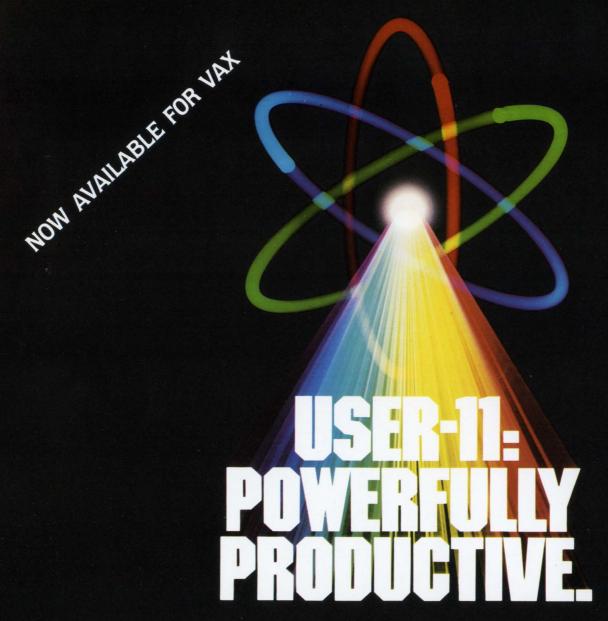
RSTS PROFESSIONAL

Volume 5, Number 3

June 1983 \$1000/issue, \$3500/year





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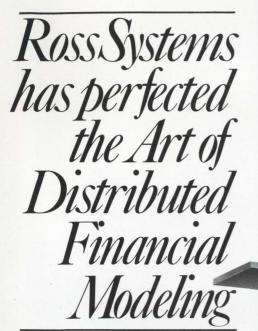
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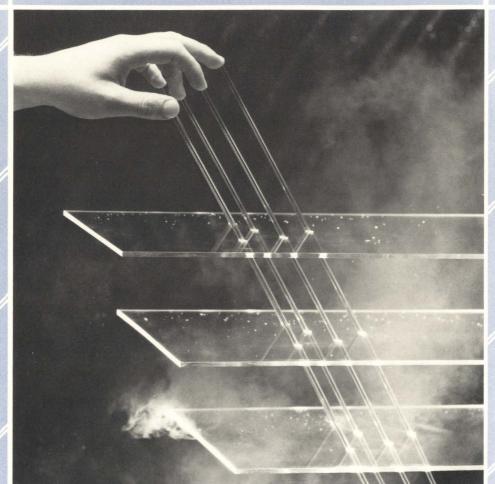
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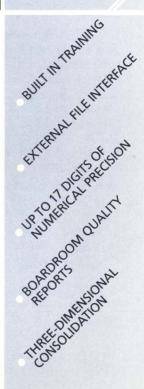
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- Act 4 A 4th Generation System Development
- Таре Сору
- ISCAN.BAS
- A Macro-Sub Program for Determining Time Between Two Dates
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From the publishers . . .

RSTS/VAX PROFESSIONAL??

Carl B. Marbach & Dave Mallery

As of the next issue, we have decided to change our flagship publication's identity. The August issue will be officially titled "RSTS/VAX PROFES-SIONAL." It will have more articles of interest to VAX people and feature several new authors who will continue to improve our editorial content.

We want to take this opportunity to explain our choice to you, our loyal RSTS crew. Like any business decision, this change of name is based on a change of realities. What realities?

Reality #1

The population is shifting. Growing, but shifting. Not exactly lemmings marching to the sea, but if you go to DEC with \$200,000 and tell them you want a computer, you get a VAX 11/750 every time.

Reality #2

Your PDP-11 is getting older. There is no 11/70 replacement (yet). You can't wait forever. Service and electricity are only going up. The J-11(11/70 chip) seems always just over the horizon and probably will emerge as a board-level machine with very different characteristics, more in keeping with DEC's overall marketing strategy. In other words, don't expect any great new exciting PDP-11s in the near future; They will slowly evolve to be smaller, faster and cheaper, but there will not be the PDP-11/90.

Reality #3

More and more of you, our readers, are moving on to the VAX/VMS and frankly, we hate to see you go. Maybe even worse, there are VAX people coming into the world without ever having known the RSTS PROFESSIONAL.

Reality #4

We thought of calling ourselves THE INTERACTIVE PROFESSIONAL, the magazine for Digital technical computer users. A technical journal is what we set out to be; and we want to be a technical journal for all of us who know and love our DEC computers. We realize that a large number of new DEC users are now VAX oriented and we need to include them in not only our readership, but to count the among our contributors.

This is not a magazine for a small clique of gurus. It can and will serve a much larger audience. The more of us there are, the more we can share, and the greater our collective clout will be.

Five years from now, RSTS will be thriving — in numbers that are unimaginable right now. However, it will be the mainstream O/S of the 'new' PDP-11 line:

> The PRO 3xx The MICRO-11 The (MACRO-11??) (sorry but your 44 is obsolete . . .)

Most of the 'hosts' in this brave new, distributed world will be VAX. RSTS will live happily in the box under your desk.

Most of us, willing or not, will be faced with VAX in our everyday work. Now is a great time to start; sharing if you already know, learning if you are new, and growing as we all must to keep up with the exciting world we work and live in.

So come with us. Start to give VAX a heart!





Publishers: R.D. Mallery, Carl B. Marbach Managing Editor: James L. Trichon Director of Advertising: Helen B. Marbach Business Manager: Peg Leiby

Assistant to the Editor: Hope Makransky

Editorial Assistant: Linda DiBiasio

Production Manager: Georgia Conrad

Subscription Fulfillment: Kathi B. Campione, Claire Hollister Steven Barsh, Margie Pitrone Connie Mahon

United Kingdom Representative: Pauline Noakes RTZ Computer Services Ltd., P.O. Box 19 1 Redcliff Street, Bristol, BS99-7JS Phone: Bristol 24181

Mid-Atlantic Representative: Ed Shaud

P.O. Box 187, Radnor, PA 19087 Phone: 215/688-7233

N.Y. & New England Representative: Jack Garland

P.O. Box 314 S.H.S., Duxbury, MA 02332 Phone: 617/934-6464

Contributors:

R. David Broom, William R. Davy, Peter Dick, Steven L. Edwards, Alexander Ehrlich, Mark Gilmore, Philip Hunt, Ken Isaacs, Michael H. Koplitz, Bob Meyer, Maury Pepper, Terry G. Ridgers, Greg Wenzel.

Cartoons: Frank Baginski, Emily Dahlstrom, Brian Hansen

Design, Typesetting & Layout: Grossman Graphics

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- DUS The set of CALLable subroutines which pre-extend file directories, reducing fragmentation.

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LETTERS to the RSTS Pro...

Send letters to: Letters to the RSTS Pro, P.O. Box 361, Ft. Washington, PA 19034-0361.

I can't let the bad words on ABLE computer and ENABLE pass without responding. We installed ENABLE here in Chile last year, with only one problem due to ABLE — the first receipt of the patches did not support our DMC-11 with COLINK. Once they fixed that we have had no problems. With three-fourths of megabyte of memory, our 11/34 ENABLEd with cache will support about 20 – 22 user terminals, some printers, and 36 total jobs. In contrast, our one megabyte 11/44 will support about 34 – 36 user terminals with the same job mix. More memory probably will not help the 34 since it appears CPU bound.

In addition, our ABLE DH/DMs have been faultless, and we are delighted. We have nothing but good things to say about ABLE, and their relations with their customers, and with their products.

However, I can not say enough bad things about the National Semiconductor NS-11 boards — they are both unreliable in themselves, and obviously are not true DEC substitutes, since DEC boards work perfectly well with ENABLE. One of our NS-11 boards can only live on the 18-bit bus, and the other won't keep its DC-DC converter working long enough to be of use.

> J. Michael Hewitt Santiago, Chile

* * *

Thank you for the fine magazine; I look forward to every issue. I find I use my RSTS PROFESSIONAL "library" just as much as I use my DEC RSTS/E manuals. I particularly enjoyed Scott Banks' articles on disk directories and I have applied his ideas to create a program that may be of interest to other RSTS DIBOL users.

RSTS is alive and well in SKI TOWN, USA!!!

David S. Williams Jr. Director of Computer Services Yampa Valley Electric Assoc. Steamboat Springs, CO Ed. Note: See Dave's article, ISCAN.BAS in the next issue.

* * *

The February 1983 issue of the *RSTS Professional* carried both an article and a letter from Alan Woloshin discussing Digital's single-board 11/44. Both are very wrong.

Here are the facts:

• Digital currently markets two Unibus PDP-11 processors: The 11/24 and the 11/44. The 11/24 is a single hex-board processor and is pinned somewhat like the 11/04 and 11/34. This is probably where the confusion arose. The 11/44 is a multi-board

processor, consisting of between 4 1/3 and 8-hex modules, depending on options.

• Each processor requires a unique backplane. Each backplane was newly introduced with its respective processor. We in no way support operation of the processor module(s) in previous backplanes.

• The 11/24 uses existing MS11-L parity memory (albeit in a new mode) or the new MS11-P 1Mb ECC memory. The 11/44 uses MS11-M or MS11-P ECC memory.

From these you can see that Mr. Woloshin's suggestion won't work. We suggest instead that:

• If you now have an uncached 11/34, and are simply swapping too much, the 11/24 may be the right solution for you. It offers 22-bit (4Mb) addressing to solve your swapping problems but uses a slightly slower CPU.

• If you now have an uncached 11/34, and need more memory and CPU, the 11/44 is the right solution for you. It also offers 22-bit (4Mb) addressing and a faster CPU.

• The speed comparison between a chached 11/34A and the 11/44 is not so dramatic. Here, the principal advantage is the larger physical memory.

The only DEC-supported method of upgrade is (CPU) box replacement. RSTS users can just move all your I/O from your existing 11/34 CPU box to your new 11/24 or 11/44 CPU box or cabinet, boot up, and reset the DEfaults. Some product lines at DEC even offer an upgrade (trade-in) program.

Atlant G. Schmidt, Principal Engineer PDP-11 Systems Engineering Maynard, MA

* * *

I applaud your editoral, *The Right Tool* (April, 1983), in which you point out how RSTS continues to lag behind other DEC operating systems. DEC, although pledging their support of the RSTS community, continues to ignore our needs or place them behind the needs of the VAX and RSX communities. Is it that we don't have enought clout?

The RSTS community must be able to present its needs and problems to DEC and expect some action on their part. I have heard estimates which claim that there are 10,000 RSTS sites across the U.S.A. (and this may be conservative). It would seem to me that 10,000 sites should be enough clout to convince DEC that RSTS is a product which must remain state-of-the-art and have their full support. What we need is a united voice.

The Independent RSTS Users Society (IRUS) can be effective in using clout to influence DEC. In addition, it can assist member sites in overcoming some of the deficiencies of RSTS by providing a means of exchanging information and experiences with other RSTS sites.

If indeed there are 10,000 sites out there, then there should be 10,000 members of IRUS. I invite each of your readers to join.

> Carlos Flores, IRUS Chairman Superior Computer Systems

Rockville Center, NY

P.S. IRUS membership dues are \$75.00 per site, per year. To join, contact Joyce Leonard, IRUS, Inc., 3657 Post Road, Warwick, RI 02886.

* * *

Your magazine is a fine instrument for the transfer of expertise, and we thank you for it.

Do you, your associates, or any of your readers know of a DIBOL program generator? I would be grateful for any informed response, yea or nea.

> Richard L. Logan, President Access Information Systems, Inc. Ashland, MA

I've just received your inquiry as to why I haven't renewed my subscription to the RSTS PROFESSIONAL.

* * *

Let me say that it's been my pleasure to subscribe to such a high-quality and *useful* publication since Volume 1, Issue 1. I don't know of any other source that can give me such a wealth of information of RSTS as the RSTS PROFESSIONAL.

However, things do change, and my office now runs 4 VAX 11/780's in addition to one PDP-11/70 (and RSTS). Furthermore, the PDP will be phased out by late summer, so in a nutshell, no more RSTS. While I find *The VAX Scene* to be an informative section, I hope that you might find it possible in the very near future to expand it into a VAX PROFESSIONAL that will rival the RSTS PROFESSIONAL in scope and quality.

> Les Hino University of Hawaii Honolulu, HI

Ed. Note: Please see our editorial on page 4. Welcome back! * * *

I am writing in response to the fellow who is looking for an escape from XQWIK. There was a contribution to the Spring 1982 RSTS symposium tape called CSORT which may be of help.

CSORT is a high performance replacement for XQWIK. It supports the same data formats and creates the same key files. However, it is ten times faster.

> Ken Harris, software consultant Greendale, WI

> > ~

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RSTS/E DISK OPTIMIZATION IN A MULTI-USER ENVIRONMENT

By William R. Davy, System Performance House, Inc.

ABSTRACT

A great deal of existing literature addresses the important matter of RSTS/E disk optimization. This article extends beyond the conventional wisdom to describe previously unpublished optimizations available for multi-user RSTS/E systems. Included are a review of common disk optimization practices, some observations about the multi-user RSTS/E environment, and how these two interact.

Most RSTS/E users are painfully aware that their systems tend to be disk bound. They perform far more disk accesses than are needed for mere data retrieval, program loading and swapping. Furthermore, disk seek time for these operations and others is longer than necessary. To minimize these problems, users are generally limited to the following methods.

• Center and pre-extend the UFDs. Some shops also preextend the MFD contiguously starting at device cluster one.

• Map files contiguously and give them the contiguous attribute where appropriate.

• Increase file clustersize so that the file is mapped in seven clusters or less, up to the maximum clustersize of 256.

• Center the swap files, run-time system files, etc. Tedious manual procedures generally limit these efforts to the few files which are perceived to be most used.

• Increase the pack cluster size to decrease FIP overhead.

• Run REORDR frequently.

Use the "new files first" attribute on the disks.

•Allocate some free space near the center of the disk. Two major focuses of this article will be to correct the misconceptions associated with the above steps and to describe other significant optimizations. It is assumed that the user is familiar with the RSTS/E file structure. There have been excellent descriptions by Mike Mayfield, Scott Banks, et al.

THE RSTS/E ENVIRONMENT

- Software Characteristics

RSTS/E systems are by definition multi-user systems: their performance problems arise under multi-user conditions. Consequently, our optimization efforts will focus on these.

The fundamental consideration of RSTS/E disk optimization is that consecutive disk accesses to the same file, UFD or MFD are statistically rare in a multi-user RSTS/E system.

For example, take a system whose file clustersize has been optimized and whose directories have been recently REORDRed. Consider what would happen if there were a number of jobs performing heavy I/O to disk files and a few jobs excercising FIP through directory accesses, etc. (i.e., a typical RSTS/E system). We would find enough idle time so that each job would receive the CPU time to queue its next disk access well before the system could process the other pending requests. For most of the time, each job would have a pending request. No given job would be able to receive two consecutive disk accesses, the RSTS/E "fairness" algorithm notwithstanding. If each job accessed a different set of files, no file would ever receive two consecutive accesses.

What about accesses to MFDs and UFDs? It is fairly well known that FIP is single-threaded. That is, it will process any operation to completion before starting another. This guarantees that two jobs will never perform FIP operations simultaneously. Although some FIP operations require multiple UFD accesses (e.g., a file lookup in a large directory), there are other jobs which do file accesses without FIP. There will be file accesses in between the UFD accesses.

Furthermore, FIP always accesses exactly one MFD/UFD block at a time. Monitor statistics show that the number of directory accesses always equals the number of directory blocks transferred. The value of this observation will become apparent when we discuss UFD optimizations.

- Hardware Characteristics

DEC and third-party vendors offer a number of disk drives with different sizes and speeds, but they all have some common characteristics.

First, total time for completion of a disk read or write operation is equal to: SEEK TIME + ROTATIONAL LATENCY TIME + TRANSFER TIME. Seek time across just one track is significant when compared to either rotational latency or transfer time. Seek time increases less than linearly as the number of tracks increases.

Rotational latency time is the time it takes for the disk to rotate the transferable data under the head(s) after the seek has been performed. The fundamental consideration discussed above shows us that it is nearly impossible to predetermine the rotational state of the disk before any given operation. Therefore, except where special circumstances warrant otherwise, it is accurate to assume an average latency time equal to one-half the maximum latency time.

From a statistical standpoint, transfer time is small compared to the other two components. Consequently, our major hope of speeding up accesses lies in reducing seek time. This is accomplished by accessing blocks on the disk which are "close" together. For now, it will suffice to say that blocks which are close together on the disk have block numbers which are near each other, and vice versa.

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CIRCLE 108 ON READER CARD

A) Determining the disk center

In many disk optimizing schemes, the UFDs, free space and certain files are centered. Some programs calculate the center of the disk to be the median numbered block. Others recognize that the entire disk may not be allocated, so that the center is better considered to be closer to the beginning of the disk, say perhaps one-third of the way from the beginning to the end.

These schemes are feeble attempts to guess the optimal "center" of the disk. A much better position can be calculated. Simply stated, the center would be the block number equal to half of the space needed for all of the files on the disk, plus some free space and the space needed for the UFDs and the MFD. An improvement upon this algorithm would be to subtract the size of all of the files not used during time sharing. (They would be placed at the end of the disk where they would not get in the way.) The other files, the UFDs, the free space and the MFD would be placed on the disk, starting at the beginning. The calculated center would then be at the center of the active files. In this article, the term "center" refers to this.

B) Pre-extending the MFD

Some installations pre-extend the MFD. Presumably, this is done to make it contiguous. In a single-user system, this can speed up MFD searches, but in a multi-user system, it practically guarantees that MFD seeks will be as slow as possible. The MFD would be entirely contained at the edge of the disk. As the fundamental consideration shows, we are unlikely to get two consecutive accesses to the MFD. A much better strategy for optimizing the MFD is described later.

C) Placing and pre-extending the UFDs.

The fundamental consideration shows that the main reason to pre-extend a UFD is not to make it contiguous, but to control its general position. The proper position for UFDs is near the center. The strategy of placing UFDs near their associated files may make some sense in a single-user system, but it is folly in a multi-user system. It will guarantee that UFDs — the most heavily accessed blocks on the disk — will be scattered as far apart as possible. It will also guarantee that if the most used files are in different accounts, they will also be scattered as far apart as possible, maximizing seek time.

The second question concerning UFDs is how much should they be extended. One strategy is to extend them to their maximum size, 112 blocks. On a system with 100 accounts, this requires 11,200 blocks of prime space on the disk, most of which will never be used.

Another strategy is to use the minimum amount necessary to hold the current directory information. This strategy is poor on two accounts. First, many systems add new files to their accounts without first deleting others. If the UFD is full, it will be extended, not necessarily in a central location. Second, when files are deleted and recreated as they are by editors, compilers, etc. they are often recreated with smaller clustersize which requires more UFD space. If the UFD was created just large enough to hold its previous contents, it will have to be extended to hold the expanded mapping information. So it is clear that some scratch space should be left in the UFD. The UFD optimization section will have more on this subject.

D) Using the optimum file clustersize

Optimum file clustersize is generally considered to be the smallest clustersize which will map the file in seven clusters or less. If the file is larger than 256*7 blocks, the optimum clustersize would be 256, the maximum clustersize. The reason for this is that the "retrieval blockettes" in the UFDs each hold pointers to seven clusters. If the file has seven clusters or less, then the minimum number of blockettes is needed to map the file.

Optimum clustersize helps performance two ways. First, since only one retrieval blockette is kept by FIP at a time, fewer "window turns" will be performed to access the file. Second, by having fewer blockettes in the UFD, it will be more compact and cached more efficiently.

There are some disadvantages to using the optimum clustersize. The space allocated to the file is a multiple of the clustersize, regardless of what is actually used. This can increase disk usage up to seven percent in a typical case, and up to 14 percent in certain pathological cases when compared to using the minimum clustersize.

Raising clustersize also makes it more difficult to place a file exactly where you want it. FIP always places files on cluster boundaries. In cases where it is desirable to pack files of different clustersize close together with no free space in between, using optimum clustersize often prohibits placing a file exactly where you want it.

On certain files, it may actually be advantageous to lower clustersize below the "optimum" value so that the files may be placed on an otherwise unattainable boundary. It may also be worthwhile to save the allocated but unused space that is wasted when using large clustersize. Run-time system and swap files are good examples.

E) Positioning files

Positioning files can increase system throughput by decreasing seek time. One of the popular optimizations is to position files accessed by a particular program as close to each other as possible, if you cannot place them on separate drives. Here is another case where a single-user optimization causes degradation of performance. As we saw when discussing the fundamental consideration, a particular job is unlikely to get two consecutive accesses to the same disk. So what is the purpose of positioning the files? The only value comes from positioning them near the files being accessed by other jobs, which will also cause them to be positioned near each other.

