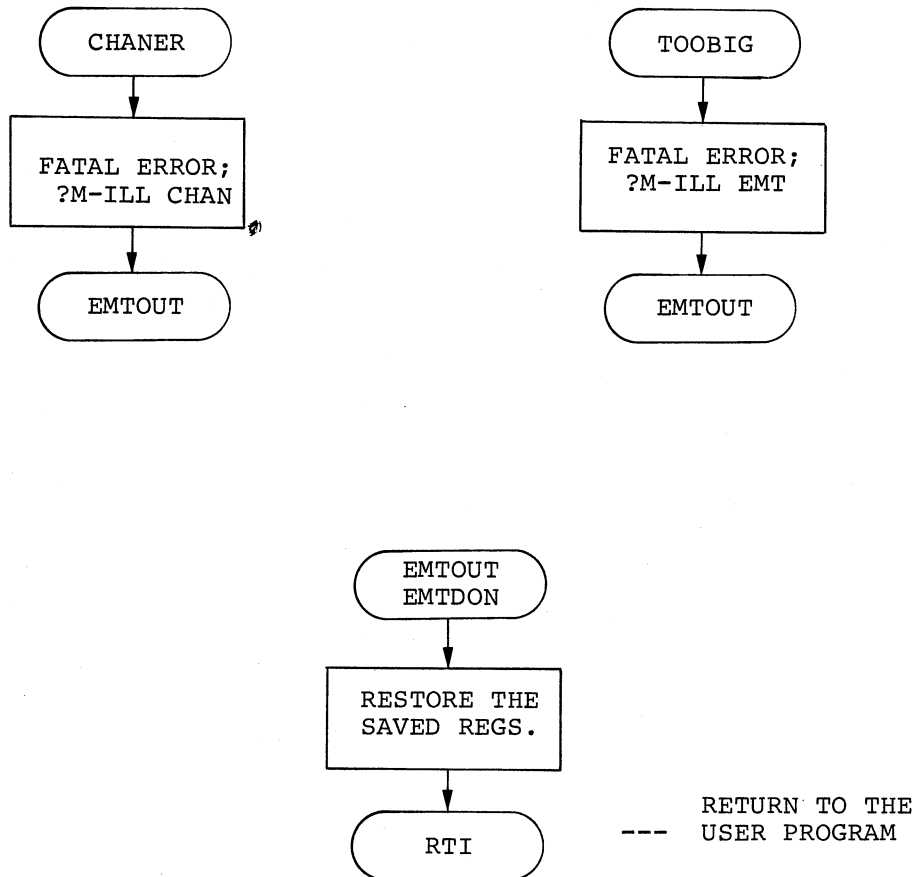
E.4 RMON (RESIDENT MONITOR) FLOWCHARTS FOR SINGLE-JOB MONITOR

# EMT DISPATCHER

The code of the EMT dispatcher is entered when an EMT instruction is executed. The EMT instruction is decoded and control passes to the appropriate code for processing.



EMT DISPATCHER (CONT.)



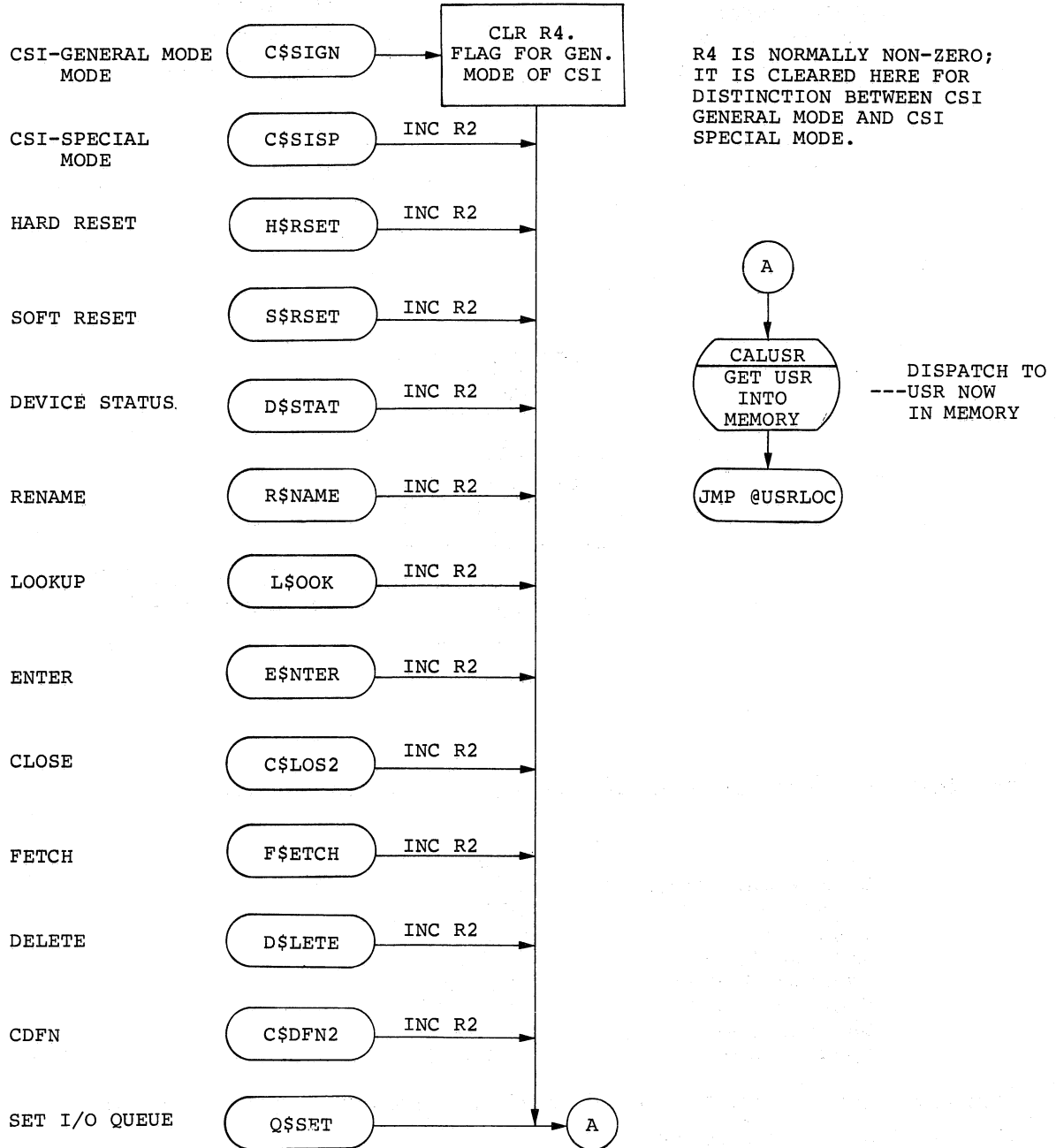
The following EMT requests are no-ops in the S/J Monitor:

Mark Time	.MRKT
Cancel Mark Time	.CMKT
Timed Wait	.TWAIT
Send Data	.SDAT
Receive Data	.RCVD
Channel Status	.CSTAT
Protect Vectors	.PROTECT
Channel Copy	.CHCOPY
Special Device	.DEVICE

Executing these requests in S/J will cause an immediate successful returns with no action taken.

USR DISPATCHER TABLE FOR EMT'S 340-357

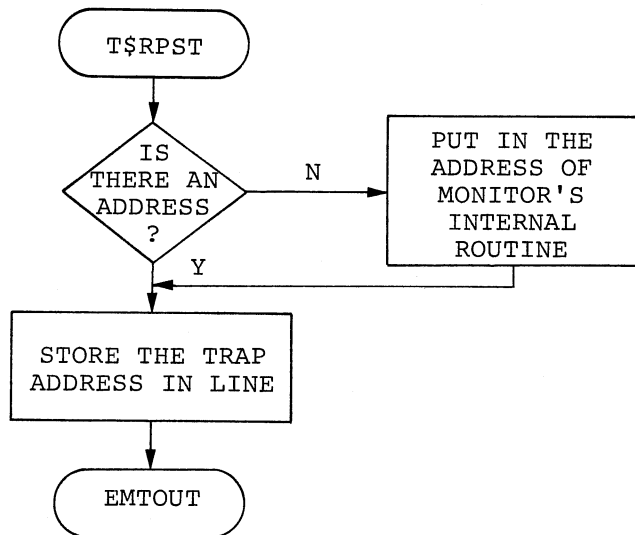
The USR Dispatch code handles dispatching those EMT's which require the USR. At each entry point, an INC R2 is performed. Thus, R2 acts as a function identifier once the USR is entered.



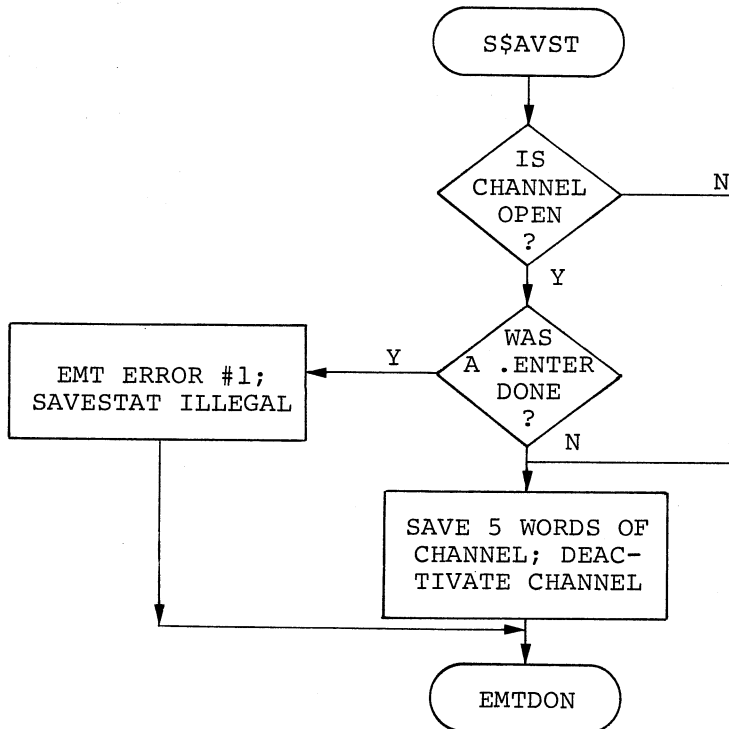
E.4.1 EMT Processors

SET TRAP/SAVE STATUS

SET TRAP ADDRESS

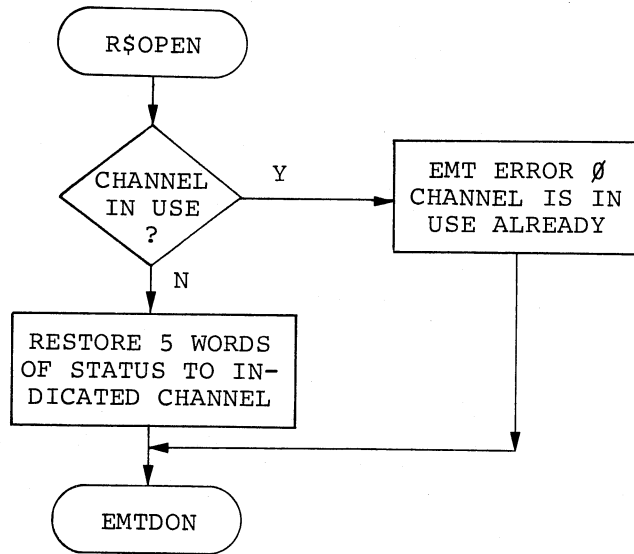


SAVESTATUS

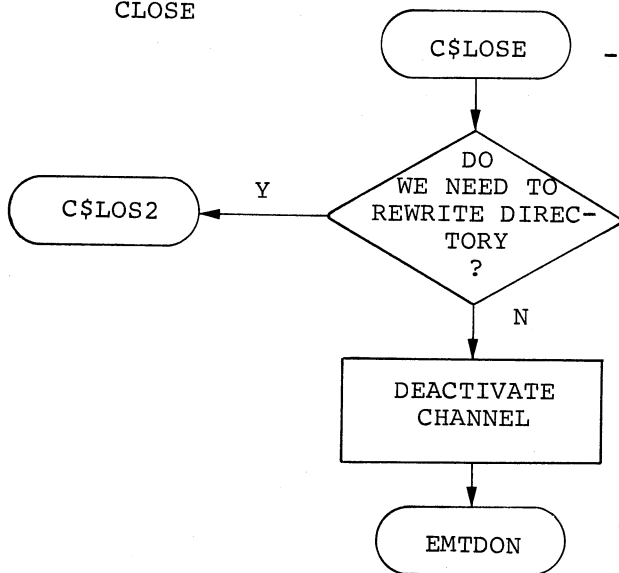


REOPEN/CLOSE/RDOVLY

REOPEN

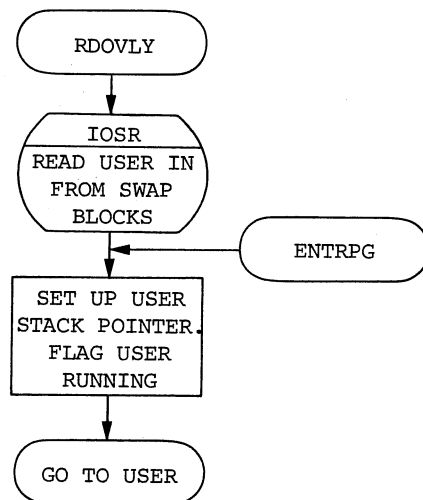


CLOSE

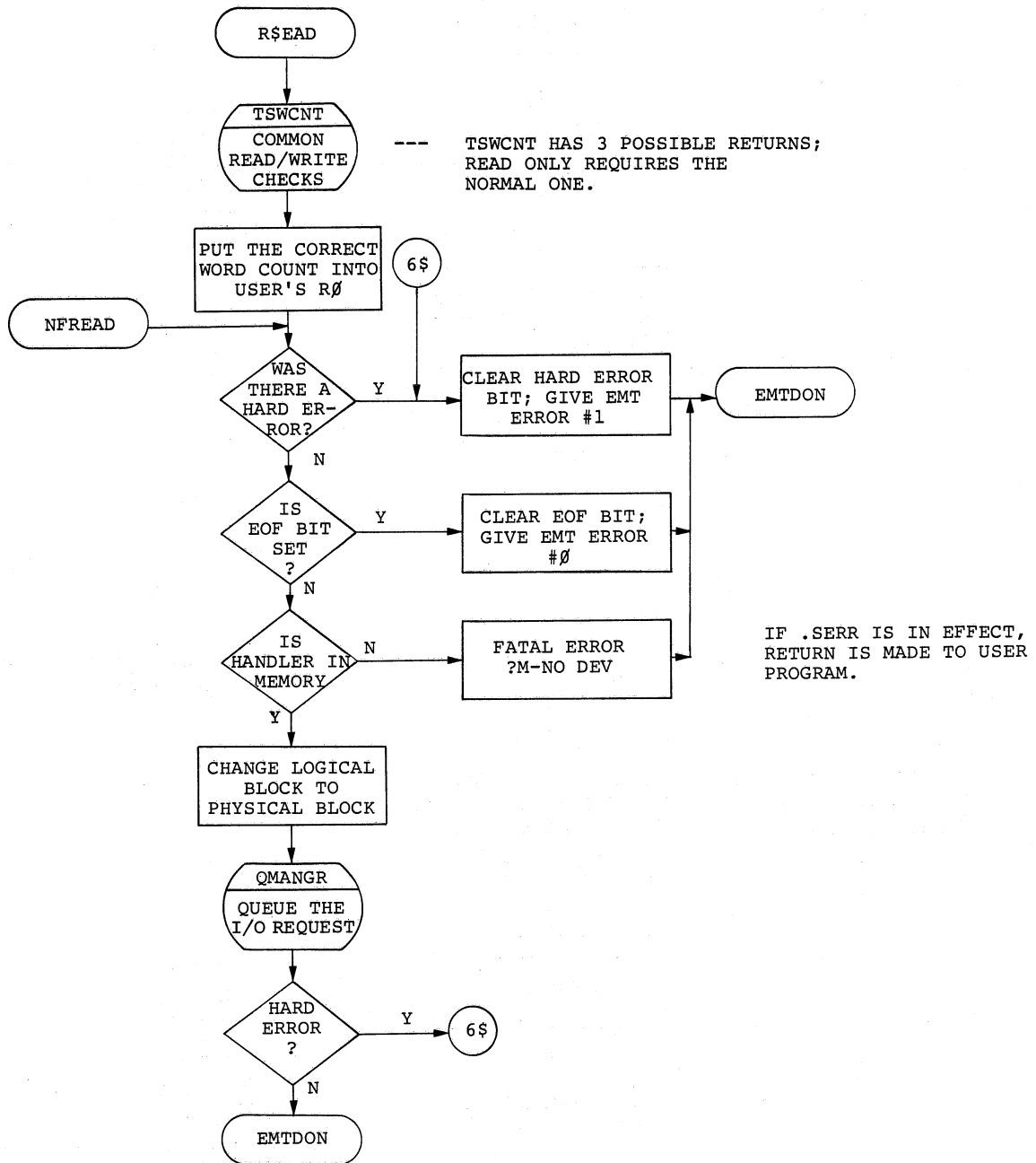


--- IF A LOOKUP WERE DONE, THE USR IS NOT REQUIRED.

RDOVLY

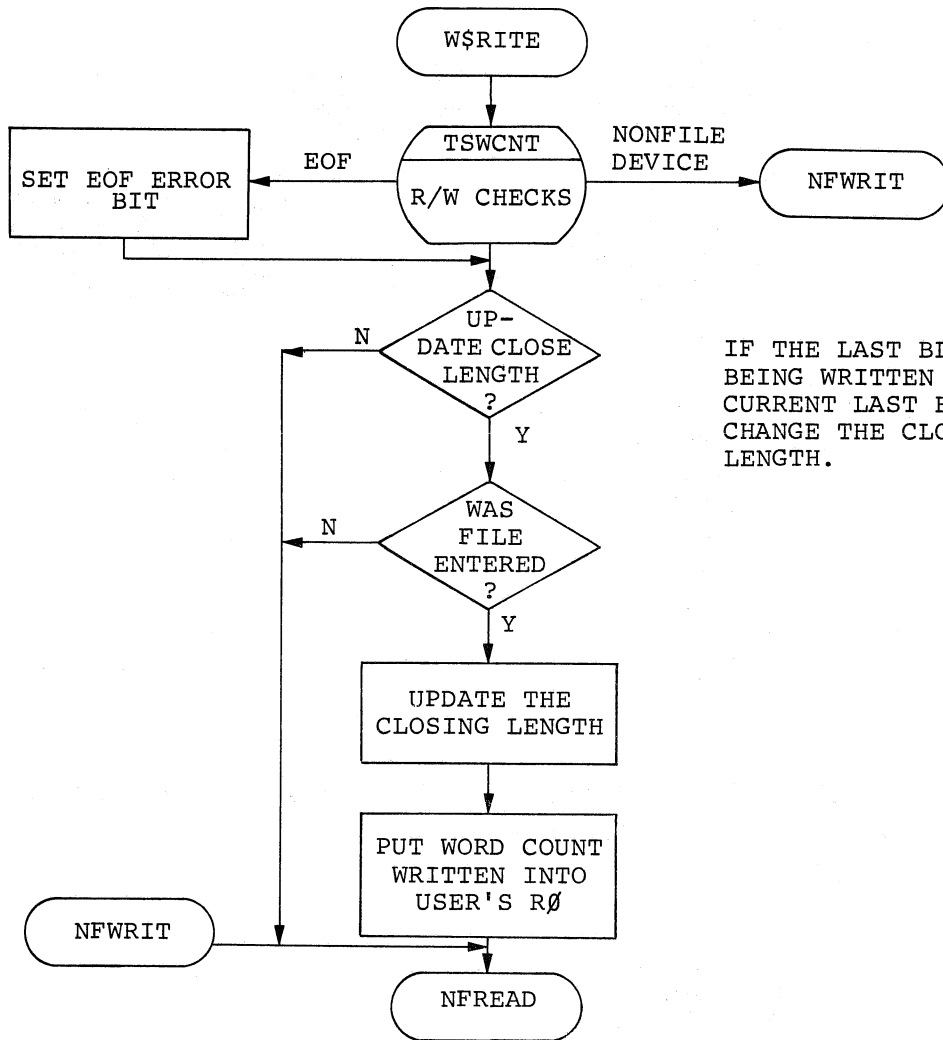


READ



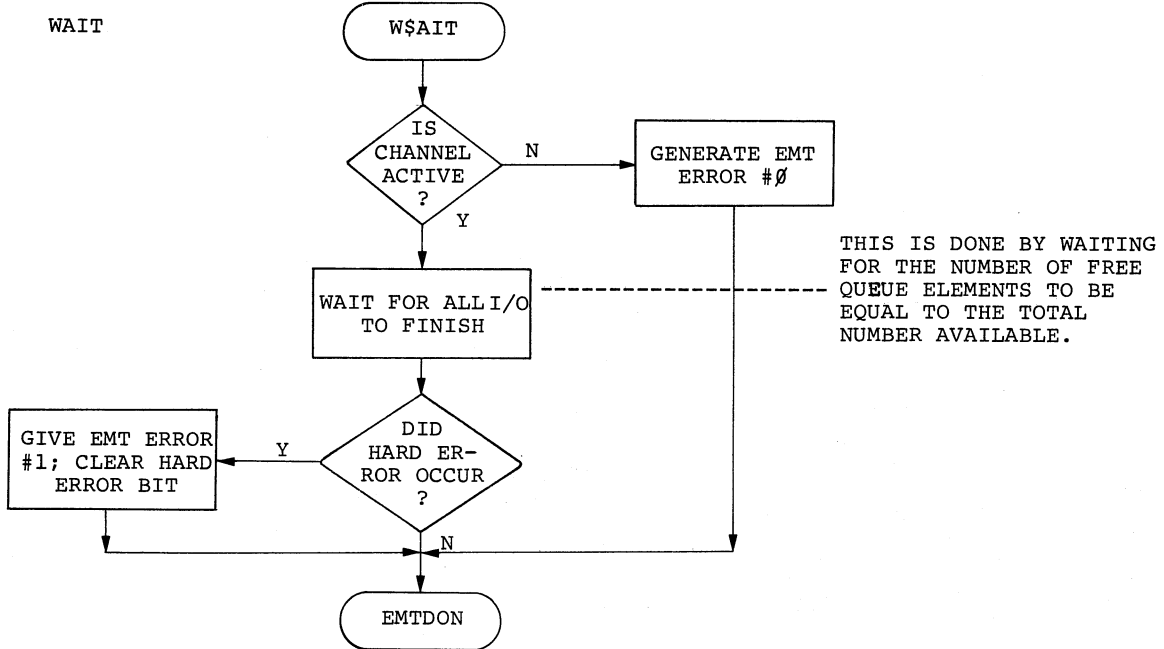


WRITE



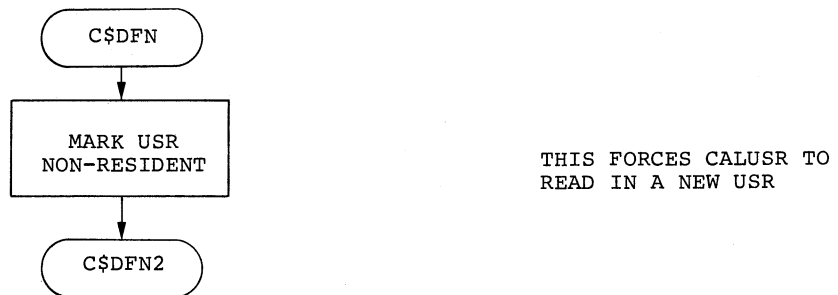
IF THE LAST BLOCK BEING WRITTEN IS > THE CURRENT LAST BLOCK WRITTEN, CHANGE THE CLOSING FILE LENGTH.

WAIT/CDFN



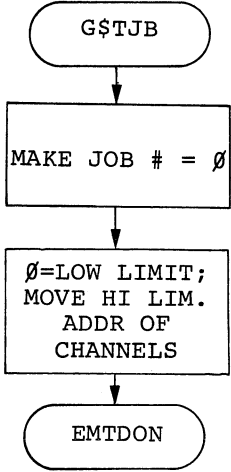
CDFN

Channel Define - the resident portion of CDFN causes a fresh copy of the USR to be read in, then enters the USR.

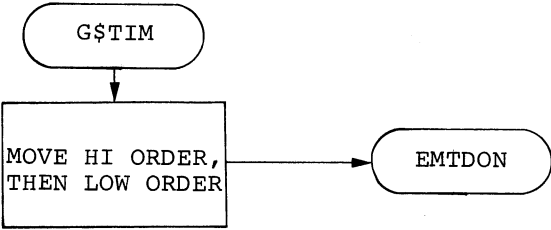


GET JOB PARAMETERS  
GET TIME OF DAY  
SET FPP EXCEPTION

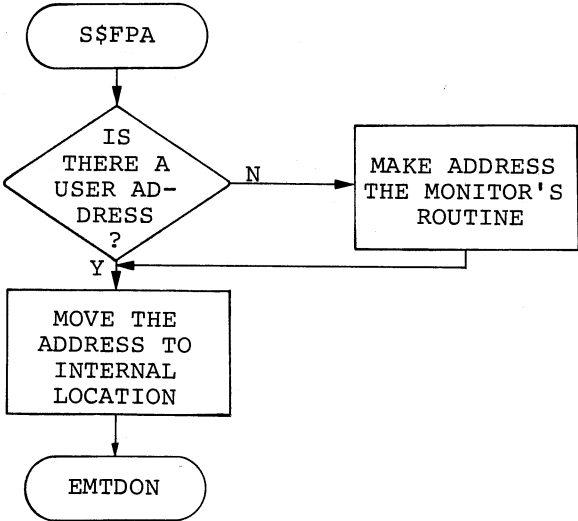
GET JOB PARAMETERS



GET TIME OF DAY

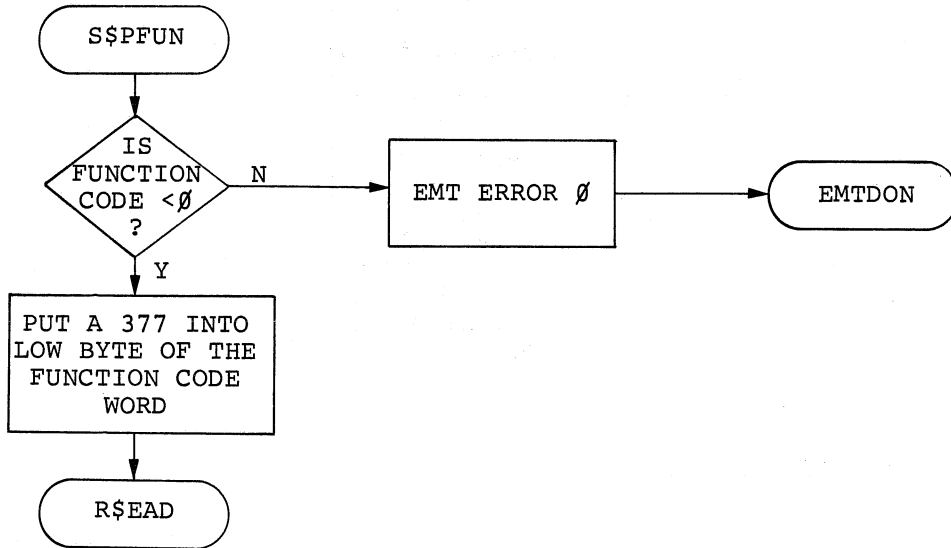


SET FPP EXCEPTION



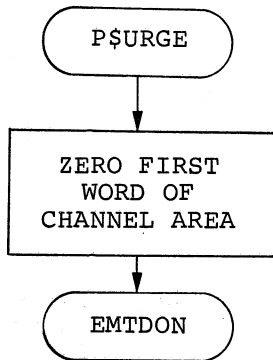
SPECIAL FUNCTIONS/PURGE  
SOFT/HARD ERRORS

SPECIAL FUNCTIONS (MAGTAPE/CASSETTE)