From the above example, we begin to see that we must position files with consideration for all of the files on the disk. The most used files will be placed closest to the center; the least used files will be placed nearest the edges. This algorithm will be expanded into a powerful file positioning strategy.

F) Making files contiguous

At this point, it is necessary to distinguish between the two types of contiguousness. Files whose block numbers are

Like DEC's.



contiguous will be referred to as being mapped contiguously. In addition, there is a RSTS/E file attribute known as the contiguous attribute, which tells RSTS/E that the file is contiguous without FIP examining the mapping. This attribute is what causes PIP and DIR to list a file as being contiguous.

There are three main reasons for mapping files contiguously. First, the file is compactly placed on the disk so that large transfers can be made with one access. Second, the entire file can be placed where desired. Finally, in cases where the file need not be extended, it can be given the contiguous attribute to help reduce window turns (FIP overhead) on files larger than 256*7 blocks. Only on singleuser systems is it necessarily true that head movement is minimized by making files contiguous.

G) Putting some free space near the center

Assuming that most systems create new or temporary files, free space becomes very active file space and should be positioned near the active files. Two questions arise concerning free space: where and how much?

Free space clearly should be positioned near the center of the disk where the most active files should also reside. Most programs, including DEC utilities, do not specify file position when they create new files. FIP searches for a file location when it is created or first extended by starting at the low numbered end of the disk until it finds contiguous space (if possible) for the specified clusters. The ramifications of this are important.

First, if there is free space between the beginning of the disk and the centered free space, it will be used. Therefore, centered free space will only be useful if there is no other free space below it — the low numbered end of the disk must be packed tightly with files.

Second, FIP has considerable overhead searching for the free space. If we have a choice between placing free space on the high or the low side of the center, we should place it on the low side so FIP will find it faster.

H) Increasing pack clustersize

Pack clustersize can be considered the minimum file clustersize. FIP maintains SATT.SYS, which is a bitmap for the pack clusters on the disk. It tells which pack clusters are in use. If those clusters are large, there are fewer of them and FIP can search the bitmap easier. Also, SATT.SYS will be smaller and there will be fewer accesses to it. FIP only keeps one block of SATT.SYS in memory at a time.

Increasing pack clustersize also results in less variation in clustersize between files. The disk becomes less fragmented and it is easier to pack files tightly, eliminating free space in front of the centered free space mentioned in the previous section.

The disadvantage of increasing pack clustersize is that it tends to waste space. Note that on a pack with pack clustersize 16, even a one block file uses 16 blocks. All storage allocations are rounded up to a multiple of 16.

I) Running REORDR frequently

Using REORDR makes systems run faster. The UFD structure can be considered a tree, and REORDR allows FIP to traverse that tree more quickly. However, while there are

many ways to organize MFD/UFDs, the one created by REORDR is almost never optimal.

POSITIONING FILES ON A MULTI-USER SYSTEM

Earlier, we defined the center of the disk. If we could somehow determine which blocks of which files/UFDs were accessed the most during a time sharing session, we could then position those files in their optimum static position by placing the most accessed blocks nearest the center.

There is no practical way to know exactly which files are used the most, but one can make some reasonable guesses by examining the nature of the system. What follows is a description of common files in what is likely to be close to their optimum order. That is, the ones mentioned first should be placed closest to the center.

A) Ordering files around the center

The swap files are frequently accessed on most systems. With the following exception, there is little reason to have more than one swap file on any given disk. The reason for having more than one swap file would be to optimize a system in which interactive jobs swap frequently with event driven jobs, the usual job count is much less than job maximum, and there is good reason to do all swapping on the same disk. Under these conditions, having SWAP3.SYS just in front of SWAP.SYS (and no other swap files) can decrease lost time by requiring slightly shorter seeks between the out-swaps and the in-swaps.

The files OVR.SYS, ERR.SYS, and BUFF.SYS are frequently accessed, if they are used. If OVR.SYS is not used, then the current SIL should be positioned near the center of the disk. Otherwise, it can go to the back edge with the other unused files.

The non-permanently resident run-time systems and resident libraries are frequently accessed. Note that in addition to being accessed at program load time, they are accessed on certain in-swaps. Under certain conditions, lost time can be decreased by positioning these files as near as possible to the swap files. Permanently resident run-time systems are not accessed once a time sharing session has started, so those files should be placed at the back edge of the disk.

MFDs and UFDs should be placed near the center, as will be discussed later.

SATT.SYS is accessed frequently on disks when files are being created or deleted. It should also be near the center.

Free space should be near the center as mentioned earlier.

Next should be frequently used files. Presumably, someone knows which files those are.

The remaining files should be positioned around the others in decreasing order of use, keeping in mind that the free space in the center will not be used until the free space toward the front of the disk is used.

B) Improvements over ordering around the center

The above rules for positioning files are effective until the system is actually used. When files are created and deleted (e.g., when they are edited or rebuilt), the new versions end up where FIP puts them, and free space is created

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elsewhere on the disk. These new creations are not likely to be well placed. Consequently, several additional steps should be taken.

First, as a general rule, files which are likely to be deleted should be placed near the center. This will force the resulting free space to be in a good location. Furthermore, if there is sufficient free space, the files likely to be deleted without being otherwise accessed, can be placed near the high numbered end of the disk. The free space generated when they are deleted will be used last; the well-positioned free space will be used first.

It may seem difficult to guess which files are likely to be deleted, but in reality it is quite easy. With few exceptions, the most recently created files are the most likely to be deleted soon. Think about it, keeping in mind that most editors, etc. do not modify old versions of files, but create new ones, deleting the old ones when done.

C) Determining the most accessed files

With the large file processor, it is relatively easy to scan the list of open files. One might conclude that monitoring this list for the most commonly opened files would reveal the most heavily accessed files. Unfortunately, such attempts will produce poor results.

Consider a typical system which runs error logging, the spooling package with batch processors, and uses the CUSPs frequently. Error logging leaves the error file open continuously and only accesses it in the event of errors. The spooling leaves a number of work files open and accesses them only occasionally. To load CUSPs, the system must open them, read them, and close them. However, all this happens so quickly that the monitoring program is quite likely to never see it happen (unless it is absolutely hogging the system). So any likely monitor program would guess the relative frequency of accesses to those files exactly backwards.

There is a large difference between a file being open and being accessed. Accesses are difficult to monitor without large overhead. Fortunately, some common sense observations can help determine which files are likely to be accessed most often.

Compiled BASIC programs and RT11 and RSX task images, can be loaded into memory with just one file access, provided that they are not overlaid. However, the ones that are overlaid (including PIP.SAV, EDT.TSK, and TKB.TSK), can be accessed many times in the course of one run. TSK files larger than 115 blocks and SAV files much larger than their core images are also overlaid. There is considerable sentiment that suggests that LOGIN should be well positioned because of its frequent use, but consider that in the course of one task build, more accesses are made to TKB.TSK than to LOGIN all day long. The same phenomenon is seen with PIP, EDT, ED2, and LBR.

Files which are accessed by EDT or compilers are typically accessed one block at a time. This implies that even if the file is only opened once, it will be accessed many times. Note that any file with a source file extension likely falls into this category. Data files, especially randomly accessed ones, are typically accessed a few blocks at a time. Object libraries, if they are used, are accessed many times per use. LOG files are not likely to be accessed more than once, sequentially, and possibly many blocks at a time (as PIP would access it).

Programs which access particular files each time they are run are most likely accessed less often than the files they access, especially if they do many accesses to those files.

D) Positioning across cylinder boundaries

Discussions about disk cylinder boundaries are rare in PDP-11 operating systems. On RSTS/E systems, the system takes care of mapping virtual blocks onto physical blocks on the disk without help from the programmer. For the most part, little is gained by positioning files across the minimum number of cylinder boundaries.

Files on most systems are accessed one or a few blocks at a time. The blocks transferred on any given access are unlikely to cross a cylinder boundary regardless of how files are placed. The fundamental consideration shows there is little to be gained from positioning such files across minimum cylinder boundaries, since they are unlikely to be accessed twice in succession.

There are disadvantages in trying to avoid crossing cylinder boundaries. Due to FIP's alignment algorithm, cylinder boundaries always straddle clusters for any clustersize greater than one. On any large disk and on any optimized disk, the pack clustersize will be greater than one. An attempt to have all files straddle minimum cylinder boundaries would create free space at many cylinder boundaries. The degradation from scattering free space would more than make up for any gains from avoiding cylinder boundaries. Furthermore, attempts to position files with large clustersize across minimum cylinder boundaries tend to position them far from where they are desired. (Explanation of this effect is beyond the scope of this article; the mathematically inclined are referred to the Chinese Remainder Theorem.)

There are benefits, though, to positioning the most useful files. Non-permanently resident run-time systems are ideal candidates for minimum-cylinder positioning. If possible, the entire file is transferred in one access. Since there are just a few of these files, it is worthwhile to straddle the minimum number of boundaries with these, if they are used often.

Similarly, frequently accessed BACs are good candidates for intra-cylinder positioning. Notice that on large disks, most of these files will miss cylinder boundaries anyway, regardless of the positioning algorithm. Ditto for unoverlaid TSKs and SAVs.

Overlaid task images are another story. They tend to be larger (harder to position), and they are accessed by a smaller number of blocks at a time, so they tend not to be accessed across cylinder boundaries. The benefit is small compared to the effort and other resulting degradations.

The swap files are large and will likely cross many cylinder boundaries regardless of where they are placed. Optimum placement can at best avoid one cylinder boundary. Once again, the benefit is very small compared to the harm resulting from the likely off-center placement.

The last "files" to be considered are the MFDs and

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UFDs. They are always accessed one block at a time, so no single access crosses a cylinder boundary. The fundamental consideration shows that it is unlikely to access the same MFD/UFD twice in succession, so there is almost nothing to be gained by avoiding cylinder boundaries with MFDs/UFDs.

OPTIMIZING MFDS AND UFDS

UFDs are divided into chunks of eight words called blockettes. There are four different types of blockettes: name, accounting, attribute, and retrieval blockettes. Blockettes are linked together by pointers and except for the requirement that they start on an eight-word boundary, can reside anywhere in the UFD. Any blockette whose first two words are zero is considered unused (free space in the UFD).

The name blockettes contain the name and extension of the file in RADIX-50. The first blockette in a UFD contains a pointer to the name blockette of the first file in the account. Each name blockette contains a pointer to the name blockette of the next file in the UFD. When FIP searches a UFD for a file, it starts with the first blockette of the UFD and follows this chain until it finds the desired file or exhausts the list.

Each name blockette has a pointer to the accounting blockette for its file. This blockette contains the last access date, the number of blocks in the file, the creation date and time, the associated run-time system, and the file clustersize.

Each name blockette also has a pointer to the first attributes blockette if the file has any attributes. The attribute blockettes contain file attributes (up to seven per blockette), plus a pointer to the next attribute blockette for the file. There is a maximum of two attribute blockettes per file.

Finally, each name blockette contains a pointer to the first retrieval blockette. Each retrieval blockette contains the starting device cluster numbers for up to the next seven clusters of the file, plus a pointer to the next retrieval blockette, if it exists.

In a sense, the internal structure of the UFDs is a tree: the name blockettes form the root of the tree, the retrieval and accounting blockettes form branches, and the attribute blockettes are leaves off of the accounting blockettes. This tree can be searched in the forward direction, that is, from the root to the leaves, but not backwards. Since all of the nodes of the tree are found only by following pointers, the nodes (blockettes) can be located in any order in the UFD. However, we will see that some orders are better than others.

A) Straight file copy

When files are copied into an empty UFD one at a time, their associated UFD blockettes are tightly packed into the UFD, one after another, starting at the first available location in the UFD. To scan the list of name blockettes (i.e., do a file lookup), FIP reads through the UFD sequentially, until it gets to the name blockette it is seeking. Straight file copy guarantees that FIP will only read a block of the UFD once when performing a file lookup.

It is a common belief that this system is efficient. It is

much better than the tangled mess that results from creating and deleting files at random. Furthermore, most of the disk structuring utilities leave their UFDs this way. Unfortunately, except for very small directories and a few pathological cases, straight file copy is never best.

B) REORDR

The other well-known UFD optimizing technique is REORDR's algorithm. REORDR puts all of the name blockettes in the requested order starting at the beginning of the UFD. The remainder of the last block of the UFD used for storing name blockettes is left empty. REORDR writes the accounting, retrieval, and attribute blockettes starting in the next block of the UFD. When possible, REORDR writes all of the non-name blockettes in the same block of the UFD so that FIP only needs the minimum number of disk accesses to retrieve them.

For large directories, the above strategy is far superior to straight file copy. To open the last file in a UFD created through straight file copy, every block in use in the UFD must be read (with the possible exception of the last block). To open the last file in a UFD created by REORDR, only the blocks containing name blockettes (usually one-third of the blocks or less) must be read to find the last name blockette. One other block — the one with the accounting, attribute, and the first retrieval blockette — also must be read. So in large directories created by REORDR, approximately onethird the UFD accesses are needed to open files as in directories created by straight file copy.

Unfortunately, REORDRed directories are not always optimal. Consider the case in which there are just a few small files in a UFD. If the UFD was created by a straight file copy, all of the directory information is in the first block of the UFD and can be retrieved with one access. If the UFD is REORDRed, the name blockettes will be in the first block, and the rest of the directory information will be in the second block, requiring two accesses to open any file!

C) DOPTERed UFDs

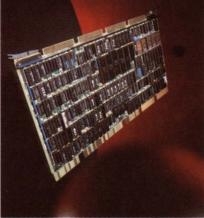
Up to this point, we have hinted that there are a number of improvements to the UFD structure possible. In developing DOPTER, the System Performance House combined the best attributes of straight file copy and REORDR, and added some new optimizations.

DOPTER has an internal routine somewhat similar to REORDR. If all the UFD information will fit into one block, it is put into the first block of the UFD. If there are more blockettes than will fit into one block, the name blockettes are separated from the others as in REORDR, but with some interesting differences.

The most obvious way to reduce the overhead of file lookups is to have the desired files as near as possible to the beginning of the name blockette list. While REORDR makes a good attempt at this with its reverse order sort on creation or access date, it certainly does not do nearly as well as a good heuristic algorithm which will accept help from the user. (Note that sorting in alphabetical order on file name or extension is basically useless for optimization.)

Since DOPTER is very good at placing the most used files at the front of the UFDs, DOPTER leaves room in the





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Associated Computer Consultants 720 Santa Barbara Street Santa Barbara, CA 93101 TWX 910 334-4907 (805) 963-9431 first block of the UFD for all of the blockettes of the first two files. Thus, the most used files can be opened with only one UFD access, a 50 percent improvement over REORDR. Except for this improvement, all of the name blockettes remain segregated from the other blockettes.

To appreciate DOPTER's most significant attribute, it is necessary to understand FIP's method of allocating free space in UFDs. When FIP needs free space, (i.e., when it is creating a new file or extending an old one), it first searches the current UFD block in memory for free space. If it finds what it needs there, FIP uses it. Otherwise, FIP searches the UFD sequentially from the beginning, examining each blockette to see if it is in use.

When FIP finds its current block full, it will search every allocated blockette before it finds a free one if the UFD was created by straight file copy. If the UFD was created by REORDR, it will search through all of the name blockettes until it comes to the small amount of free space left at the end of the name blockettes (if any). If that space has been used (which it will be after a few new files are created), FIP must search to the end of the allocated blockettes, just as in straight file copy.

The solution is to leave some free space near the beginning of the UFD. Since FIP only accesses one block at a time in the UFDs, there is little to be lost with this strategy. DOPTER leaves blocks two through nine empty. Thus, FIP finds free space with a minimum number of reads. In a large UFD, this strategy can cut file open overhead by 80 percent.

There is another important benefit from leaving free space near the start of the UFD. In a typical UFD, files are both created and deleted in somewhat random order. If an otherwise tightly packed UFD has a few files deleted, there will likely be a few free blockettes scattered throughout the UFD. When FIP reallocates the space, it scatters the blockettes throughout the UFD, causing FIP to perform many disk accesses to do file lookups, creations, etc. FIP generates what is commonly referred to as a tangled directory.

When free space is left at the front of the UFD, files still are deleted wherever they are. However, when new files are created, they are built as they would be by straight file copy. As we have seen, straight file copy is not best, but it is far better than the alternative.

D) DOPTERed MFDs

MFDs are similar in structure to UFDs, except that the name blockette entries may be for UFDs as well as files. However, MFDs provide special opportunities for optimization.

The first cluster of an MFD must be at device cluster one. MFDs are normally extended contiguously. However, this is a poor strategy because it guarantees that some of the most frequently accessed blocks on the disk are at the very edge. There are several ways to improve the situation.

An MFD may be pre-extended anywhere. That is, clusters two through seven may be placed on any cluster boundary, preferably near the center. Not only does DOPTER place the other clusters near the center, but it only uses the first block of the first cluster. Thus, instead of building the entire MFD at the edge of the disk, it only uses one block there. All the others are where they belong. Once again, the fundamental consideration shows us that we have lost very little by removing the contiguity of the MFD because con-

secutive accesses are seldom made to an MFD. (Note: all blocks of the MFD are available if they are needed, but that is seldom necessary.)

E) Additional MFD/UFD optimizations

Many files have associated attributes which are unused. The prime example is task images created by the task builder. The task builder uses RMS I/O to create its task images, but the attributes are never used. EDT Version 1 put attributes on its output files, but in the case of most source files, they were unneeded by the compilers or editors. The unneeded attributes take up space in the UFDs and create extra overhead for FIP. DOPTER effectively eliminates supernumerary attributes by recognizing source file and task image extensions. The same can be done with PIP.

Another UFD optimization can be made with files with the contiguous attribute. Once such a file is opened, RSTS/E knows the mapping for the entire file from having read the first retrieval blockette. UFD accesses can be minimized by putting the rest of the retrieval blockettes for large contiguous files at the very end of the UFD, away from the useful information in the UFD. In this manner, they will never be accessed, not even when searching for free space. This strategy assumes that there is sufficient free space in the UFD, which is almost never a problem.

OTHER TOPICS

A) DLA versus DLW

When initializing a RSTS/E volume, RSTS/E can be made to keep either the date of last access (DLA) or the date of last write (DLW). In addition to the arbitrary needs of the installation, there are several trade-offs when considering one against the other.

The date of last access/write is stored in the accounting blockette. If the DLW option is chosen, this blockette only needs to be rewritten when a file is modified. In the case of read-only files, of which executable programs are a major example, a physical UFD access can be saved by using DLW instead of DLA. Note that all writes to UFDs generate physical accesses — they are not cached.

On the other hand, if REORDR is used to sort the files in the UFD in reverse order by access date, a better order is likely to occur if DLA is used instead of DLW. If DLW is used, some frequently accessed programs/files can have very old last access dates and thus be put at the end of the UFD. With DLA, they would be put near the front.

On balance, DLW is better with all small directories because only a block or two will ever be read to search the name blockette list for any file. The savings from not writing the last access date on read-only files will more than compensate for the slightly longer lookup path after running REORDR. DLA is better only on large accounts where heavily used read-only programs/files will be placed at the front of the UFD by REORDR.

B) New Files First

It is easy to overrate the advantages of using New Files First (NFF). The arguments for using NFF are as follows.

1). It is convenient to have the most recently created files at the front of the directory listing. (This has nothing to do with performance.)

. . . continued on page 30

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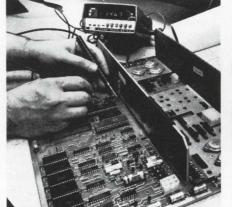
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H-1580

THE RSTS PRO VISITS THE NCC

By Carl B. Marbach

Los Angeles is busy. Even at night (or early morning) the airport is busy and choked with traffic. There are no rental cars and the taxi ride to Anaheim requires the national debt to be raised substantially. From the East coast it takes twelve hours of travel to make the trip and when you finally roll into bed in the plastic motel it seems only marginally worthwhile. Morning comes early in California and arriving at the Anaheim Convention Center it looks like a large conference with the routine exhibitors, coffee in styrofoam cups, sore feet and lots of exciting equipment. The press rooms are small, but there are two of them. After the obligatory cup of coffee in the first, I wander into the seccond. There are tables of typewriters for the reporters and . . . five tables of DECmate II's! Ten in all, complete with two letter quality printers and two DECpersons to help all the unfortunate ones who do not know how to use this terrific machine. So, the National Computer Conference is off to (for me) a great start; I feel at home.

The DECmate II might be the most hidden secret DEC has. It is a really fine machine. With its PDP-8 based word processing software (among the best in the industry), it also has CP/M capability which means it can run a large number of microprocessor based software packages. These packages include professional time and billing, spreadsheet planning, many financial planning languages, BASIC compilers and many, many others. The price is getting more competitive all the time and if you haven't tried one yet just find your nearest DIGITAL computer store and take a good look. Note that the DECmates are being sold by the DEC computer stores; they also will be selling MICRO-11s and Vax 11/730s. They will NOT be selling Rainbow 100s and Professional 300's for the time being, so that while supplies are short the DEC stores will not be competing with the independent DEC dealers. But, when they are finally making more than they can sell . . .

The National Computer Conference is a massive event, taking up the whole Anaheim Convention Center, the Disneyland Hotel and several large (and hot) tents in the parking lot. It used to be known as the Fall Joint Computer Conference (The Fall Joint for short, but that means something else now) and the Spring Joint Computer Conference, each in its own season. The show was combined into one event when it became too expensive for exhibitors to display at both shows, and when travel to two shows was too much for people. The price we pay for one show is the crowd. Exhibitors love it; attendees put up with it. When you attend, be prepared.

For most of us it is a Howard Johnsons for computers only there are more than 27 flavors; there are hundreds. All the majors are there and it is very informative to look and learn about the computer systems we don't use. When you run out of systems, you can start with the peripherals. All the old ones you knew about and many new ones you haven't thought about yet. A new letter quality printer for \$995 list (20 CPS), A VT100 emulating terminal, a page size (66 lines) CRT, 100 MB disks in 5½ inch size, 400 + MB disks in slightly larger configurations, streaming tape drives, cartridge tape drives, laser printers, supplies, desks, and, of course, software.

It is an impressive show. I mentioned that in general and was told that "without the micros, it wouldn't be nearly as exciting." Right. VisiCorp (the VisiCalc people) introduced VisiWord and VisiSpell on the xxx personal computer (it will be on the Rainbow 100 in the future). While the word processing isn't that impressive, the Speller is dynamite. DECword, WORD-11 and DECmate people take note: it is much better. It will even make a very good guess at what you were TRYING to spell.

DEC was not outdone by anyone. The DECbooth was filled with new and old DEC products. The Personal computer line was well represented, as was the MICRO-11 running all PDP-11 operating systems (even RSTS) and the VAX 11/730 with its see through panels showing the inside of this well packaged VAX.

The personal computers were well represented with their color graphics really making the booth look good. The telephone management system was alive and well on the Pro 350 and will be an impressive product; this is due out in the fall and will be the subject of a complete article in Personal and Professional (THE magazine for DEC personal computers). DEC also introduced IVIS, Integrated Video Information System which is a Pro 350 connected to a laser disc video recorder. Between the color graphics, the TV picture and the Pro 350 controlling the video disk it is an amazing array of gadgetry that will tie together two pieces of electronic equipment which DEC hopes will make a training tool par excellence. Can you imagine having the computer control the video disk bringing up random sections of a training manual depending on the user's response to the prior video section?

VAX users were not left out. Enter the new VAXstation 100. This gigantic (18"??) display is connected to a MICRO box that houses a Motorola 68000 microprocessor. It emulates up to four VT-100 "windows," has bit-mapped graphics with a resolution of 1000 x 800, and can put a graph in three quarters of the picture and still give you complete VT-100 functionality in the remaining quarter. Oh yes, it has an optional mouse. Want to see a VAX that looks like a LISA?

The NCC is a dizzying display of equipment and services; I have not even touched on the academic sessions which appear to be of high quality. There is more to do than you can in one week. Whatever time you can spend is not only informative but also good for the soul. We all need to be rejuvenated once in a while. This could be the place.

COMPILE.BAS: COMPILATION AID FOR BASIC PLUS TWO

By Alexander Ehrlich, Arel Software Ltd., Yavneh, Israel

1. INTRODUCTION

While working on a large data processing project, I found that a lot of my nice BASIC PLUS programs were growing into not so nice BASIC PLUS 2 programs. Gradually my account filled up with .ODL, .OBJ and .TKB files. Also, whenever I ran a compilation, my terminal was tied up for at least several minutes. COMPIL.BAS was written to solve these and other problems.

COMPIL.BAS helps the BASIC PLUS 2 programmer by doing the following:

- 1) Generates the auxiliary files that are necessary for a compilation.
- After the compilation, deletes any auxiliary files that are no longer needed.
- Performs all the commands needed to run a compilation.
- Runs the compilation in "background" so as not to tie up the user's terminal. See DETACH option below.

2. USING COMPIL.BAS

The programmer uses COMPIL by typing the following: RUN COMPIL

(The program responds with "PROGRAM NAME?") program name / option list

The extend part of the program name will be assumed to be .B2S unless explicitly stated.

The following can be in the option list, separated by commas:

- CHAIN The program will be compiled with the chain option
 - 2) NOCHAIN The program will be compiled with the no chain option
 - 3) LINE The program will be compiled with the line option
 - 4) NOLINE The program will be compiled with the no line option
 - 5) DEBUG The program will be compiled with the debug option
 - 6) XFILE The program can use file channels 13 and 14
 - 7) DETACH The compilation will run in "background" on a pseudo keyboard. This will enable the programmer to log on again and continue working (or at least play PAC-MAN) until his or her compilation is finished. The programmer will receive a message when the compilation is finished and any error messages will be stored in the .ERR file.

If DETACH is not specified, then the compiler commands will be forced into the job stream automatically.

3. NOTES TO THE SYSTEM MANAGER

Although COMPIL is an aid for BASIC PLUS 2 programming, it is written in BASIC PLUS. Place the .BAC version of COMPIL into a convenient account (such as $\{1,2\}$) with a protection code of < 232 >. COMPIL uses privileged commands.

The default option specified at line number 65 is CHAIN. Substitute any options that you prefer, separated by commas.

Examine carefully lines 210 through 250. These lines generate the command and overlay files for our installation and might have to be modified.

0	EXTEND		
0 !	PROGRAM	- COMPIL.BAS<232	>
	DESCRIPTION	- BASIC PLUS 2 C	OMPILATION AID
	PROGRAMMER	- Alexander Ehrl Arel Control S Yavneh, Israel	ystems Ltd.
	DATE	- 1-October-82	
0 !			
		MAIN PROGRAM	
0	CR\$ = CHR\$(13		
	OPTION.LIST\$(ST\$(7%), OPTIONS%(7) 1%) = "DETACH"	\$), SYS1\$(30\$) !OPTION LIST
		2%) = "CHAIN" 3%) = "NOCHAIN"	
	OPTION.LIST\$(
	OPTION.LIST\$(6%) = "DEBUG"	
	DUMMY\$ = SYS(7%) = "XFILE" CHR\$(6%) + CHR\$(9%)	
	KEYBOARD\$ = A	SCII(MID(DUMMY\$, 2%	, 1%)) * .5 ! OBTAIN KEYBOARD NUMBER
5	DEFAULT.OPTIO	N\$ = "CHAIN"	
7	GOSUB 100		! OBTAIN PROGRAM NAME ! OPTIONS
0	GOSUB 200		I CREATE COMMAND AND OVERLAY I FILES
0 1		NON DETACL DUN	COMPILATION ON USER'S KEYBOARD
			COMPILATION ON USER'S REIBOARD
	IF NOT OPTION THEN GOSUB		CREATE INSTRUCTION TEXT
	GOSUB	400	FORCE TEXT TO KEYBOARD
	GOTO	99	
0 1	DETAC	H - RUN COMPILATION	IN BACKGROUND
	PRINT "		
		CHR\$(6%) + CHR\$(7%))	
	GOSUB 500		OBTAIN PASSWORD
	GOSUB 600		I RUN PSEUDO JOB
	GOSUB 700		I SEND MESSAGE TO USER'S TERMINAL
	DUMMY\$ = SYS()	CHR\$(6%) + CHR\$(5%))	I LOGOUT
9	GOSUB 32000		
1 00			G100
	OBTAI	N PROGRAM NAME AND C	PTIONS

105 1	
1	READ PROGRAM NAME AND OPTIONS
1	PRINT PRINT "PROGRAM NAME " ;
\ \	INPUT LINE PROGRAMA\$ PROGRAMA\$ = CVT\$\$(PROGRAMA\$, 132\$)
1	GOTO 32000 IF PROGRAMA\$="" OPTIONS%(I%) = 0% FOR I% = 1% TO 7%
1	IF INSTR(1%, PROGRAMA\$, "/")=0% THEN PROGRAM\$ = PROGRAMA\$ OPTIONS\$ = DEFAULT.OPTION\$+"," ! DEFAULT = "CHAIN"
110	GOTO 120 PROGRAM\$ = LEFT(PROGRAMA\$, INSTR(1\$, PROGRAMA\$, " ")-1\$)
110	OPTIONS\$ = RIGHT(PROGRAMA\$, INSTR(1\$, PROGRAMA\$, "/")+1\$)+","
120 ! !	CHECK OPTIONS
1	COMMA.LOC\$ = INSTR(1\$, OPTIONS\$, ",")
~	GOTO 140 IF COMMA.LOC%=0% OPTION\$ = LEFT(OPTIONS\$, COMMA.LOC%-1%)
`	GOTO 130 IF LEFT(OPTION.LIST\$(1\$), 3\$)=LEFT(OPTION\$,3\$) FOR 1\$ = 1\$ TO 7\$ PRINT "ILLEGAL OPTION: "; OPTION\$
1	GOTO 105
130 \ \	OPTIONS\$(I\$) = -1\$ OPTIONS\$ = RIGHT(OPTIONS\$, COMMA.LOC\$+1\$) GOTO 120
140	IF (OPTIONS\$(2%) AND OPTIONS\$(3%)) OR
0	(OPTIONS#(4#) AND OPTIONS#(5#)) THEN PRINT "BOTH NO AND YES ON SAME OPTION"
145	GOTO 105 OPTION\$ = ""
145	OPTION\$ = OPTION\$ + OPTION.LIST\$(I\$) + "," IF OPTIONS\$(I\$) FOR I\$ = 2\$ TO 6\$
1	<pre>OPTION\$ = "/" + LEFT(OPTION\$, LEN(OPTION\$)-1\$) IF OPTION\$<>""</pre>
150 !	
1	CHECK THAT PROGRAM SOURCE FILE EXISTS
	PROGRAM\$ = PROGRAM\$ + ".B2S" IF INSTR(1\$, PROGRAM\$, ".")=0\$! DEFAULT EXTEND IS .B2S
160	ON ERROR GOTO 170 OPEN PROGRAM\$ FOR INPUT AS FILE 1≸
1	CLOSE 1% ON ERROR GOTO 0
~	<pre>PGM\$ = LEFT(PROGRAM\$, INSTR(1\$, PROGRAM\$, ".")-1\$) GOTO 190</pre>
	RESUME 180
1	PRINT PRINT "UNABLE TO FIND THE PROGRAM FILE -" ; PROGRAM\$
,	ON ERROR GOTO 0 GOTO 105
190	RETURN
!	
200 I 1	G200
200 I I I I	
200 I I I	G200
200 I I I I I I 210	G200 CREATE COMMAND AND OVERLAY FILES OPEN FGM\$+".ODL" AS FILE 1≸ I CREATE OVERLAY FILE
200 I I I I I 210	G200 CREATE COMMAND AND OVERLAY FILES OPEN FGM\$+".ODL" AS FILE 1\$ I CREATE OVERLAY FILE PRINT #1\$, CHR\$(9\$) ; ".ROOT USER"
200 I I I I I 210	G200 CREATE COMMAND AND OVERLAY FILES OPEN FGM\$+".ODL" AS FILE 1\$ I CREATE OVERLAY FILE PRINT #1\$, CHR\$(9\$); ".ROOT USER" PRINT #1\$, "USER:"; CHR\$(9\$); "FCTR SY:"; PGM\$; "-LIBR" PRINT #1\$, CHR\$(9\$); ".FCTR LB:BP2COM/LB" PRINT #15, CHR\$(9\$); ".ROT"
200 I I I I I I I I I I I I I I I I I I I	G200 CREATE COMMAND AND OVERLAY FILES OPEN PGM\$+".ODL" AS FILE 1\$! CREATE OVERLAY FILE PRINT #1\$, "USER:"; CHR\$(9\$); ".ROOT USER" PRINT #1\$, "USER:"; CHR\$(9\$); ".FCTR SY:"; PGM\$; "-LIBR" PRINT #1\$, "LIBR"; CHR\$(9\$); ".FCTR LB:BP2COM/LB" PRINT #1\$, CHR\$(9\$); ".END" CLOSE 1\$
200 I I I I I 210 \ \ 220 \	G200 CREATE COMMAND AND OVERLAY FILES OPEN FGM\$+".ODL" AS FILE 1\$ I CREATE OVERLAY FILE PRINT #15, CHR\$(95) ; ".ROOT USER" FTRT SY:" ; PGM\$; "-LIBR" PRINT #15, "UIBR:" ; CHR\$(95) ; "FCTR SY:" ; PGM\$; "-LIBR" PRINT #15, CHR\$(95) ; ".END" CLOSE 1\$ I CREATE COMMAND FILE PRINT #15, FOM\$; "=" ; FOM\$; ".ODL/MP" IF NOT OPION\$5(75)
200 I I I I I 210 \ \ 220 \ \	G200 CREATE COMMAND AND OVERLAY FILES OPEN PGM\$+".ODL" AS FILE 1\$! CREATE OVERLAY FILE PRINT #1\$, CHR\$(9\$); ".ROOT USER" PRINT #1\$, "USER:"; CHR\$(9\$); "FCTR SY:"; PGM\$; "-LIBR" PRINT #1\$, "LIBR"; CHR\$(9\$); ".FCTR LB:BP2COM/LB" PRINT #1\$, CHR\$(9\$); ".END" CLOSE 1\$ OPEN FGM\$+".CMD" AS FILE 1\$! CREATE COMMAND FILE PRINT #1\$, FOM\$; "="; PGM\$; ".ODL/MP" IF NOT OPTIONS3(7\$) THEN PRINT #1\$, "UNITS=12" PRINT #1\$, "ASG=SY:5:6:7:8:9:10:11:12"
200 I I I I I 210 \ \ 220 \ \ \	G200 CREATE COMMAND AND OVERLAY FILES OPEN PGM\$+".ODL" AS FILE 1\$ I CREATE OVERLAY FILE PRINT #1\$, CHR\$(9\$); ".ROOT USER" PRINT #1\$, "USER:"; CHR\$(9\$); ".FCTR SY:"; PGM\$; "-LIBR" PRINT #1\$, "LIBR"; CHR\$(9\$); ".FCTR LB:BP2COM/LB" PRINT #1\$, CHR\$(9\$); ".END" CLOSE 1\$ OPEN FGM\$+".CMD" AS FILE 1\$ I CREATE COMMAND FILE PRINT #1\$, FOM\$; "="; FOM\$; ".ODL/MP" IF NOT OPTIONS\$(7\$) THEN PRINT #1\$, "UNITS=12" PRINT #1\$, "ASG=SY:5:6:7:8:9:10:11:12" GOTO 250
200 I I I I I 210 \ 220 \ 220 \ 220 \ 230 \	CREATE COMMAND AND OVERLAY FILES OPEN FGM\$+".ODL" AS FILE 1\$ I CREATE OVERLAY FILE PRINT #15, CHR4(95); ".ROOT USER" PRINT #15, "USER:"; CHR4(95); ".FCTR SY:"; PGM\$; "-LIBR" PRINT #15, "LIBR:"; CHR4(95); ".FCTR SY:"; PGM\$; "-LIBR" PRINT #15, CHR4(95); ".END" CLOSE 15 OPEN FGM\$+".CMD" AS FILE 1\$ I CREATE COMMAND FILE PRINT #15, PGM\$; "="; PGM\$; ".ODL/MP" IF NOT OPTIONSX(75) THEN FOINT #15, "ASG=SY:5:6:7:8:9:10:11:12" GOTO 250 PRINT #15, "UNITS=14" I OPTION=XFILE PRINT #15, "ASG=SY:5:6:7:8:9"
200 I I I I 210 \ 220 \ 220 \ 220 \ 230 \ \	CREATE COMMAND AND OVERLAY FILES OPEN FGM\$*".ODL" AS FILE 1\$ I CREATE OVERLAY FILE PRINT #15, CHR4(95); ".ROOT USER" PRINT #15, "USER"; CHR4(95); ".FCTR LB:BP2COM/LB" PRINT #15, "LIBR:"; CHR4(95); ".FCTR LB:BP2COM/LB" PRINT #15, "LIBR:"; CHR4(95); ".CND" CLOSE 15 OPEN FGM\$*".CMD" AS FILE 1\$ I CREATE COMMAND FILE PRINT #15, "CMM\$; "="; PGM\$; ".ODL/MP" IF NOT OPIIONS(75) THEN PRINT #15, "ASG=SY:5:6:7:8:9" PRINT #15, "ASG=SY:5:6:7:8:9" PRINT #15, "ASG=SY:5:6:7:8:9" PRINT #15, "ASG=SY:10:11:12:13:14"
200 I I I I 210 \ 220 \ 220 \ 220 \ 230 \ \	CREATE COMMAND AND OVERLAY FILES OPEN FGM\$+".ODL" AS FILE 1\$ I CREATE OVERLAY FILE PRINT #15, CHR4(95); ".ROOT USER" PRINT #15, "USER:"; CHR4(95); ".FCTR SY:"; PGM\$; "-LIBR" PRINT #15, "LIBR:"; CHR4(95); ".FCTR SY:"; PGM\$; "-LIBR" PRINT #15, CHR4(95); ".END" CLOSE 15 OPEN FGM\$+".CMD" AS FILE 1\$ I CREATE COMMAND FILE PRINT #15, PGM\$; "="; PGM\$; ".ODL/MP" IF NOT OPTIONSX(75) THEN FOINT #15, "ASG=SY:5:6:7:8:9:10:11:12" GOTO 250 PRINT #15, "UNITS=14" I OPTION=XFILE PRINT #15, "ASG=SY:5:6:7:8:9"
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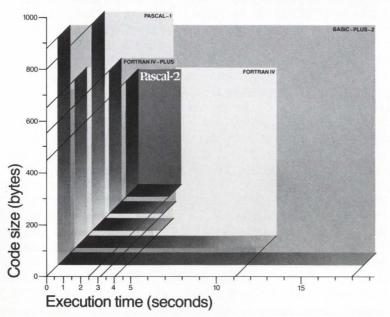
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00 !	G400
	FORCE TEXT TO KEYBOARD
10	T.PNT% = 1%
20	<pre>DUMMY\$ = SYS(CHR\$(6\$) + CHR\$(-4\$) + CHR\$(KEYBOARD\$) + RIGHT(TEXT\$, T.PNT\$))</pre>
	IF RECOUNT<>O IADDITIONAL TEXT THEN T.PNT\$ = (LEN(TEXT\$)-RECOUNT) + 1\$ GOTO 420
90	RETURN
00 !	G500
	OBTAIN PASSWORD
0	DUMMY\$ = SYS(CHR\$(6\$) + CHR\$(14\$) + STRING\$(20\$,0\$))
	CHANGE DUMMY\$ TO SYS1% PROG% = SYS1%(7%)
	PROJ\$ = SYS1\$(8\$) PASS\$ = RAD\$(SYS1\$(9\$) + SWAP\$(SYS1\$(10\$)))
	+ RAD\$(SYS1%(11%) + SWAP%(SYS1%(12%)))
0	RETURN
0 1	G6 00
	RUN PSEUDO JOB
)	ON ERROR GOTO 620
	PK.NUM# = 1%
5 1	FIND AVAILABLE PSEUDO KEYBOARD
	FIND AVAILABLE PSEUDO KEYBOARD OPEN "PK"+NUM1\$(PK.NUM\$)+":" AS FILE 2\$, MODE 0\$
	GOTO 640
	IF PK.NUM%=10% THEN PRINT "NO AVAILABLE PSEUDO KEYBOARDS FOR THE JOB"
	ERROR.IND\$=1\$ GOTO 690
D	PK.NUM% = PK.NUM% + 1%
	RESUME 615
0	FIELD #2\$, 128\$ AS MSG\$
D !	SEND COMMANDS TO BACKGROUND JOB
	<pre>DUMMY\$ = FNSEND\$("HELLO "+NUM1\$(PROJ\$)+","+NUM1\$(PROG\$)+ ";"+PASS\$, "Ready")</pre>
	DUMMY\$ = FNSEND\$("BP2", "BASIC2")
	DUMMY\$ = FNSEND\$("OLD "+PROGRAM\$, "BASIC2")
	DUMMY\$ = FNSEND\$("COM "+PGM\$+".OBJ"+" "+OPTION\$, "BASIC2") IF NOT ERROR.IND\$
	<pre>DUMMY\$ = FNSEND\$("TKB @"+FGM\$+".CMD", "Ready") IF NOT ERROR.IND\$</pre>
	DUMMY% = FNSEND%("EXIT", "Ready") IF ERROR.IND%
	DUMMY# = FNSEND#("UNSAVE "+PGM\$+".CMD", "Ready")
	<pre>DUMMY\$ = FNSEND\$("UNSAVE "+PGM\$+".ODL", "Ready")</pre>
	DUMMY# = FNSEND#("UNSAVE "+FGM\$+".OBJ", "Ready") IF NOT ERROR.IND#
)	CLOSE 2%
	RETURN
1	G700
	CENT MECCACE TO HOPPIC TERMINA
	SEND MESSAGE TO USER'S TERMINAL
	IF ERROR.IND\$=-1\$
D	THEN MESSAGE\$ = "PROGRAM: "+PROGRAM\$+" COMPILATION ERRORS "
)	+"-CHECK "+PGM\$+".ERR" ELSE IF ERROR.IND\$=0\$
	+ "-CHECK "+ PGM\$+". ERR "