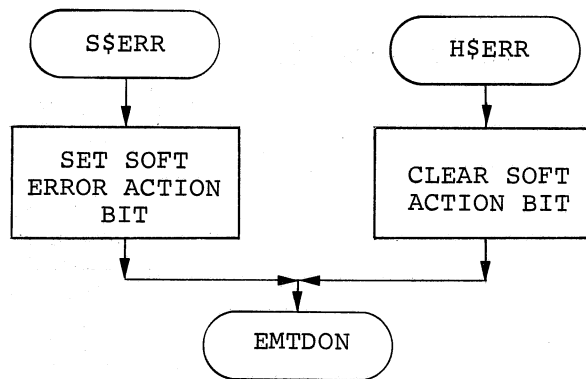


SPECIAL FUNCTIONS/PURGE  
SOFT/HARD ERRORS

PURGE

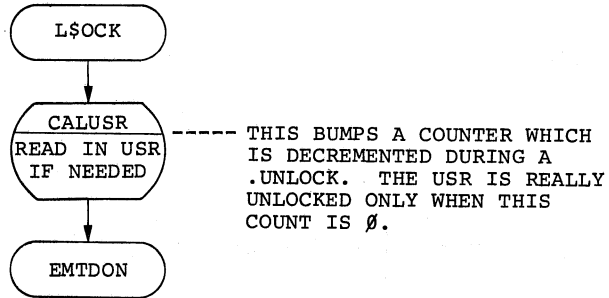


SOFT/HARD ERRORS

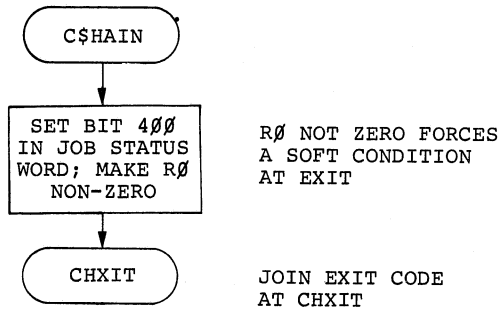


LOCK USR/CHAIN/UNLOCK USR

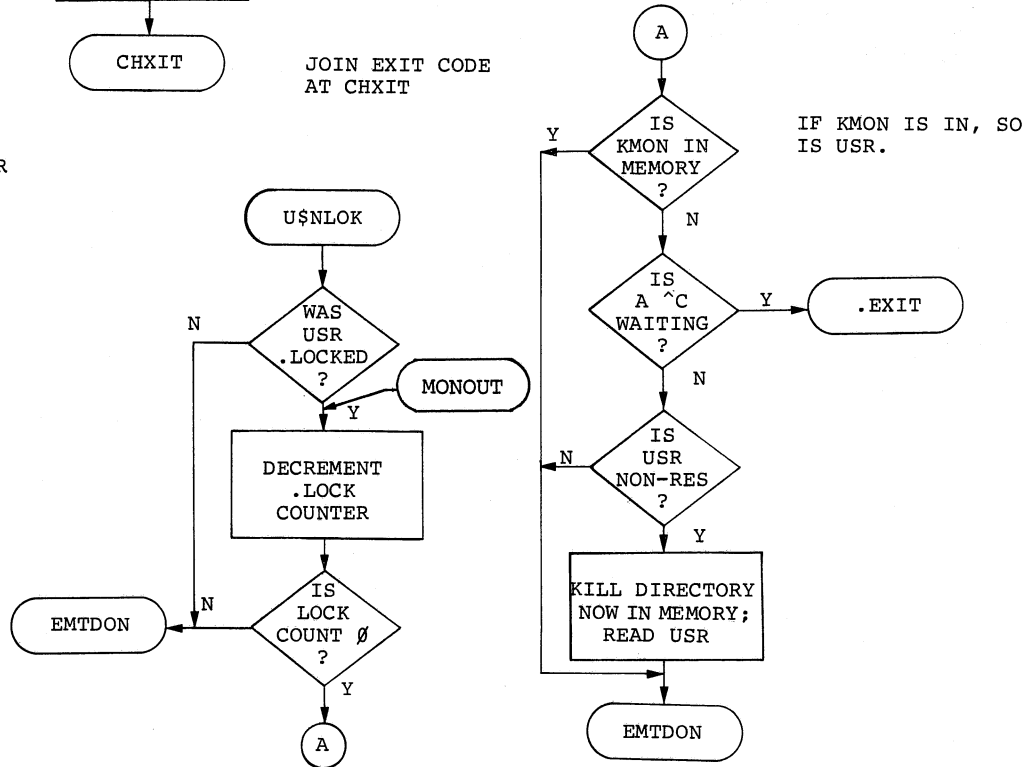
LOCK USR



CHAIN

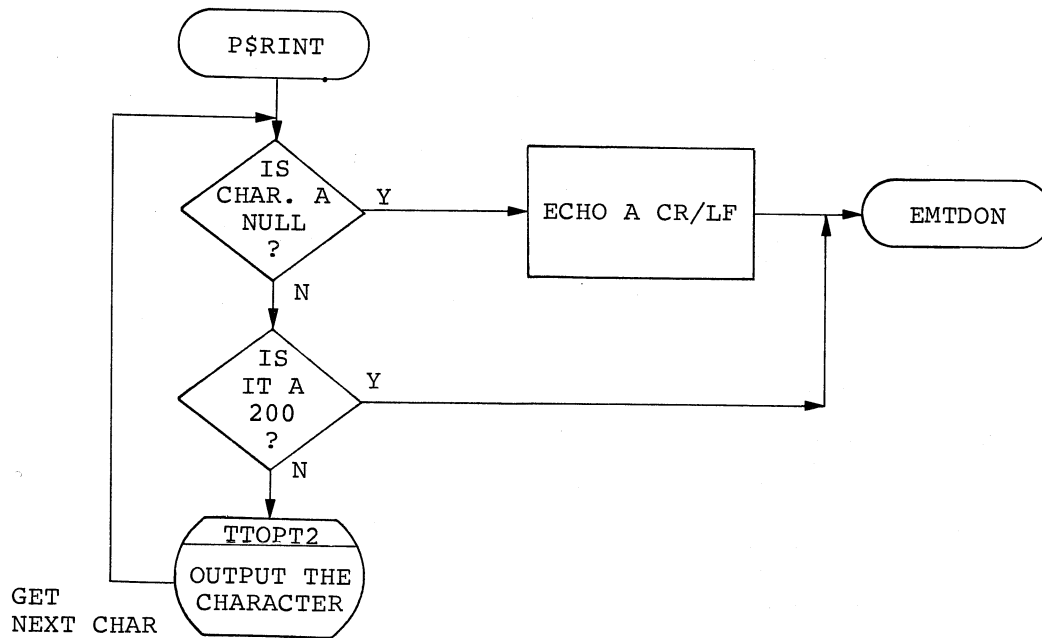


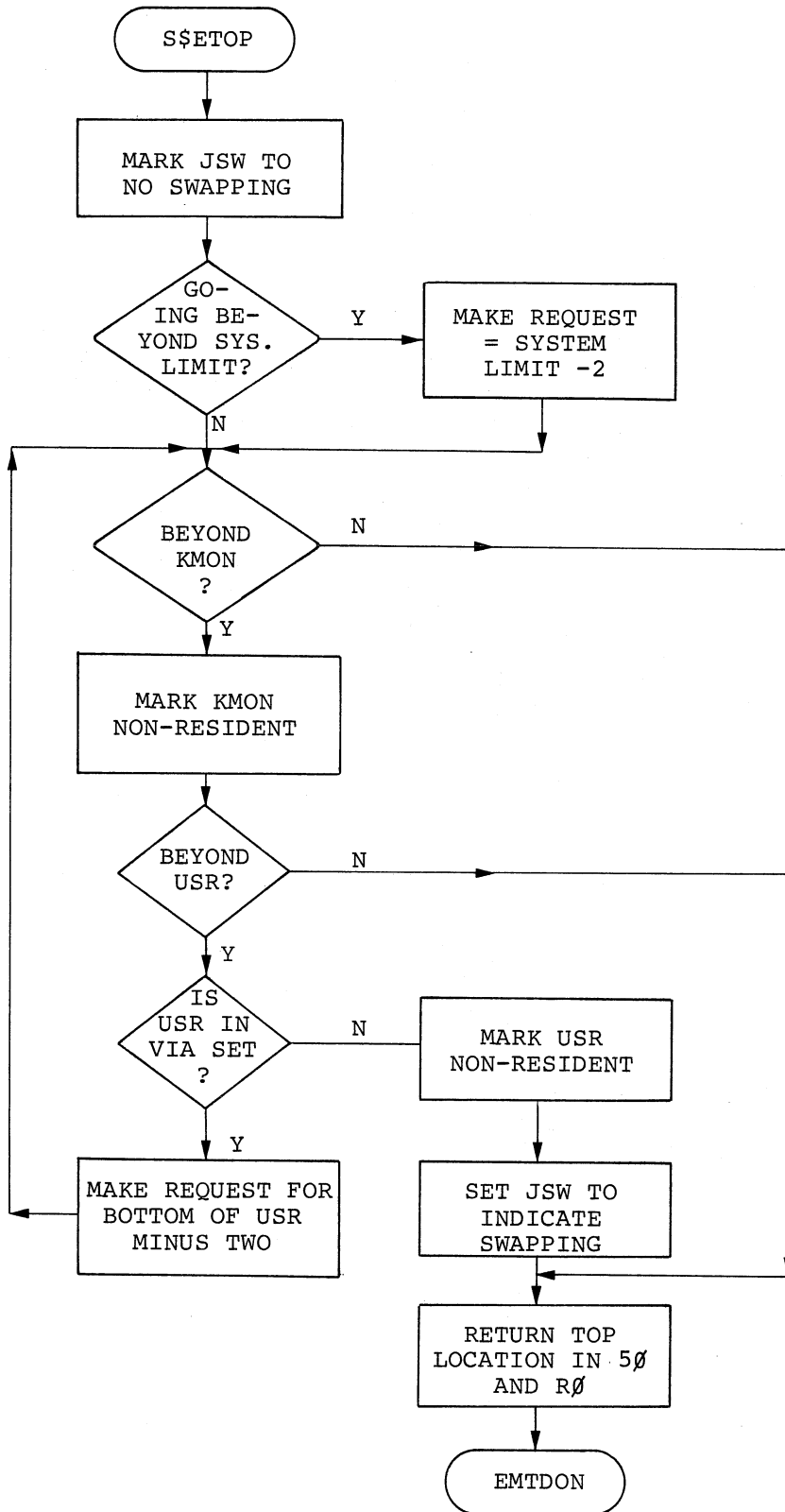
UNLOCK USR



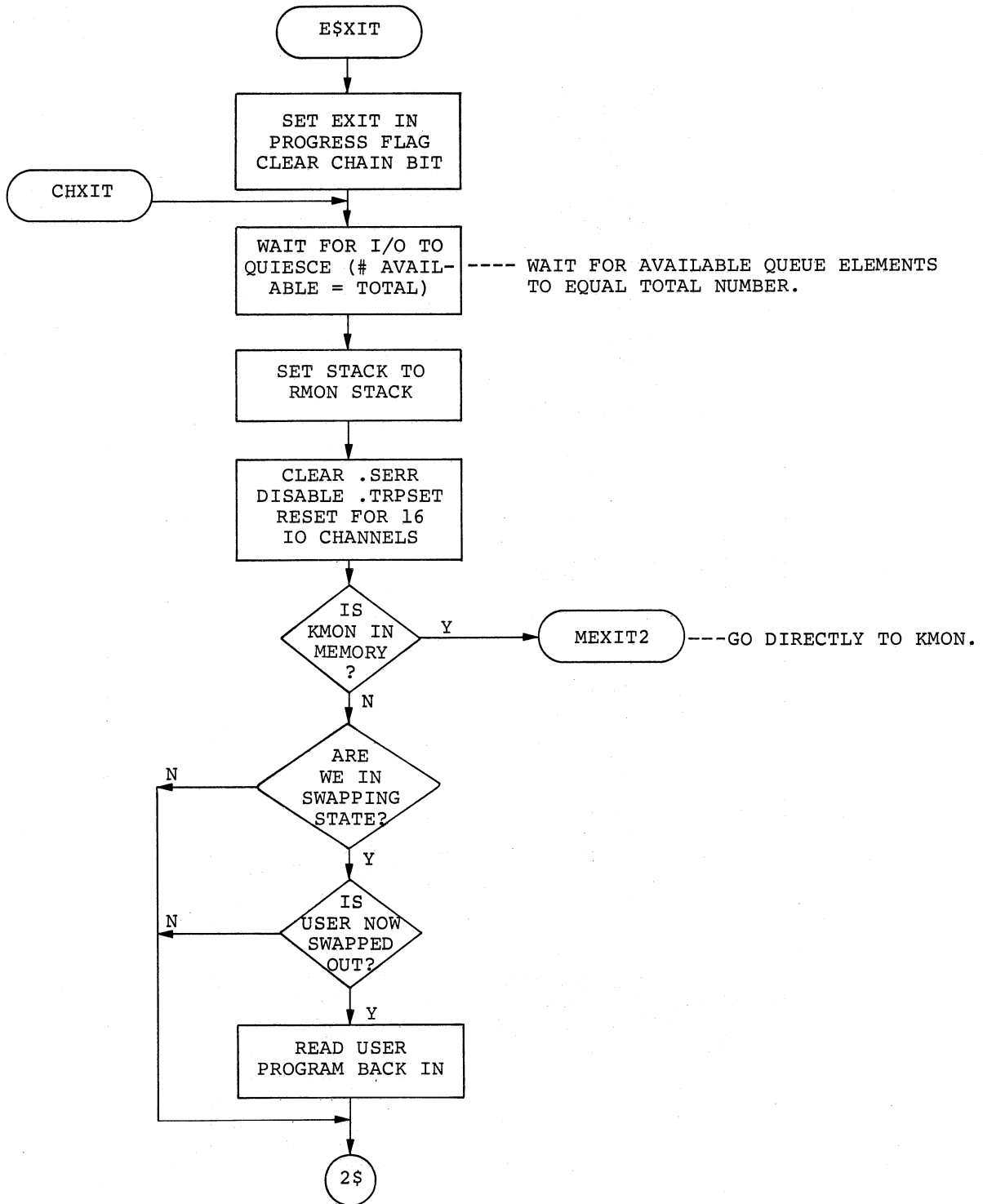
PRINT

PRINT - Causes a line to be output to the console terminal.

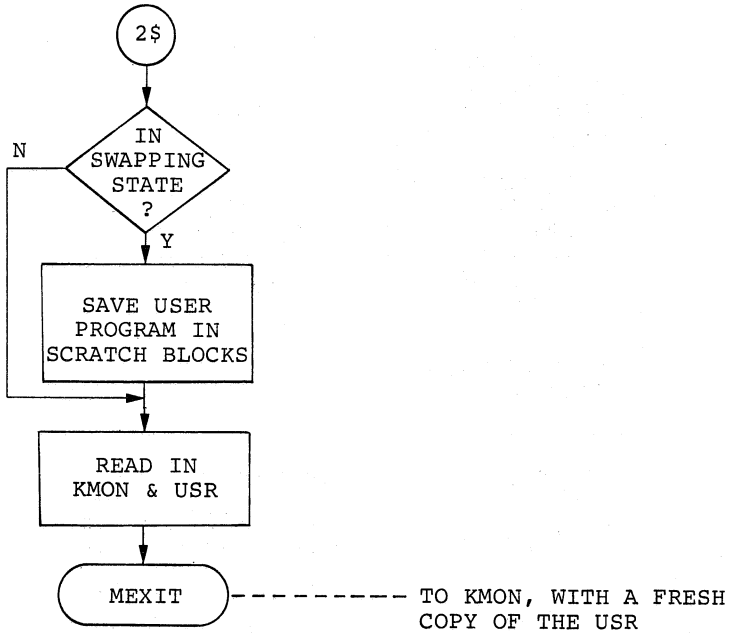




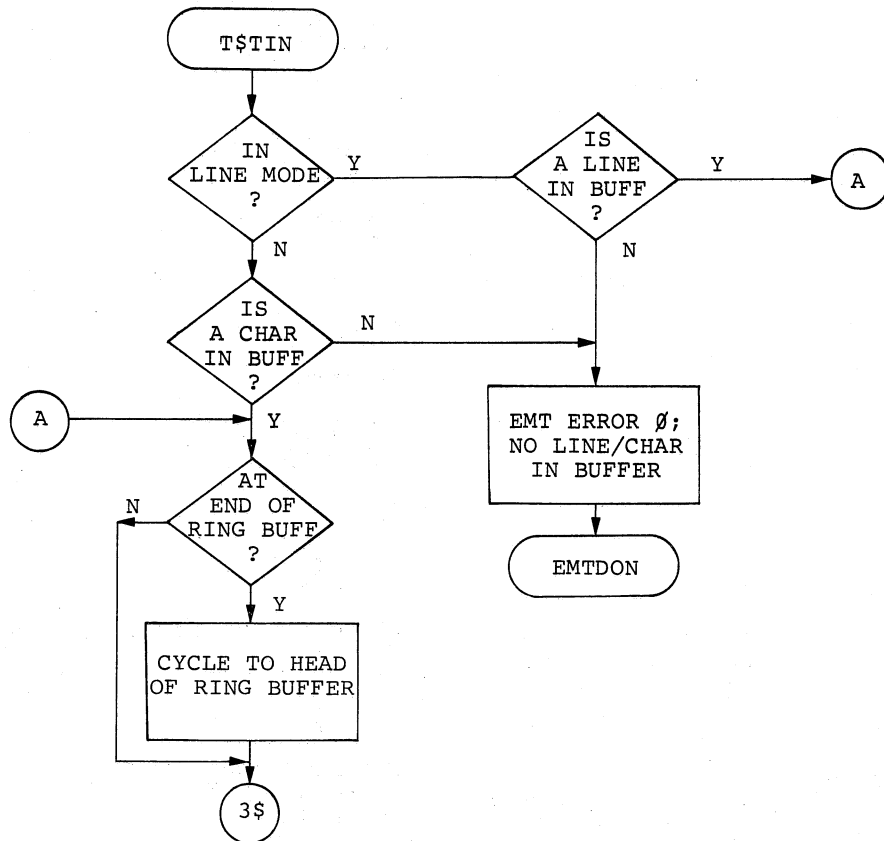
EXIT



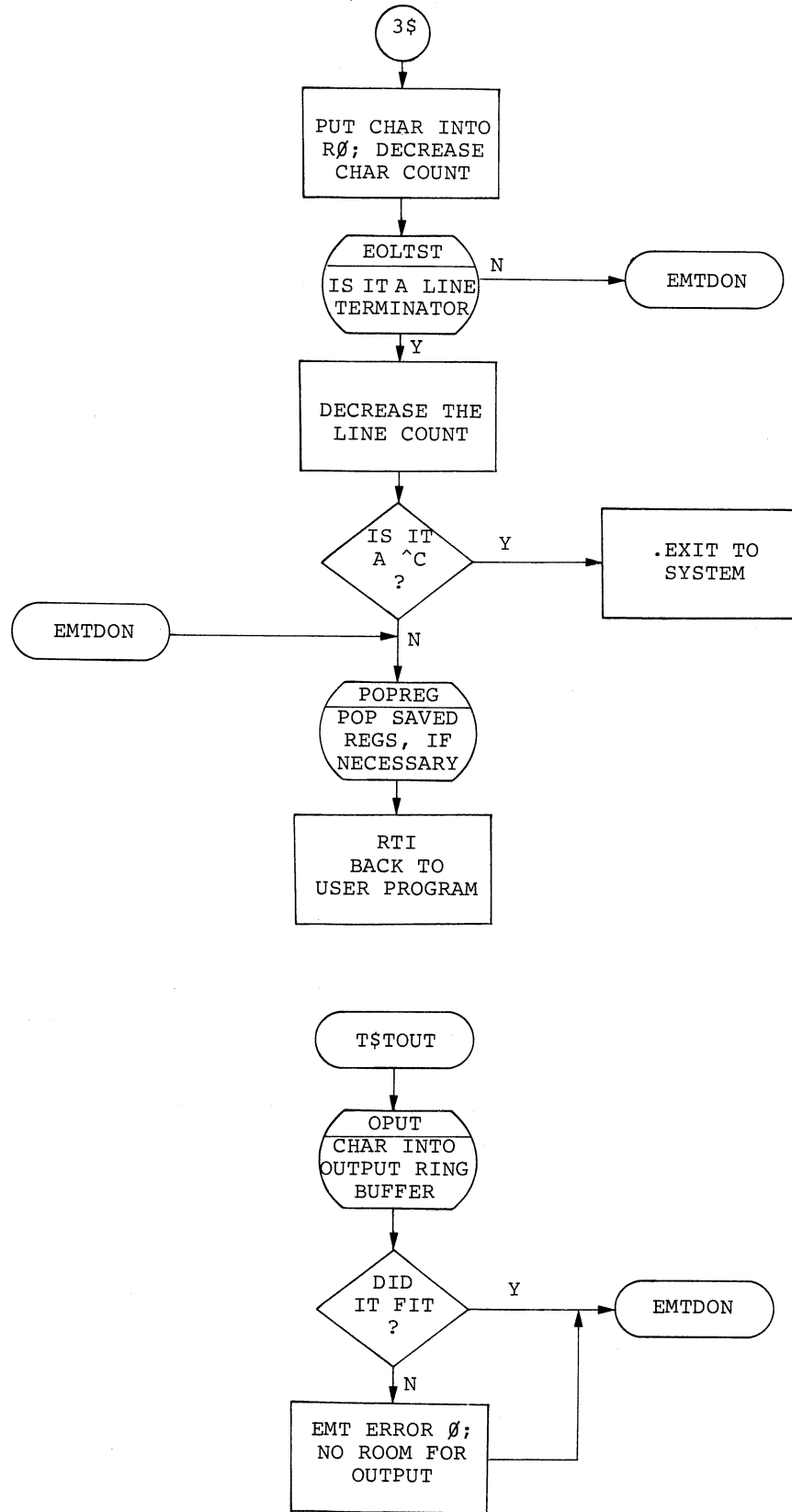




TTYIN



TTYIN (CONT.)/TTYOUT

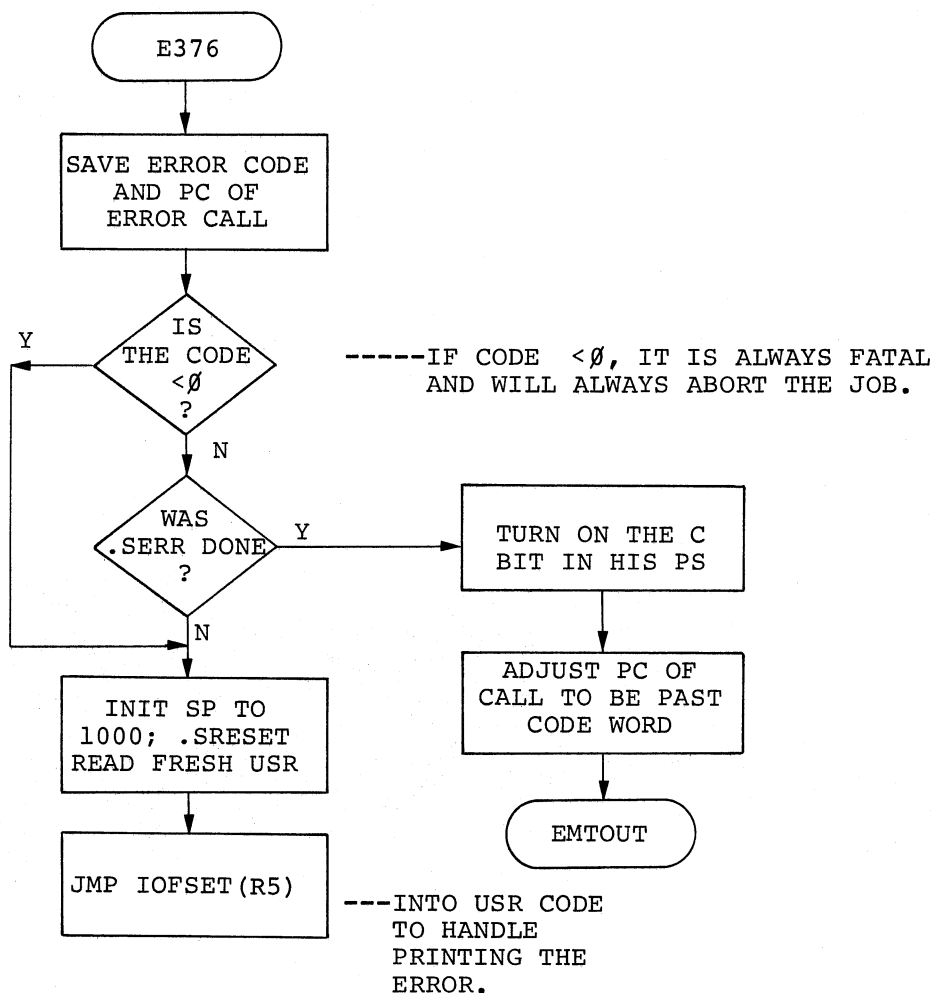


FATAL ERROR PROCESSING

EMT 376 is reserved for reporting fatal monitor errors. When a fatal error condition is encountered, a call of the form:

EMT 376  
code

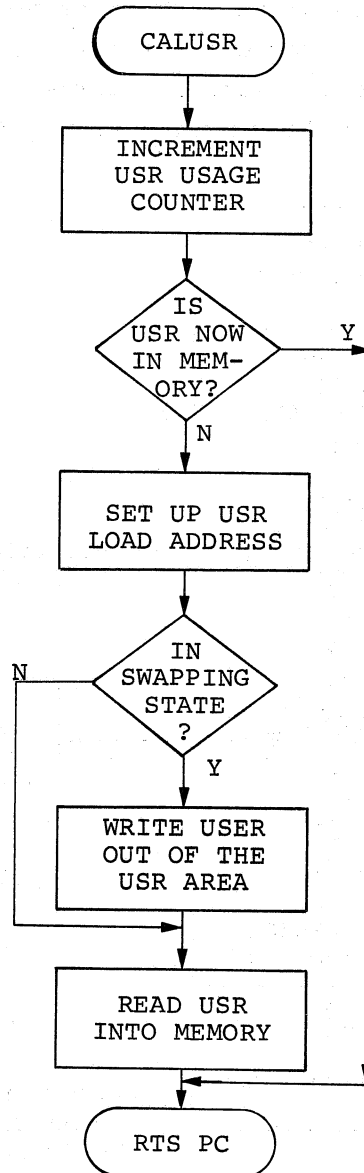
is executed. This indicates to RT-11 that a fatal error has occurred. The normal response is to print a ?M-error message and then abort the job. However, if a .SERR request has been done, no message will occur and control will pass to the user's program. The error bit (C bit) will be set and byte 52 will contain the negative of the error code.



# CALUSR

CALUSR is used to ensure that the USR is in memory for a USR type request. It will handle the situation where the user program must be written to scratch blocks before the USR is read in. Entry is made at CLUSR2 when an error has occurred and the error processing code in the USR buffer is required.

EITHER 'NORMAL' VALUE  
OR WHAT THE USER HAS  
PUT IN LOCATION 46.



#### E.4.2 Clock Interrupt Service

The interrupt service for the clock is primitive. The clock vector is set up such that the interrupt routine is always entered with the C bit = 1. At the interrupt routine, the code is:

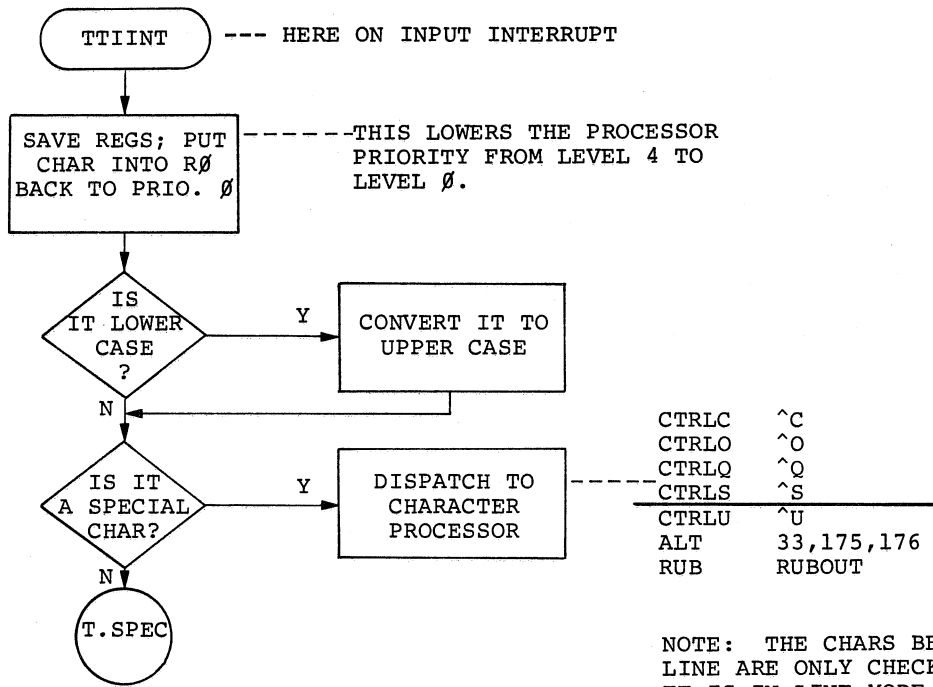
```
ADC $TIME+2
ADC $TIME
RTI
```

Since the C bit is 1, \$TIME+2 is incremented by the ADC. When the low order word goes from 177777 to 0, the C bit remains on and \$TIME is then incremented. No 24 hour wrap around is provided.



### E.4.3 Console Terminal Interrupt Service

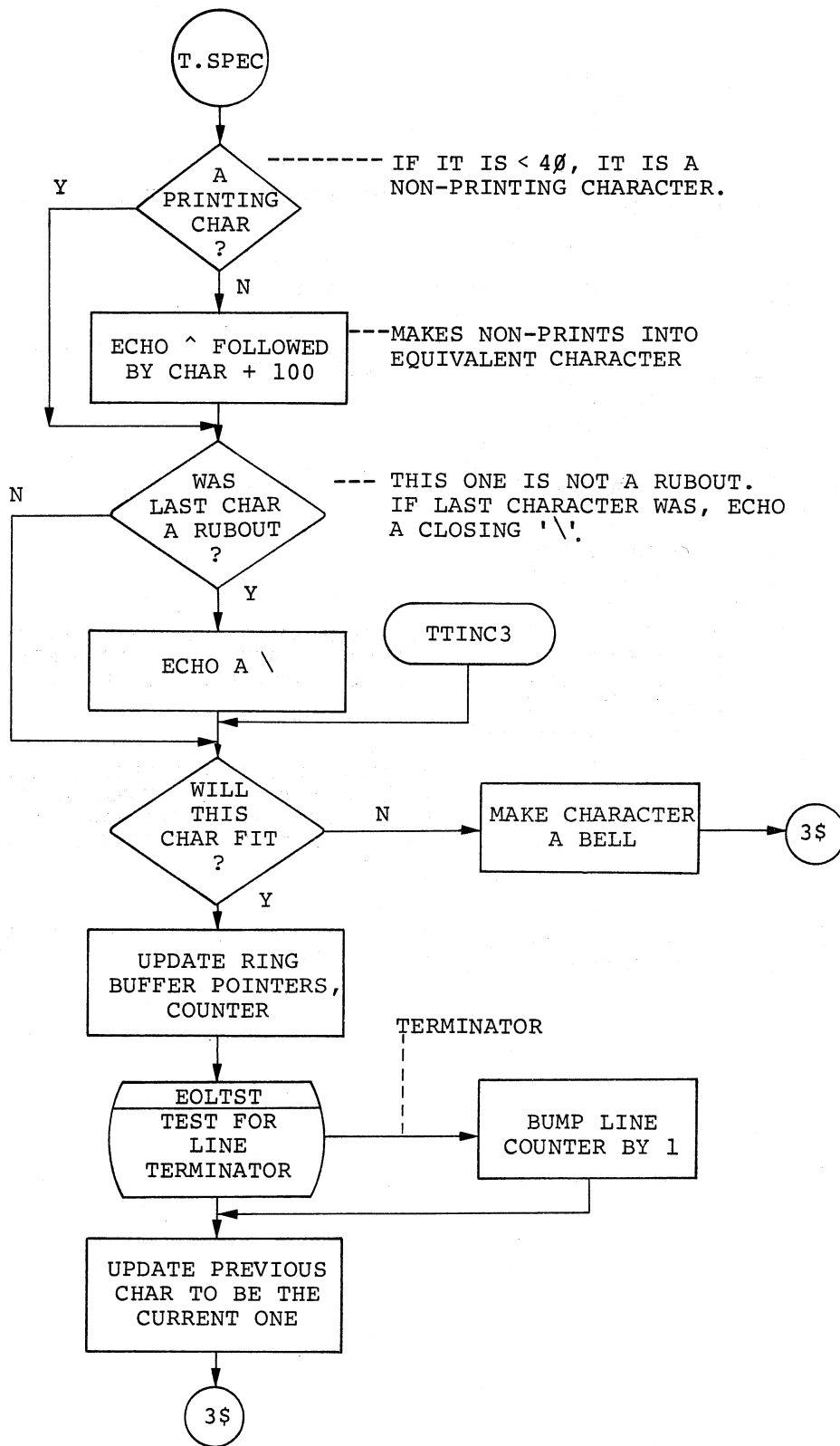
TT INPUT INTERRUPT SERVICE



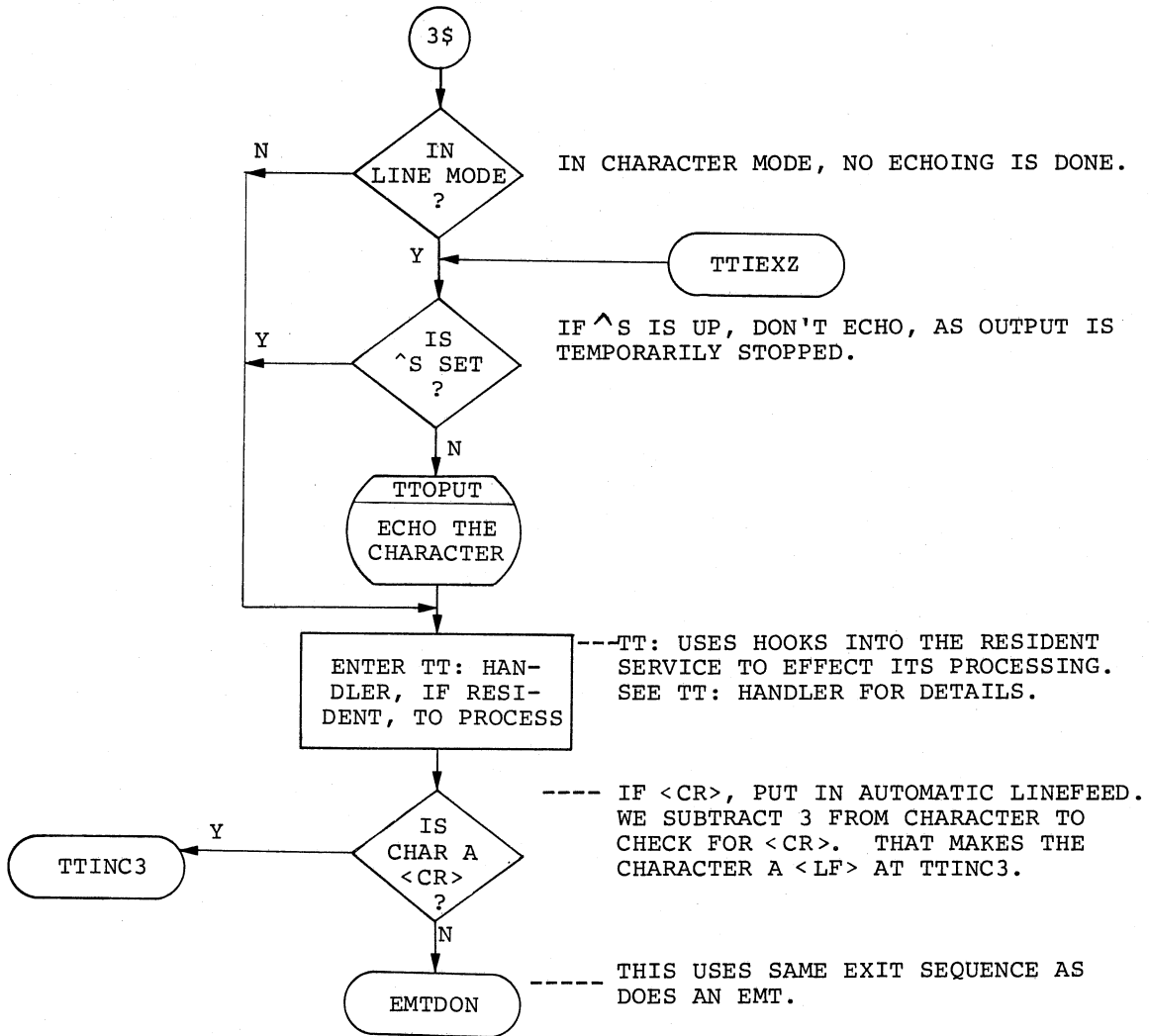
NOTE: THE CHARS BELOW THE LINE ARE ONLY CHECKED WHEN TT IS IN LINE MODE. IN CHARACTER MODE THEY ARE NOT ACTED UPON.