720 790	DUMMY\$ = SYS(CHR\$(6\$) + CHR\$(-4\$) + CHR\$(KEYBOARD\$) + MESSAGE\$+CHR\$(13\$)+CHR\$(10\$)) RETURN	& & &	OR INSTR(1\$, OUT1\$, "?")<>0\$ & & OR INSTR(1\$, OUT1\$, "\$")<>0\$ & & ! SET ERROR INDICATOR IF & ! STRING CONTAINS AN ERROR & ! MESSAGE & &
		å	THEN OPEN PGM\$+".ERR" AS FILE 3%
1000 !	FNSEND	& &	\ PRINT #3\$, OUT1\$ & \ ERROR.IND\$ = -1\$ & \ CLOSE 3\$ & &
1	SEND AND RECEIVE STRINGS FORM PSEUDO TERMINAL	& &	1090 FNEND
1	SEND.STRING\$ - STRING TO BE SENT	α &	! &
1 1 1	END.STRING\$ - END OF STRING TO BE RECEIVED	& & & &	å 32000 IF NOT OPTIONS≸(1≸) ! NOT DETACH å THEN PRINT *END OF COMPIL" å
1000 \ \ \ \	DEF FNSEND\$(SEND.STRING\$, END.STRING\$) ERR.COUNTER\$ = 0\$ OUT1\$ = "" LST MSG\$ = SEND.STRING\$+CHR\$(13\$) IF INSTR(1\$, SEND.STRING\$, "HELLO")<>0\$ THEN REC.NUMBER\$ = 1\$ ELSE REC.NUMBER\$ = 10\$	& & & & & & & &	32767 END
1010 I 1 1	SEND STRING TO PSEUDO KEYBOARD	& & &	
	ON ERROR GOTO 1030	u	
1020	PUT #2%, RECORD REC.NUMBER%, COUNT LEN(MSG\$) GOTO 1050	& &	
1030 \ \ \	SLEEP 2 ERR.COUNTER≸ = ERR.COUNTER≸ + 1≸ RESUME 1020 IF ERR.COUNTER≸<10≸ RESUME 1050 I SMALL BUFFERS ARE FULL	& & & &	
1050 I I	RECEIVE STRING FROM PSEUDO KEYBOARD	& & &	
'	ON ERROR GOTO 1070	4	
1060	INPUT LINE #2#, OUT\$	ě	
1062 \	OUT3\$ = LEFT(OUT\$, RECOUNT) OUT1\$ = OUT1\$ + OUT3\$	& &	
1064	GOTO 1080 IF INSTR(1%, OUT1\$, END.STRING\$)<>0% ! CHECK FOR END OF STRING	& &	
1	GOTO 1060 RESDME 1060	& &	HANSEN
1070	RESUME 1060	å	
1080	IF INSTR(1\$, OUT1\$, "Error")<>0\$ OR INSTR(1\$, OUT1\$, "error")<>0\$	& &	"You haven't been the same since the boss said you were a programming star."
			이 같은 것 같은

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CIRCLE 57 ON READER CARD

THE HISTORY OF RSTS

By Peter Dick, Silver Programs, London

The history of RSTS starts in the second World War. There was a requirement to be able to train pilots without having them actually go up in aeroplanes. Project Whirlwind was the name given to the first electronic computer that was able to operate in real time; this meant it could be programmed to behave like an early flight simulator. The connection between Project Whirlwind and RSTS is that one of the people on that project was Ken Olsen, and the fact that the Whirlwind computer was a 16-bit computer.

Ken Olsen stayed with MIT until 1957 when he left and formed Digital Equipment Corporation. Folklore tells us that Ken Olsen started with three other people, one of whom was his brother, using their garage to assemble boards. The truth of course is slightly more boring, and that from the first day of operations they had the backing of a quarter of a million dollar loan.

The big breakthrough for computing was in 1961 when President Kennedy decided that the American people would land a man on the moon before the end of the decade. This decision gave computing in general an enormous source of funds; i.e., NASA.

In 1963 the PDP 6 was announced and on this machine it was possible to do timesharing. In 1964 BASIC was invented by Dartmouth College, New Hampshire — BASIC, of course, standing for Beginners All-purpose Symbolic Instruction Code. 1968 saw TSS/8 being released by Digital which was timesharing on the PDP 8 that in fact* gave birth to RSTS. 1970 saw the release of the hardware, the PDP 11, which has become the most successful minicomputer in the history of computing.

RSTS 11 Version 1, 1970. It was never released and it was only tried inhouse.

RSTS Version 2A, 1971. Again never released only tried in-house.

RSTS Version 2B, June 1971. This is the date from which Digital claims that RSTS was introduced. There are no known users.

RSTS Version 3A, February 1972.

It was supported on the 11/20 which was the only PDP 11 available, and was Pre-SYSGENed.

It obviously worked on an unmapped machine, job max 8K words, swap max allowed for eight users. It was interesting to note that the 20K words were both BASIC and RSTS together.

RSTS Version 3B, May 1972. There are no known users.

RSTS Version 3C, June 1972. At long last we have the first recorded user, Bob Branton, at Southeastern State College. This user was in fact responsible for documenting the first complaint about RSTS, that under Version 3C the monitor was unable to write to a disk file from a terminal while reading that file from another terminal. Very simply, this was because there was no such thing in those days as UPDATE mode.

RSTS Version 4A, October 1972. Support was increased to include the newly announced 11/40 and 11/45. SYSGEN was under DOS and it was interesting to note that the CUSPS were built from paper tape at 10 cps. There were only three manuals in those days, the BASIC PLUS Language Manual, RSTS Users Manual, and the System Managers Manual. RECORD I/O was optional, and in fact there was a separate sub option for UPDATE mode which the book told us took an extra 75 words. There was talk of BADS for future releases. Hardware was fixed disks or RKO5's. There was no memory management. The garbage collector was FIP overlayed and it wrote out the program and had to read it back in again. PI was not write protected; i.e., if you were tired of it being 3.142 you could make it 3.1,3,16,189. Maximum program size was 8K words but this could be PATCHed to 9K if you knew Rosemary Phillimore. Swap max was increased from the original eight users to 16. The program SHUTUP achieved the system closing down by PEEKing at an odd address which caused the machine to crash.

RSTS version 4B, contained no new code and was in fact just a fully



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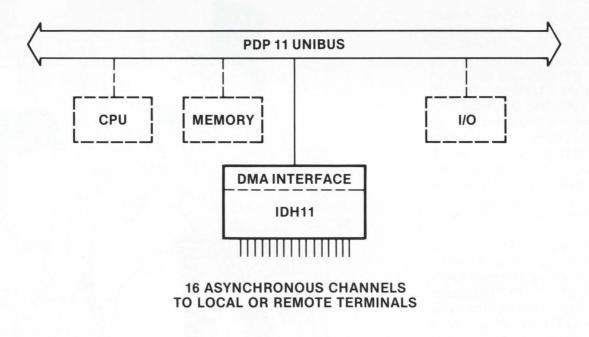
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PATCHed version of 4A. It was released in 1975 and marks the end of RSTS 11.

RSTS/E Version 5A/21, July 1973. It no longer supported the 11/20. SYSGEN was under DOS but at least you had batch command files to help. Job max was still 8K words for users but 16K for privileged users. Swap max had again been increased to allow for 32 users. The maximum machine size was 128K words, of which 124 was usable. Commercial Extensions were available which included Sort, IAM, QUE, SPOOL, Decimal Arithmetic package. Hardware included the RPO3, which was the 40M byte disk drive, DH11s and 1200 LPM printers. APPEND as a command was introduced.

RSTS/E Version 5B/24, November 1974. It allowed for CCLs as long as you created them during SYSGEN time. EDIT and EDITCH were introduced. Auto answer was introduced for SYSGEN generations and so were multiple SWAP files. There was a hidden option for overlapped seek with multiple disks. Among the new BASIC PLUS commands were CVT\$\$, STRINGS\$, STATUS, BUFSIZ. It was an extremely untidy release and so four months later . . .

RSTS/E Version 5C-01, March 1975. This in fact was a PATCHed version of Version 5B and contained no new commands.

RSTS/E Version 6A-02, August 1975. Support was now for the 11/34, 11/35, 11/40, 11/45, 11/60 and 11/70. The swap max had been increased to 63 users and maximum memory on the 11/70 was now 2M bytes. The RPO2 and the RPO4 were added, as was the concept of multiple run time systems: i.e., COBOL, FORTRAN, SORT11 were all added. XBUF was introduced. Pseudo keyboards. STATS was a hidden option. The VT50 was the replacement for the original VT05 and TU16's were available.

RSTS/E Version 6B-02, February 1977. SYSGEN was now under RT11. New hardware included RK05Fs, RK06, RP05, RP06, DZ11, LA180. Contiguous files were allowed. CCLs could be added at run time. EXTEND mode was introduced, Echo control mode 8. Among the run time systems allowed were BASIC PLUS 2, FORTRAN, COBOL, RPG2, APL, RMS-11, RT11 and of course yet another new BACKUP package.

RSTS/E Version 6C-03, February 1978. TECO and VTEDIT were released as part of a DECUS supported package. DECNET was introduced, DATA-TRIEVE, DIBOL/DECFORM. PIP.SAV replaced the old PIP and PIP EXTEND. Sexy lights were available. Also the program STATUS from Martin Minnow, which is still not supported by Digital. Hidden options included the words GET BLOCK and SPEC %.

RSTS/E Version 7.0. December 1979. Supported CPUs now included the 11/24, 11/34, 11/35, 11/40, 11/44, 11/45, 11/60, 11/70. RDC was introduced. SYSGEN was now under RSTS during time sharing and data caching was introduced. The RSX monitor allowed 31K word programs. Maximum memory on an 11/70 was now 4M bytes. Resident libraries were introduced. XBUF for line printers. SAV/RES for backup. A new version of FIP called large file FIP was created, and by now they had of course stolen the switch panel so that sexy lights were no longer possible.

RSTS/E Version 7.1, Februrary 1982. This now uses separate I & D space to relieve the small buffer problem. DECNET route through is allowed. TRACE, BREAK and DUMP are all hidden options. RSTS/E Version 7.2, August 1982, supports new UDA50 disks; i.e., the RA80. The typical machine is now an 11/44, 1M byte of memory, RA80, TS11, DZ11E for under 50,000 pounds including a general licence.

RSTS/E Version 8.0, April 1983. Support for the recently announced PDP Micro 23. BASIC Version 2 now allows 31K words and no overhead for the use of RMS.

It is fair to say that the number of users has always been underestimated; Digital's own figues are as follows:

> December 1973 — 150 users August 1975 — 1,200 users July 1977 — 3,100 users December 1979 — 5,000 users

It is now thought that there are well over 10,000 licenced users and at least an equal number of unlicenced users!!!.

RSTS has a future, as Digital has yet to implement its greatest secret weapon, separate I & D space in terms of the application programs. It is interesting to note that RSX11M PLUS Version 2 allows the programmers to have separate 31K word areas for both I & D; in other words it should be possible under RSTS to have a total 62K word user area. RSTS remains the key commercial time sharing operating system that Digital sells.



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QUEUING TO A SPECIFIC DEVICE QUEDEV.B2S

by Terry Ridgers, DuPont Canada, Corunna, Ontario

QUEDEV is written and compiled in BASIC + 2. It can be called from any BASIC + 2 main line program. The main line file name .ODL must contain QUEDEV i.e., .FCTR MAIN-LIBR-* (QUEDEV-LIBR, etc.). QUEDEV when called and when passed valid arguments will

queue files to any of the LPn:, BAn: and RJn: devices.

- QUEDEV uses five passing arguments, namely:
- 1) SV.FIL\$ full file name as a string or as a string variable.
- 2) SV.DEV\$ Device name:
 - must be of the form "LP", "BA", or "RJ".

3) SV.UNIT% - Unit number

- i.e., if sent to LP1: SV.UNIT% must = 1%.
- SV.CPY% Number of copies to be sent to the line printer (valid only for the line printer)
- 5) SV.DEL% Delete list file after printing to LP:
 - 1% yes
 - 0% no
 - (valid only for line printer)

NOTE: For unit number, number of copies and delete, the values passed must be integer values. If real values are passed then incorrect number of copies etc., may occur.

Examples of possible calls from the Main Program are:

1) CALL QUEDEV ("SY: (1,91)TEST.LST", "LP",0%, 2%, 1%) Will queue 'SY:(1,91) TEST.LST" to LP: for two copies and will delete "SY:(1,91)TEST.LST" after listing the last copy.

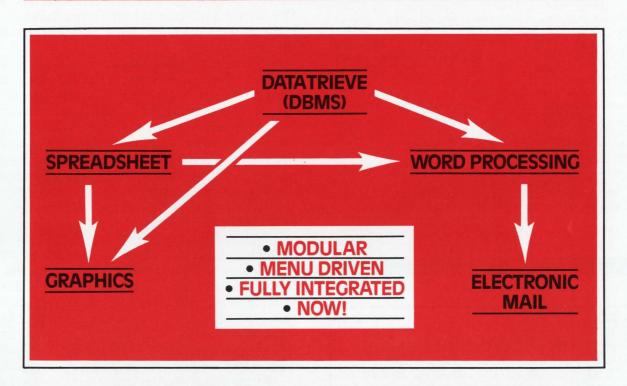
- 2) SV.CTL\$ = "DR1:(18,5)BATCH.CTL"
 - CALL QUEDEV(SV.CTL\$, "BA", 1%, 0%, 0%)

This call will queue "DR1: (18,5) BATCH.CTL" to BA1: and it ignores the number of copies and will not delete the file regardless of what is passed as the last two arguments.

1 :	SUB QUEDEV (SV.FIL\$,S	V.DEV\$,SV.UNIT\$,SV.CPY ! SUBPROGRA	(,SV.DEL%) QUEUES FILES TO VARIOUS DEVICE	S&
		1	SV.FIL\$ - FILE NAME TO QUEUF	
		1	SV.DEV\$ - DEVICE NAME	ő:
		1	IE: BA:,LP:,RJ:	&
		1	SV.UNIT# - DEVICE UNIT #	å
		1	SV.CPY# - NO OF COPIES	å
			APPLIES TO LP: ONLY	8
			SV.DEL% - DELETE FILE AFTER	å
			1 - YES	8
		i	0 - NO	&
		1	APPLIES TO LP:	å
		1	ONL Y	å
10				&
	1			&
	I SUB PROGRAM DESCRI	PTION AND MODIFICATION	LAYOUT	å
	!			8
	SUB PROGRAM:	QUEDEV		8
	1 50D PROGRAP.	COLDE.		R
	1			å
	! PROGRAMMER :	T. RIDGERS		å
	1	DUPONT		&
	1	ST. CLAIR RIV	ER WORDS	&
		CORUNNA, ONT. NON 1GO		& &
		NON TOO		8
	I SUB-PROGRAM DESCRI	PTION:		å
	!	TUTS SUB PROCEAM	CAN BE CALLED FROM ANY BASIC +	8
	2 PROGRAM TO QUEUE	FILES TO VARIOUS DEVI		&
	1 .			&
	! MODIFICATION LOG:			&
	!			α &
	! VERSION	DATE	DESCRIPTION	&
	1	al	ODTO TNAL DELEASE	8
	! 001	04-Aug-82	ORIGINAL RELEASE	& &
	999	DD-MMM-YY	******	8
	1		***********************	&
	1 .		XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX	&
	!			&
		*****************************	***************************************	&

1 MAIN SUB-PROGRAM 1000 - 9999 IF SED\$(SV.DEV\$,1\$,2\$) <> "LP" THEN ! ONLY LF: DEVICE CAN SPECIFY SV.DEL1\$ = 0\$! ONLY LF: DEVICE CAN SPECIFY COSUB 10000 ! QUEUE FILE SV.FIL\$ TO EVICE ELSE SV.DEL1\$ = 0\$! INITILIZE TO NOT DELETE COSUB 10000 ! QUEUE SV.FIL\$ TO DEVICE BUT MEXT X\$! INITILIZE TO NOT DELETE COSUB 10000 ! QUEUE SV.FIL\$ TO DEVICE BUT NEXT X\$! SUBROUT COSUB 10000 ! QUEUE SV.FIL\$ TO DEVICE BUT NEXT X\$! SV.DEL1\$ = 1\$ SV.DEL1\$ = 1\$! SV.DEL1\$ IS ALSO 1\$ COSUB 10000 ! QUEUE SV.FIL\$ TO DEVICE BUT SUBEXIT ! SUBROUT IF SUBEXIT ! SUBROUT IF SUBROUTINE - QUEUES FILE TO DEVICE SEG\$(FSS\$(SV.FIL\$,1\$),5\$,12\$) +! SYS CALL TO FIP SUBROUTINE - QUEUES FILE TO DEVICE SEG\$(FSS\$(SV.FIL\$,1\$),5\$,12\$) +! FILE NAME STRING SCAN TO GET SEG\$(FSS\$(SV.FIL\$,1\$),5\$,12\$) +! IN RADIX SO FORMAT SUBROUTINE - QUEUES FILE TO DEVICE CHH\$(0\$) +! AS UNIT SEG\$(FSS\$(SV.FIL\$,1\$),23\$,26\$) +! INT AD AD OF ORMAT	1000!******						\$
IF SED & (SV. DEV. 1, 1, 2, 2) (> "LP" THEN SV. DEL. 1 * O' SV. DEL. 1 * O' SV. DEL. 1 * O' DELETE SV. DEL. 1 * O' LESE SV. DEL. 1 * O' REXT X5 (GOUD 10000 (REXT X5 (GOUD 10000) NEXT X5 (GOUD 10000) NOT 10001) NEXT X5 (GOUD 10000) NOT 10001) NOT USED (HEX 0000) NOT 10001) NOT USED (HEX 0000) NOT 10000) NOT USED (HEX 0000) NOT USED (HEX		JB-PROGRA	м	1000 -	- 9999	9	& &
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<pre> NEXT XS SV.DEL15 = 15 SV.DEL75 = 15 SV.DEV5 SV</pre>	\ FOR]	K% = 1% T	O SV.CPY#	- 1%	!	! NON-DELETE COPIES TO LP: &	
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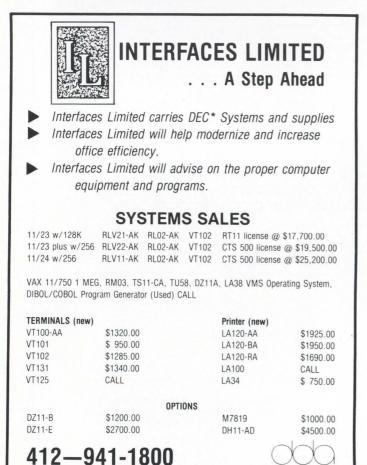
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CIRCLE 9 ON READER CARD



CIRCLE 174 ON READER CARD

RSTS/E DISK OPTIMIZATION ...

... continued from page 18

2). Recently created files are frequently accessed, so putting them at the front of the list decreases open time.

Closer inspection shows these arguments against using NFF.

1). In accounts with more than 31 files, it requires an extra UFD access (a physical write access), to create a new file. All of the name blockettes must be read to make sure that the file does not already exist. Then the pointers in both the first and the last name blockette must be rewritten. If NFF is not used, only the last name blockette must be rewritten. Note too, that these are physical accesses to disk that are not cached.

2). Recently created files are very likely to be deleted, resulting in two side effects. First, two name blockettes must be rewritten when the file is deleted, and they are likely to be in two separate UFD blocks. Without NFF, it is likely that they will be in the same UFD block. Second, in many environments such as word processing and development systems, most files are created, opened only once for reading, and then deleted. The main advantage of NFF, quick file opens, is minimized.