TT INPUT INTERRUPT SERVICE (CONT.)



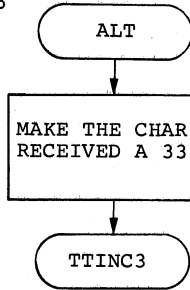
TT INPUT INTERRUPT SERVICE (CONT.)



ALTMODE/CONTROL O, S, Q

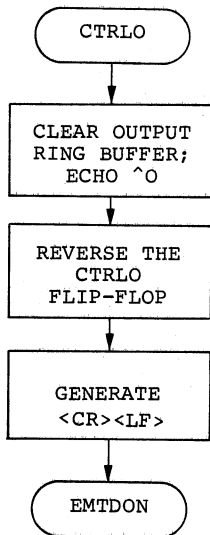
These routines are entered when any of the corresponding special characters are struck.

ALTMODE - 33,175,176

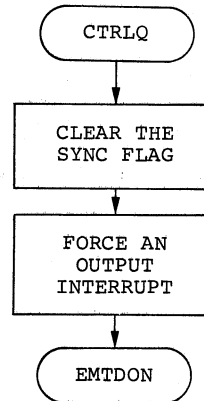
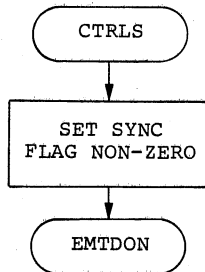


----- ALL POSSIBLE CODES ARE CONVERTED TO 33. AT SUBROUTINE TTOPUT, A 33 ABOUT TO BE ECHOED IS CHANGED TO A \$.

CONTROL O

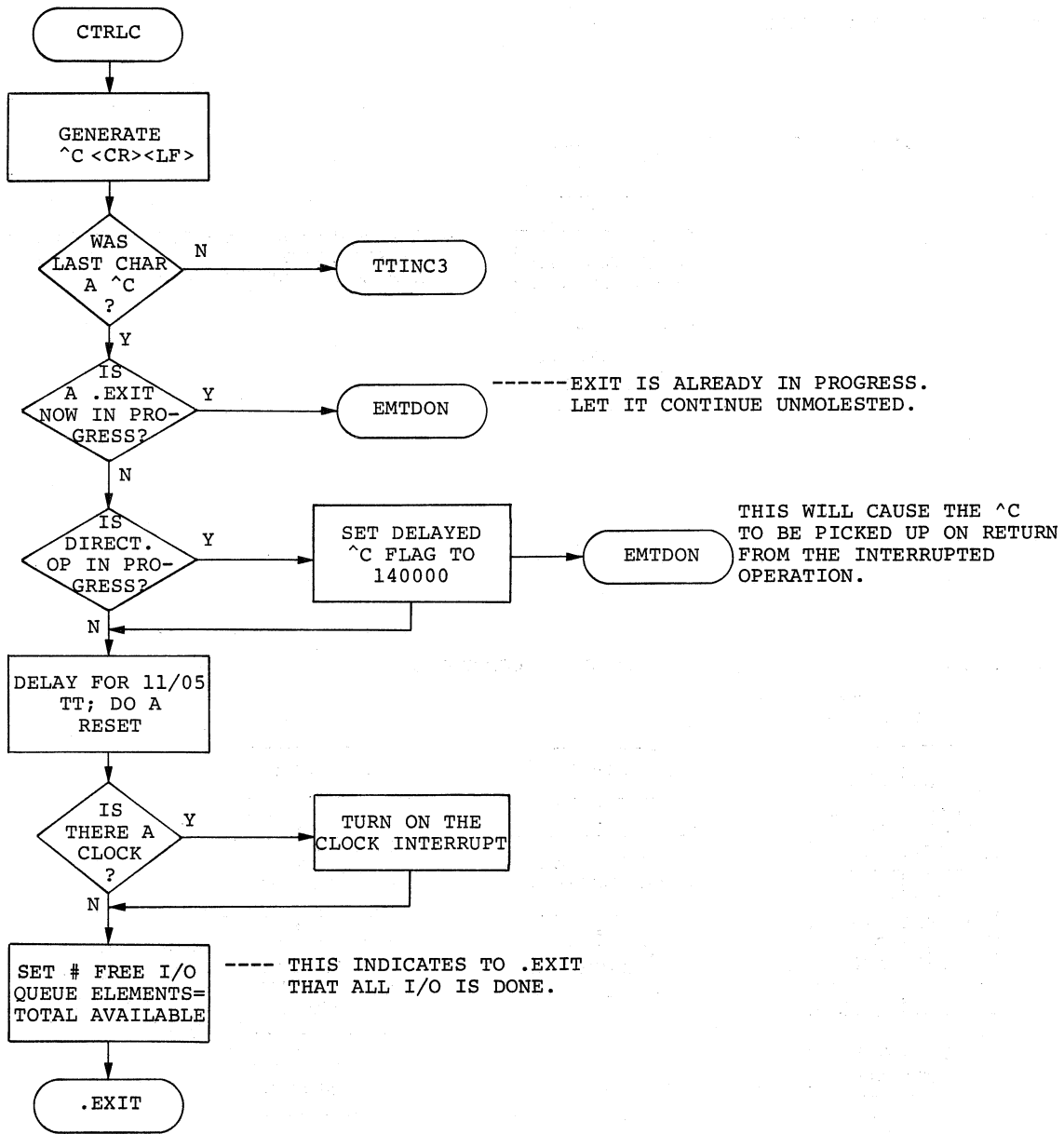


CONTROL S, CONTROL Q



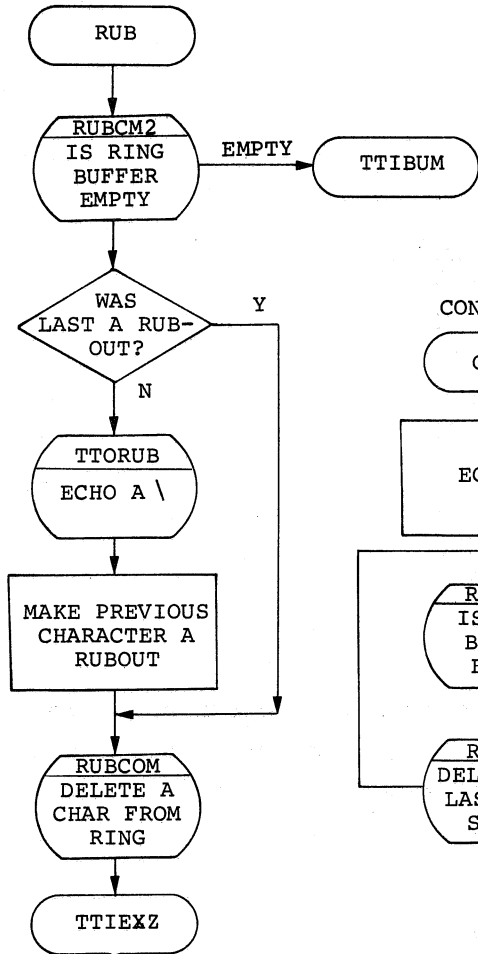
-----DO THIS BY CLEARING THEN SETTING THE OUTPUT INTERRUPT BIT

CONTROL C

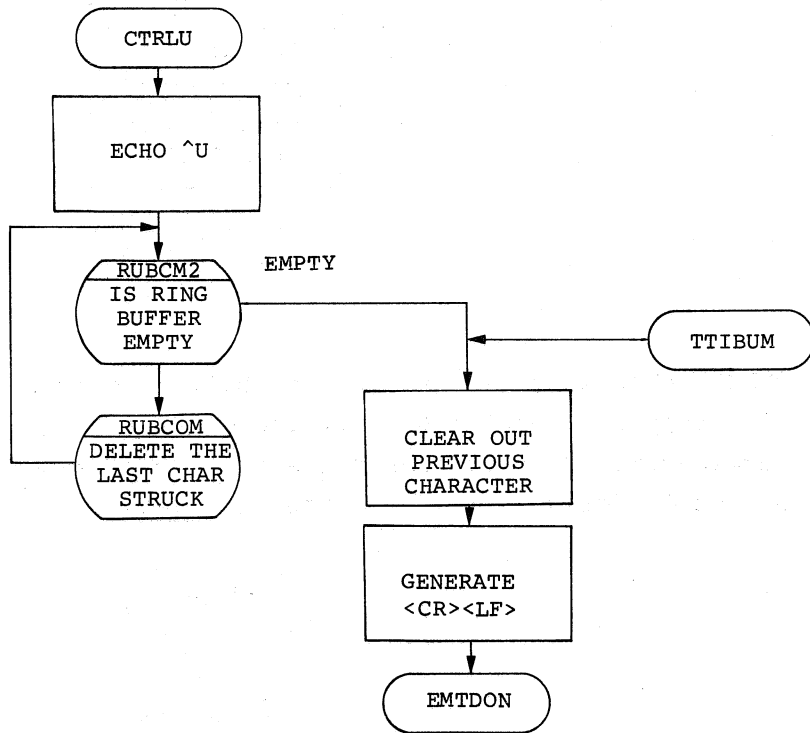


RUBOUT/CONTROL U

RUBOUT

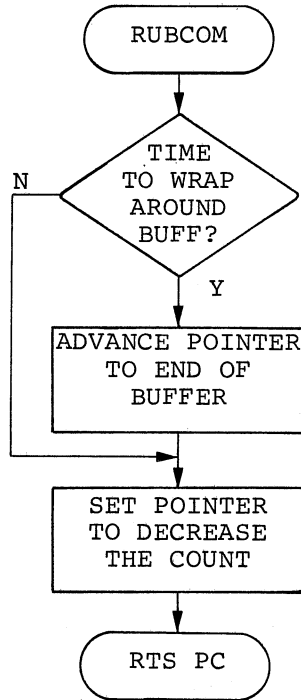


CONTROL U

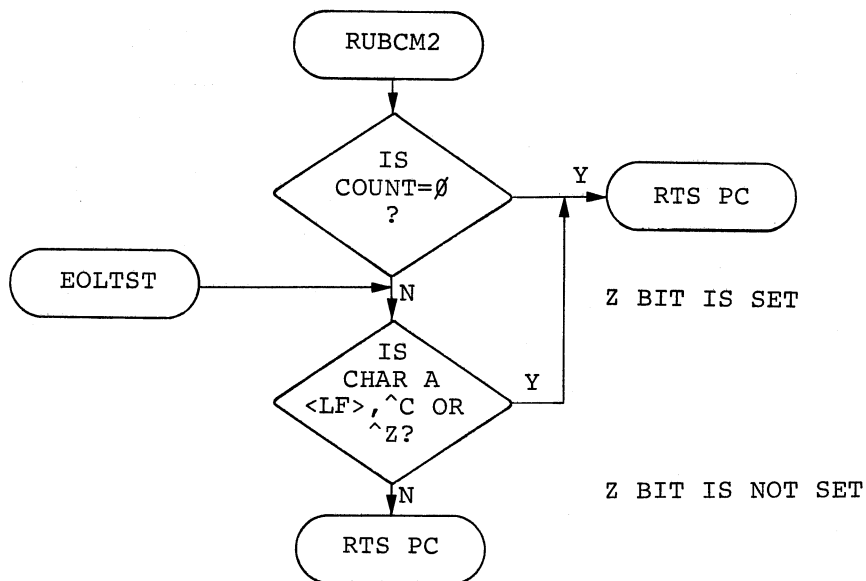


RUBCON/RUBCM2

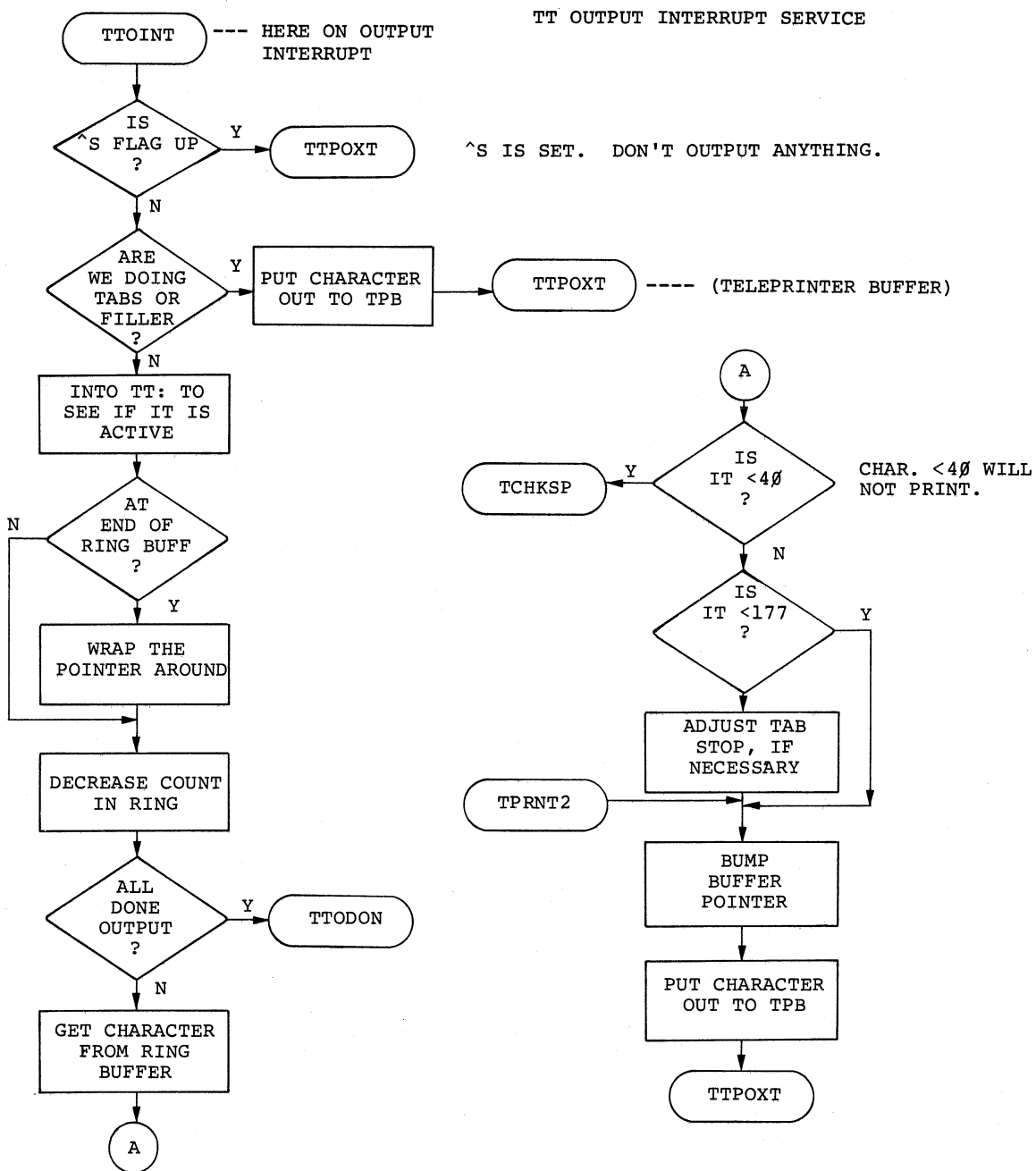
RUBCOM will update the input ring buffer pointers when a character is to be deleted.



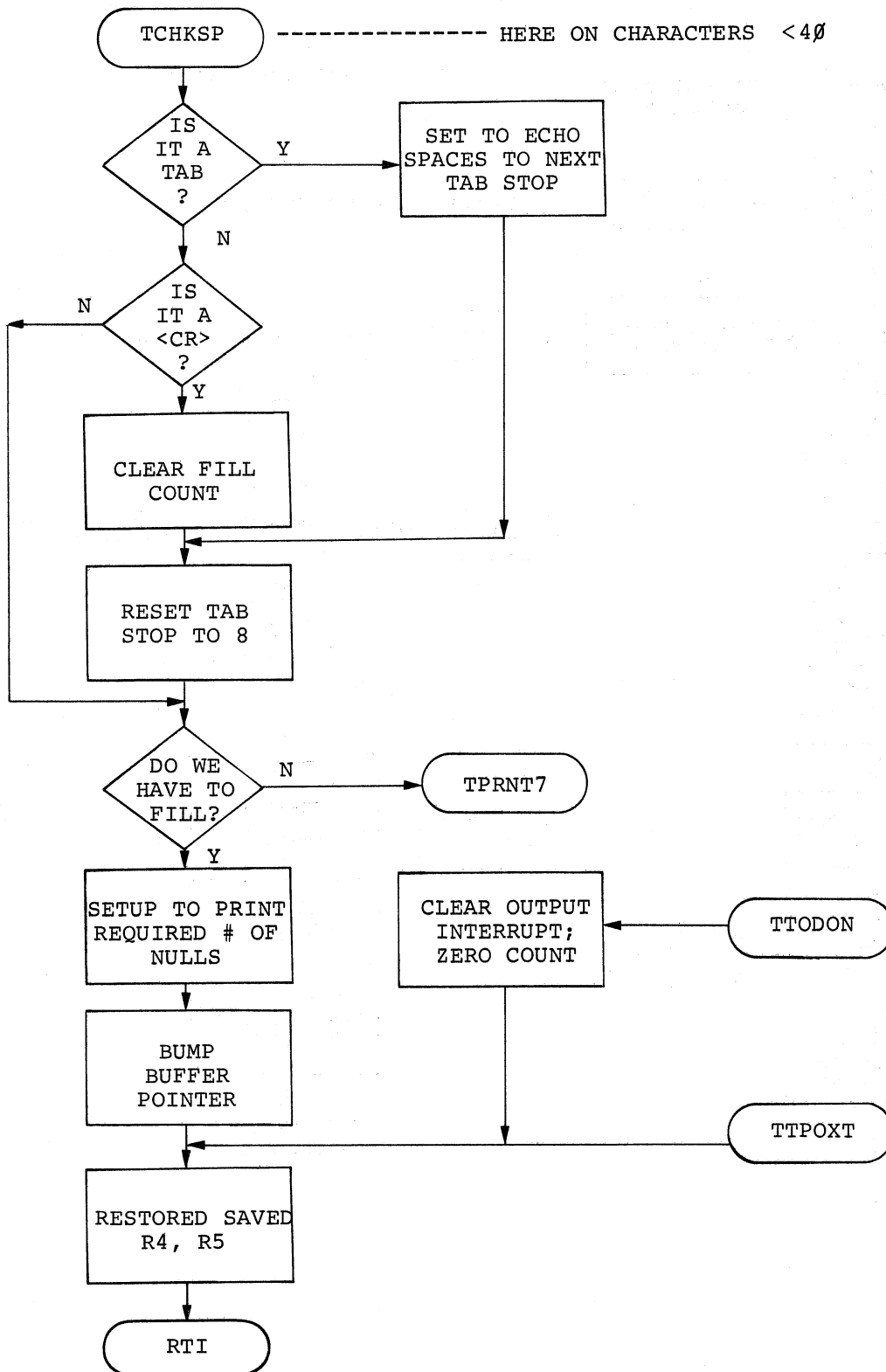
RUBCM2 checks to see if the ring buffer is empty. The buffer is empty if either the count = 0 or if the character to be deleted is a line terminator. This routine falls into routine EOLTST. The zero condition is returned if the buffer is empty.



TT OUTPUT INTERRUPT SERVICE



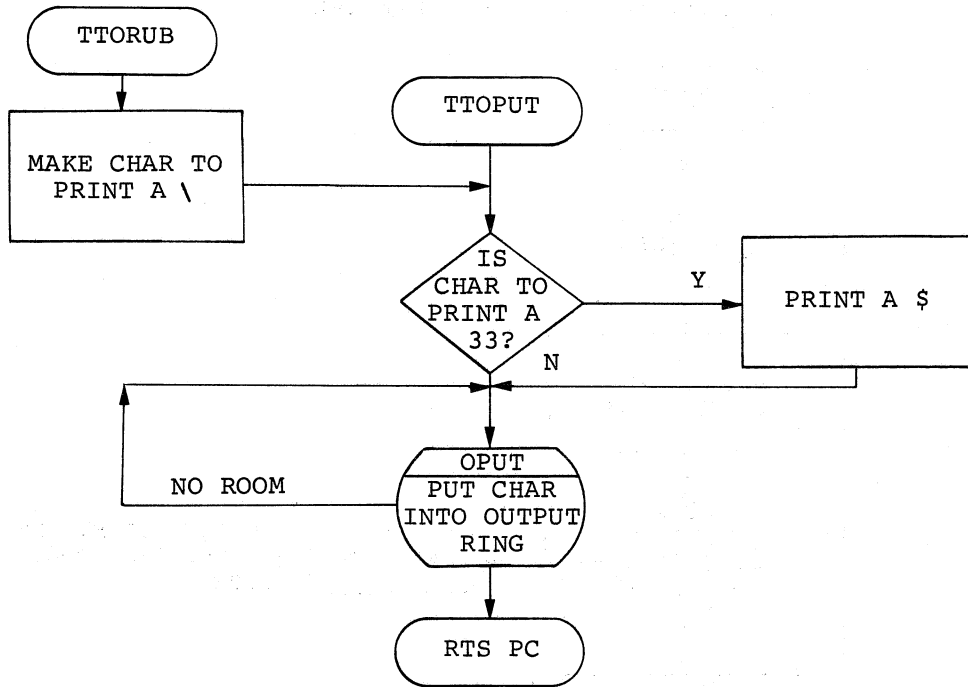
TT OUTPUT INTERRUPT SERVICE (CONT.)





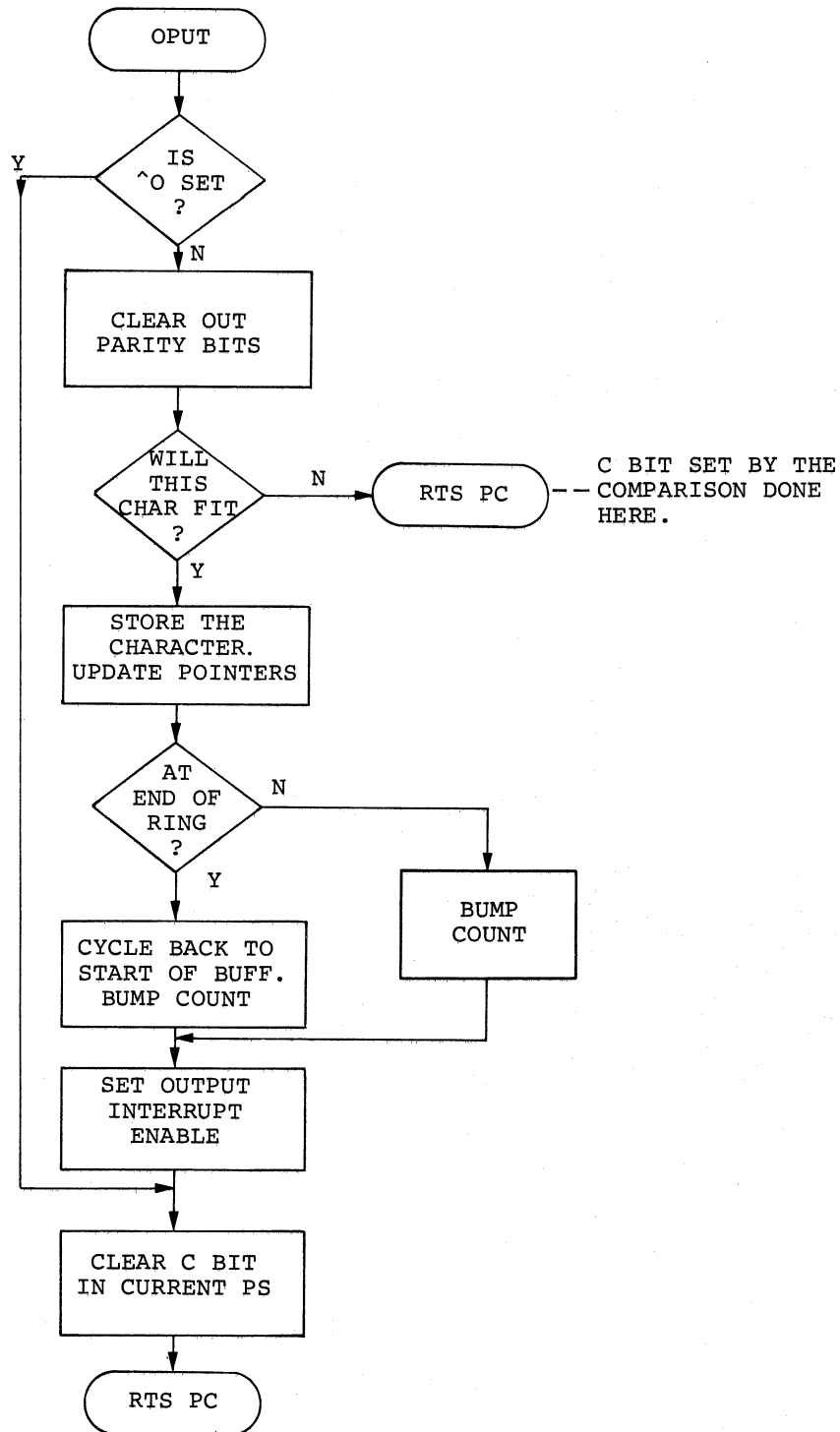
TTORUB/TTOPUT

TTORUB and TTOPUT handle the printing of ALTMODE and RUBOUT. They print a \$ for ALTMODE and \ for RUBOUT.



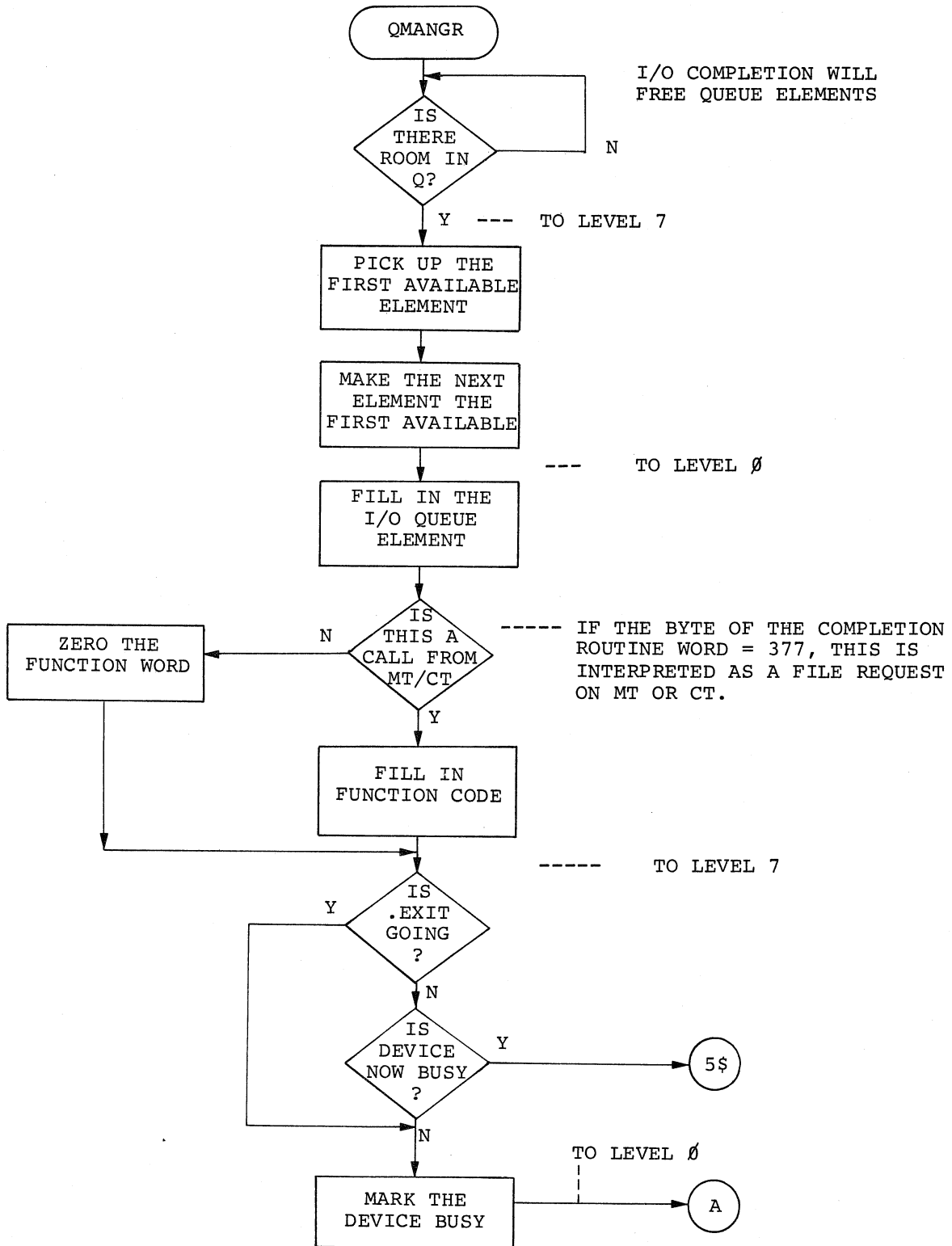
OPUT

OPUT actually puts the output character into the ring buffer. It updates the ring pointers and sets the interrupt enable bit. If the buffer is full, it returns with the C bit set.

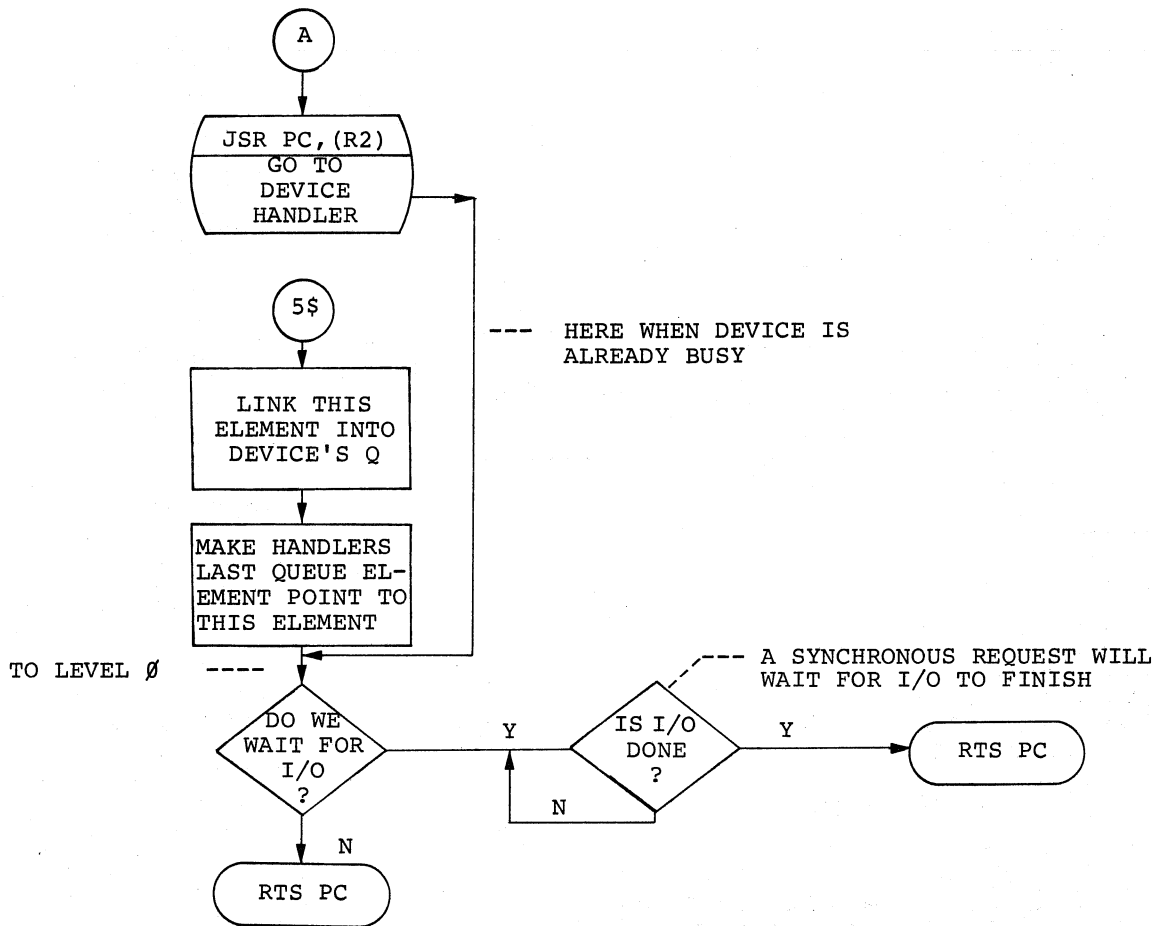


E.4.4 I/O Routines

I/O QUEUE MANAGEMENT ROUTINES

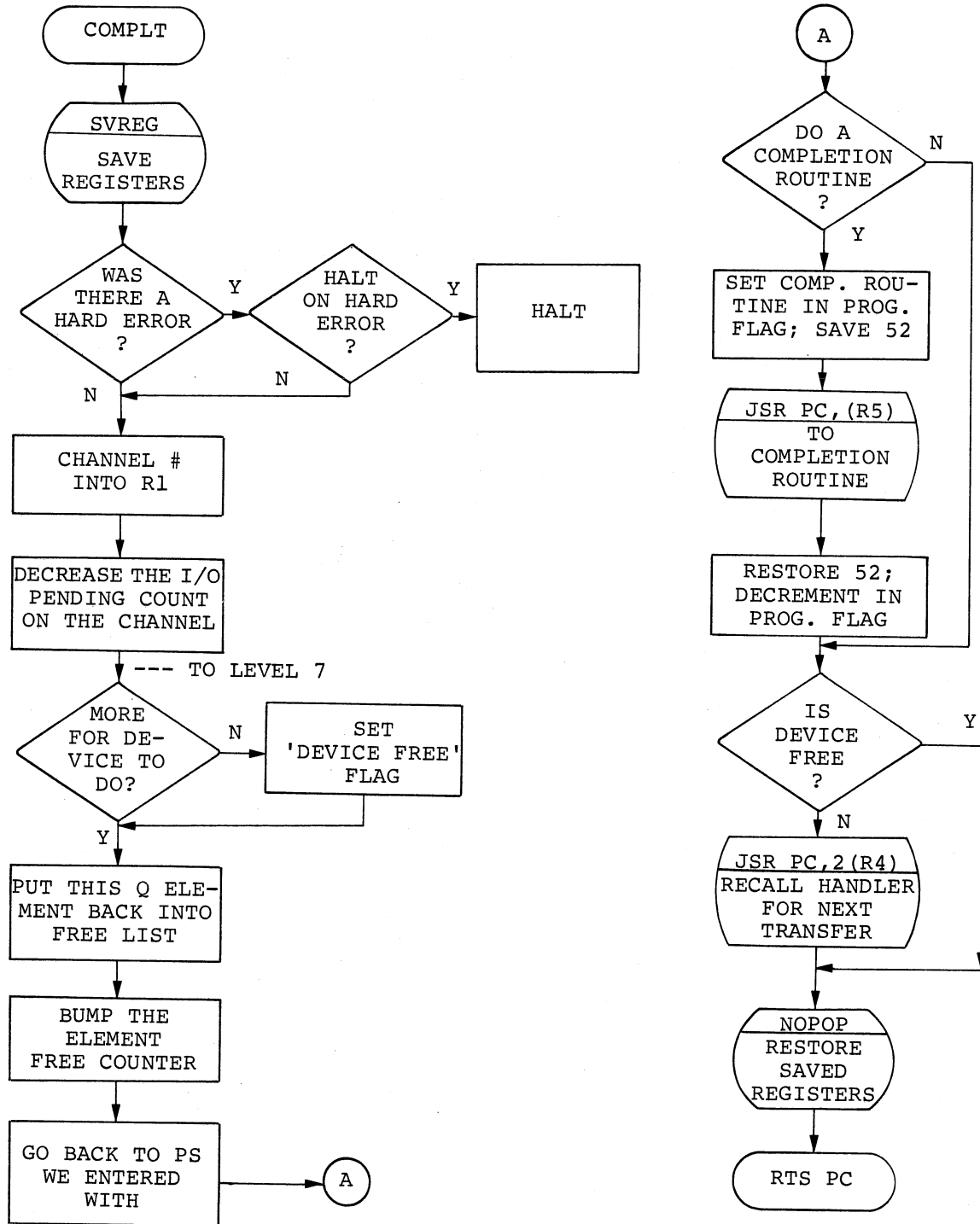


I/O QUEUE MANAGEMENT ROUTINES (CONT.)



I/O QUEUE COMPLETION

COMPLT is entered when an I/O transfer finishes.