BENCHMARKS

Typical RSTS/E systems are difficult to objectively measure. System load varies from minute to minute and hour to hour. Monitor statistics are easy to gather but difficult to interpret (e.g., what does it mean if directory accesses decrease: less activity, a slower system, or a more efficient UFD structure?). One way to eliminate uncontrollable variables is to run a single job on an otherwise unused system and measure wall clock time. Unfortunately, single-user benchmarks are poor indicators of multi-user performance.

At the System Performance House, we have developed a set of 12 programs which run simultaneously, simulating a multi-user environment. They are heavily disk I/O bound programs which perform a mix of file creations, opens, closes, lookups, logins, logouts, swaps, run-time system loads, and random file accesses. The programs run for a fixed length of time and measure the number of each type of operation they can perform during the allotted time. Although they overstate the true performance differences from changes in disk I/O efficiency, they are very sensitive to small performance changes.

These benchmark programs vividly demonstrate the performance gains generated by the optimizations described above. The benchmark programs showed 50 percent more throughput on this arrangement than they did on the identical system generated by another commercially available "disk structuring" program. This increase was due to both faster disk accesses and fewer directory accesses.

REFERENCES

1. Banks, Scott. "RSTS Disk Directories," RSTS Professional, Vol. 1, No. 1; Vol. 2, No. 1; Vol. 2, No. 3; and Vol. 2, No. 4.

2. Mayfield, Mike. "RSTS/E Disk Internal Structures," Proceedings of the Digital Equipment Computer Users Society, April, 1978.

WATCH and RINGME

By Maury Pepper and Greg Wenzel, STS, Inc., St. Louis, MO

We found that after submitting a task build to ATPK, programmers were frequently using SYSTAT to see if the job had terminated. The two programs included here were designed to increase productivity and ease the anxiety of those waiting for a job to finish.

Any job can be WATCH'd. When the job disappears from the job table (i.e., when it is killed or logged out), RINGME will send three bells (beep, beep, beep) to the terminal which originated the WATCH. We set up a CCL for calling WATCH

UT CCL WA-TCH = [1,2]WATCH.BAC;PRIV 0 The syntax is:

WATCH jn watch job number jn

or WATCH/jn watch job number jn

or WATCH program will prompt for job number WATCH will determine which terminal you are on and then spawn a job to run RINGME — passing the terminal number and job number to RINGME in core common. Both programs should be saved with a <232> protection code. RINGME will run "logged-out" and detached. It's left as an unnecessary exercise for the purists to have RINGME log itself into the originator's account.

1 EXTEND 10 !

END Program: WATCH.BAS Author: Greg Wenzel/Maury Pepper Date: 28-Jan-83

100 JOBNO\$ = CVT\$\$(RIGHT(SYS(CHR\$(7\$)),7\$),-1\$) 1 Get Job # From CC IF JOBNO\$ = GOSUB 1000 Request Job IF JOBNO\$ = " Exit If No Job # Set Error Trap GOTO 32767 & ON ERROR GOTO 32767 JOBNOS = VAL(JOBNOS) Validate Job # GOTO 32767 IF JOBNOS < 15 OR JOBNOS > 645 & 110 KB\$ = MID(SYS(CHR\$(6\$)+CHR\$(26\$)+CHR\$(0\$)+CHR\$(0\$)),4\$,1\$) 8 ! Determine Calling Keyboard 2 I\$ = SYS(CHR\$(6\$)+CHR\$(-10\$)+'SY:[1,2]RINGME.BAC') 120 & ! Convert Program File Spec For Job Spawn & M\$ = SYS(CHR\$(6\$)+CHR\$(24\$)+CHR\$(0\$)+CHR\$(0\$)+MID(1\$,5\$,8\$)+ 130 8 CVT\$\$(JOBNO\$)+KB\$+' '+MID(1\$,23\$,4\$)) * ! Spawn "RINGME" 140 GOTO 32767 I Exit & INPUT "Watch Job # "; JOBNO\$ 1000 ! Request Job # \$ RETURN 2 32767 END EXTEND 10 Program: RINGME.BAS å Author: Greg Wenzel/Maury Pepper 1 Date: 28-Jan-83 2 100 J\$ = SYS(CHR\$(7\$))I Get Core Common & 110 IS = ASCII(MID(J\$,35,15)) 1 Determine KB # For Broadcast å 120 I\$ = SYS(CHR\$(6\$)+CHR\$(-13\$)+CHR\$(255\$)+CHR\$(-1\$)+CHR\$(-8\$)) Ł ON ERROR GOTO 140 I Lower Priority to -8 2 130 I\$ = SYS(CHR\$(6\$)+CHR\$(26\$)+CHR\$(CVT\$\$(J\$))+CHR\$(0\$)) GOTO 130 SLEEP ! Loop until SYS call fails 1 & ON ERROR GOTO 32767 140 M\$ = SYS(CHR\$(6\$)+CHR\$(-5\$)+CHR\$(I\$)+CHR\$(7\$)+CHR\$(7\$)+CHR\$(7\$)) & ! Broadcast 3 Bells to Terminal &

BAC mac can do it all!

32767

END

&

BAC into RTS / BAC into MAC / BAC into BAS



Western Distributor: Telecom Computer Systems, Inc. P.O. Box 03285 Portland, Oregon 97203 503/286-5122 BACmac is a unique software tool, running under RSTS/E, which provides the following conversions:

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translation from Basic-Plus into Macro source code, which compiled under RSTS runs faster than Basic-Plus

■ translation from Basic-Plus into Macro source code which may be compiled under RSTS for execution under RT11 — a migration facility

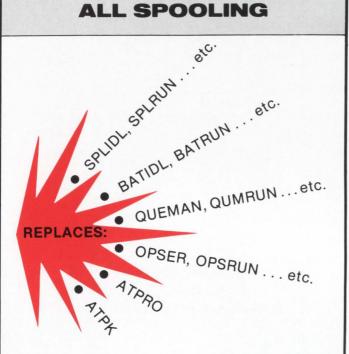
■ translation from Basic-Plus into a RUN-TIME-SYSTEM. Now you can write an RTS in Basic-Plus. The ideal solution to memory thrashing due to "multi-copy" applications programs.

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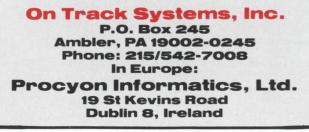
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CIRCLE 11 ON READER CARD

MEMO A COMPUTERIZED NOTE FILE

By Mark Gilmore, Data Processing Dept.

California State University at Long Beach, Long Beach, CA

This program is intended to be an aid to people who are constantly logged in and working on their RSTS systems, and who are also subject to frequent interruptions from bosses, users, or their own brilliant thoughts. MEMO will allow you to store your gems of wisdom (or your boss' demands) where they won't be forgotten (or can be "filed" indefinitely).

Memo numbers, once assigned, are never re-used (even though the associated memo may be deleted). This allows you to keep some sort of record based on memo number if you wish to make paper copies of everything. Memo numbers are assigned starting with number one and incrementing by one.

The program creates a file called MEMO.DAT on the disk specified by the variable SY.MEMO.FILE\$ in line 1110. Change this variable to specify the disk you want the memo file created on. Once created, the file can be moved to any other disk on the system. The program will check all the disks, looking for a memo file that it has read/write access to. This disk scan routine is the only reason that MEMO must be privileged, and privileges are dropped immediately after this scan is completed. The program will not currently handle systems with more than 20 disks. To allow more, change the dimension statement in line 1010.

The MEMO.DAT file is built with a linked list structure. Records are added to the file as needed, but they are never deleted. Instead, they are added to the list of free (available) records. This means that the file can grow to almost any size (more memos than you'll want to keep anyway) but that it can never shrink (until someone writes a program to compress the file — this is left as an exercise for the reader).

It is possible to remove privileges by placing disk names to be scanned in a DATA statement. These names should be placed in the DISKS\$ array and the variable DISKS% should be set to the number of disk names. Do this in line 1100 before the GOTO 15000, and delete lines 15010 — 15020.

If the disk scan feature is not wanted at all (the file specified in SY.MEMO.FILE\$ is the only possibility) delete lines 15010 - 15060 and make the following changes:

15010		MEMO.FILE\$ = SY.MEMO.FILE\$
	\	OPEN MEMO.FILE\$ FOR INPUT AS FILE #4%
	\	IF (STATUS AND 1024%) = 0% THEN 1110
		ELSE
		PRINT "Don't have write access to"; MEMO.FILE\$
	\	GOTO 32767
		! If we can't write on the file forget it.
32010		IF ERL = 15010 THEN RESUME 20050
		! Memo file didn't exist — create it.

The PROMPT command allows the user to change the input prompt. This allows a little "customizing" for each user, and could be used to store an important note that the user wants to be reminded of the next time the program is run. The prompt may also be set to print escape sequences to allow the use of some of the special functions available on video terminals (note that this may have undesirable effects if you find yourself on a different terminal type than you anticipated). This feature may be disabled by removing the loop in lines 6020-6030 which converts CHR\$(155)s in the prompt.

For video terminals, it may be advantageous to modify the PROMPT command to allow the user to specify several "set-up" sub-prompts. These could include (1) a set-up string to be printed when the program is first run; (2) an input prompt (available now); (3) a string to print before the response to a command; and (4) a string to reset the terminal at the end of execution. For example, this would allow the VT100 user to define the lower part of his screen as a scrolling region for program responses, print a double-width "Function?" at the top of the screen, move to the scrolling region to see the response, and reset the terminal when he exits.

Another feature of the program is that the instructions are printed automatically the first time the program is run. This is possible quite simply due to the fact that the first time the program is run there is no MEMO.DAT file. This feature could be considered for inclusion in any program that creates a data file (or anything else unique to the first run).

While this program has been tested, the disclaimer in the program text should be taken seriously. I am interested in hearing about problems or suggested changes, but I will not promise to do anything about them. Finally, if you find ways to make this program more efficient, please feel free to do so. The inefficiencies are there so that you may have the enjoyment of finding and fixing them. (Thus the larger the inefficiency, the greater my generosity in leaving it for you to find.)

A LIST OF FUNCTIONS

		1	CHOICE %	Users choice of function
		1	CRLF\$	Carriage return/Line feed combination
Add	Adds a memo to your memo file.	1	CUR.MEMO\$	Stores memo number currently listing
	fielde a memorie se jear memorie me	1	DELNUMS	Memo number to delete
		1	DEV\$	Device name for disk lookup
Delete	Removes a previously added memo.	1	DEVCNT\$	Table of maximum unit numbers on syst
Delete	Removes a previously added memo.	1	DEVNAMS	Device name table
		1	DEVOKB\$	Number of disk devices
		1	DISKS\$	Index into DISKS\$ table
List	Lists all memos in the file.	1	DISKS\$ (20)	Table of disk names
	Prints entry date and subject.	1	F%	Flag for output file specified
	Finits entry date and subject.	· · · · · · · · · · · · · · · · · · ·	F.LINK\$	FIELD - Link to next record in list
			FIRST%	Holds first record num in FNFIRST.REC
-		1	FIRST.RUN%	Flag to print help for first-time use
Type	Prints the contents of one or more m	nemos.	FOUND [®]	Contains result of SEARCH function
	Several options are available:	1	FQB\$ (30)	Returned values from SYS call
	Several options are available:		FR.REC%	Free record found by FNUNLINK
	Type Prints a memo. Pro	ogram asks which one.	FREE%	FREE - constant to link pointers
	5.	- grann anna rinnan - rinn	HEAD%	Flag for header printed in SEARCH
	Type 10 Prints memo 10.		I%	Scratch variable
	Type * Prints all memos.		INIT\$	Field used to initialize MEMO.DAT fil
	Type Fints an mentos.		JUNK%	Scratch variable Last record number (in functions)
	Type STRING Prints all memos w	vith the specified STRING	LAST\$ LINKAFTER\$	Record to link new data after (FNLINK
	51		M%	Flag for memo number specified in TYP
	in the subject field		M. NUM\$	Memo number to add to file in ADD
			MAX.MEMO\$	FIELD - Highest memo number used
			MEM. NUMS	Memo number from command
	An output file may be specified as f	follows:	MEMO \$	The memo to be placed in the file
			MEMO, DAT\$	FIELD - MEMO data field (the memo tex
	Type * > Filename		MEMO, DATES	FIELD - The date the memo was filed
			MEMO.FILE\$	Name of memo file
		in the second	MEMO.LINE\$	A line of data to add to the memo
Search	Lists memos with a specific string in	their subject fields.	MEMO.NUM\$	FIELD - Memo number in disk record

More	Appends more information to an existing memo.
Prompt	Changes the current prompt to something else.
Help	Gives some basic help to the user.
Exit	Leaves the program.

A memo number may be specified with the Type, Delete and More commands.

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text)

	MEMO. SEQ\$	FIFID - Next av	ailable memo number					
	MEMO.SPEC\$	Memo number/string/* in TYPE FIELD - unused area in MEMO base rec (END) - constant to link pointers				12000, 2000, 3000, 21000, 5000,	9000, 6000, 7000, 11000, 8000	
	I MEMO.UNUSED			2000	,	Add a memo to the file		
	! NEW.PROMPT\$ The user-spec:					Add a memo to the life		
	I NXT.PTR\$	NULL.LINK\$ Used to indicate end of linked list NXT.FTR\$ Pointer to next record NXT.REC\$ Next record number (FNNEXT.REC\$) NXT.REC\$ Pointer to next record OLCCALL\$ Save previous value of CALL\$ P.NUM\$ Number of memo to print		2005		IF A.NUM% <> 0% THEN		
					1	PRINT "I choose the memo numb GOTO 1200	er."	
				2020		SUB.ADD% = 0%	1 Normal entry.	
	I PART. OF. THE. NEW	. PROMPT \$	A line of new prompt		1	PRINT		
	PREV.REC% PRINT.FILE\$	Output file for	number (FNPREV.REC%) TYPE function		1		! Get the memo subject.	
	PROMPT\$	PROMPT.LEN\$ FIELD - Length of PROMPT\$ in file PRV.PT\$ Pointer to previous record PTR\$ (4) FIELD - List pointers (free/used; start/end) R.LINK\$ FIELD - Link to prev rec (reverse link) R.TYPE\$ FIELD - Record type (free/used) REC\$ Record number passed to functions		2030		PRINT "Now type your memo. End with	Ctrl-Z."	
	I PRV. PTR \$			2030	>			
	! PTR\$ (4)							
				2035		GET #1%		
	! REC%				1	FIELD #1%, RECOUNT AS MEMO.LIN		
	! REQUEST\$! SEARCH\$	SEARCH\$ String to search memo subjects for			1	MEMO\$ = MEMO\$ + MEMO.LINE\$ NEXT	! Get the memo data ! End on Ctrl-Z error	
	I START%	SUB.ADD\$ True if ADD being used as subroutine SUB.SEARCH\$ True if SEARCH being used as subroutine SUB.TYPE\$ True if TYPE being used as subroutine SUB.SECT\$ Subject of memo		2040		RETURN IF SUB.ADD%	! Get out if subroutine.	
	1 SUB.SEARCH%					GET #4%, RECORD 1%	! Get the main pointers.	
					1	LSET MEMO.SEQ\$ =	! Get memo sequence number.	
	SY.MEMO.FILE\$ Name of memo fi TEMP.MEM\$ Temp storage fo		le on system disk r SEARCH			CVT\$\$(CVT\$\$(MEMO.SEQ\$)+1\$) PUT #4\$, RECORD 1\$! Increment the seq. num. ! And replace the pointers.	
	! THIN.AIR\$	THIN.AIR\$ Record number to create in FNUNLINK\$ TODAYS.DATE\$ Obvious TYPE\$ Record type - FREE or USED		2050		CALL\$ = FNUNLINK\$(FREE\$, 0\$)	1 Get a free record.	
	1 TYPE%			2050	\ GOTO 2090 IF CALLS	GOTO 2090 IF CALLS = 0%	0% ! Complain if none free.	
	I TYPE.FILE\$	File to TYPE me FIELD - Anythin				! Get the free record. ! Store the message number.		
	I USED%	USED - pointer	to link pointers		1	LSET MEMO.DATE\$ = CVT\$\$(TODAYS.DATE		
	I VER.DATE\$ I VERSION\$	Date of current Current version			1		! Trim down the memo.	
	1	FUNCTIONS			1	PUT #4, RECORD CALL% CALL% = FNLINK%(USED%, CALL%, 0%)	! Write the record to the file. ! Link the record in the used queue.	
		Returns first r	and the selfest		1		! If more memo, repeat.	
	I FNFIRST.REC%	Returns last re	cord in a list	2060		PRINT		
	I FNL INK% I FNNEXT.REC%		Links a record into a list Returns next record in a list			PRINT "Memo "; CVT\$\$(M.NUM\$); " filed." GOTO 1200		
	I FNPREV.REC%		s record in a list	2090		PRINT		
	I FNUNLINKS		a from a fist	2090	١	PRINT "? No FREE record available"		
1010	DIM PTR\$(4%), FQB%(3) \ USED% = 0%	0), DISKS\$(20)				PRINT GOTO 21000	! Should always be a free ! record available.	
	\ FREE% = 2%			3000	1		! Bad error if not.	
	\ START\$ = 1\$ \ ND\$ = 2\$				1	Delete a memo		
	<pre>\ NULL.LINK\$ = CHR\$(0% \ CRLF\$ = CHR\$(13%) +</pre>				·			
	\ FIRST.RUN% = 0%			3010	1	PRINT CALL2% = 0%		
	<pre>\ TODAYS.DATE% = SWAP\$(CVT\$\$(MID</pre>	(SYS(CHR\$(6%)+CH	R\$(-3%)),27%,2%)))		1	PRINT "Delete memo number: "; UNLESS DELNUM% &		
1100	PROMPT\$ = "Function? "				1	INPUT #1%, DELNUM% UNLESS DELNUM%	! Get memo num if not given.	
1100	<pre>\ VERSION\$ = "1b "</pre>				1	GOTO 3050 IF DELNUM% < 1% GET #4%, RECORD 1%	! Exit if bad number. ! Check against high number.	
	<pre>\ VER.DATE\$ = "Nov 82" \ TYPE.FILE\$ = "_KB:ME</pre>	MO.KBD"			1	GOTO 3050 IF DELNUM\$ >= CVT\$\$(MEMO	.SEQ\$) ! Exit if number too high.	
	\ SY.MEMO.FILE\$ = "_DB	0:MEMO.DAT"						
	<pre>\ ON ERROR GOTO 32000 \ OPEN TYPE.FILE\$ FOR</pre>	TNDIT AS FTIF #1	*	3020	1	CALL% = FNFIRST.REC%(USED%) CALL% = FNNEXT.REC%	! Start at the beginning.	
	\ GOTO 15000	INFOI AS FILE #1				UNTIL MEMO.NUM\$ = CVT\$ OR CALL\$ = 0\$	\$(DELNUM%)	
			! Set up some defaults ! and go find the data file.		1		! Find the record or end. ! Get out if end.	
1110	FIELD #4%, 2% AS	PTR\$(1%),	! Used start pointer			CALL2% = FNUNLINK%(USED%, CALL%)	! Unlink the record.	
1110	2% AS	PTR\$(2%),	! Used end pointer		1	GOTO 3020 UNLESS CALL2% = 0% OR CALL2% <> CALL%	! Repeat for more records. !unless there was an error.	
	2% AS	PTR\$(3%), PTR\$(4%),	! Free start pointer ! Used end pointer	3030		IF CALL2% = 0% THEN		
		PROMPT.LEN\$, MAX.MEMO\$,	! Length of prompt ! Highest record counter	3030		PRINT "? Couldn't unlink record	d ";CALL%	
	2% AS	MEMO.SEQ\$,	I Memo sequence number			ELSE PRINT "? Deleted record ";CALL:	2%;" not ";CALL%	
	8% AS 490% AS	UNUSED.FOR.NOW\$ BUF\$		3040		PRINT		
			! FIELDS for record 1 ! (main pointer block)	3040	1	GOTO 1200		
		D I THEA	I Forward link	3050		IF CALL2% <> 0% THEN		
	2% AS	AS F.LINK\$, AS R.LINK\$,	! Reverse link			PRINT "Memo "; DELNUM%; " dele ELSE	ted."	
		R.TYPE\$, MEMO.NUM\$,	! Record type ! Memo number &			PRINT "Can't find memo ";DELNUM%		
	2% AS	MEMO.DATE\$, MEMO.UNUSED\$,	! Memo date & ! Unused				! Print the results.	
	500% AS	MEMO.DAT\$! Memo data	3060		GOTO 1200		
			! FIELDS for data records	5000	1	List any pending memos.		
	\ GET #4, RECORD 1% \ IF CVT\$%(PROMPT.LEN\$) > 0% THEN			1	LISE any pending memos.		
		BUF\$, CVT\$%(PROM	PT.LEN\$))	5005		IF A.NUMS <> 0% THEN		
1120	PRINT				`	PRINT "Show some guts - list t GOTO 1200	hem all."	
	<pre>\ PRINT "M E M O \ PRINT RIGHT(SYS(CHR\$</pre>	"; (6)+CHR\$(9)+CHR\$	(0)),3)				I Goto first used record.	
	\ PRINT "Version ";VER	SION\$;" ";VER.D	ATE\$	5010		CALL% = FNFIRST.REC%(USED%) GOTO 5050 IF CALL% = 0%	! If noneno memos.	
	<pre>\ PRINT \ GOTO 11000 IF FIRST.</pre>	RUN%				PRINT PRINT "Memo", " Date", "Subject"		
						PRINT PRINT		
1200	PRINT \ CLOSE 1\$			5020		CUR.MEMO\$ = MEMO.NUM\$ + "" ! Copy the memo number.		
	\ OPEN TYPE.FILE\$ AS F	<pre>\ OPEN TYPE.FILE\$ AS FILE #1\$! Re-open keyboard for input \ A.NUM\$, P.NUM\$, DELNUM\$ = 0\$! Zero memo nums (Type, Delete, More) \ SUB.ADD\$, SUB.SEARCH\$, SUB.TYPE\$ = 0\$</pre>			1	<pre>PRINT CVT\$\$(MEMO.NUM\$),</pre>		
	\ SUB.ADD%, SUB.SEARCH							
	<pre>\ PRINT PROMPT\$; \ INPUT LINE #1, REQUEST\$ REQUEST\$ = CVT\$\$(REQUEST\$, 85 +4\$) ! Dump leading spaces and CR/LF. IF LEFT(CVT\$\$(REQUEST\$, -1\$), 1\$) = "T" THEN 13000</pre>				/	UNTIL MEMO.NUM\$ <> CUR	. MEMO\$	
						OR CALLS = 0%	! Find next memo.	
			1 Check fancy type functions.		1	GOTO 5020 IF CALL\$ <> 0\$! And print it.	
1210	MEM.NUM% = VAL(RIGHT	(REQUEST\$, INSTR(5030		PRINT		
	I Get memo number if given.					PRINT GOTO 1200	! All done.	
1000			19)	5050	,	PRINT		
1220	REQUEST\$ = LEFT(CVT\$ \ GOTO 1200 IF REQUEST	\$ = ""	! No function - try again	5050		PRINT "No memos on file."	I Nothing to mint	
	<pre>\ CHOICE\$ = INSTR(0\$, \ ON CHOICE\$ GOTO</pre>	"ADELMPTHS", REQ	UEST\$) + 1%		1	GOTO 5030	! Nothing to print. ! SEARCH func comes here, too.	