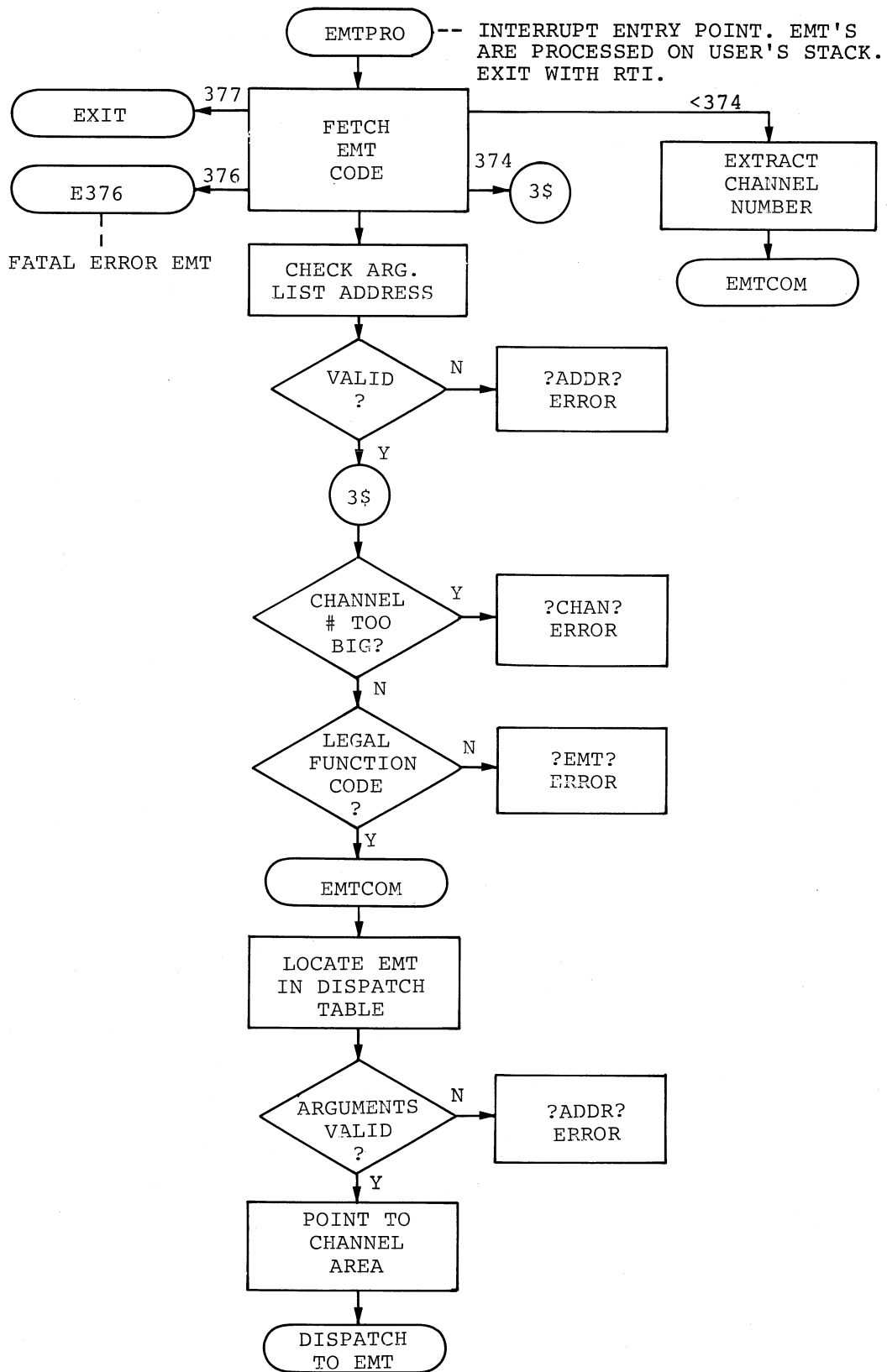
E.5 RMON (RESIDENT MONITOR) FLOWCHARTS FOR  
FOREGROUND/BACKGROUND MONITOR



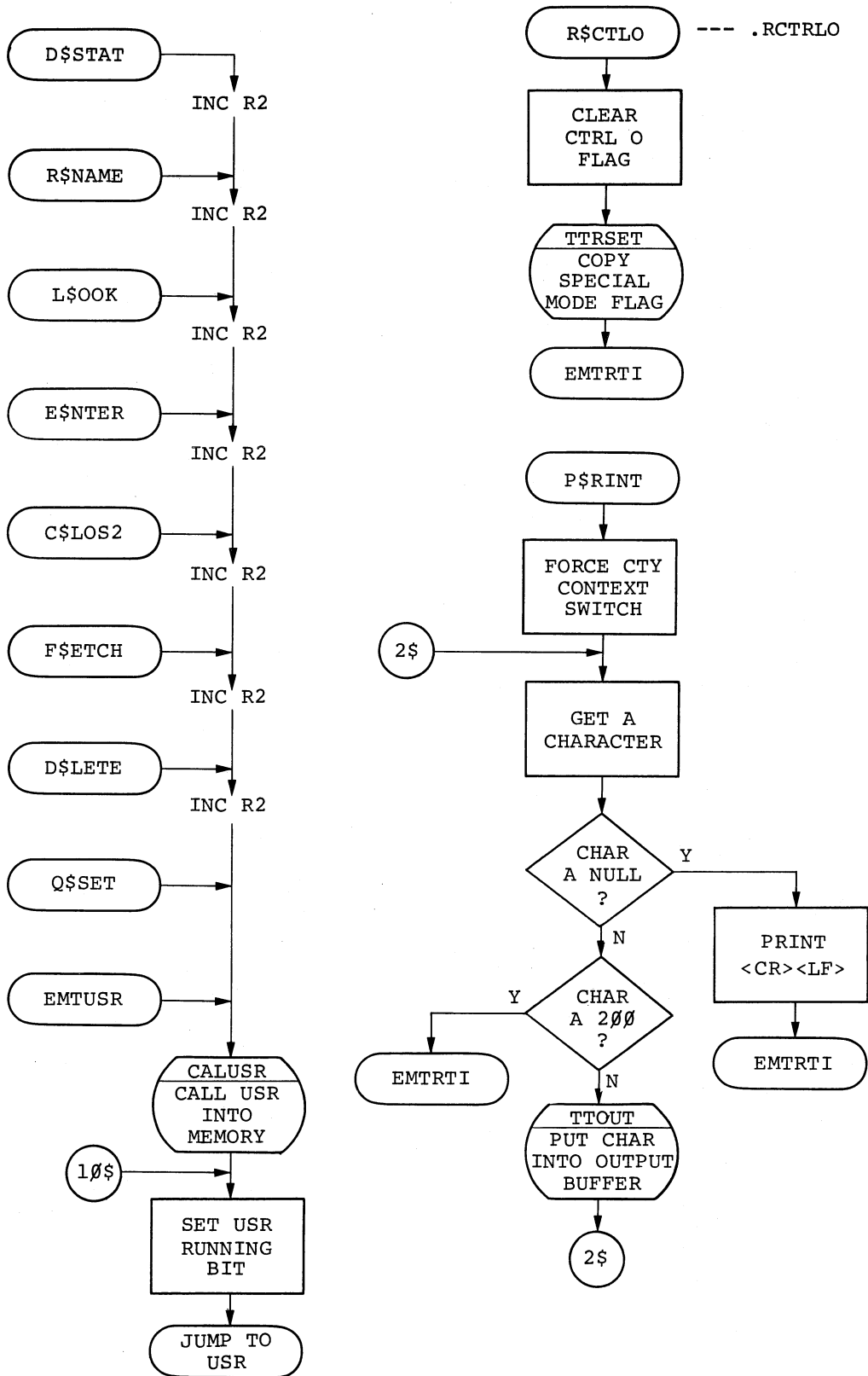


E.5.1 EMT Processors

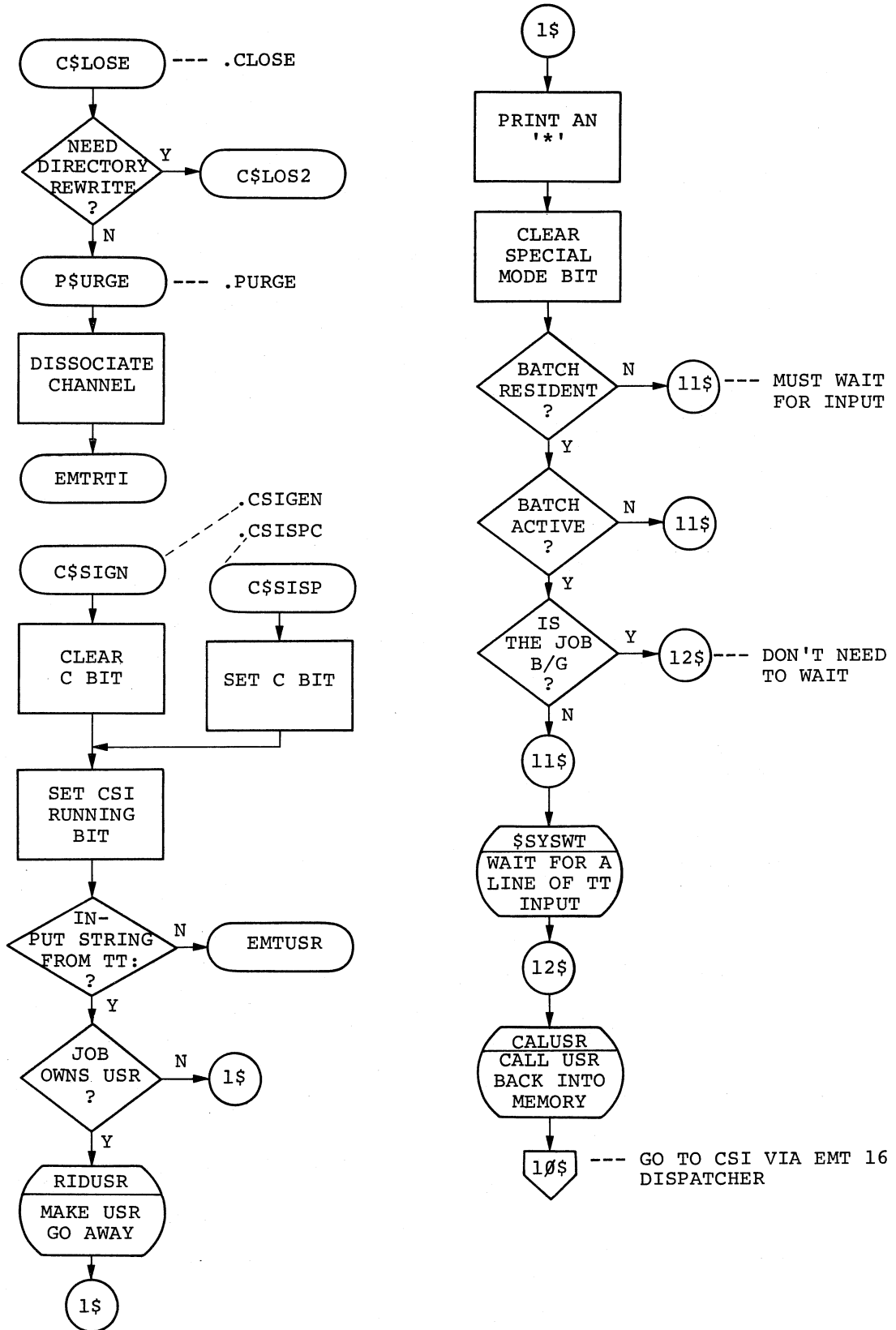
EMT DISPATCHER



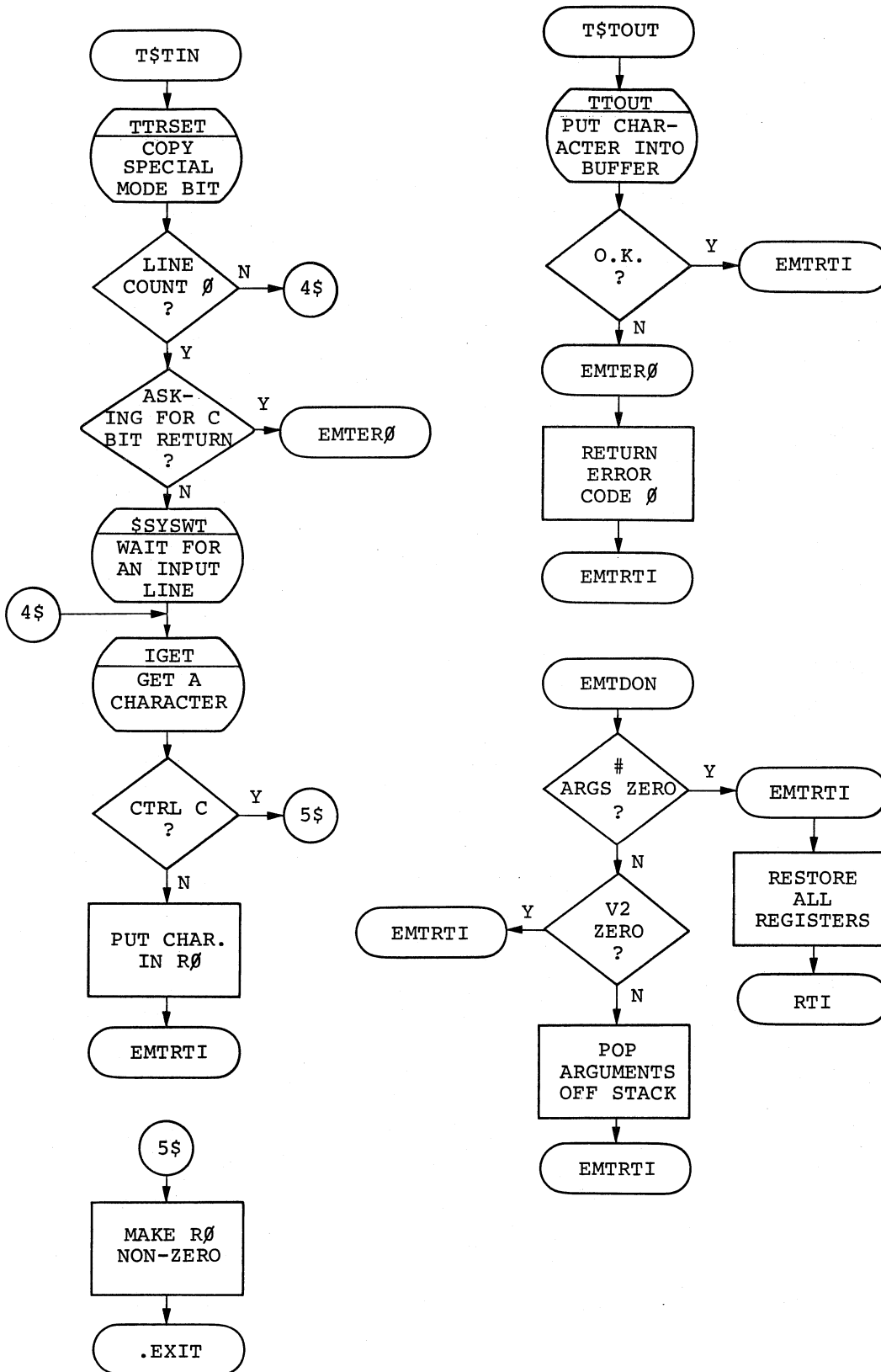
EMT 16 DISPATCH, .RCTRL0, .PRINT



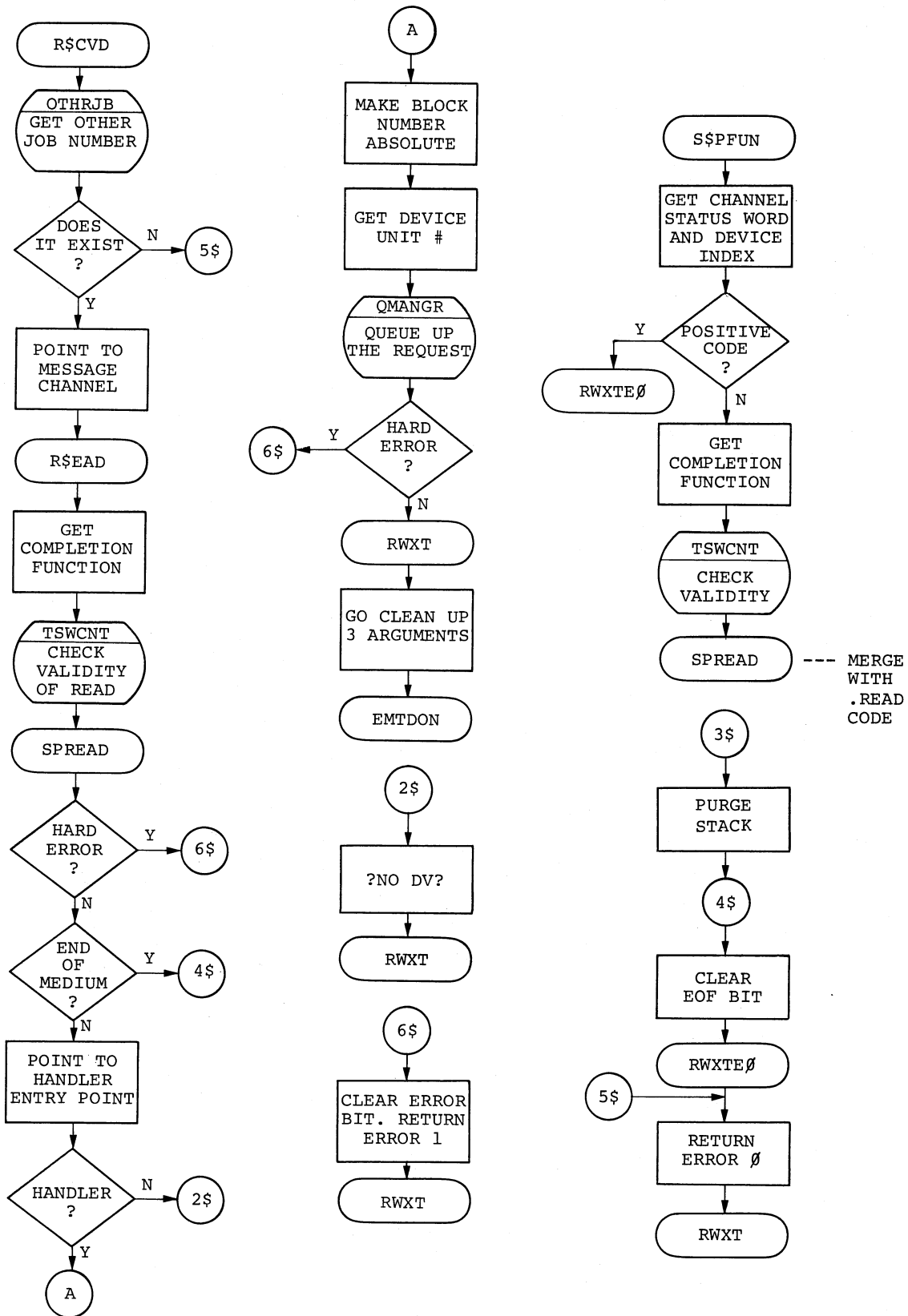
.CLOSE, .PURGE, .CSISPC, .CSIGEN

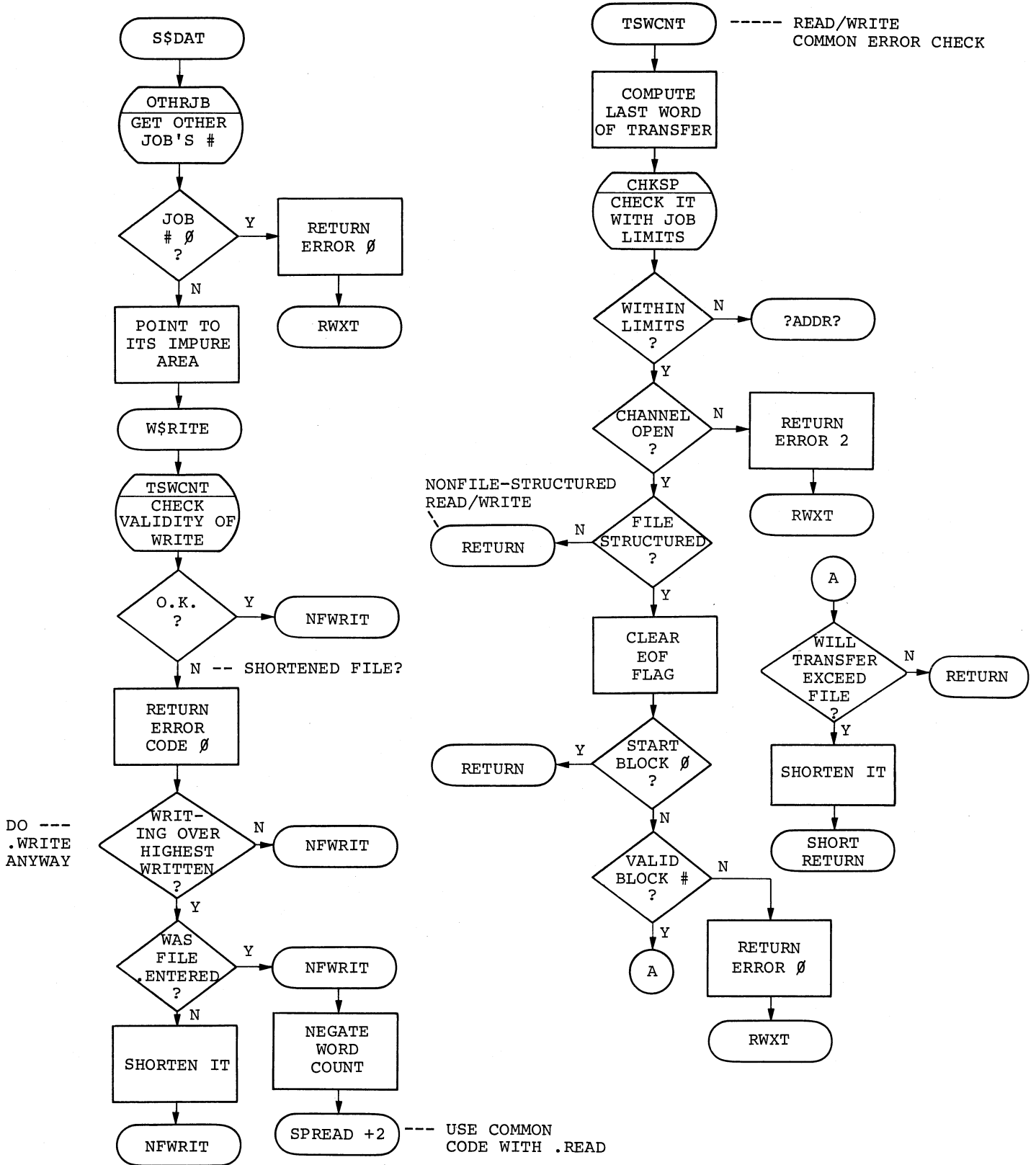


.TTYIN, .TTYOUT, EMT RETURN

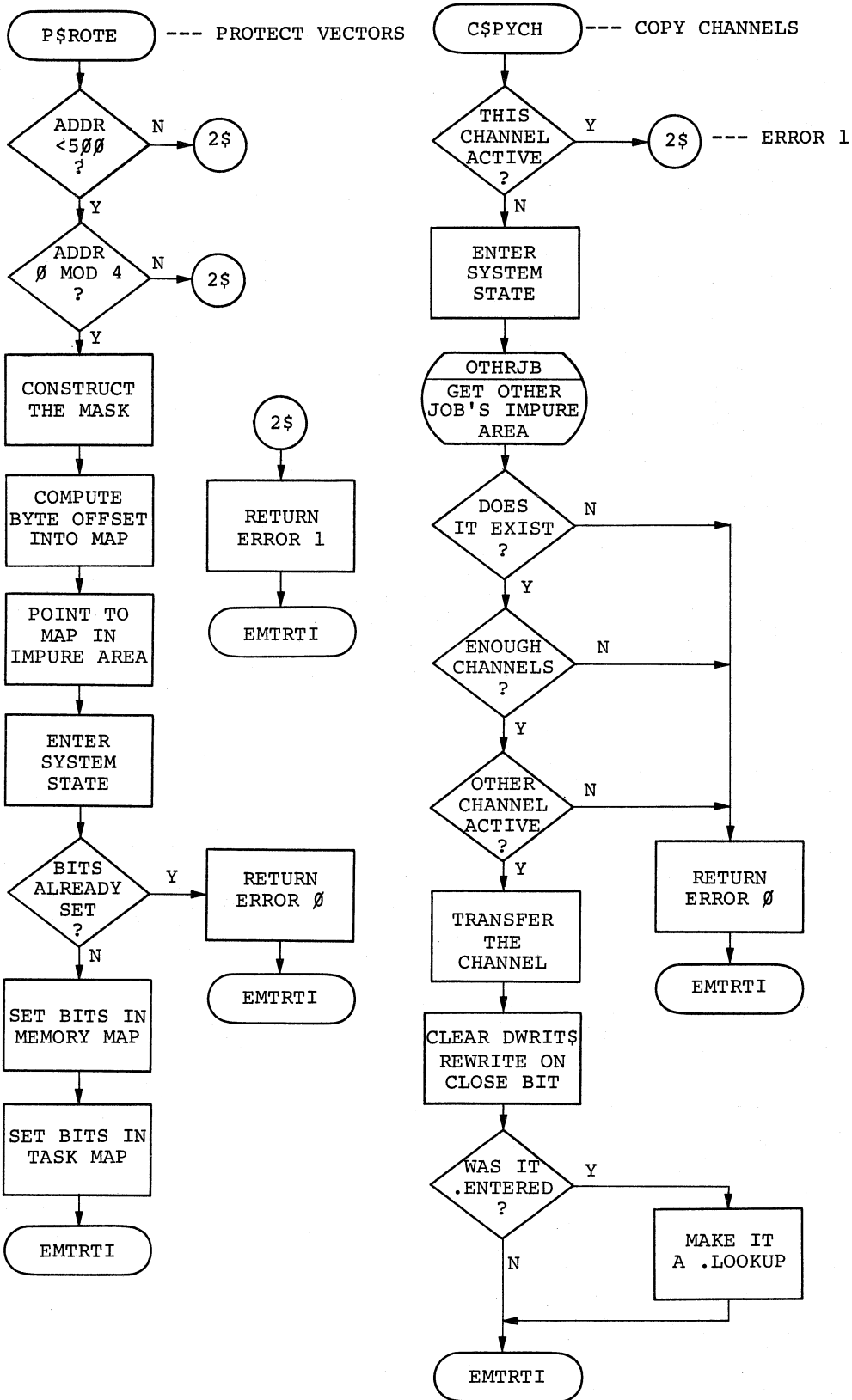


.READ, .RCVD, .SPFUN



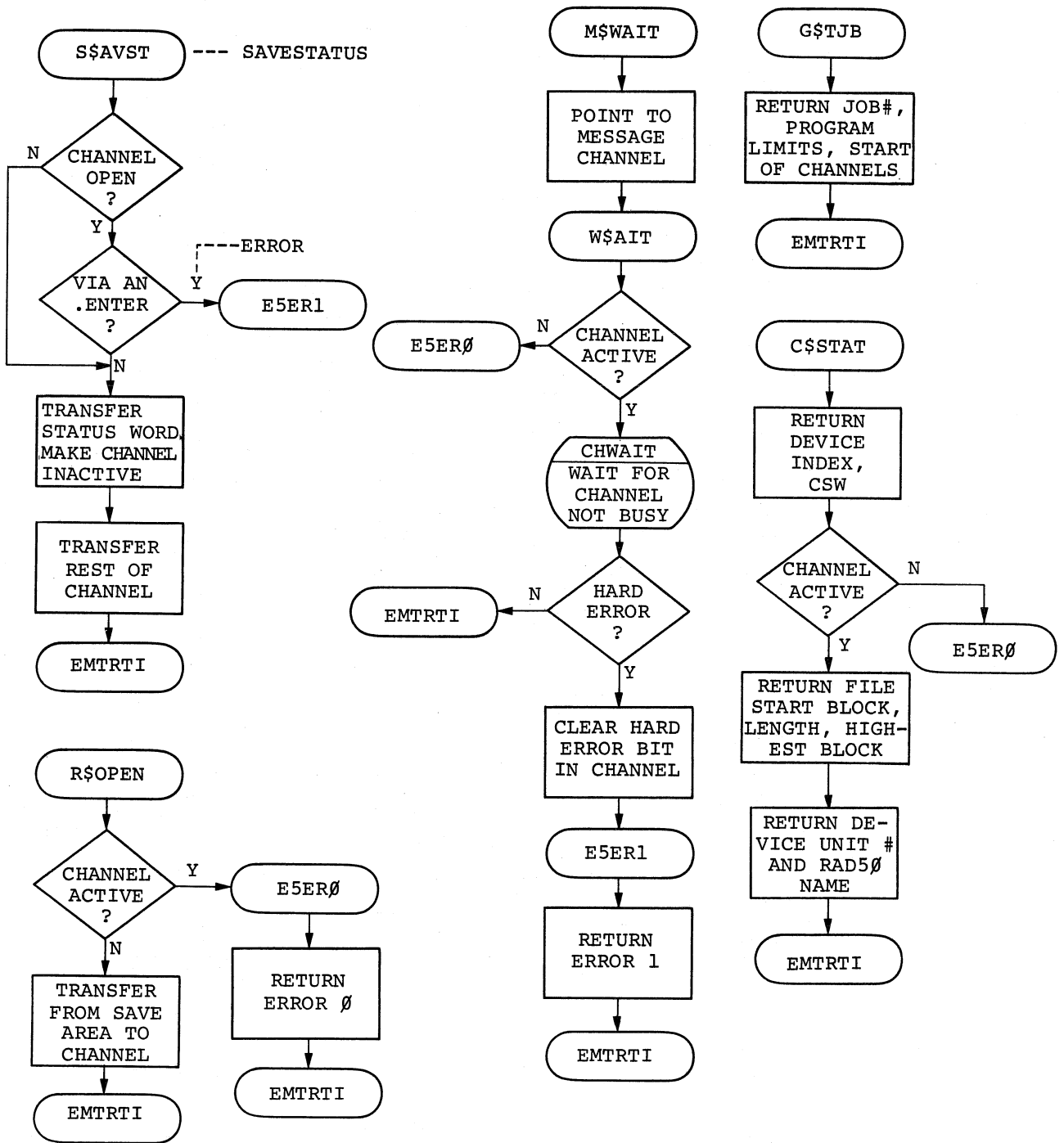


.PROTECT, .CHCOPY

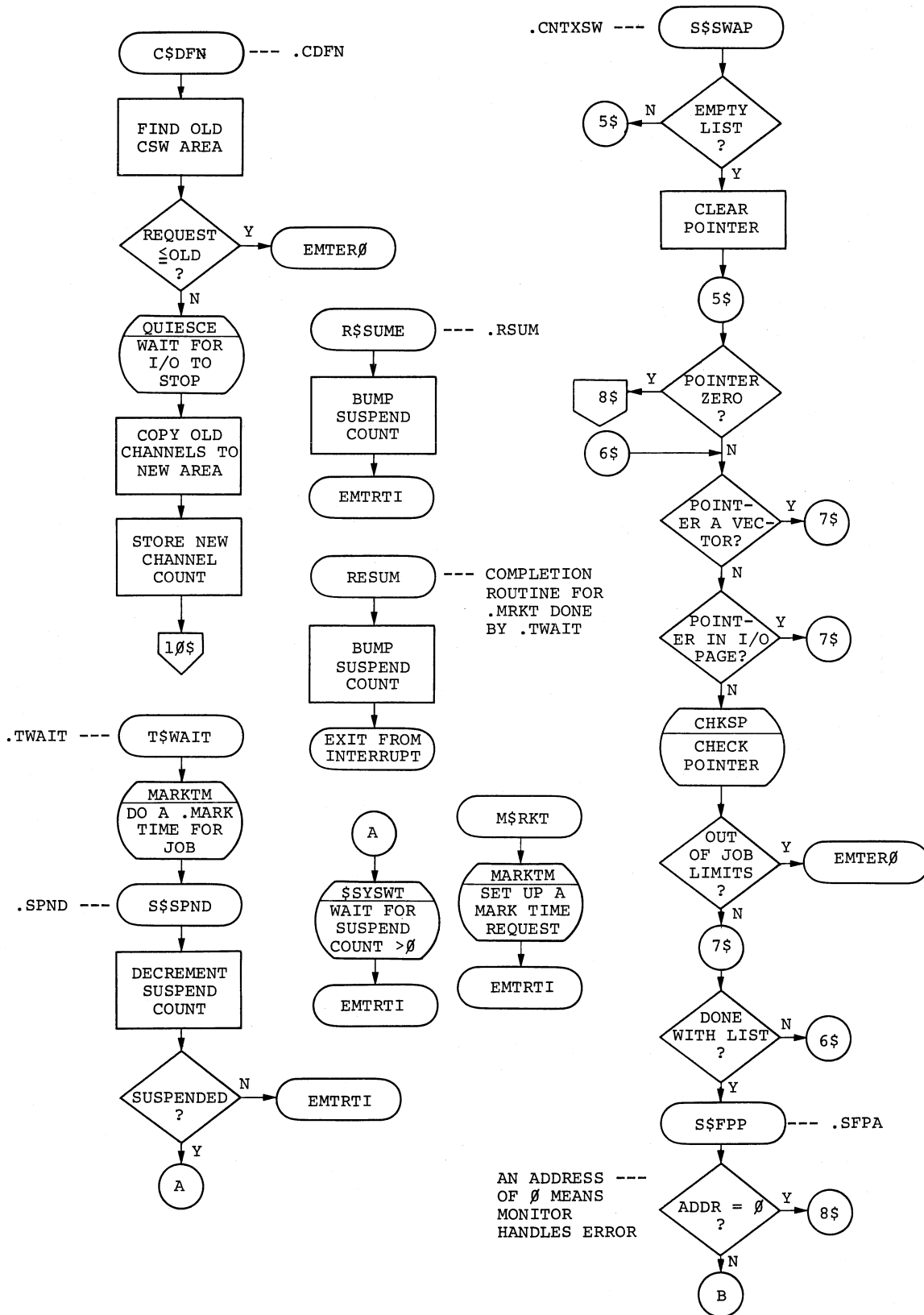




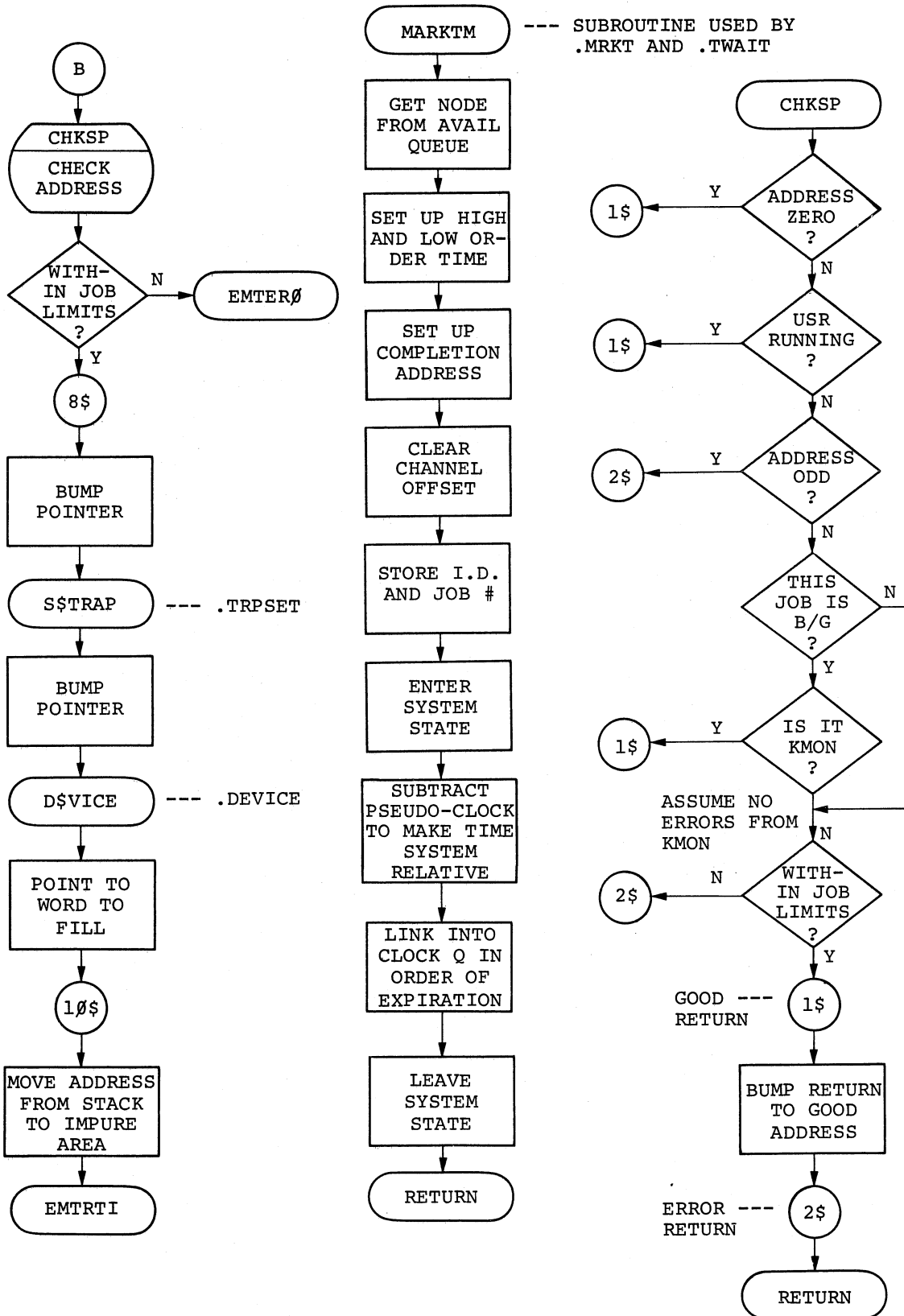
.SAVESTATUS, .REOPEN,  
 .M\$WAIT, .W\$AIT,  
 .GTJB, .C\$STAT



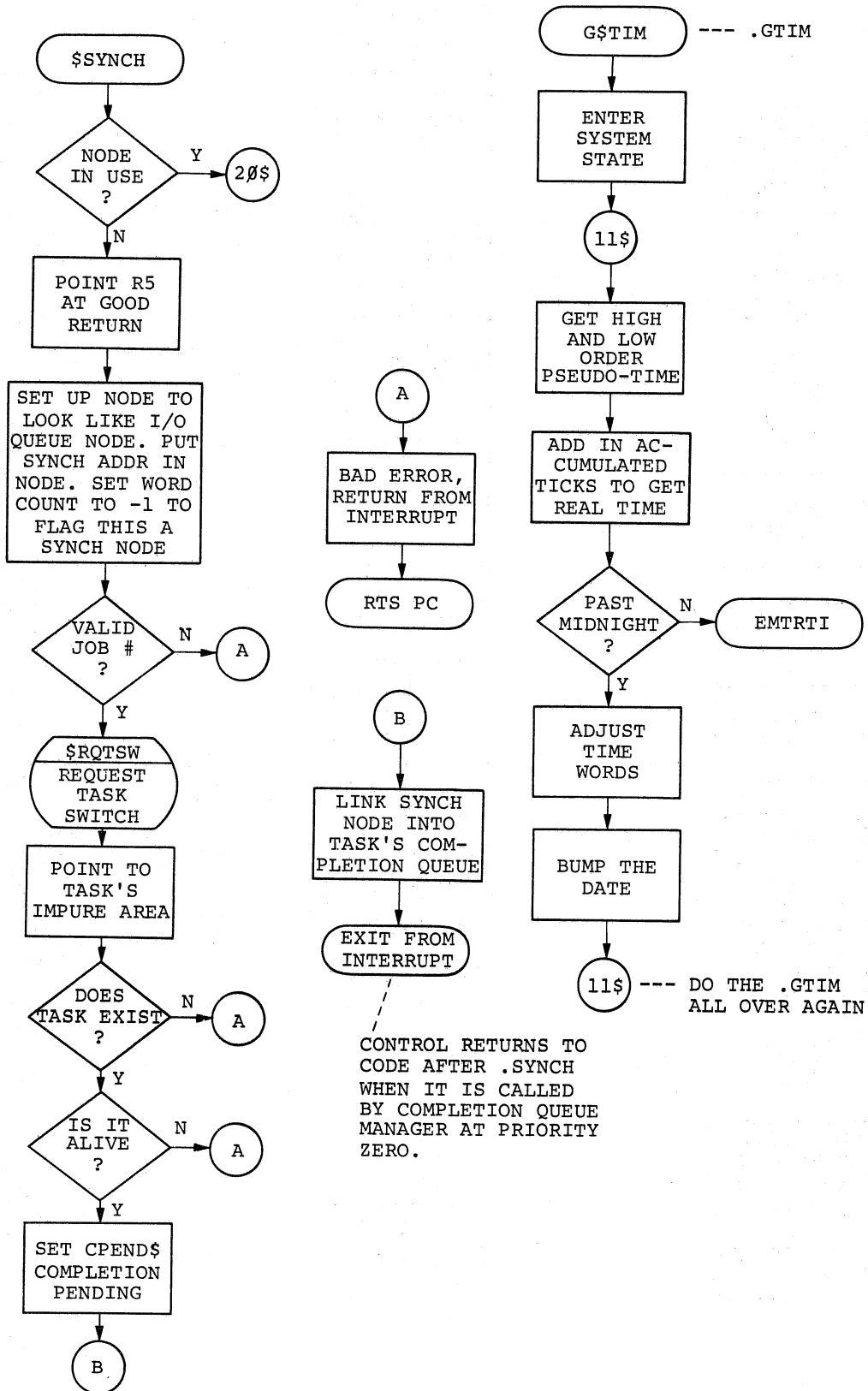
.CDFN, .TWAIT, .SPND, .RSUM,  
 .CNTXSW, .SFPA, .TRPSET, .DEVICE



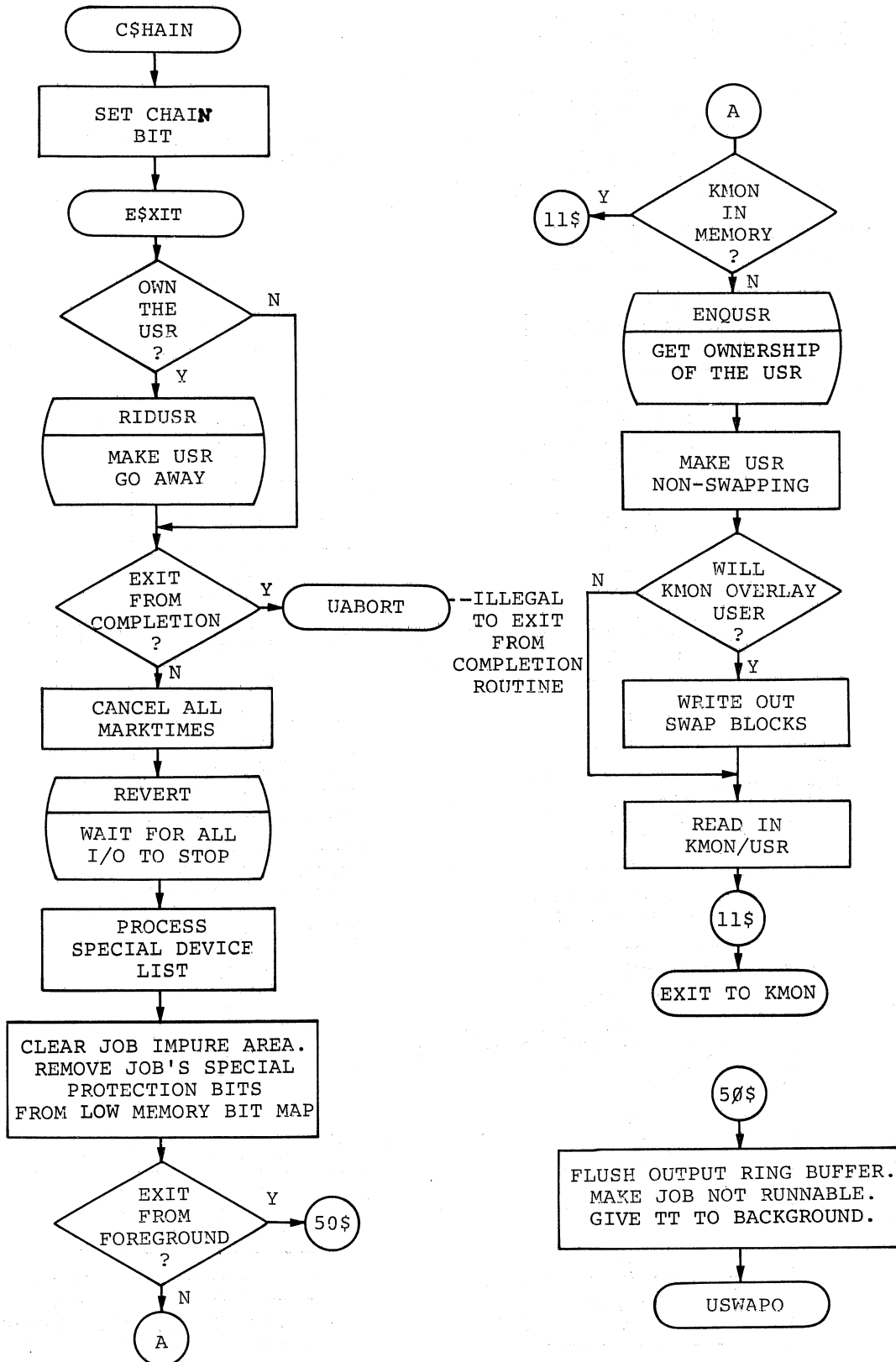
.TRPSET, .DEVICE (CONT.)



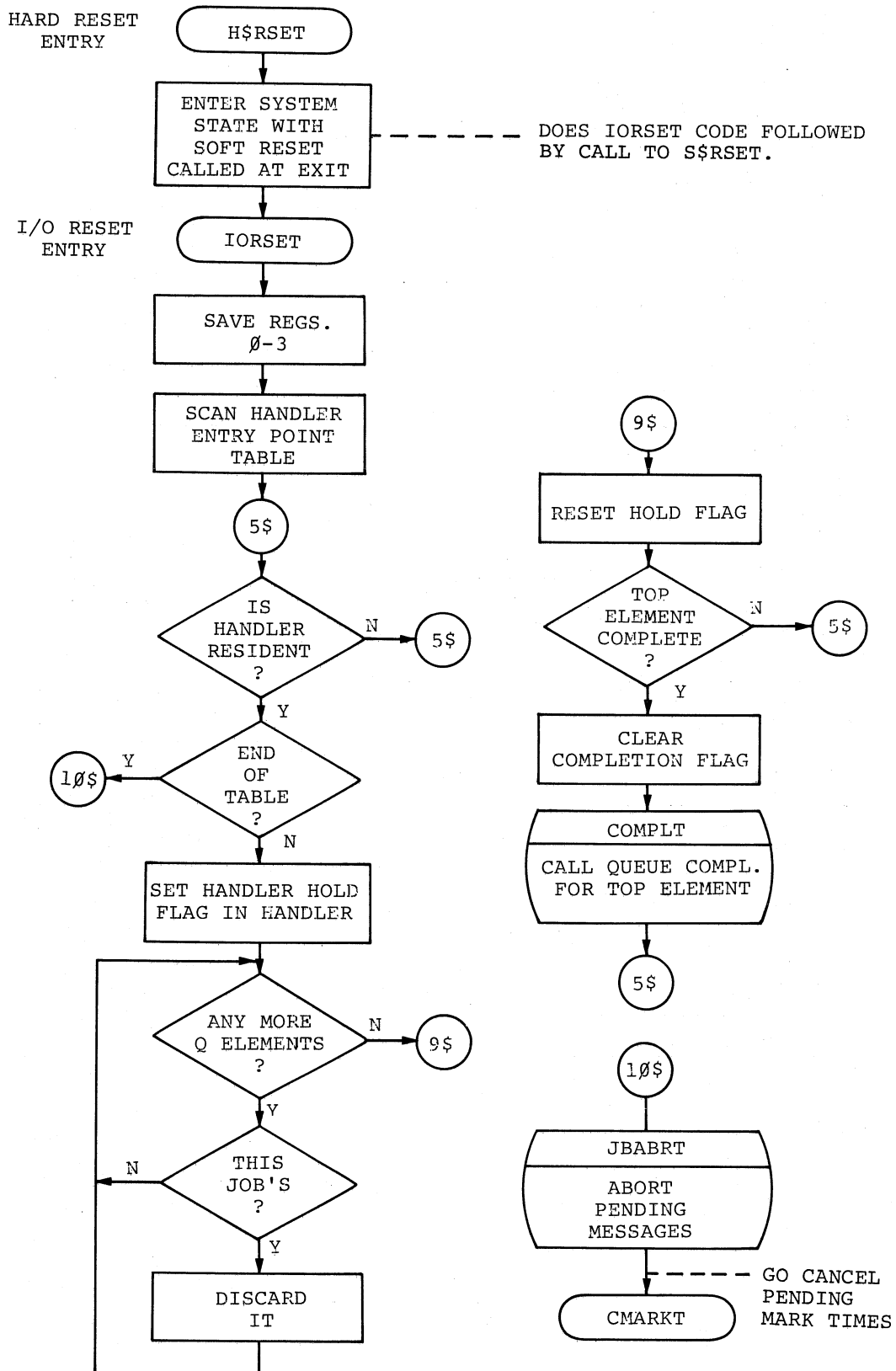
.SYNCH, .GTIM



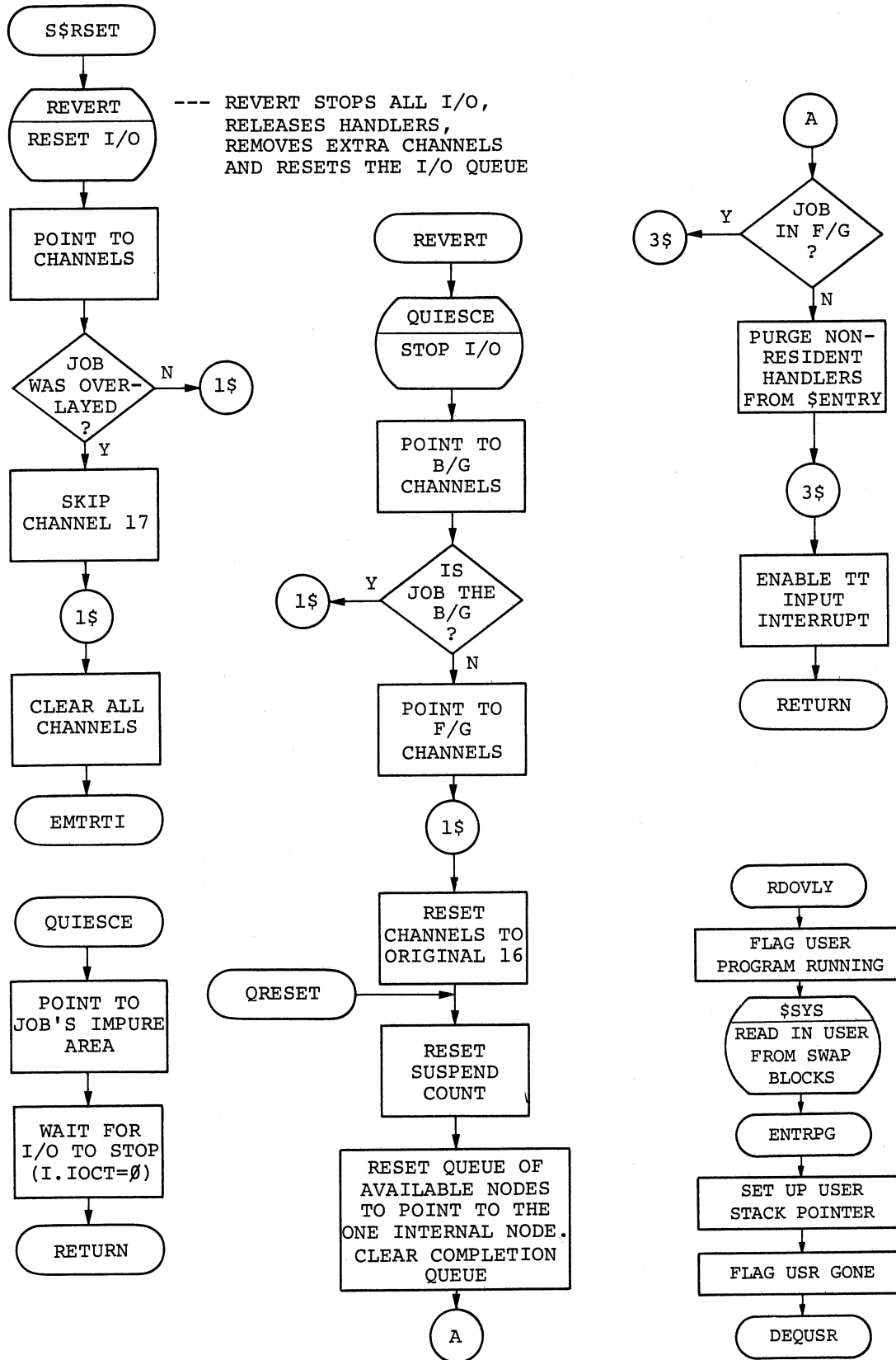
.EXIT  
.CHAIN



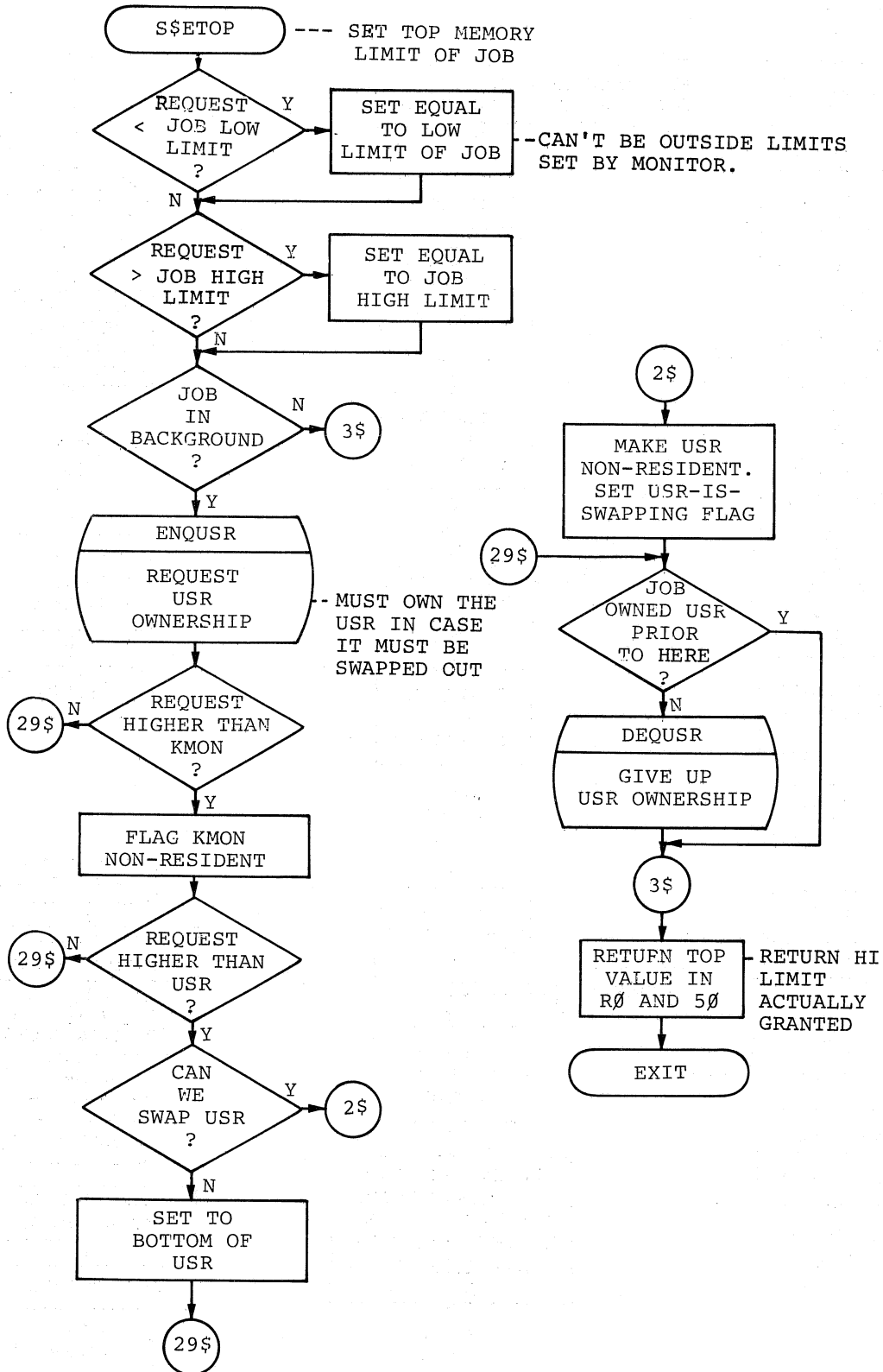
HARD AND SOFT RESET



HARD AND SOFT RESET (CONT.)/RDOVLY

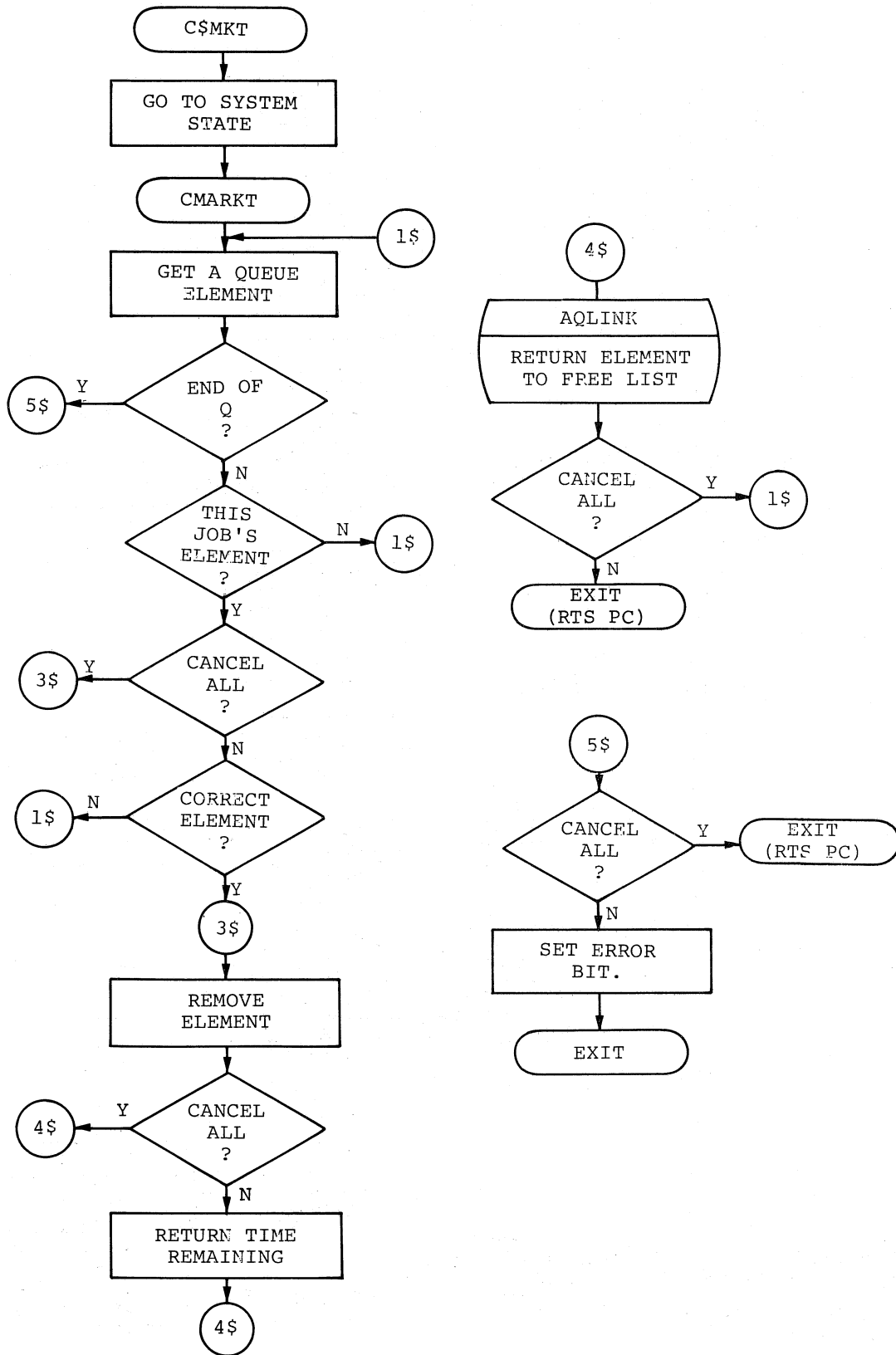


.SETTOP

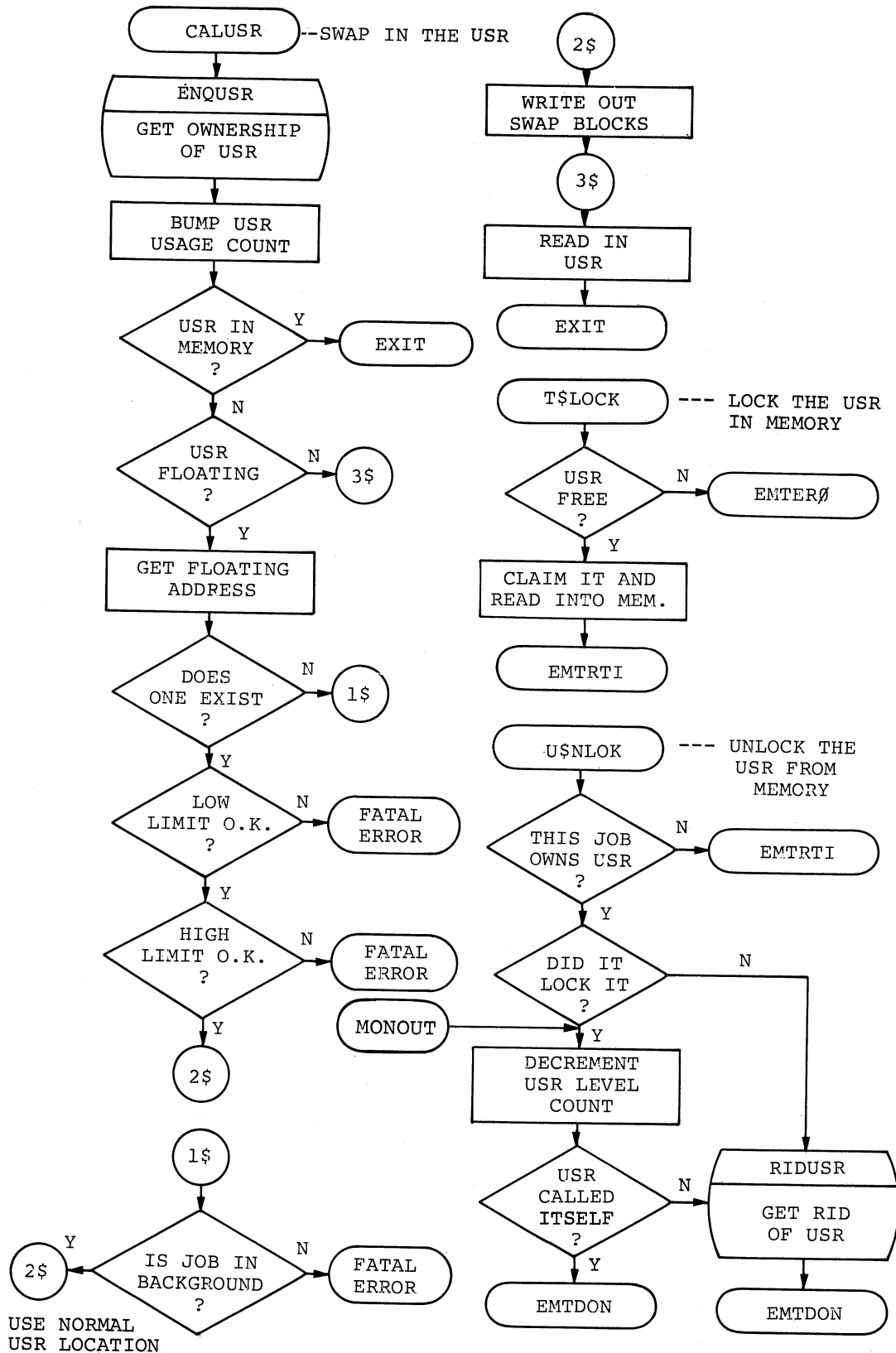




CANCEL MARK TIME



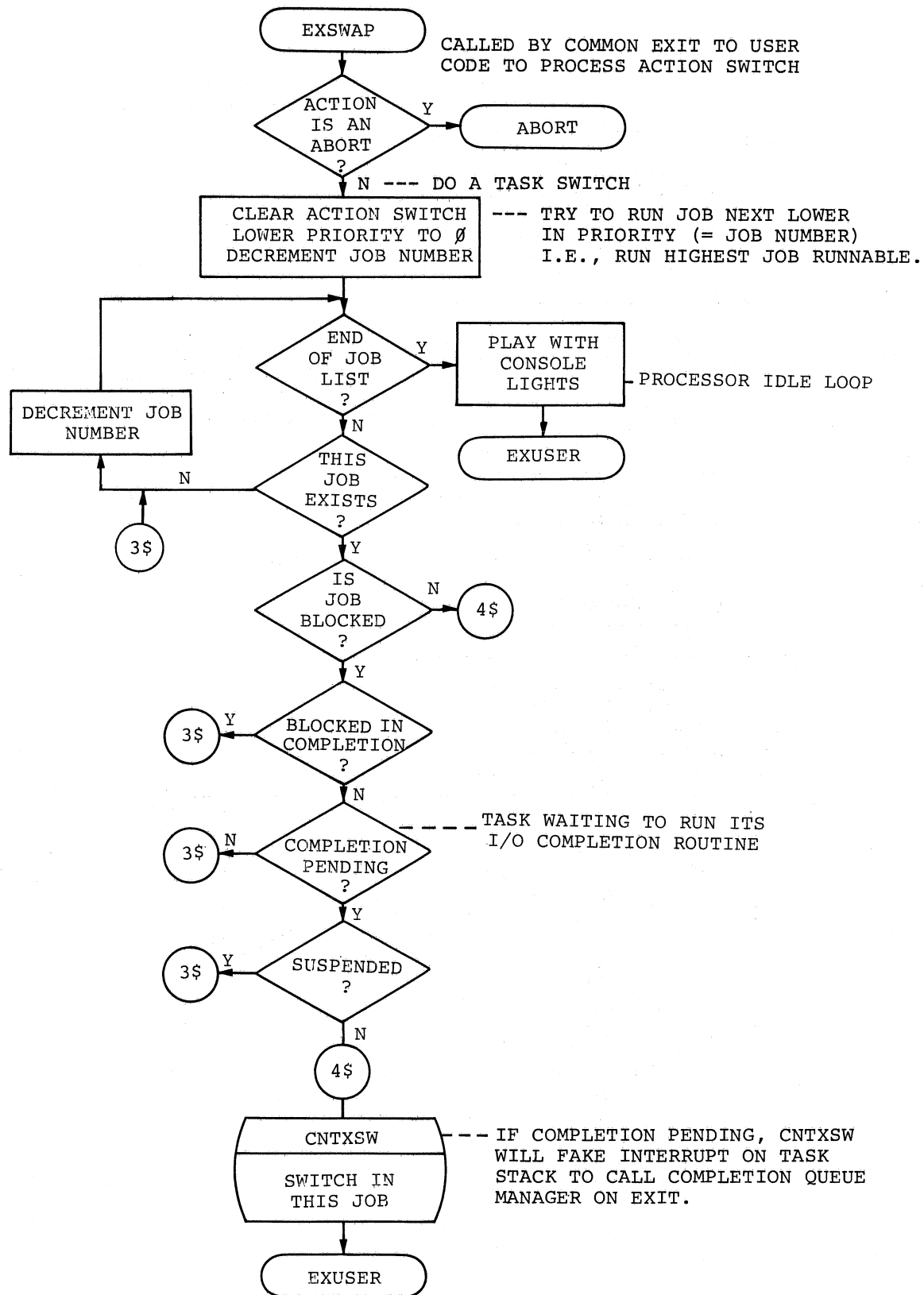
SWAP IN USR, LOCK/UNLOCK USR



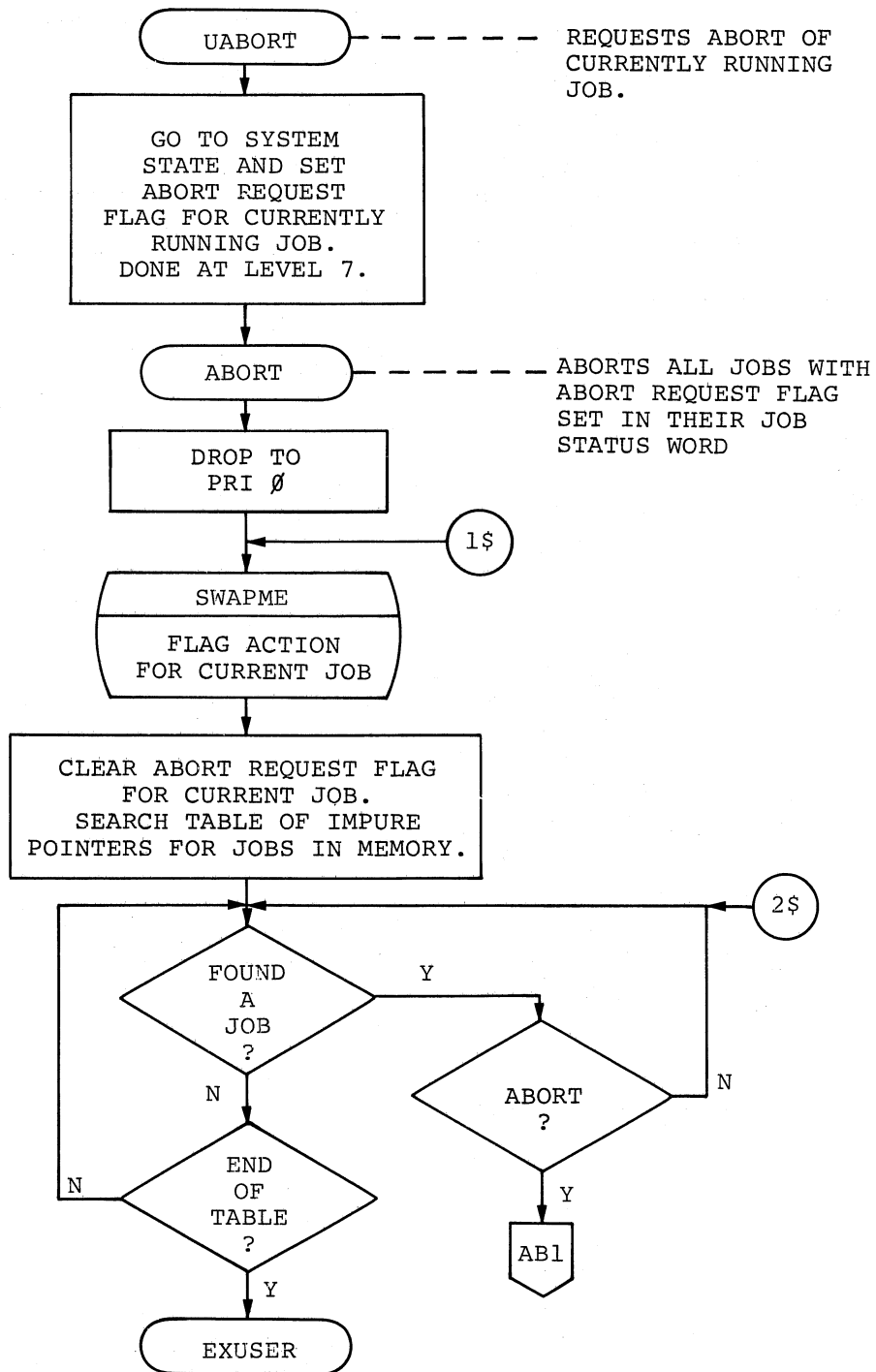
E.5.2 Job Arbitration, Error Processing



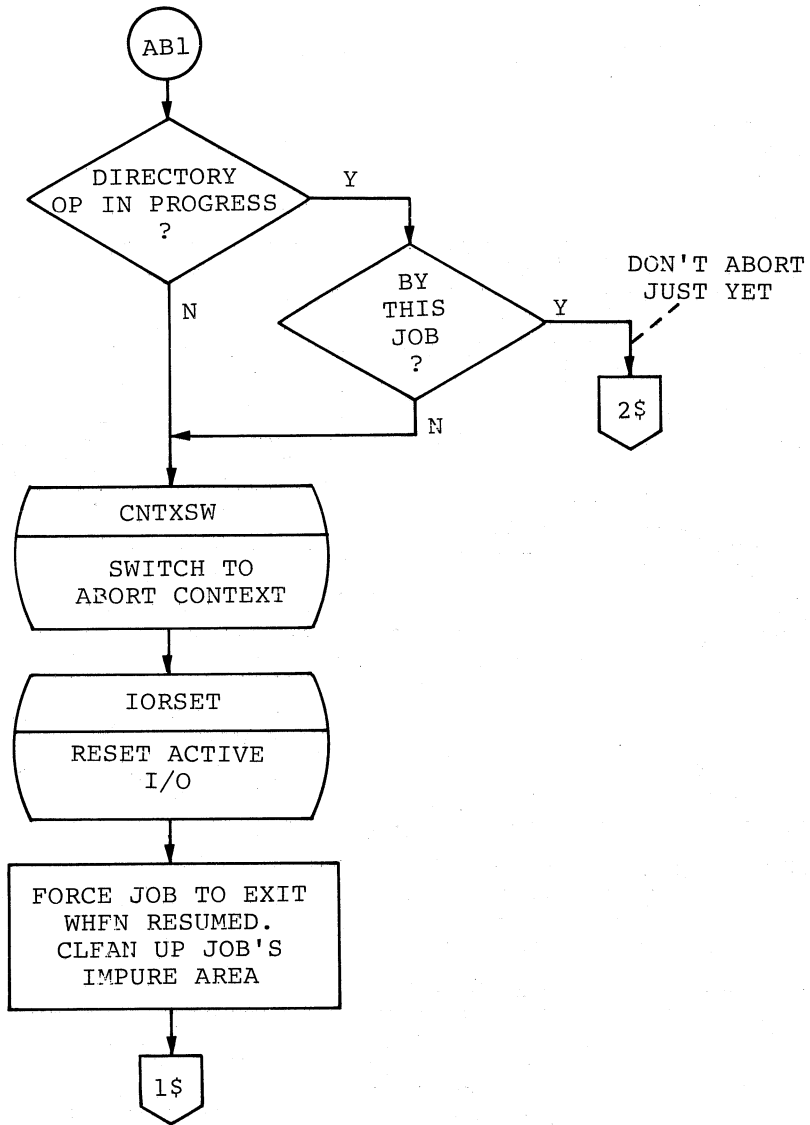
SCHEDULER



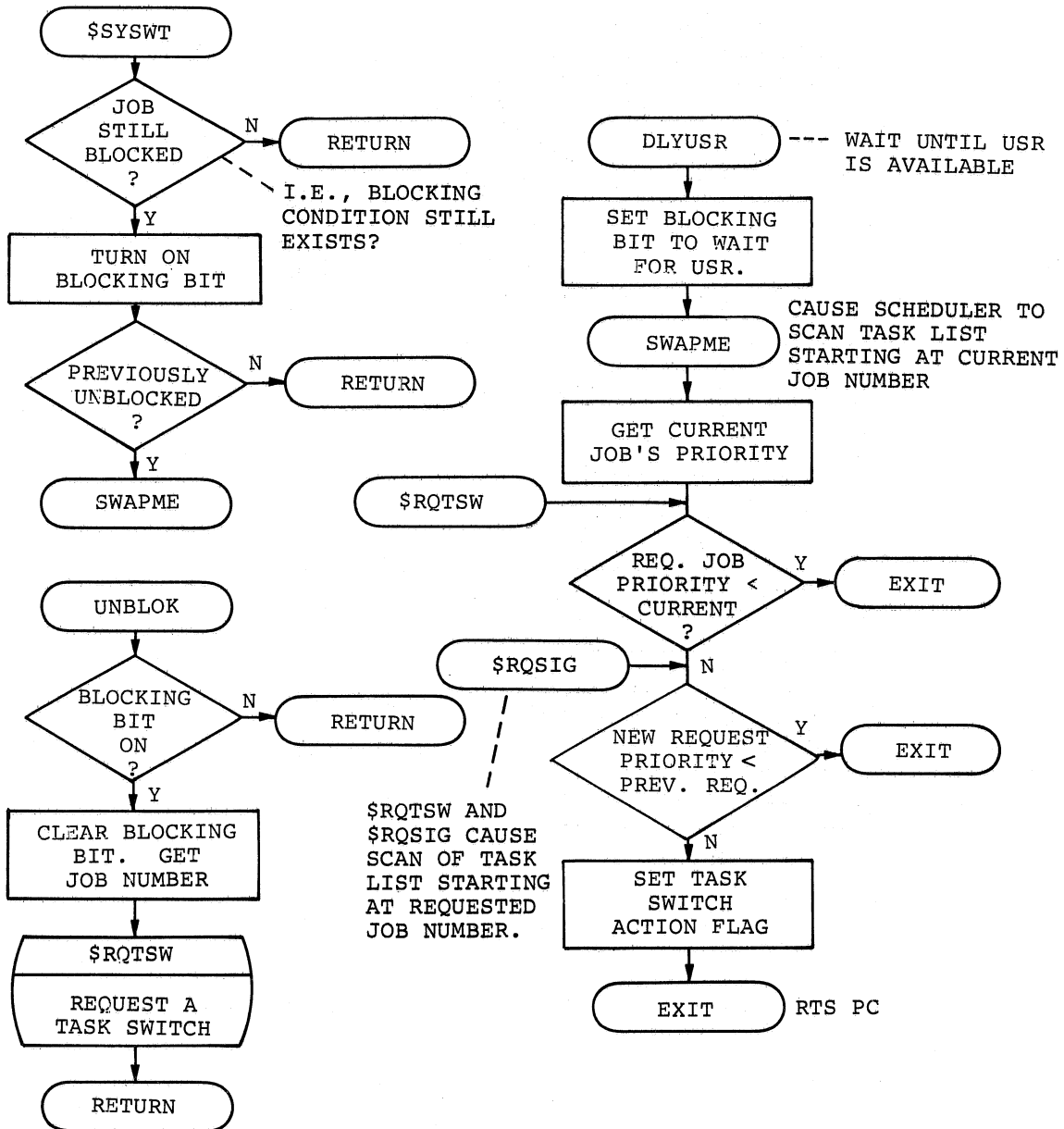
JOB ABORT



JOB ABORT (CONT.)

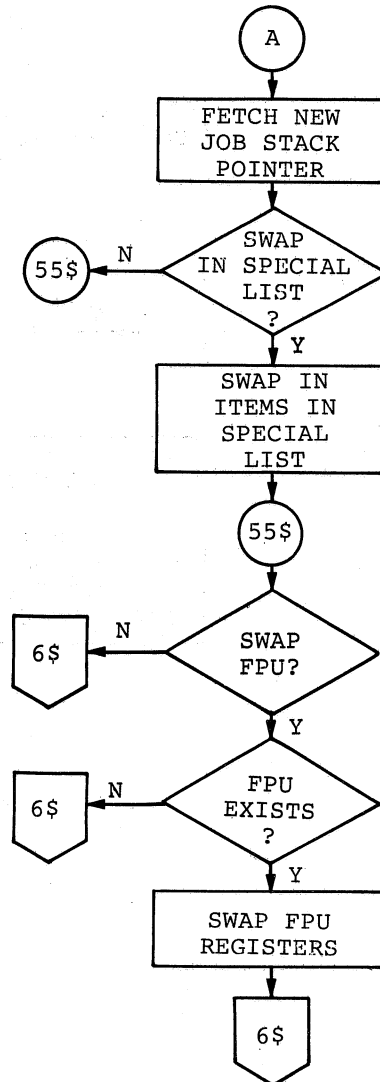
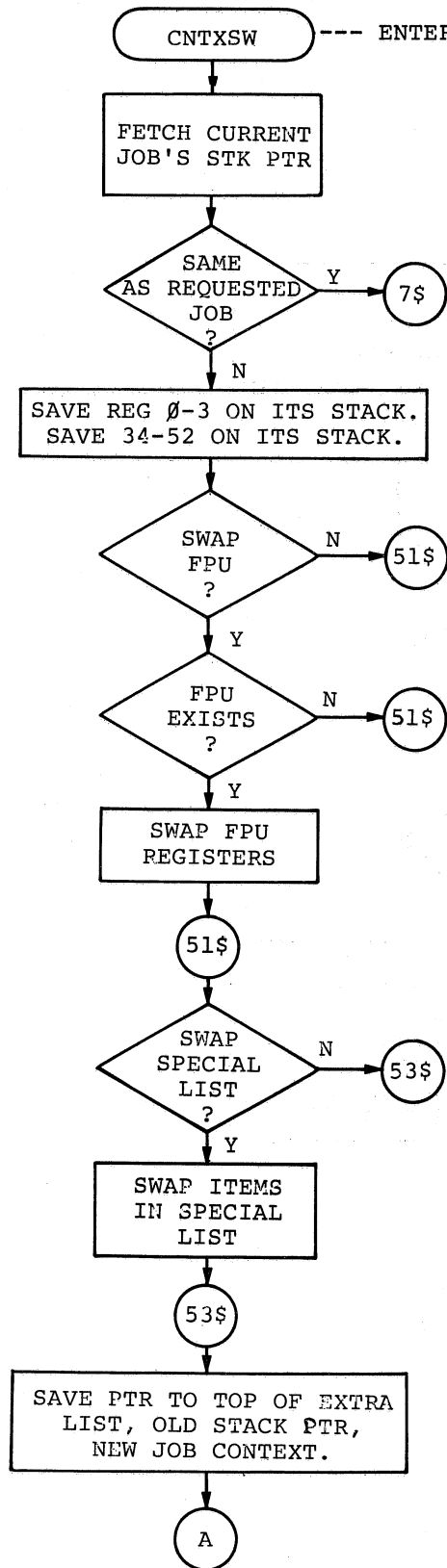