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6000	I Change the prompt to something else.
6005	IF A.NUMS <> 0% THEN PRINT "Change the prompt in a memo? Ridiculous!" GOTO 1200
6010	PRINT \ PRINT "Enter your new prompt, ending with Ctrl-Z." \ PRINT \ NEW.PROMPT\$ = ""
	\ WHILE 1\$ \ GET #1\$ \ FIELD #1\$, RECOUNT AS PART.OF.THE.NEW.PROMPT\$
	<pre>NEW.PROMPT\$ = NEW.PROMPT\$</pre>
6000	\ NEXT
6020	GET 44\$, RECORD 1\$ NEW.PROMPT\$ = LET(NEW.PROMPT\$,490\$) FOR 1\$ = 1\$ TO LEN(NEW.PROMPT\$)
	<pre>\ IF ASCII(MID(NEW.PROMPT\$, I\$, 1\$)) = 27\$ THEN NEW.PROMPT\$ = LEFT(NEW.PROMPT\$, I\$-1\$)</pre>
	+ CHR\$(155\$) + RIGHT(NEW.PROMPT\$, I\$+1\$) I Set parity on escapes.
6030	NEXT I\$ \ LSET BUF\$ = NEW.PROMPT\$ + STRING\$(407\$-LEN(NEW.PROMPT\$),0\$) \ LSET PROMPT.LEN\$ = CVT\$\$(LEN(NEW.PROMPT\$)) \ PUT #4\$, RECORD 15 \ PROMPT\$ = NEW.PROMPT\$ \ GOTO 1200
7000	I I Print a memo
7010	9 PRINT
	\ SUB.TYPE\$ = 0\$
7012	PRINT "Which memo? "; UNLESS P.NUM\$ \ INPUT 41\$, P.NUM\$ UNLESS P.NUM\$ \ IF SUB.TYPE\$ THEN IF F\$ THEN
	OPEN PRINT.FILE\$ FOR OUTPUT AS FILE #1\$ ELSE
	OPEN TYPE.FILE\$ FOR OUTPUT AS FILE #1% 1 Open the file if specified.
7015	GOTO 7050 IF P.NUM\$ < 1\$ I Ask for number - quit if bad. GET 04\$, RECORD 1\$ I Check against high number. GOTO 7050 IF P.NUM\$ >= CVT\$\$(MEM0.SEQ\$)
	<pre>\ CALL\$ = FNFIRST.REC\$(USED\$) & CALL\$ = FNNEXT.REC\$ UNTIL MEMO.NUM\$ = CVT\$\$(P.NUM\$) OR CALL\$ = 0\$</pre>
	\ GOTO 7050 IF CALL\$ = 0\$! Find matching memo or end. \ Complain if none.
7020	GOSUB 7100 IF INSTR(0\$, MEMO.DAT\$, CHR\$(3)) I Print header if first rec.
	<pre>\ PFINT #1\$, CUT\$\$(RIGHT(MEMO.DAT\$,INSTR(0,MEMO.DAT\$,CHR\$(3))+1\$),128\$); \ Call\$ = FNNEXT.RC\$ \ GOTO 7020 IF NEMO.NUM\$ = CVT\$\$(P.NUM\$) UNLESS CALL\$ = 0\$ RETURN IF SUB.TYPE\$</pre>
	PRINT POPULATION
7050	PRINT 01%, "Memo "; P.NUM%; " not found." PRINT 01%, "Memo "; P.NUM%; " not found." \ RETURN 01% \ GOTO 1200
7100	PRINT 015 PRINT 015
	<pre>\ PRINT #1\$, "Memo "; P.NUM\$, "Date "; DATE\$(CVT\$\$(MEMO.DATE\$)) \ PRINT #1\$, "Subject: "; \ PRINT #1\$, LEFT(MEMO.DAT\$, INSTR(0\$, MEMO.DAT\$, CHR\$(3\$)) -1\$) \ PRINT #1\$ \ RETURN</pre>
8000	1 SEARCH Function
8005	IF A.NUMS <> 0\$ THEN PRINT "Why search if you know the memo number?" \ GOTO 1200
8010	PRINT \ SUB.SEARCH\$ = 0\$
	\ HEAD\$ = 0\$ PRINT "Search for: "; \ INPUT LINE #1\$, SEARCH\$ \ CALL\$ = FWFIRST.REC\$(USED\$) \ GOTO 505 IF CALL\$ = 0\$ I Exit 1f no memos.
8015	SEARCH\$ = CVT\$\$(SEARCH\$, 32\$ + 4\$)
8020	<pre>I\$ = INSTR(0\$, MEMO.DAT\$, CHR\$(3\$)) \ GOTO 8030 IF I\$ = 0\$ TEMP.MEM\$ = CVT\$\$\$ (LEFT(MEMO.DAT\$, I\$), 32\$ +4\$) \ FOUND\$ = INSTR(0\$, TEMP.MEM\$, SEARCH\$) \ GOTO 8030 UNLESS FOUND\$ \ RETURN IF SUB.SEARCH\$ \ IF HEAD\$ = 0\$ THEM</pre>
8025	<pre>PRINT #1\$, CVT\$\$(MEMO.NUM\$), DATE\$(CVT\$\$(MEMO.DATE\$)), LEFT(MEMO.DAT\$, INSTR(0\$, MEMO.DAT\$, CHR\$(3))-1\$);</pre>
8030	CALL\$ = FNNEXT.REC\$ \ GOTO 8020 IF CALL\$ > 0\$

```
8050
                          RETURN IF SUB. SEARCHS
                          PRINT
                   \ GOTO 1200
 9000
                   1
                           Append to a memo
9010
                          PRINT
                         PRINT "Add to which memo? "; UNLESS A.NUM$

INPUT #1, A.NUM$ UNLESS A.NUM$ I Get the memo number.

GOTO 9070 IF A.NUM$ < 1$ I See if too small.
                        ! Find the first record of
! ...the memo to append to
! ...or end of the list.
! Get out if not found.
! While reading the same memo.
                        GOTO 9070 IF CALL$ = 0$

WHILE MEMO.NUM$ = CVT$$( A.NUM$)

AND CALL$ <> 0$

OLDCALL$ = CALL$

CALL$ = FNNEXT.REC$

NEYT
                        CALLS = FNNEXT.RECS I Get the next record.

NEXT I Now have last record of memo.

MEMO$ = CRLF$ + "***** Appended "+DATE$(0$)+" ******-CRLF$
                                                                                                            ! Set up append notice.
! Flag add as a subroutine.
                        SUB.ADD% = 1%

      CALL$ = FNUNLINK$( FREE$, 0$)
      1 Get a free record.

      GOTO 9050 IF CALL$ = 0$
      1 Complain if none free.

      GET 4#$, RECORD CALL$
      1 Get the free record.

      LSET MEMO.DMT$ = CYT$$(NOW$)
      1 Store the message number.

      LSET MEMO.DAT$ = LEFT(NEMO$,500$)
      1 Store the date.

      LSET MEMO.DAT$ = LEFT(NEMO$,500$)
      1 Store some memo data.

      MEMO$ = RIGHT(MEMO$,500$)
      1 Write the record to the file.

      CALL$ = NENIN$(USE$, CALL$, 0LDCALL$) I Link the record in the used queue.

      IF LEN(MEMO$) > 0$ THEN
      1 If more memo, repeat.

      OLDCALL$ = CALL$
      01000

9030
                                   GOTO 9030
                        PRINT
PRINT "Memo "; A.NUM%; " updated."
SUB.ADD% = 0%
9040
                  \ GOTO 1200
                  PRINT

SUB.ADD$ = 0$

PRINT "? No FREE record available"
9050
                         PRINT
                                                                                                              ! Should always be a free
                        GOTO 1200
                                                                                                               ! record available.
! Bad error if not.
                  PRINT
\ PRINT "Memo "; A.NUM≸; " not found."
\ PRINT
\ GOTO 1200
9070
11000
                 .
                                  HELP Function
                        PRINT

PRINT "The MEMO program is designed as a computerized 'filing-box'."

PRINT "Its purpose is to allow people who spend a good deal of time"

PRINT "ising the computer system to use the computer to keep track of"

PRINT "ideas or messages."
11010
                        The TYPE command has several options available. They are:"
                                                                                                            Type a memo. Program asks which one."
Type memo number 10."
Type all current memos."
Type all memos with the"
specified string in the subject."
                         PRINT " Specified String in the subst

PRINT " You may have your memos typed to a file by including"

PRINT " filename in the TYPE command as follows:"

PRINT " TYPE • > FILENAME Puts all memos into the"

PRINT " file specified by FILENAME."
                          PRINT
                         PRINT "Type the RETURN key to continue.";
INPUT LINE #1,A$
PRINT
                         PRINT "
                                                  A list of commands follows:"
                  \ GOTO 12020
12000
                                   List the legal functions.
                  PRINT 

\ PRINT "I humbly apologize, but I do not recognize this command." 

\ PRINT "My limited repertoire is as follows:"
12010
12020
                        PRINT
PRINT "Add
PRINT "Delete
PRINT "List
PRINT "Type
PRINT "Search
PRINT "More
PRINT "Prompt
PRINT "Help
PRINT "Exit
PRINT "Exit
                         PRINT
                                                                         Add a memo to your memo file."
Remove a previously added memo."
List memos on file."
Print the contents of a memo."
                                                                        Frint the contents of a memo."
List memos with a specific subject."
Append more information to an existing memo."
Change the current prompt to something else."
Gives more details about this program."
Leaves the program."
                         PRINT
                        PRINT "All commands may be abbreviated to their first letter. The Type, "
PRINT "Delete and More commands may be followed by a memo number."
GOTO 1200
```

13000	! ! !	The magic routine to handle out co-ordinate a type by string se					ero all pointers and prompt length. To need to put the prompt tring in yet. Program init an handle it. &
13010		REQUEST\$ = CVT\$\$(REQUEST\$, 128\$+4\$) F\$ = INSTR(0\$, REQUEST\$, ">") IF F\$ THEN	! Look for file indicator.			<pre>\ PUT #4%, RECORD 1% \ FIRST.RUN% = 1% \ GOTO 1110</pre>	
	1	PRINT.FILE\$ = RIGHT(REQUEST\$, REQUEST\$ = LEFT(REQUEST\$, F%-1	F\$+1\$) \$) ! Trim file out of command.	20050		OPEN MEMO.FILE\$ FOR OUTPUT AS F: \ GOTO 20010	LE #4\$! Re-init a bad memo file.
13020	1	REQUEST\$ = CVT\$\$(REQUEST\$, 128%) MEMO.SPEC\$ = ""	I Look for memo number/search.	21000		! ! Exit !	
	,	M% = INSTR(0%, REQUEST\$, " ") IF M% THEN MEMO.SPEC\$ = RIGHT(REQUEST\$, M		21010		CLOSE 1,4 GOTO 32767	
13025	`	P.NUM% = VAL(MEMO.SPEC\$) GOTO 7000 UNLESS F%	<pre>! Get memo number (string = err). ! If memo number and no file,</pre>	25000			
	1	SUB.TYPE# = 1# GOSUB 7012	<pre>! just print the memo. ! Print single memo to file.</pre>			! ! Remove a record from the sp	ecified linked list.
	1	GOTO 13060					d list indicator. er to be removed.
13030 13040		IF MEMO.SPEC\$ = "*" THEN 13100 SUB.SEARCH\$ = 1\$! Print all the memos.			! (va !	lid for USED list only)
13040	1	CALL\$ = FNFIRST.REC\$(USED\$) CALL\$ = 1\$ FOUND\$ = 0\$! Start at the first record.			I This routine is the real br structure of this program. Without being UNLINKED from	No record should be used somewhere using this
	1	SEARCH\$ = MEMO.SPEC\$ IF F% THEN	! Get the search string.			<pre>! routine. It performs the or ! you are likely to find in !</pre>	
		OPEN PRINT.FILE\$ FOR OUTPUT AS ELSE				IF TYPE = FREE THEN	
		OPEN TYPE.FILE\$ AS FILE #1%	! Do file or keyboard.		3	REMOVE THE FIRST RE I IF THERE WAS A RECO	CORD FROM THE FREE QUEUE RD TO REMOVE
13050	١	GOSUB 8015 IF FOUND\$ > 0\$ AND CALL\$ > 0\$ THEN P.NUM\$ = CVT\$\$(MEMO.NUM\$)	<pre>! Look for the record. ! If found, go print it.</pre>		1	I ELSE	THE FREE QUEUE LINKS
	1	SUB.TYPE% = 1% GOSUB 7020	I Set the subroutine flag.			I CREATE A NE I I IF TYPE = USED	W RECORD OUT OF THIN AIR
	1	GOTO 13050	! And look for another one.			I IF TIPE = USED I THEN I REMOVE A RECORD FRO	M THE USED QUETIE
13060		PRINT #1% PRINT #1%				PLACE IT ON THE FRE	
	1	CLOSE #1% OPEN TYPE.FILE\$ AS FILE #1% SUB.SEARCH\$, SUB.TYPE\$ = 0\$ P.NUM\$, FOUND\$ = 0\$ GOTO 1200		25010		DEF* FNUNLINK\$(TYPE\$, REC\$) \ GOTO 25050 IF TYPE\$ = USED\$ \ FR.REC\$ = FNFIRST.REC\$(FREE\$) \ GOTO 25030 IF FR.REC\$ = 0\$	I Check for used record. I Get the first free record. I If none, do something else.
13100	! ! !	Print all the memos.				<pre>NXT.REC\$ = F.LINK\$ + "" GET #4\$, RECORD 1\$ LSET PTR\$(FREE\$+START\$) = NXT.RE LSET PTR\$(FREE\$+ND\$) = NXT.REC\$ IF CVT3\$(NXT.REC\$) = 0\$</pre>	! Save link to next free rec. ! Go to the main link pointers.
13110	1	CALL\$ = FNFIRST.REC\$(USED\$) IF CALL\$ = 0\$ THEN PRINT "No memos to print." GOTO 13060	! Get the first memo. ! Unless there wasn't one.			<pre>PUT #4\$, RECORD 1\$ IF CV1\$\$(NXT.REC\$) > 0\$ THEN GET #4\$, RECORD CV1\$\$(NXT.F</pre>	<pre>! And save it. ! If there was a next record EC\$) !grab it.</pre>
13115		IF F% THEN OPEN PRINT.FILE\$ FOR OUTPUT AS	FILE #1%		,	LSET R.LINK\$ = NULL.LINK\$ PUT #4%, RECORD CVT\$%(NXT.F	! Null out the reverse pointer EC\$) !and replace the record.
		ELSE OPEN TYPE.FILE\$ AS FILE #1%				GOTO 25090	! And exit gracefully.
13120	、	CUR.MEMO\$ = MEMO.NUM\$ + "" P.NUM\$ = CVT\$\$(CUR.MEMO\$)	I Get current memo number.	25030	,	GET #4%, RECORD 1% \ THIN.AIR\$ = MAX.MEMO\$ + "" \ LSET MAX.MEMO\$ =	I Go for the big data. I Get the next to create.
	1	SUB.TYPE# = 1# GOSUB 7020	! Print a memo.		,	CVT\$\$(CVT\$\$(MAX.MEMO\$) +1 > PUT #4\$, RECORD 1\$	I Save the new big data.
	1		! Do another if more.		1	<pre>\ LSET F.LINK\$ = CVT\$\$(0\$) \ LSET R.LINK\$ = CVT\$\$(0\$) \ LSET R.TYPE\$ = CHR\$(FREE\$)</pre>	<pre>! Null the forward !and reverse links. ! Initialize as free record.</pre>
15000	1	Check out all the disks for the	memo file		;	LSET MEMO.NUM\$ = CVT\$\$(0\$) LSET MEMO.DATE\$ = CVT\$\$(TODAYS.D	! Might as well do memo num. ATE\$)! Put in the date for laughs. \$) ! And empty out the data fieid.
15010	1	CHANGE SYS(CHR\$(6\$) + CHR\$(-3\$)) TO DEVCNT\$ = FQB\$(5\$) OR SWAP\$(FQB\$(6\$)	FQB\$)			<pre>FR.REC\$ = CVT\$\$(THIN.AIR\$) GOTO 25090</pre>	! Keep track of it. ! Exitstage left.
	1	CHANGE SYS(CHR\$(6\$) + CHR\$(-12\$)) TC DEVNAM\$ = FQB\$(5\$) OR SWAP\$(FQB\$(6\$) DEVOKB\$ = FQB\$(9\$) OR SWAP\$(FQB\$(10\$! Get :)				
15020	1	FOR IS = 0% TO DEVOKES-2% STEP 2% DEV\$ = CVT\$\$(SWAP\$(PEEK(DEVNAM\$	+1\$)))	25050		<pre>FR.REC\$ = FNFIRST.REC\$(USED\$) FR.REC\$ = FNNEXT.REC\$</pre>	! Get the first used record.
	1	FOR J% = 0% TO PEEK(DEVCNT%+1%) DISKS% = DISKS% + 1%				UNTIL (FR.REC% = RE OR (FR.REC% = 0	
	1	DISKS\$(DISKS\$) = "_" + NEXT J\$	DEV\$ + NUM1\$(J\$) + ":"		1	GOTO 25090 UNLESS FR.REC%	! Get out if not there. ! Save forward link.
	/	! Build	through all disk names each disk name in sequence for all units of this disk			<pre>PRV.PTR\$ = R.LINK\$ + "" IF CVT\$\$(NXT.PTR\$) <> 0\$ THEN UNU\$\$ CNNEY" PEC\$\$</pre>	! Save reverse link. ! If there was a next record
		! Incre ! Build	ment disk unit counter and store complete disk name		,	JUNK\$ = FNNEXT.REC\$ LSET R.LINK\$ = PRV.PTR\$ PUT #4\$, RECORD JUNK\$! Get it ! Update reverse link ! And replace the record.
15030		! Until CALL\$ = SYS(CHR\$(6\$) + CHR\$(-21))	we have all of them	25060		IF CVT\$%(PRV.PTR\$) <> 0% THEN	! If there was a previous record.
15030			with privileges - drop them.			JUNK\$ = FNPREV.REC\$ LSET F.LINK\$ = NXT.PTR\$ PUT #4\$, RECORD JUNK\$	<pre>! Get it ! Update forward link ! And replace the record.</pre>
	1	MEMO.FILE\$ = DISKS\$(I\$) + "MEMO OPEN MEMO.FILE\$ FOR INPUT AS FI	LE #4% ! Try to open file.	25065		IF CVT\$%(NXT.PTR\$) = 0%	! If no next record.
	1	IF (STATUS AND 1024\$) = 0\$ THEM ELSE PRINT "Don't have write	access to ";MEMO.FILE\$,	THEN GET #4%, RECORD 1% LSET PTR\$(USED\$+ ND\$) = PR	
15060		NEXT 1%		2506*	1	PUT #4%, RECORD 1%	! Update used end ptr.
		MEMO.FILE\$ = SY.MEMO.FILE\$! Didn' GOTO 20050 ! And g	t exist - default to system. To create it.	25067		IF CVT\$\$(PRV.PTR\$) = 0% THEN GET #4%, RECORD 1%	! If no previous record. ! Get main pointers.
20000	!	Initialize a new memo file			1	LSET PTR\$(USED\$+ START\$) = PUT #4\$, RECORD 1\$	
	!			25070		IF FR.REC% <> 0% THEN	! If a record was removed REC\$, O\$)!place it on the free queue.
20010	1	FIELD #4%, 512% AS INIT\$ LSET INIT\$ = STRING\$(10%,0%) + CUT\$\$(24) + CUT\$\$(14)		25090		FNUNLINK# = FR.REC#	! Return the record number.
		+ CVT\$\$(2\$) + CVT\$\$(1\$)			1	FNEND	