BLOCK A TASK/UNBLOCK A TASK  
REQUEST TASK SWITCH

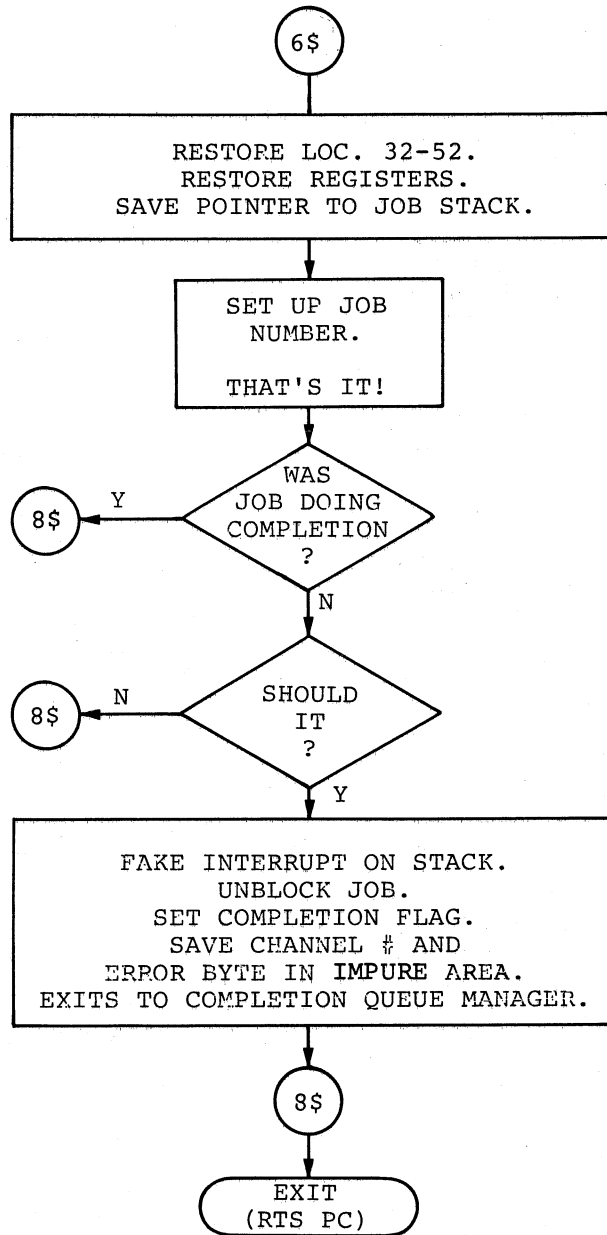




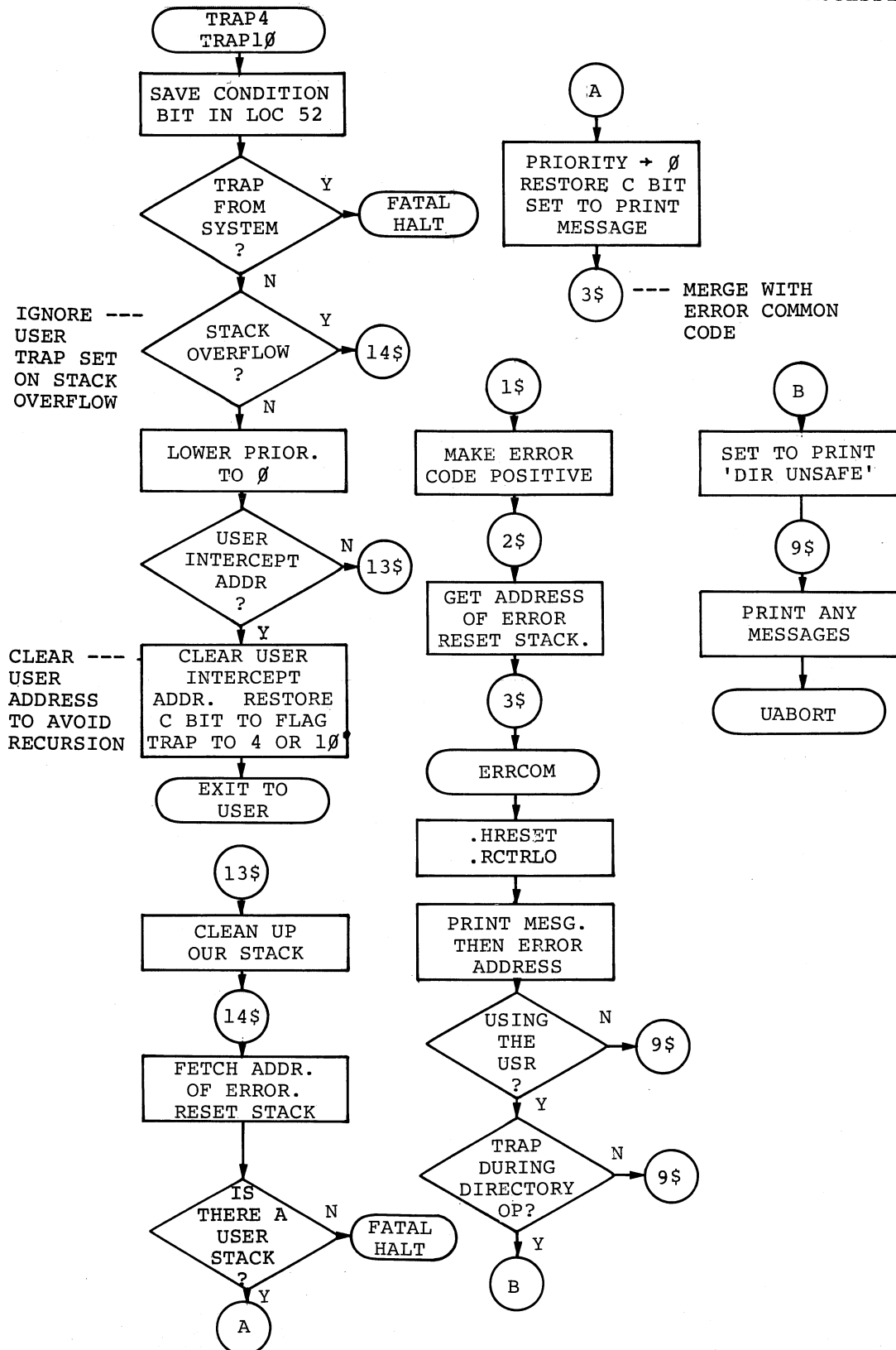
CHANGE CURRENT CONTEXT



CHANGE CURRENT CONTEXT (CONT.)



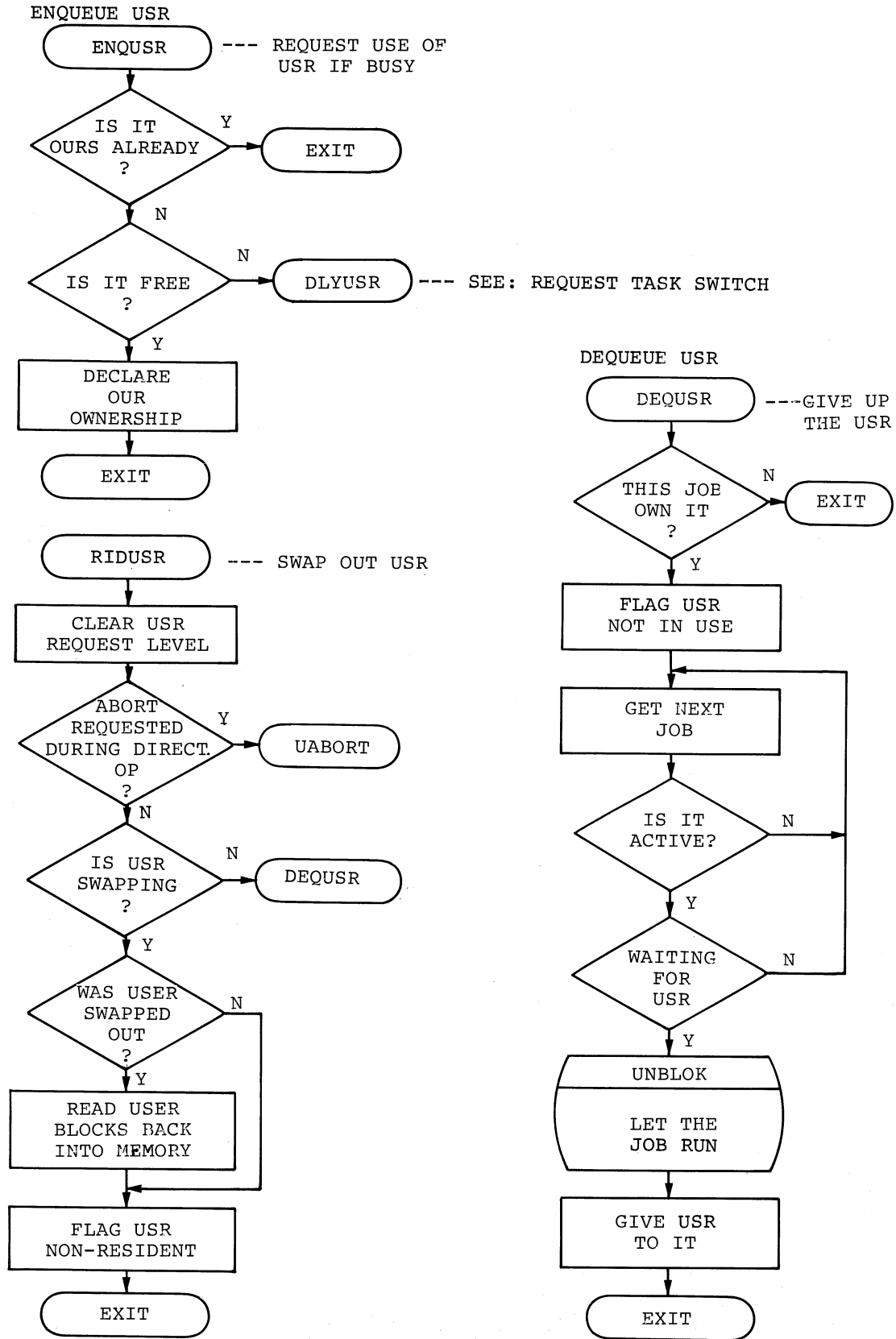
ERROR PROCESSING



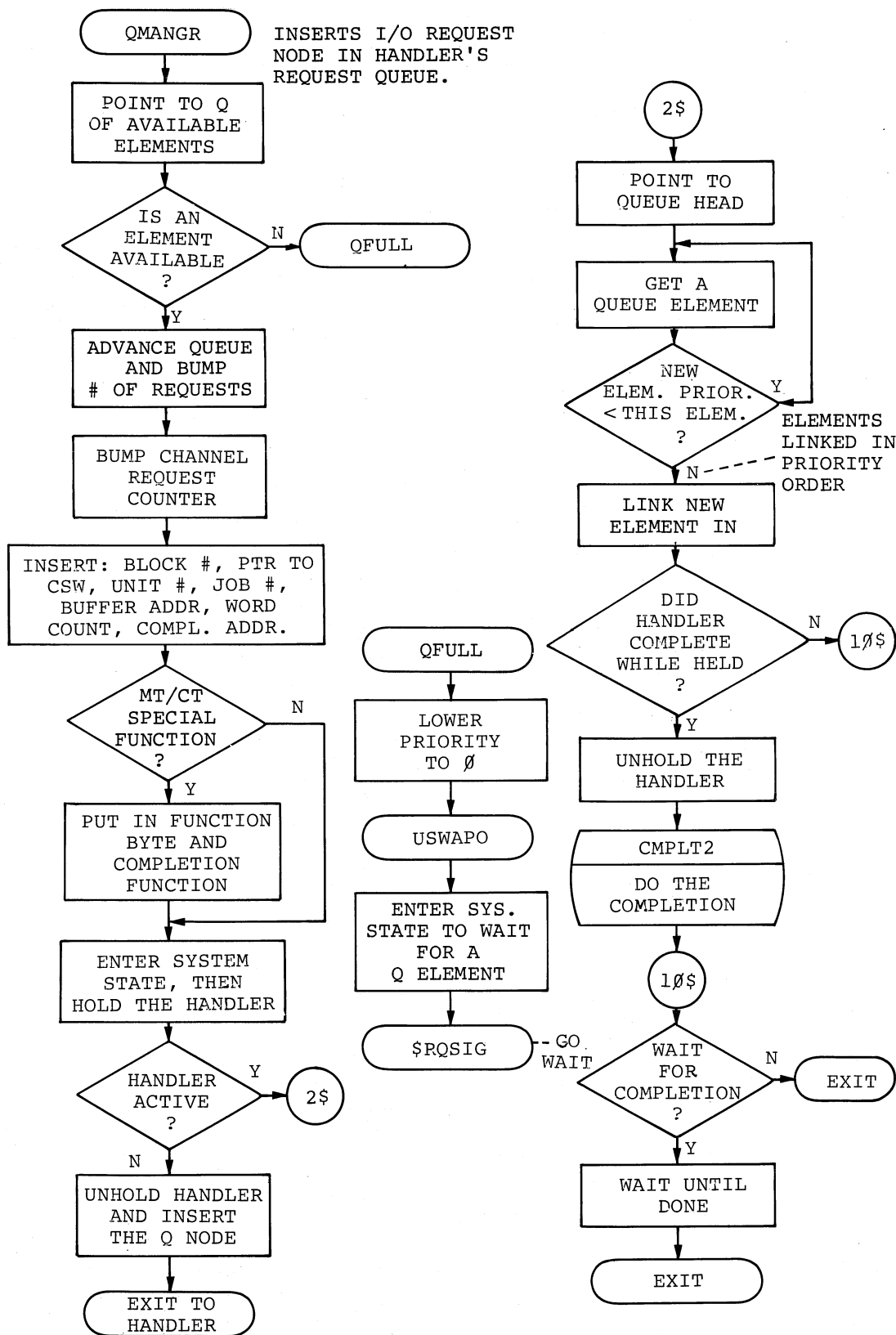


E.5.3 Queue Managers (I/O, USR, Completion)

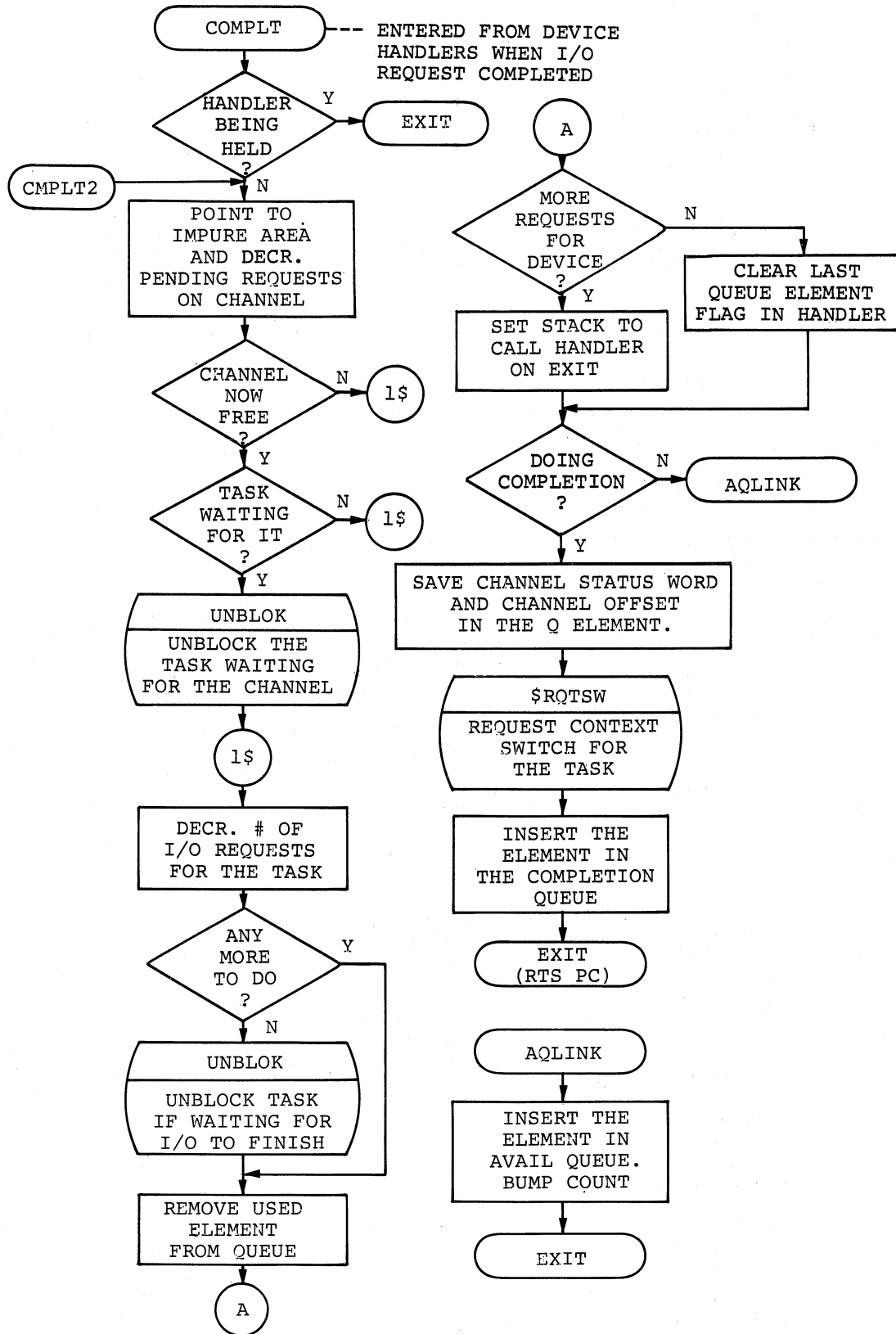
ENQUEUE/DEQUEUE USR



I/O QUEUE MANAGER

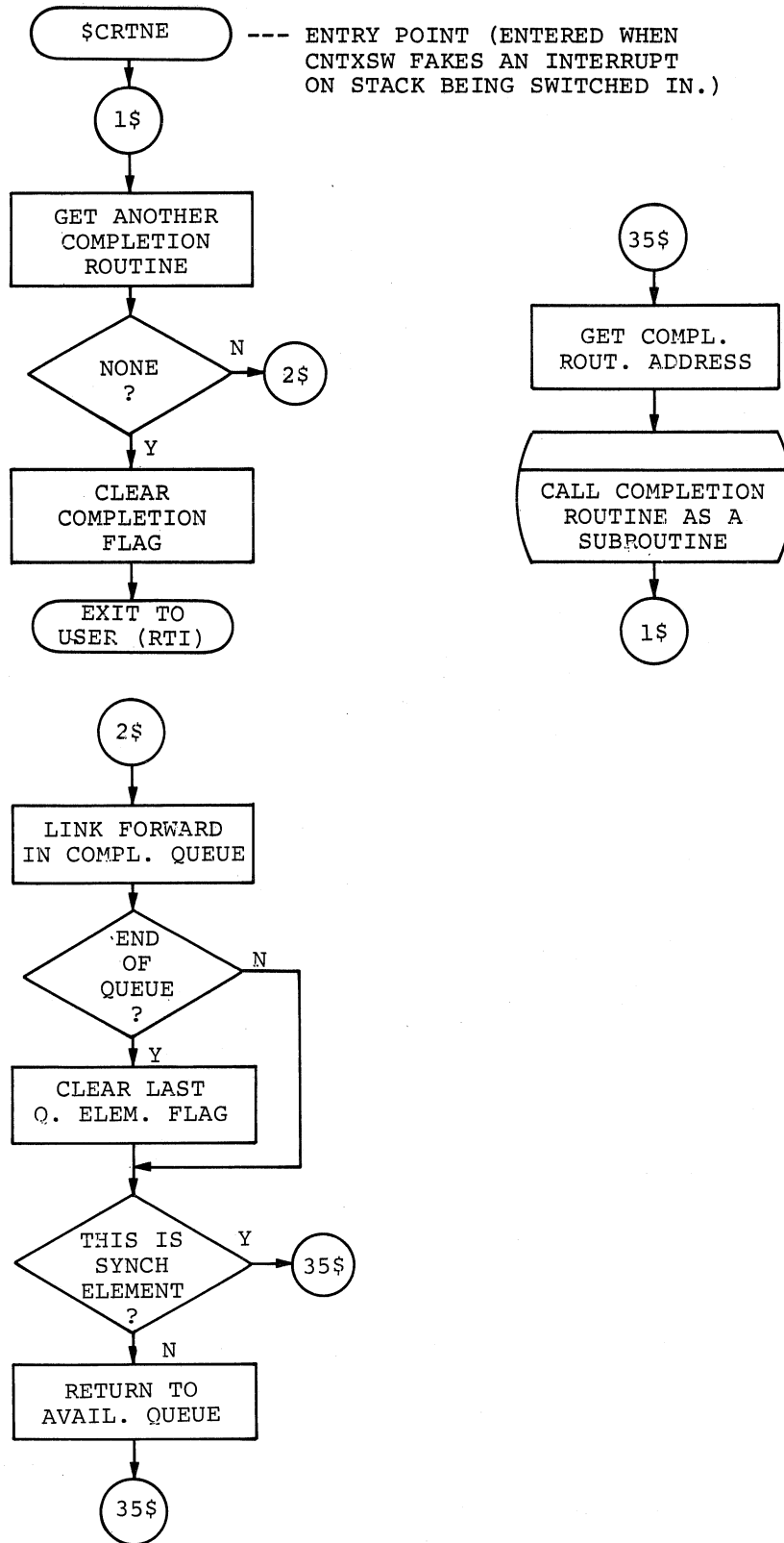


QUEUE COMPLETION





COMPLETION QUEUE MANAGER



(

1.

5.

(

(

(

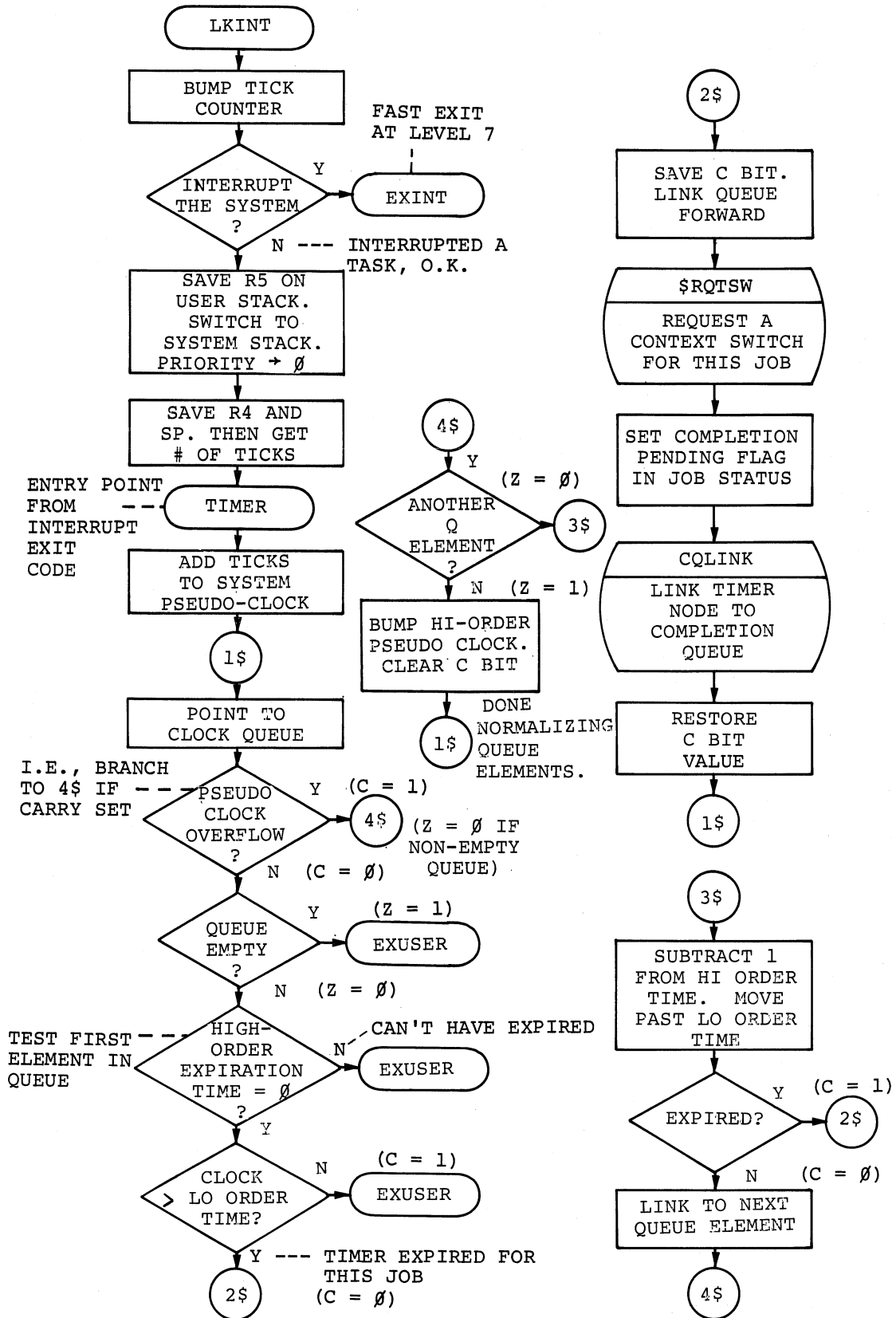
4.

4.

(

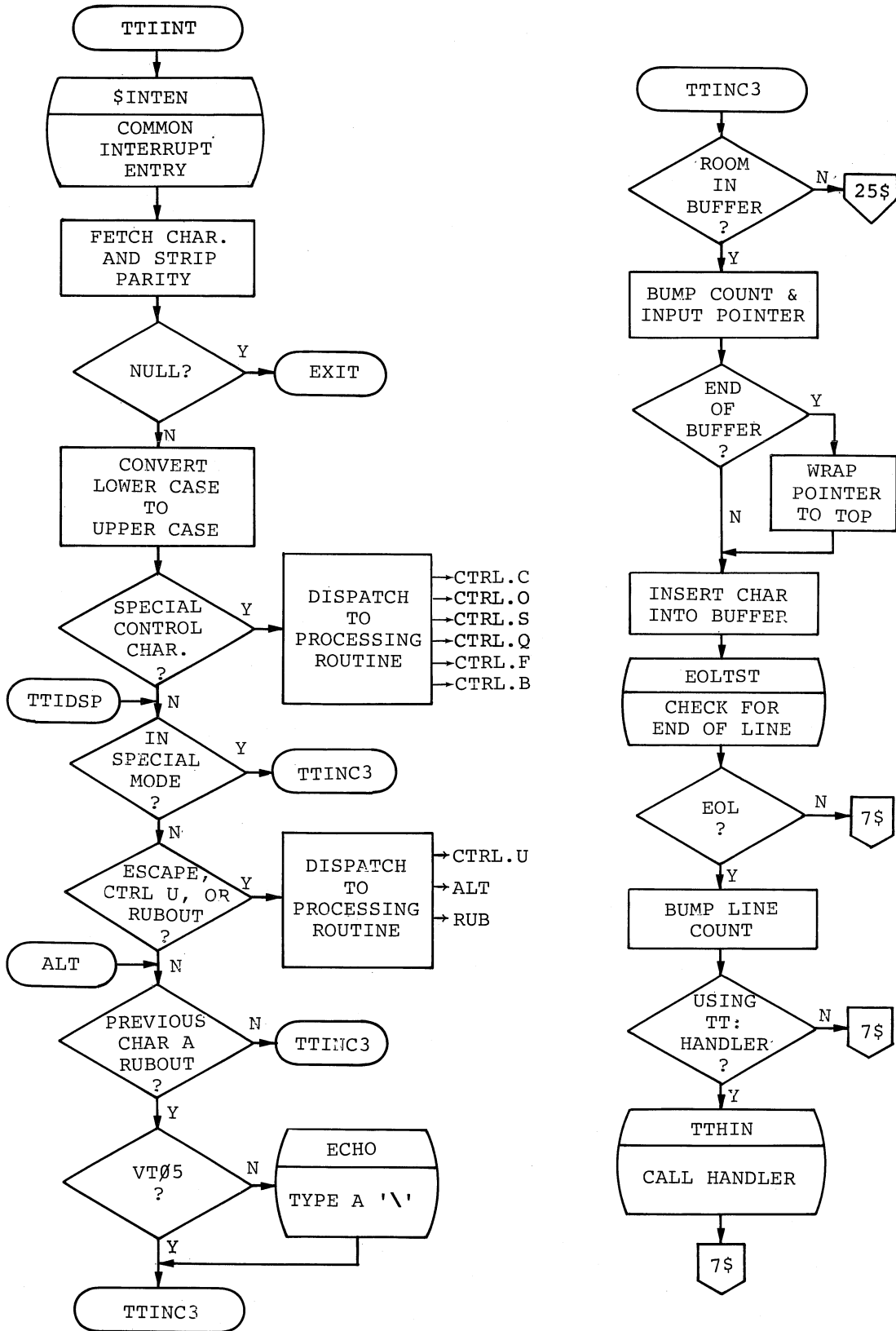
E.5.4 Clock Interrupt Service

CLOCK INTERRUPT HANDLER

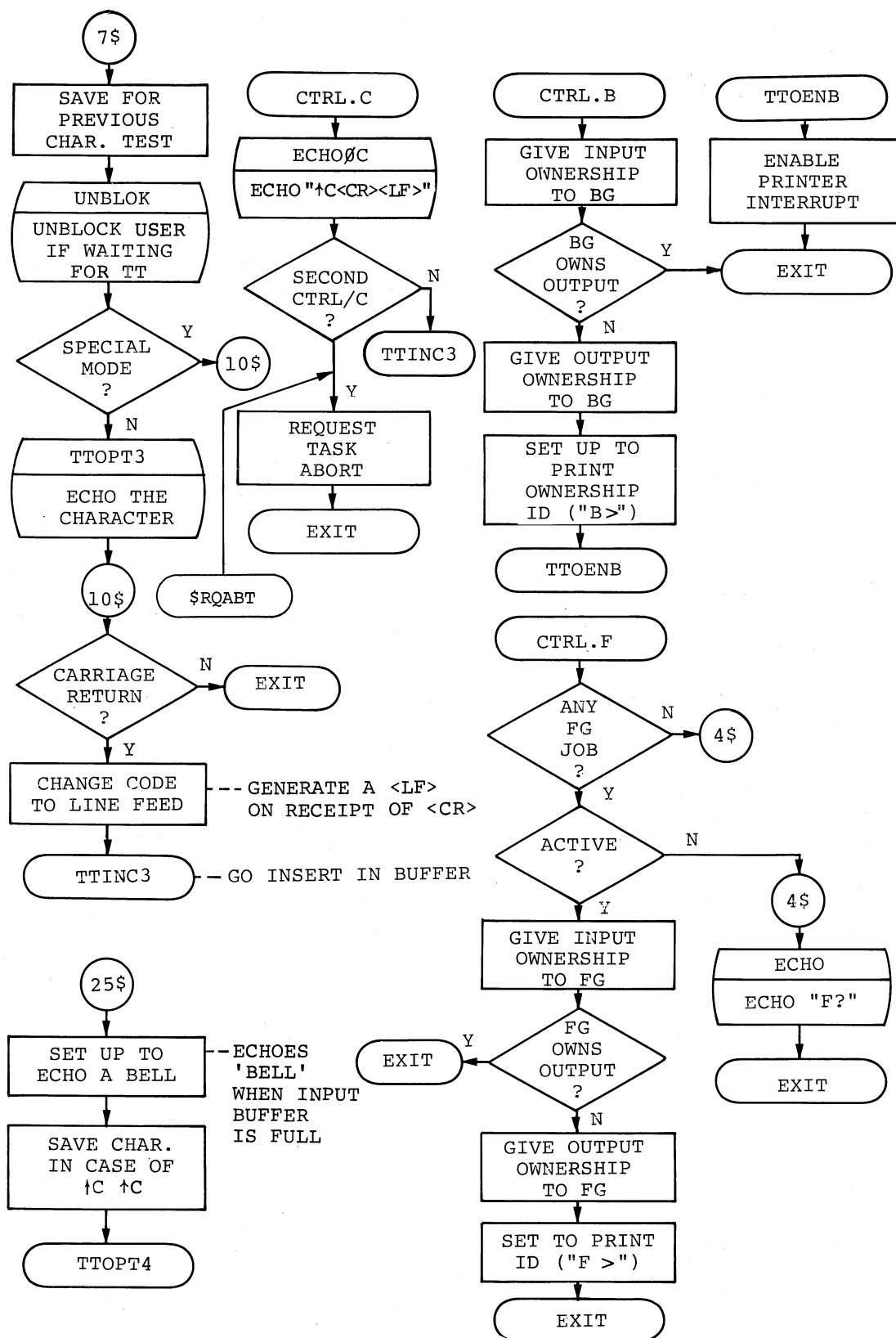


E.5.5 Console Terminal Interrupt Service

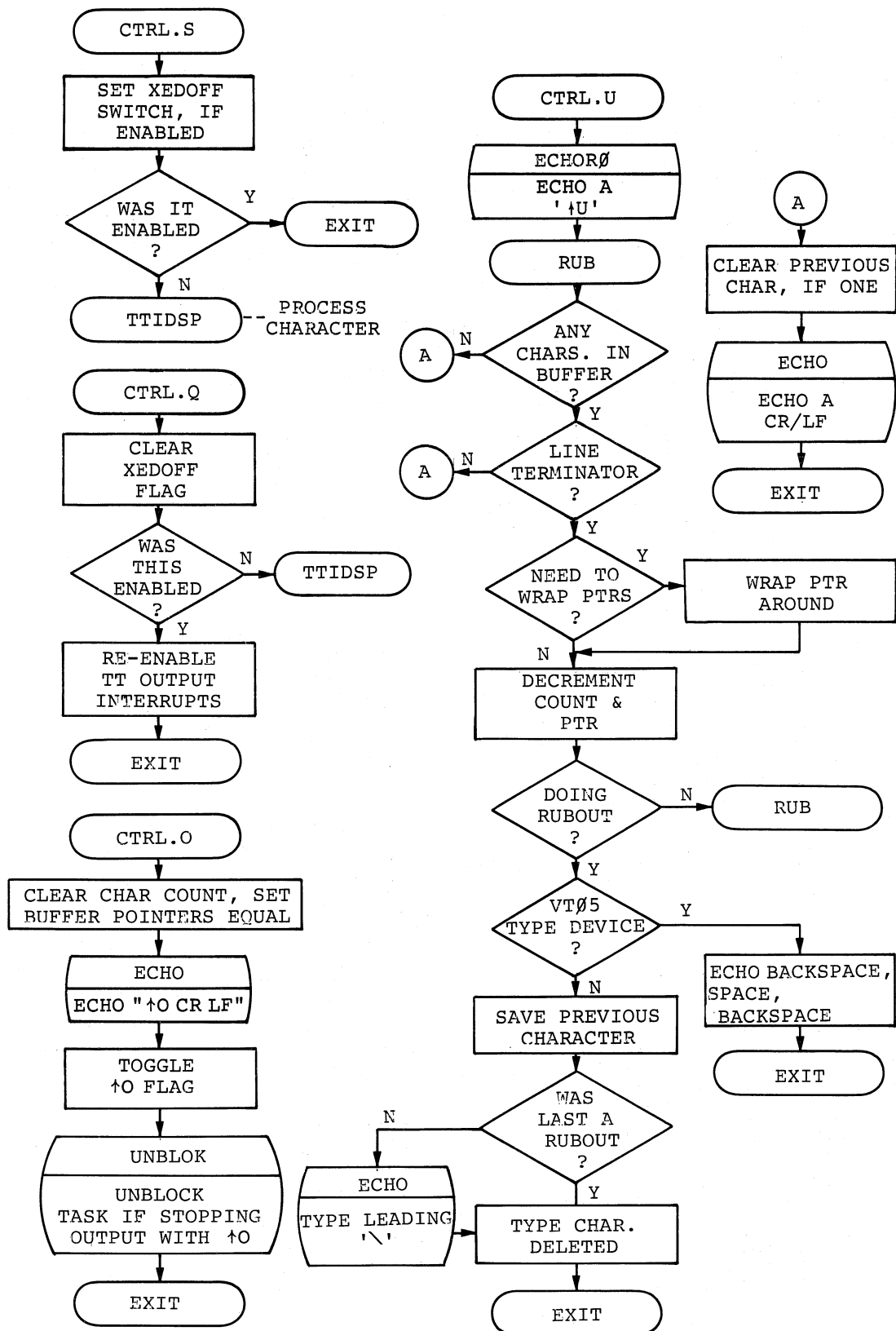
TT INPUT INTERRUPT ROUTINE



TT INPUT INTERRUPT ROUTINE (CONT.)

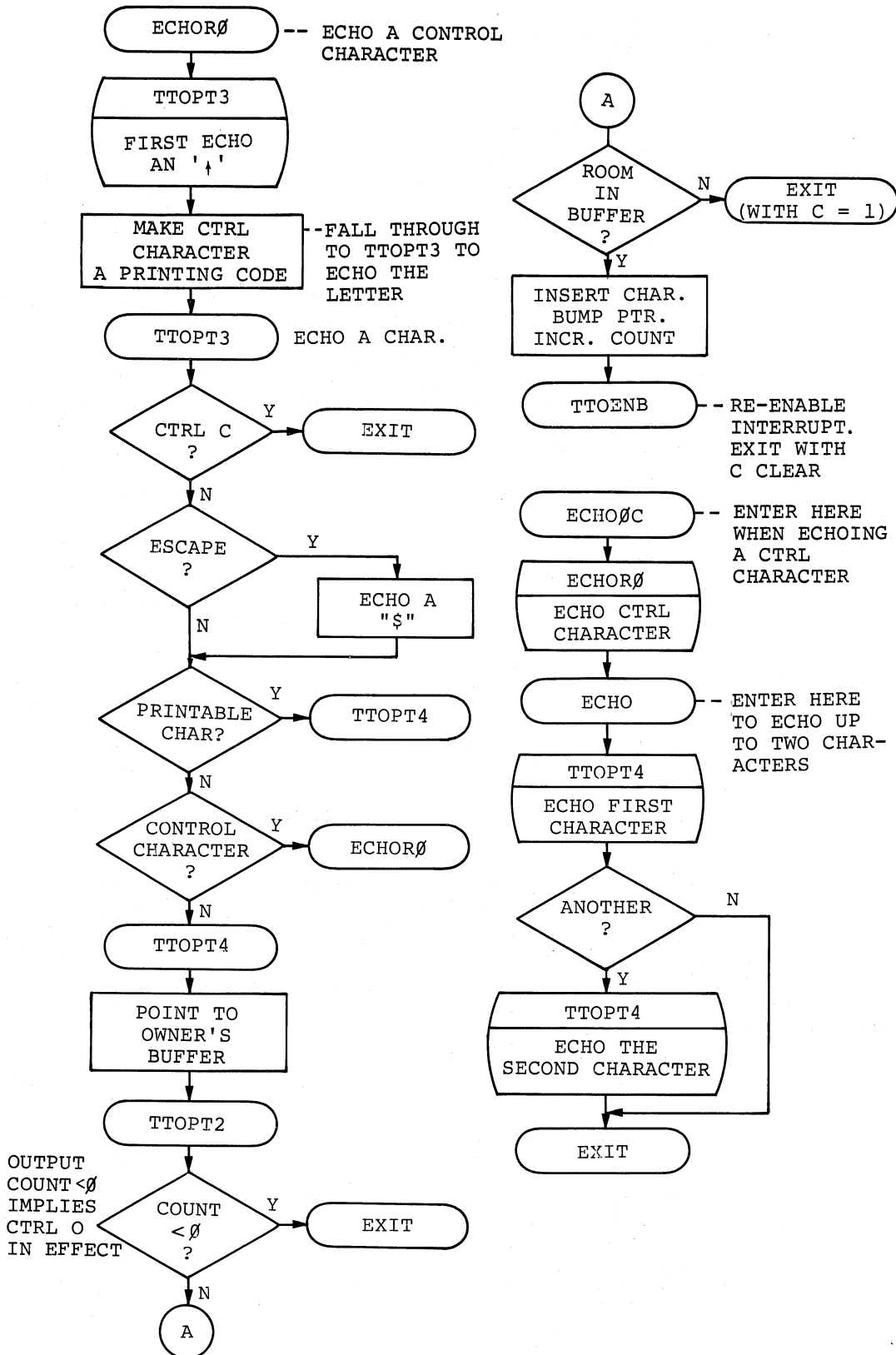


TT INPUT INTERRUPT ROUTINE (CONT.)

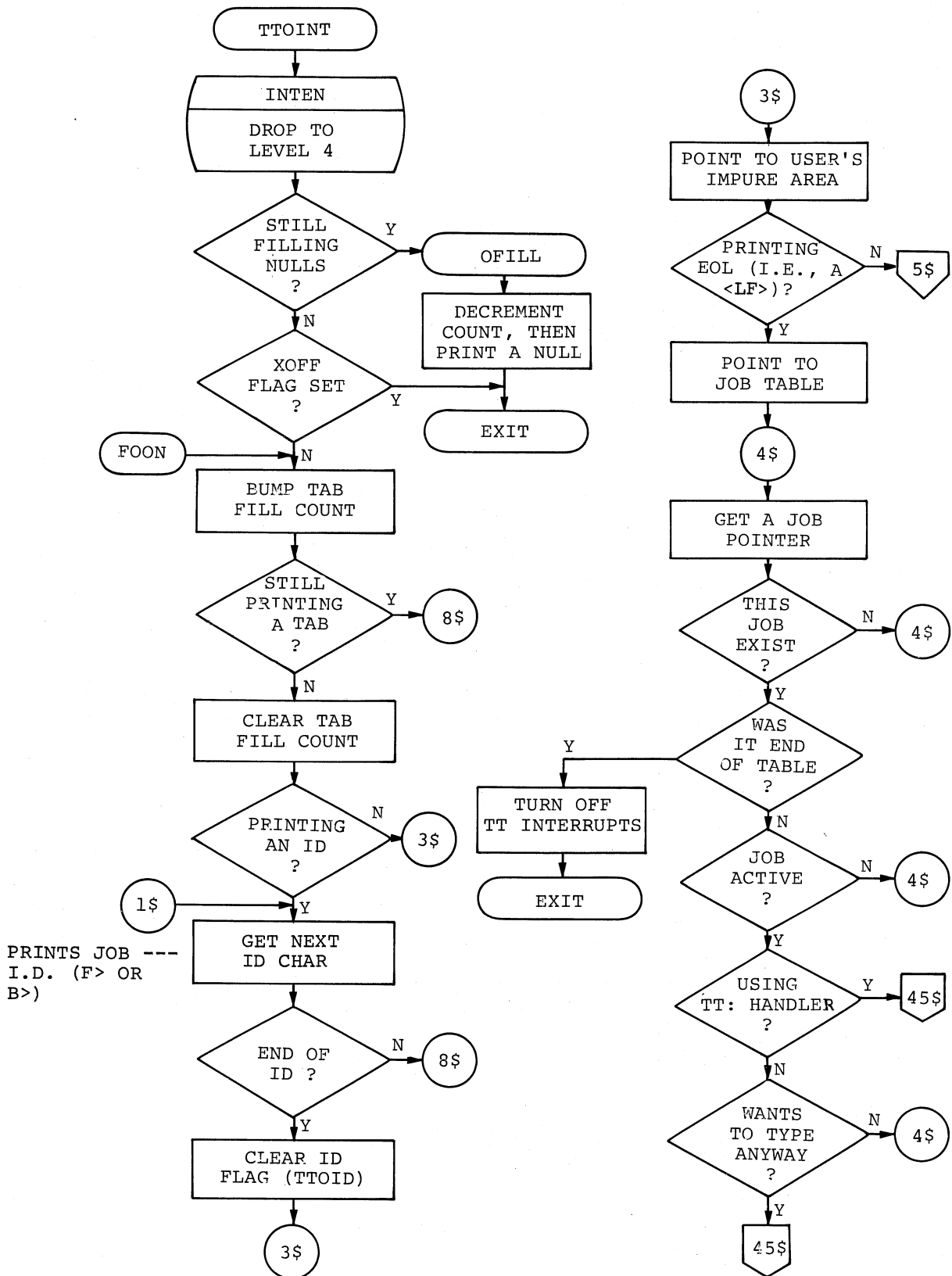




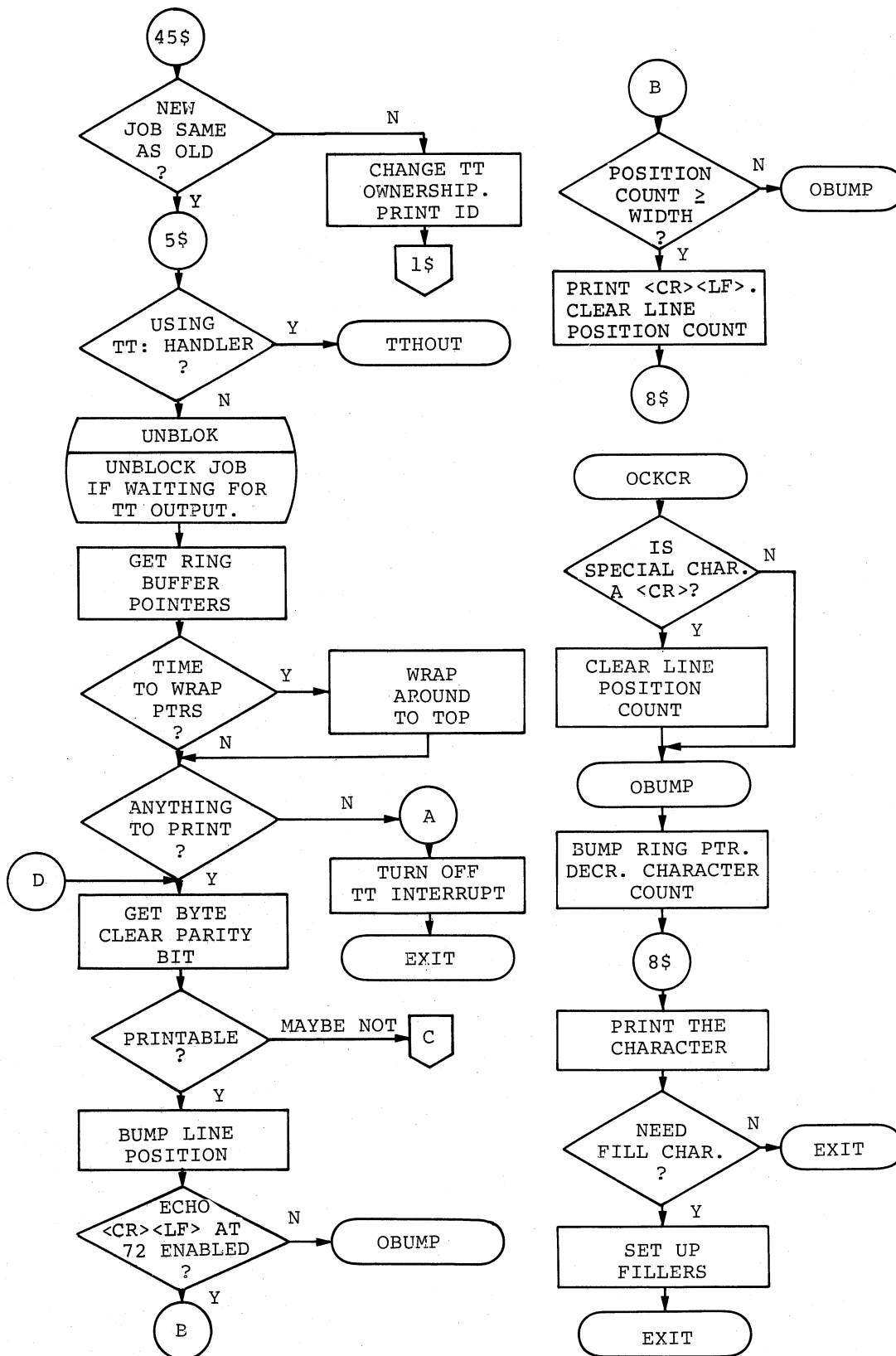
TT ECHO SERVICE



TT OUTPUT INTERRUPT ROUTINE

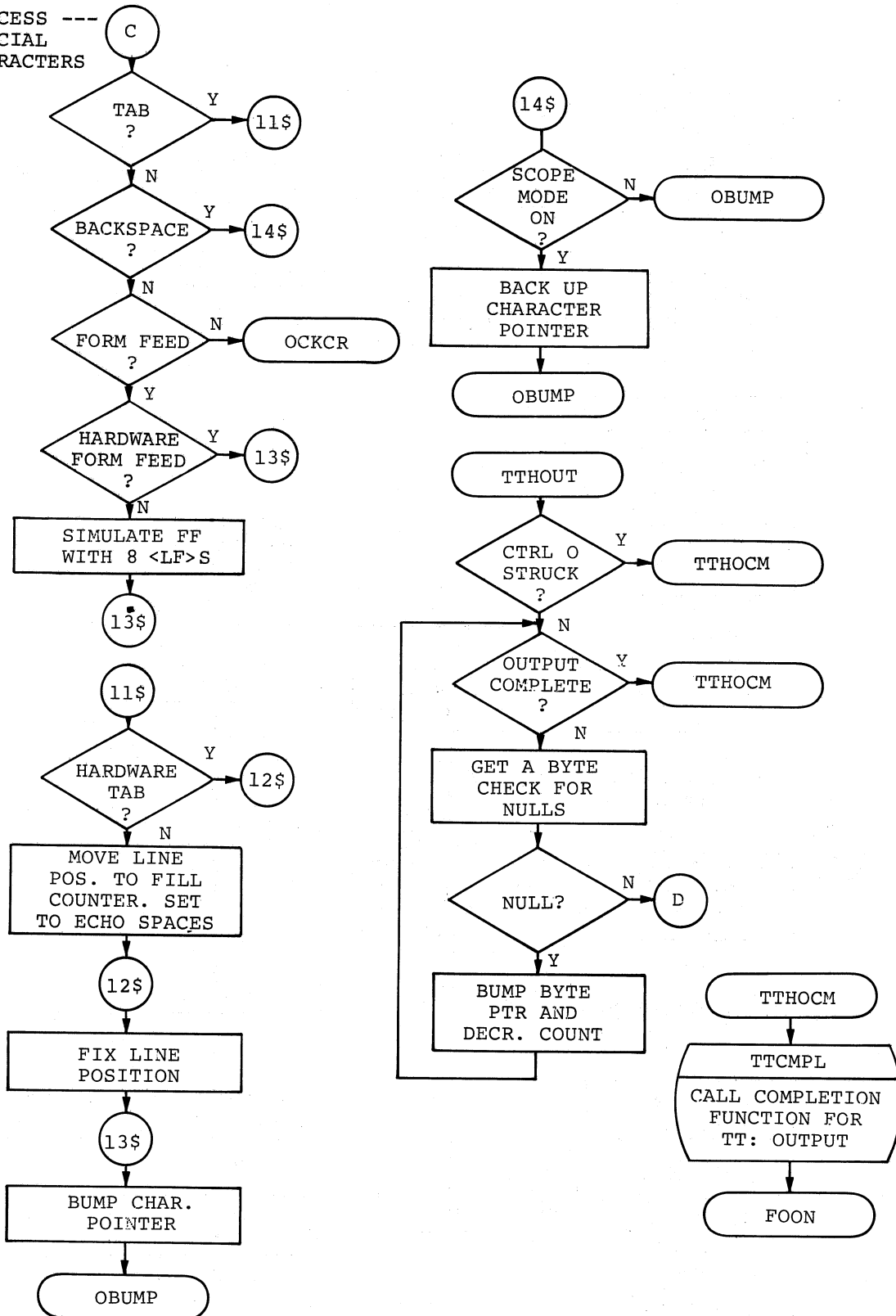


TT OUTPUT INTERRUPT ROUTINE (CONT.)



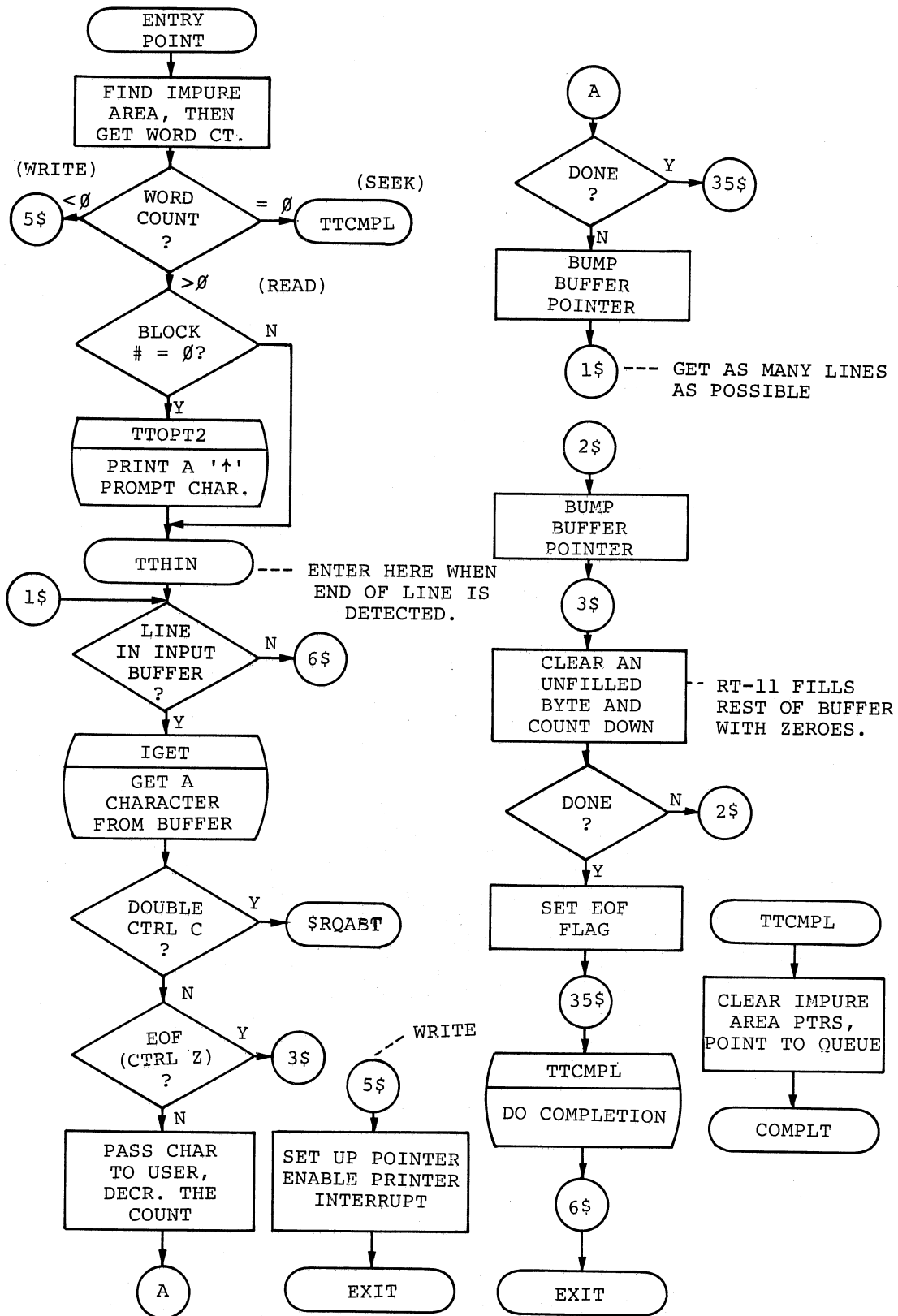
TT OUTPUT INTERRUPT ROUTINE (CONT.)

PROCESS ---  
SPECIAL  
CHARACTERS

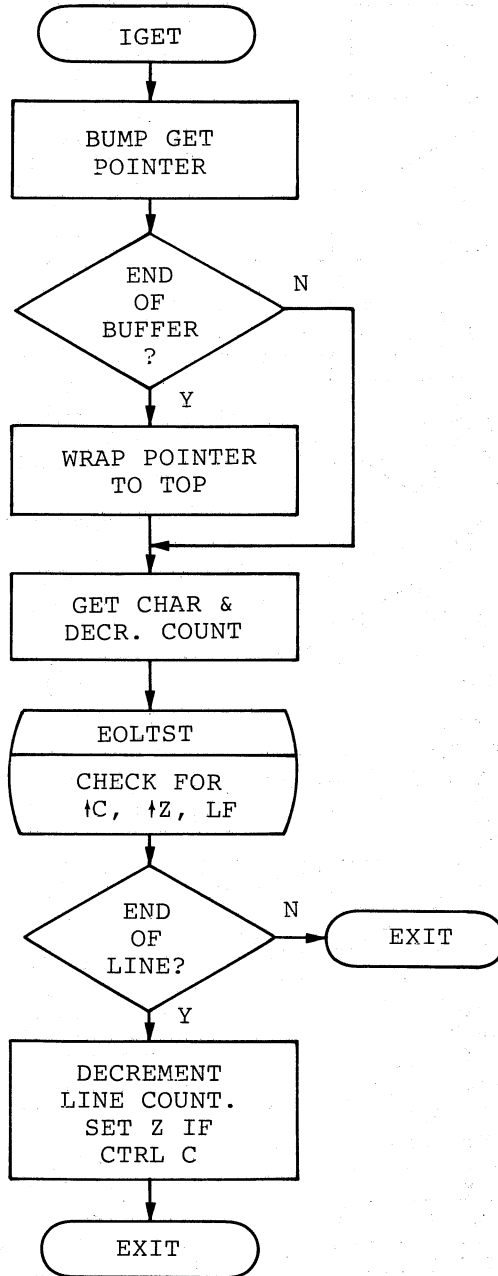


E.5.6 Resident Device Handlers (TT, Message)

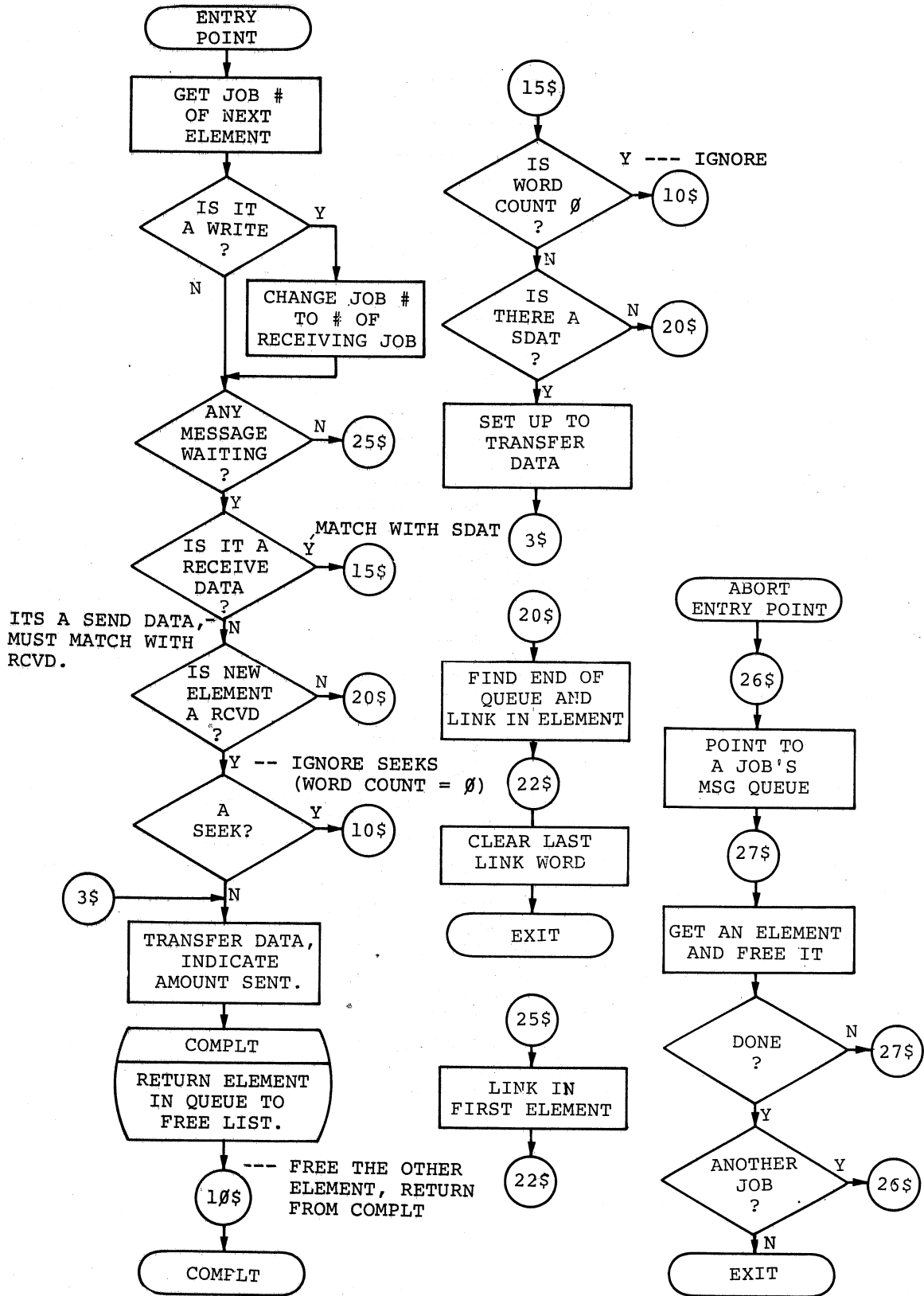
TT: RESIDENT HANDLER



TT: RESIDENT HANDLER (CONT.)



MESSAGE HANDLER





ENTRY POINT INDEX

Page numbers marked by an asterisk indicate the flowchart of the entry point.

\$CRTNE, E-135*	CTRLS, E-89*
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\$INTEN, E-122*	
\$RQABT, E-141*, E-148	
\$RQSIG, E-126*, E-133	D\$DTAT, E-105*
\$RQTSW, E-126*	D\$LETE, E-66*, E-105*
\$SYNCH, E-114*	D\$STAT, E-66*
\$SYSWT, E-126*	D\$VICE, E-113
	D, E-4, E-6*
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C\$DFN2, E-66*, E-72	E, E-4, E-5*
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C\$LOS2, E-66*, E-69, E-105*, E-106	E5ER0, E-111*
C\$LOSE, E-69*, E-106*	E5ER1, E-111*
C\$MKT, E-119*	ECHO, E-143*
C\$PYCH, E-110*	ECHO0C, E-143*
C\$SIGN, E-66*, E-106*	ECHOR0, E-143*
C\$SISP, E-66*, E-106*	EMTCOM, E-104*
C\$STAT, E-111*	EMTDON, E-65*, E-68, E-69, E-70,
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QSET, E-43\*  
QUIESCE, E-117\*

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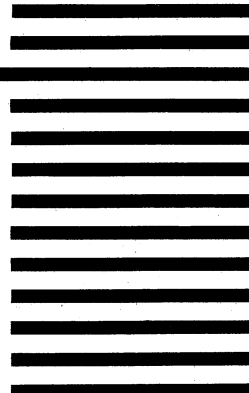
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