25100	! ! FNL INK%		
	! ! Add a record to the	e linked list	specified.
	I TYPES Fro	ee or Used lis	t indicator
	I LINKAFTERS Rea		link record after
	! (1) !	n USED 11st on	ily - zero means end)
25110	DEF* FNLINK\$(TYPE\$, REC\$ \ GOTO 25150 IF (TYPE\$ = U	, LINKAFTERS)	VARTER > Of)
	\ LAST\$ = FNLAST.REC\$(TYPE) \ IF LAST\$ > 0\$	\$)	! Move to the last record. ! If there is a last record
	THEN GET #4%, RECORD LAS	CT.	I Read the last record.
	<pre>\ LSET F.LINK\$ = CVT \ PUT #4\$, RECORD LAG</pre>	\$\$(REC%) ST%	! Update the forward link. ! And replace the record.
	\ GET #4, RECORD 1% \ LSET PTR\$(TYPE\$ + 1)	NDS)	
	= CVT\$\$(RE) PUT #4\$, RECORD 1\$	C\$)	! Update last ptr = new record.
	\ GOTO 25130		
25120	IF LAST\$ = 0\$ THEN		I If there was no last record I Get the main pointers.
	GET #4\$, RECORD 1\$ \ LSET PTR\$(TYPE\$+ST. = CVT\$\$(R)	ART\$)	I Init the start and
	<pre> LSET PTR\$(TYPE\$+ND)</pre>	\$)	! end pointers.
	\ PUT #4%, RECORD 1%		, end pointer of
25130	GET #4%, RECORD REC%	K\$! Get the record to link. ! Null the forward link.
	<pre>\ LSET R.LINK\$ = CVT\$\$(LA: \ LSET R.TYPE\$ = CHR\$(TYP)</pre>	ST\$) E\$)	<pre>Null the forward link. 1 Old LAST is reverse link. 1 Add record type.</pre>
	V PUT #4%, RECORD REC%		! Put the record in the file.
25140	FNLINK\$ = REC\$		I Return installed record #.
25150	CALLS = FNFIRST.RECS(U	SED%)	I Get the first record in list.
	<pre>\ CALL\$ = FNNEXT.REC\$ UNTIL CALL\$ OP</pre>	S = LINKAFTERS	I Find the record to add to.
	\ GOTO 25190 IF CALL\$ = 0	\$.	I Exit if not there. I Save pointer to next record.
	<pre>\ NXT.PTR\$ = F.LINK\$ + "" \ PRV.PTR\$ = R.LINK\$ + "" \ LSET F.LINK\$ = CVT\$\$(R)</pre>		I Update forward pointer.
	V PUT #4%, RECORD CALLS	- NULL. LTNKS	I Restore the record.
	<pre>\ GET #4%, RECORD CVT\$%(\ PRV PTR* - R LINK* + ""</pre>	NXT.PTR\$)	I Get the next record.
	<pre>\ LSET R.LINK\$ = CVT\$\$(R \ PUT #4\$, RECORD CVT\$\$(</pre>	EC\$)	! Install new pointer. ! Restore the record.
25160	GET #4%, RECORD REC%		I Get the record to link in.
	<pre>\ LSET F.LINK\$ = NXT.PTR\$ \ LSET R.LINK\$ = PRV.PTR\$</pre>		! Set up the links.
	<pre>\ LSET R.TYPE\$ = CHR\$(TY \ PUT #4\$, RECORD REC\$ \ COTO 25180</pre>	PE%)	
25190	\ GOTO 25140 REC\$ = 0\$		
20190	\ GOTO 25140		
25200	1		
	I FNNEXT.RECS		
	Return the next re Works with the fre	e or used list	h a linked list. t, depending on
	! which list was use ! ! F.LINK\$ Mu		ltered after last
	I F.LINK\$ Mu	record n	read. This contains ward link used here.
	FNNEXT.RECS Re	turns next red	cord number (or zero ext record). If there
		is a ner	kt record, it is read.
25210	DEF* FNNEXT.RECS		
	<pre>\ NXT.REC\$ = CVT\$\$(F.LINK \ IF NXT.REC\$</pre>		<pre>! Read forward link for next. ! If there is a next rec. !then read it.</pre>
05000			1then read it. 1 Return the record # (or zero).
25220	FNEXT.REC% = NAT.REC%		I NECULAI CHE LECOLU P (OL 2010).
25300	I FNPREV.REC%		
	I I Return the previou	s record numbe	er in a linked list.
	! Works with the fre ! which list was use	e or used list	t, depending on
	I F.LINK\$ Mu	st remain unal	ltered after last read. This contains
	I I FNPREV.REC≸ Re	the back	kward link used here. s record number (or zerô
		if none). If there is a s record, it is read.
	i		
25310	DEF* FNPREV.REC% > PREV.REC% = CVT\$%(R.LI)	NK\$)	I Get the backward link.
	\ IF PREV.REC%		! If there is one !read the record.
25320	FNPREV.REC# = PREV.REC		! Return record # (or zero).
	\ FNEND		
25400	! I FNLAST.REC%		
	! ! Return the last re	ecord number i	n a linked list.
	I TYPES L.	ist type free	or used.
	I F.LINK\$ M	ust remain una record	ltered after last read. This contains

	the backward link used here.
	f FNLAST.REC% Returns last record number (or zero if none). If there is a
	if none). If there is a last record, it is read.
	1
25410	DEF* FNLAST.REC%(TYPE%)
	<pre>\ GET #4\$, RECORD 1\$! Get the main links. \ LAST\$ = CVT\$\$(PTR\$(TYPE\$+ND\$)) ! Get END pointer for TYPE.</pre>
	\ IF LAST ! If there is a last record
	THEN GET #4\$, RECORD LAST\$! then read it.
25420	FNLAST.REC% = LAST% ! Return record number.
	\ FNEND
25500	1
	I FNFIRST.REC%
	Return the first record number in a linked list.
	! TYPES List type free or used.
	I FNFIRST.REC% Returns first record number (or zero
	i if none). If there is a first record, it is read.
	1
25510	DEF" FNFIRST.REC\$(TYPE\$)
	<pre>\ GET #4\$, RECORD 1\$! Get the main links. \ FIRST\$ = CVT\$\$(PTR\$(TYPE\$+START\$)) ! Get START pointer for TYPE.</pre>
	\ IF FIRST\$! If there is a last record
	THEN GET #4%, RECORD FIRST\$! then read it.
25520	FNFIRST.REC% = FIRST% ! Return record number.
	\ FNEND
32000	
5	1 Handle those errors
	1
32010	IF ERL = 15050\$
	THEN RESUME 15060 ! Keep checking disks to find the memo file.
32030	IF ERR = 11% AND ERL = 6010% THEN FIELD #1%, RECOUNT -1% AS PART.OF.THE.NEW.PROMPT\$
	NEW. PROMPT\$ = NEW. PROMPT\$
	+ PART.OF.THE.NEW.PROMPT\$
	! Handle the control-Z at the end of the prompt.
32040	IF ERR = 11% AND ERL = 2035% THEN
5	FIELD #1%, RECOUNT -1% AS MEMO.LINE\$
	<pre>\ MEMO\$ = MEMO\$ + MEMO.LINE\$ \ RESUME 2040</pre>
	! Handle the control-Z at the end of the memo.
32050	IF ERR = 11% AND ERL = 2020% THEN
	RESUME 1200 1 Control-Z no good to end memo subject.
	I CONTROL-2 NO BOOD TO BUD memo subject.
32070	IF ERR = 28\$ OR (ERR = 11\$ AND ERL = 1200\$) THEN
	CLOSE 1\$,4\$
	RESUME 32767 I Exit on control-C
	! or control-Z at prompt.
32080	IF (ERL = 3010\$)
32000	OR (ERL = 7010%)
	OR (ERL = 8010 THEN RESUME 1200 I Head back to the prompt if any problems.
32090	RESUME 1220
	I Bad number given on input (Type, Delete)
32100	IF ERL = 13025% THEN
	RESUME 13030 ! If memo number was a search string
32110	IF (ERL = 13040\$) OR (ERL = 13115\$)
	OR (ERL = 13210\$) THEN
	PRINT
	V PRINT PRINT.FILE\$; "? Looks like GARBAGE to mel" V PRINT
	N RESUME 1200 ! Complain about filename.
32120	IF ERL = 9010% THEN RESUME 1200
	! Bail out if bad memo number in MORE.
32600	IF ERL > 25000\$
32000	AND ERL < 32000% THEN
	PRINT PRINT "MEMO file format error"
	RESUME 32710
	! Fix the error in the file.
32610	IF ERL = 32710\$
	THEN RESUME 21000 ! Quit if error in init file answer.
32620	IF ERL = 11010\$ THEN RESUME 1200
	! They got tired of the instructions.
32700	PRINT
32100	\ PRINT "? Unexpected Error"
	<pre>\ PRINT RIGHT(SYS(CHR\$(6)+CHR\$(9)+CHR\$(ERR)),3%); \ PRINT " in line "; ERL</pre>
	\ RESUME 21000
	! Can't expect everything
327 10	PRINT
	<pre>\ PRINT "Re-initialize the file? "; \ INPUT LINE #1, A\$</pre>
	<pre>\ A\$ = LEFT(CVT\$\$(A\$,-1\$),1\$)</pre>
	<pre>\ IF A\$ = "Y" THEN GOTO 20050 ELSE PRINT "? Bad MEMO file - can't continue"</pre>
	CLOSE 1,4
32767	END

Meet Eddie. He's learning everything about the business.

Everything.

LOCK-11 is a system security and management package for RSTS.

LOCK-11 gives you absolute control of access by keyboard or user-I.D.

LOCK-11 provides an optional MENU environment that keeps non-privileged users where they belong. **LOCK-11** offers the system manager powerful surveillance utilities that actually improve thruput. **LOCK-11** is very well documented, supported and enhanced regularly.

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HOW TO USE SHAREABLE DATA

By Ken Isaacs, Digital Management Group, Ltd., Willowdale, Ontario, Canada

OVERVIEW

With the advent of RSTS/E V7.0, there was considerable support given to resident libraries. They are usable for code or data sharing between many concurrently active jobs. We also have been provided with all the tools and raw data from Digital with which to manipulate these new features. What has not been clear, however, is "HOW DOES ONE SHARE DATA EASILY BETWEEN PROGRAMS WRITTEN IN A HIGH LEVEL LANGUAGE?" Relax, sit back and read. I will present the simple steps you too can take. I will briefly discuss:

- why use resident common areas,
- cookbook instructions on how to make one,
- elephant traps to avoid.

INTENDED AUDIENCE

This document is intended for those of you who know and love RSTS/E and plan to do great things with your system. The following discussion does get rather detailed. I have tried to make it as clear as possible so that it is not necessary to know lots about machine language, task builders and the like. My objective is that you will learn how to share data between concurrently active programs using resident common areas.

PREREQUISITE READING

Although there is no mandatory prerequisite reading, I suggest that you become familiar with:

• one high-level language (BASIC+2 or DIBOL) and what it generates in terms of object code (PSECTs, Global symbols, etc.),

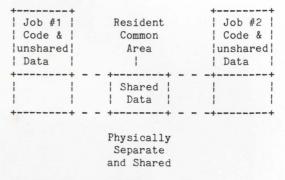
- a little (just a little, mind you) MACRO,
- the Task Builder utility.

The following reference manuals will be handy:

- Task Builder Reference Manual,
- MACRO-11 Language Manual,
- RSTS/E System Manager's Guide for UTILTY,
- RSTS/E Programmer's Utility Manual for MAKSIL.

DEFINITION OF COMMON AREAS

A common area is a portion of a program's address space that is used to store/retrieve information that is used between modules. Many modules in one program can access this area of common information. The term, resident common area applies to a special case of common area. A resident common area is one that is separately loaded by the system, and is connected to by one or more programs. Thus it is possible for several jobs to easily share their data.



WHY USE RESIDENT COMMON AREAS?

The addressing limitations of PDP-11s in commercialapplications are well known by this time. Many applications use:

- complex program overlay structures,
- many detached/chaining jobs,
- virtual arrays,
- scratch pad files,
- large amounts of interjob message sending,

• frequently used smallish size (1-40 blocks) tables that are expensive to put in files,

- any of the above,
- all of the above.

I cannot realistically do away with the requirement for a large amount of program segmentation. What I can inform you of is an easy to use mechanism for sharing data between many small jobs. Your application design will have to ensure that all individual programs do not become too large and unwieldy.

SAMPLE USES OF RESIDENT COMMON AREAS

It is hard to visualize the use of a resident common area unless it is put into an "application" example. The following are two examples of existing applications that could use resident common areas:

- high speed table lookup,
- processing of the same data by several jobs.

HIGH SPEED TABLE LOOKUP

One application involves high volumes of data entry with minimal verification. The longest time spent in verification is looking up the product code to see if it exists. Additionally, the verification table is updated periodically and the changes must be in effect at once. As a result, this table is stored in a file. There are only a couple of hundred product codes. This would fit neatly into an array in a resident common area and allow any product additions or deletions to take effect immediately.

+-						+						+-							+
ł	J	ob	#	1		1						1	Jo	b	#2	2			1
+•						+						+-							+
+	-	-	-	-	-	-+	-	+-			+	-+	-	-	-	-	-	-	-+
1						1		ł	Product	Table	1	1							1
+	_	_	-	-	_	-+	-	+-			+	-+	-	-	-	-	-	-	-+

++	<-	-	++
Slot #2			User Pgm 1
++	<-	-	++
++			
1 02 1 1			

| Slot #n +-----

Product Table is part of each program's virtual address space

PROCESSING BY SEVERAL JOBS

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The other application involves several on-line programs and several detached programs processing the same area of information. The detached programs are responsible for handling one function common to accessing the data. For example, the first detached program copies from a master file to a common scratch pad where the online programs manipulate the data. The second detached program copies the data back to the master file. The third and fourth programs provide some application dependent functions common to each set of records. A large resident common area could be used to store the scratch pad and the 'overlay'' feature that I discuss later could be used to address the pertinent portion of the scratch pad.

The smoothest path between RSTS/E and VAX / VMS just got smoother: there's a major new release of

ROSS/V

ROSS/V has always provided:

- the fastest way to bring up RSTS/E applications on the VAX.
- the only way to do RSTS/E development on the VAX.
- an extensive subset of RSTS/E monitor calls and standard RSTS/E features, like CCLs, DOS-formatted magtape, and RSTS/E-style file update mode.

Now, in Version 3, ROSS/V supports:

- the "hidden" RSX run-time system (with 32 KW job size).
- resident libraries.
- job spawning and detached jobs.
- spooling to VMS print and batch queues.
- mailbox send/receive for communication with VAX-11 BASIC and other native mode applications.

How ROSS/V works:

ROSS/V is written in VAX-11 MACRO, and RSTS/E monitor calls are performed in VAX native mode. The rest of your PDP-11 code (in applications, run-time systems, TKB, etc.) is executed directly in the PDP-11 microcode that's present in every VAX. ROSS/V runs under VMS, not in place of it. Thus, some users may be working under the RSTS/E subsystem provided by ROSS/V while others are concurrently using any of the other VAX/VMS capabilities.

Call or write for the new ROSS/V technical summary, which describes all of ROSS/V's features.

ROSS/V

Evans Griffiths & Hart, Inc. 55 Waltham Street

Lexington, MA 02173 (617) 861-0670

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Spokane, WA 99202 (509) 484-3400

CIRCLE 176 ON READER CARD

Each program only maps the part of resident common area that it is currently using.

OTHER POSSIBLE USES

Another use of a large resident common area is that of a "file handler." There are several sites that use file handlers so that it is not necessary to build the file access code into each program (i.e., RMS). As a result, they use a file handler and inter-job messages to access the data. A file handler could use a resident common area in two ways:

 store the file (or portions of it) in its own private resident common area. This amounts to a more sophisticated data caching technique. It also points out another interesting characteristic of resident common areas and that is they do not have to be shared; they can be used by one program exclusively.

• passing the records back and forth between the file handler and

the rest of the application.

Scratch Pad in Resident Common Area

| Server Pgm 1| | Cont

| Control Area | | Slot #1 |

COMMON AREA COOKBOOK

I will now discuss the steps necessary to build a resident common area suited to your requirements.

• The first thing I must do is address some language dependencies.

RPM CAN DRAMATICALLY INCREASE THE PERFORMANCE OF YOUR SYSTEM

You can double or triple your system's performance without adding costly hardware. RPM provides everything you need to optimize your system and keep it running at peak performance.

RPM analyzes your system performance automatically, identifying problem areas. But, RPM doesn't stop there It identifies the cause of the problems and makes suggestions for correction in plain English.

AUTOMATIC PERFORMANCE ANALYSIS

Instead of dumping columns of cryptic numbers, RPM gives you a plain English report that describes how your system is performing. It tells you where you have problems, what caused the problems, and how to fix them.

This report analyzes each of the resources that are common problem areas. Any resources that are not being used optimally are identified. The program or files that caused the problem are then identified and suggestions are made for correcting the problem.

DETAILED PROGRAM ANALYSIS

In addition to identifying problem programs, RPM can analyze the operation of individual programs, identifying problem areas.

RPM's detailed program analysis breaks the operation of the program down into CPU usage, input/output and system calls. Usage and directory overhead counts are displayed by channel and by system call.

File P	rocessor (FIP)	Usage			
The file processor (FI least one job 67% of t are waiting to use FIF time, it is waiting fo use. The disk informa 91% Directory.	he time. In ad Although the r information f	dition, an a file process rom the disk	verage of 2.7 sor is in use 72% of the ti	other u 67% of t me it is	sers he in
	ting to the dir	ectory struc	ture.		
7% Disk Allocation					
Reading or wri 2% Monitor Overlays	ting informatio	n about free	blocks on the	disk.	
	r overlays into	memory.			
0% Miscellaneous.					
	ache informatio			data.	
The five programs that OE047A 275.2 PAYREC 213.1 TKB 158.0	use the most F	IP resources	are:		
BP2COM 110.3					
LOGIN 78.7 FIP usage can be decre	and he antimin	ing the alus	tonging of fug	augent lu	unod
files (see section 4.4 where possible, reorde	.1 of the RPM L	ser's Guide)	, using contig	uous fil	es
PM> EXAMINE JOB 5 EVERY	5 MINUTES				
ob: 5 Program: *ALL*	CPU: 54%	Sample Time:	297 Seconds		
han File	Count Ovrhd	Chan	File	Count	Ovrh
0 KB2:	3	1 * Closed		4	
3 DM1: [1, 3] CSPCOM. OLB		4 DM1: [5, 1		39	4
6 DM1 . 15 11 EXAMPL STR	9 91	7 DM1 - 15 1	IFYAMPI TSK	799	

	1: [1,	3] TKB	. TSK		278					1011	10	
EMT	Count	Ovrhd	EMT	Count	Ovrhd	EMT	Count	Ovrhd	EMT	Count	Ovrhd	
CALFIF	9 149		. READ	1441	2	. WRITE	460	133	. CORE	7		
. TTECH	1 1		. TTRST	r 1		. DATE	1		. NAME	2		
. RTS	1		. LOGS	12		. CLEAF	2 2		. CCL	1	5	
. FSS	42		. UUO	71		RSX	1		CLSFQ	62	28	
OPNFQ	58	535	CREFQ	2	226	DLNFQ	1	23	RSTFQ	2		
LOKFQ	22	259	CRTFQ	1	64	CRBFQ	1	175	UU. ATH	R 70	250	
UU. NAM	1 1	3										

EXTENDED PERFORMANCE STATISTICS

RPM adds extended performance statistics to your RSTS/E monitor. This monitor extension captures information about overall system performance plus information on individual programs and files. Using this information, you can improve system performance by minimizing disk head movement, reducing file processor (FIP) waits, reducing swapping, and optimizing disk cache operation.

DYNAMIC PLOTTING

In addition to comprehensive reports, **RPM** can generate a wide variety of graphs. It can plot curves, draw bar charts, and plot histograms using any combination of system information. By plotting one variable against another, you can immediately see correlations between variables. This is espescially useful for determining critical resource usage points.

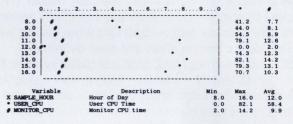
IDENTIFY PROGRAMS BY RESOURCE USAGE

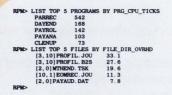
With **RPM**, you can identify the programs that will provide the most benefit from optimization. You can determine which programs are used most often and which programs use the most of critical resources.

Once identified, these programs can be optimized. Overall system performance can be increased by making changes only where they will do the most good.

	01234567	.890
DLO: SEEK DIST 1	*****	47.1
DLO: SEEK DIST 4	*******	22.9
DLO: SEEK DIST 16	*******	16.5
DLO: SEEK DIST 64	*****	10.8
DLO: SEEK DIST 100	**	2.7
PM> HISTOGRAM DLO		
RPM> HISTOGRAM DLO	_DISK_USAGE 01234567	. 8 9 0
	01234567	
DLO:DISK_USAGE_0	01234567	2.1
DLO:DISK_USAGE_0 DLO:DISK_USAGE_10	01234567 •••	2.1
DLO:DISK_USAGE_0 DLO:DISK_USAGE_10 DLO:DISK_USAGE_20	01234567	2.1 2.0 3.1
DLO:DISK_USAGE_0 DLO:DISK_USAGE_10 DLO:DISK_USAGE_20 DLO:DISK_USAGE_30	01234567	2.1 2.0 3.1 3.6
DL0: DISK_USAGE_0 DL0: DISK_USAGE_10 DL0: DISK_USAGE_20 DL0: DISK_USAGE_30 DL0: DISK_USAGE_40	01234567	2.1 2.0 3.1 3.6 17.6
DLO:DISK_USAGE_0 DLO:DISK_USAGE_10 DLO:DISK_USAGE_20 DLO:DISK_USAGE_30 DLO:DISK_USAGE_40 DLO:DISK_USAGE_50	01234567	2.1 2.0 3.1 3.6 17.6 42.2
DL0: DISK_USAGE_0 DL0: DISK_USAGE_10 DL0: DISK_USAGE_20 DL0: DISK_USAGE_20 DL0: DISK_USAGE_40 DL0: DISK_USAGE_50 DL0: DISK_USAGE_50 DL0: DISK_USAGE	01234567	2.1 2.0 3.1 3.6 17.6 42.2 19.8
DLO:DISK_USAGE_0 DLO:DISK_USAGE_10 DLO:DISK_USAGE_20 DLO:DISK_USAGE_30 DLO:DISK_USAGE_40 DLO:DISK_USAGE_50	01234567	2.1 2.0 3.1 3.6 17.6 42.2

NO PLOT USER_CPU, MONITOR_CPU BY HOUR





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• Secondly, I'll talk about simple resident common areas.

Lastly, I'll describe overlaid resident common areas.

LANGUAGE DEPENDENT BACKGROUND

Your choice of high-level language will affect the remainder of the discussion. This is a result of the differing ways that the languages generate common data references. The remainder of this discussion centers on the use of BASIC+2, although it may also be applied to F4P and DIBOL.

BASIC + 2 shares data through MAPs or COMmons. The output from the compiler contains program sections (PSECTs). Each MAP/COMmon has one PSECT that has the same name as the MAP/COMmon name. For example:

100	Map	(ADATA)	
		First\$	= 10%
	,	Second\$	= 10%

produces a PSECT named ADATA. The task builder (TKB) builds tasks by combining all references to a PSECT within a segment. This characteristic is important when dealing with overlaid resident common areas and their definitions.

DIBOL, on the other hand, produces global symbols for each of the COMMON variables. This makes using the overlaid resident common areas extremely easy to use in DIBOL, in that the task builder resolves all offsets. However, I haven't run into very many DIBOL shops. If there are those of you who are interested, feel free to give me a call.

FORMS OF SHARING

I have experimented with two different forms of resident common areas:

• non-overlaid resident common area (mapped using a single static window into the resident region),

• overlaid resident common area (mapped using a dynamic window(s) into a LARGE resident region),

and found several differing characteristics between the two forms. The appropriate use of either must be determined at design time.

NON-OVERLAID RESIDENT COMMON AREA

The non-overlaid resident common area is the easiest form to use. The fundamental sequence to generate one is:

1. Define it — determine which PSECTs and/or global symbols are the ones to be shared and design the resident common area.

2. Create allocation for it — using a quicky MACRO program (3-4 lines usually, examples to follow).

3. Build it — using the task builder with resident library switches. (See the examples and Chapter 7 of the RSTS/E Task Builder Manual.)

4. Prepare for installation — run the output task through MAKSIL.

5. Install the resident common area mapped Read/Write.

To use a non-overlaid resident common area, you take the following steps:

1. Create your high-level language program(s).

2. Task build using the new resident common area (see Options in TKB manual, Chapter 3.2).

3. Run your programs.

The first set of steps is done only once. After that, using it is as easy as falling out of bed.

DEFINE A COMMON AREA

Let's assume that we have decided to share two record/common areas. We will call them RESRC1 and RESRC2. Their lengths are 1000 bytes and 2500 bytes respectively. We will call the common area RESDAT and use account [1,53] for the development. These names and sizes are examples only and can be anything that you want them to be.

RESDAT Resident Common			
++	Base	+	350
RESRC2 Mapped area			
2500 bytes long			
++	Base	+	100
RESRC1 Mapped area			
1000 bytes long			
++	Base		

where: BASE is on an even 8k byte boundary.

CREATE IT

The first step is to write the MACRO program, thusly:

.TITLE RESDAT		; Looks pretty on listings.
.IDENT /01KI/		; For those of you who keep track of ; versions.
.RADIX 10)	: So that you do not have to count in ; octal.
.PSECT F	RESRC1,RW,	, D,GBL,REL,OVR
		; The attributes match those
		; generated by B+2 and are
		; described in TKB manual (section
		; 4.2.1). The name must match the
		; COMmon/MAP name that you plan
		; to use in the program.
.BLKB 1	1000	; Reserve 1000 bytes
.PSECT F	RESRC2, RW,	D,GBL,REL,OVR
.BLKB 2	2500	; Reserve 2500 bytes
.END		; Required by MACRO

BUILD IT

Next, you assemble this module (MAC RESDAT = RESD AT) and task build it as follows:

TKB

TKB > RESDAT/-HD/PI,,RESDAT = RESDAT TKB > / ENTER OPTIONS: TKB > PAR = RESDAT:0:0 TKB > STACK = 0 TKB > //

Notice the following points about the task build.

• The task image has NO header and IS position independent.

- NO header is mandatory for resident libraries.
- Position independence is recommended.

• No map was called for. You may if you want, but there is nothing much there.

• A symbol file WAS called for. This is required for MAKSIL.

• The PAR option specified no fixed address range (0:0). This is the preferred way, because the task builder can then fit it in more easily to your program later.

• The stack was set to zero. Remember that this is a shareable thing. Stacks exist in your own job area.

PREPARE FOR INSTALLATION

The next step is to run MAKSIL and create the resident library, sort of like this:

RUN \$MAKSIL

AKSIL V07 . . .

Resident Library Name ? RESDAT

Task-Built Resident Library Input File < RESDAT. TSK > ?

Include symbol table (Yes/No) < Yes> ?

Symbol Table Input File < RESDAT.STB> ?

Resident Library Output File < RESDAT.LIB> ?

RESDAT built in 2 K-words, 0 symbols in the directory RESDAT.TSK renamed to RESDAT.TSK < 40>

INSTALL IT

Finally, you add the resident common area, using UTILTY. Up to this point, all functions may be performed under a non-privileged account. This is the ONLY step that must be done using a privileged account.

RUN \$UTILTY

UTILTY V07 . . .

ADD LIBRARY [1,53]RESDAT < 0 > / ADDR:???/RW
EXIT

Two points to notice:

1. The use of an explicit protection code.

The default protection code is <42> when you add a library. That implies Read access for everyone, but no Writing. You must supply a valid protection code in the ADD statement despite the protection code that is on the resident library disk file. Now, you set any protection code that you want to protect it at the user, group, or system level. Even privileged programs will get "?Protection violation" if they try to access a library that does not allow them write access. I suggest that you use a protection code of <0>.

2. The use of the /RW switch.

The default mapping of the resident library is Readonly. You must specify that the resident common area is mapped Read/Write or the memory management hardware might take it to mind to zap you, if RSTS/E doesn't get you first.

FIRST STEP FOR USAGE

You now have a resident common area ready to use. The following BASIC + 2 example will show just how easy it is to really use it.

Мар	(RESRC1)! Use the same PSECT or ! MAP name, please Cust.ID\$(199) = 5%! 200 customers
Мар	(RESRC2)! Second record Prod.Code\$(99) = 5% Prod.Desc\$(99) = 45%

1000

I

MAIN - LINE STARTS HERE

The program itself is standard BASIC + 2 code. The two stipulations you have are that:

• The MAP names must match the PSECT names in the MACRO module.

• The length of your MAPs CANNOT exceed the allocated space without incurring some problems. I'll leave it up to you to discover them. They quickly become self evident.

SECOND STEP FOR USAGE

In the command file, you add ONE line to use the resident common area.

		Sample Command File
		SY: PROGA, SY: PROGA=SY: PROGA/MP
		UNITS=12
		ASG=SY:5:6:7:8:9:10:11:12
		RESLIB=LB:BASICS/RO
Additional line	->	RESCOM=[1,53]RESDAT/RW
		EXTTSK=1024
		11

Use TKB to build your program. There should be no error messages issued by TKB.

If there are, you've messed up somewhere. You should check the map file and make sure that the starting address of the first PSECT (RESRC1 in this example) is starting on a high-up 8kb boundary. Since the starting addresses are in octal, the thing to remember is that 8kb is 20000 (octal) bytes. So you should see a line in the memory allocation synopsis that looks like this:



RESRC1(RW, D, GBL, REL, OVR) 120000 001750 001000.

```
This is the starting address-+
```

Now, run your program. This is the end of the simpler version of resident common.

OVERLAID RESIDENT COMMON AREA

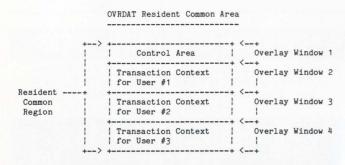
The term "over-laid resident common area" is loosely used. The data is not physically overlaid, only the Active Page Registers are modified to map the correct section of data into your virtual address space. All the data resides physically in memory. I use the term overlaid common because it is mapped using the "auto-load overlay mechanism." The concept of building overlayable resident common areas came to me from the combination of two references in the task builder manual. Knowing how to use them has come from a lot of hacking. You can check the TKB references yourself:

• using the .NAME directive to create loadable segments that contain no executable code (Chapter 5, pp. 5-11 to 5-13),

• using the auto-load/overlay mechanism to dynamically remap your window in a resident region (Chapter 7).

SAMPLE USE OF OVERLAID COMMON

Let's assume that we have an application that is a transaction processor. We want to have several common modules that perform a given function for the whole application. For example, a screen manager to handle the screen traffic, several file handlers to handle the RMS traffic and a transaction processing manager to schedule all and sundry activities. As a result, we have many common areas of the same layout, one for each of the currently active users. It is possible to use a scratch pad file or send around a lot of large messages. Both of these approaches involve significant overhead (one on disk accesses and the other on small buffers). Instead let's design a resident common area that looks like this:



The usage of the different areas and sizes of each may be different. This example uses two types of areas and two different sizes. The control region would contain the information used in controlling the remainder of the areas (e.g., areas in use, areas free) and perhaps some application data (e.g., current business date, application status). Each of the context areas would contain the same format for the data area but each would contain data unique to a given user. Then, when a program was servicing user #1, it would "overlay" into the transaction context area for that user, and so on for each user being serviced. In this fashion, a program could easily service many users.

Now, I realize that the example of the application is rather complex, but I have made it so to deliberately show possible uses. There are many more complex and many less complex. Actual uses I will leave to the reader (that's you).

MAKING IT

The fundamental sequence to generate an overlaid resident common area is similar to generating a non-overlaid resident common area:

1. Define it — determine what data is to be shared and design the overlay structure.

- for BASIC + 2 use MAPs/PSECTS
- for DIBOL use variables/Global symbols
- 2. Create it using several MACRO programs.
- 3. Build it using an overlay descriptor file.

4. Prepare for installation — by running the output task through MAKSIL.

5. Install the resident common area mapped Read/Write.

To use an overlaid resident common area, you take the following steps:

1. Create your high-level language program(s) with the following characteristics:

- specifically defined MAPs (more on that later),
- calls to NAMEd segments to remap the areas as required.

2. Build your program using the new resident common area.

3. Run your program.

As before, the first set of steps is done only once.

The three points that are CRITICAL to the successful completion of this portion are:

1. setting up the MACRO program correctly and sneakily to overcome TKB "features,"

2. specifying the correct overlay description so that your data does not get muxed ip, and

3. coding your program so that you do not tamper with the wrong things.

DEFINING A RESIDENT COMMON AREA

Using the previously described example, we will design an overlaid resident common area with a zero length root and two types of segments.

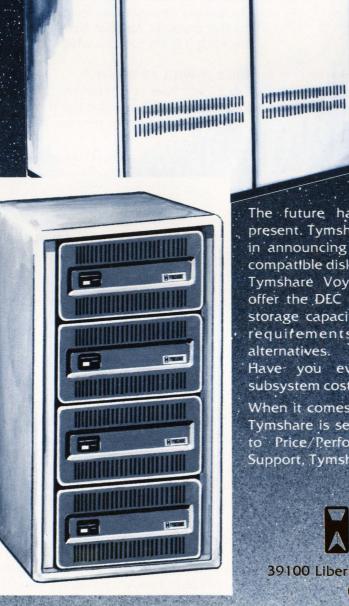
1. The first segment type is the control area and it exists once. The size of the segment is 1000 bytes and it contains data about the availability of the other three slots and who owns them, etc.

2. The second segment type is the transaction context area and it is duplicated three times, once for each active user. The size of this area is 8192 bytes. It contains three major sections:

- a screen buffer containing the current data from the entire screen,
- a workspace that contains the current records and other essential variables for this transaction,
- and a journaling area for record updates.

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CREATE IT

First we write a MACRO program to define and allocate the space for the control area. It looks something like this (explanations follow details, like in a mystery):

```
.TITLE OVRCTL ; Keeps it separate from the
; other example.
.IDENT /01KI/
.RADIX 10 ; for decimal phreaks
.PSECT DUMMY,RW,D,GBL,REL,CON
.BLKB 448
.PSECT DUMMY,RW,D,GBL,REL,CON
.BLKB 448
.PSECT DUMMY,RW,D,GBL,REL,CON
.BLKB 104
.END
```

Additionally, we describe a MACRO program to allocate the space for the transaction context area, thusly:

```
.TITLE OVRTXN ; Pretty, eh?
.IDENT /01KI/
.RADIX 10
.PSECT DUMMY,RW,D,GBL,REL,CON
.BLKB 448
.PSECT DUMMY,RW,D,GBL,REL,CON
.BLKB 448
. repeat for a total of 18 times to
. get 18*448 bytes (or 8054)
.
.PSECT DUMMY,RW,D,GBL,REL,CON
.BLKB 78
.END
```

The last MACRO program is required to give a you a zero length PSECT to use in BASIC + 2. It's the simplest.

.TITLE NULROT .IDENT /01KI/ .PSECT OVRDAT,RW,D,GBL,REL,OVR .END

Looking at these examples, the first question that comes to mind is "What the @#%1@#?". More lucid questions are:

• Why repeat the PSECT names with all those unusual numbers?

Why have a zero length PSECT?

CREATION EXPLAINED

Answers: First of all, I don't know "What the @#%1#?"

• TKB gets very upset about allocating empty space in overlaid segments that have the GBL and NODSK attributes. What this boils down to is that an overlay segment in an overlaid resident common has:

- the GBL and NODSK attributes set automatically for each overlaid segment. NODSK states that there is no loading of data/code that occurs when the segment is invoked. So in practice, the autoload mechanism will map the window but will not load from disk. This is exactly what we want.
- to have space allocated so that we can access the required addresses when we reference the data. Combine this with the fact that TKB does not allow

pre-allocation of empty space and we have . . .

- an interesting opportunity, AND
- a tacky solution. TKB will not complain IF the PSECT has the CONcatenate attributes AND each of references does not exceed 448 bytes (340 octal words). Therefore, you will notice that the attributes for DUMMY are not OVeRlaid.

• A zero length PSECT is only required for BASIC+2. BASIC+2 generates a PSECT with the same name as your MAP. The PSECT is then the common point of reference. TKB will not store PSECT names from overlay segments in the symbol table. Therefore, when you build a resident common you get all PSECTs from the root, but none from the overlay segments and as a result it must be faked. The solution is to define a zero length root with one PSECT name and use TKB's inbred intelligence to overcome the small discrepancies in length (amidst many diagnostic messages).

An aside for DIBOL users, you do not have to go through the dummy PSECT route. You may define your common variables as global symbols at the correct offset. The naming convention of the global symbols is such that the five character variables have a dollar sign appended or preceding it (I am not sure which). Experiment.

BUILD IT WITH AN ODL FILE

If I still have your attention at this point . . . let's get onto the ODL file. Remember this is the second critical point (but not as messy as the previous one):

	.ROOT	NULROT-*!(AREAO, AREA1, AREA2, AREA3)
	.NAME	CTRL00,GBL
	. NAME	SLOTC1,GBL
	. NAME	SLOT02,GBL
	. NAME	SLOT03, GBL
AREAO:	.FCTR	CTRL00-OVRCTL
AREA1:	.FCTR	SLOT01-OVRTXN
AREA2:	. FCTR	SLOT02-OVRTXN
		SLOTO3-OVRTXN
	. END	

and we need the accompanying command file:

```
SY:OVRDAT/-HD/PI,OVRDAT,OVRDAT=SY:OVRDAT/MP
PAR=OVRDAT:0:0
STACK=0
GBLREF=CTRL00,SLOT01,SLOT02,SLOT03
//
```

Task build this arrangement. There should be no errors. Errors that I have encountered are:

• TKB — FATAL — I/O ERROR ON OUTPUT FILE — this very informative message says clearly that the preallocated space in one of your PSECTs in an overlay segment is too big. Reduce the size of the .BLKB directive.

STEPS 3 & 4

You then run it through MAKSIL as in the previous case (just substitute OVRDAT for every occurence of RESDAT).

You must now add the library using UTILTY. The notes mentioned in the non-overlaid common also apply here.

This is the end of creating an overlaid resident common. Now, aren't you glad that it needs be done only once?

Software Packages for VAX/VMS and RSTS/E.

Resource Accounting, Auditing and Billing VAX/VMS and RSTS/E.

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RABBIT-4 is a stand-alone data file security system that operates under RSTS/E Version 7. RABBIT-4 prevents access to classified or confidential data fields by unauthorized personnel. Even priviledged users may be excluded from data files secured by RABBIT-4. Coupled with normal security measures, RABBIT-4 will ensure file integrity while monitoring all file access attempts.

RABBIT-4 "locks-up" designated data files from would-be intruders. All attempted file violations are logged and available for analysis. Computer operations and security management may be immediately notified of also easy to install. After loading the RABBIT-3 tape or disk, just answer a few questions to tailor the system to your needs. In less than an hour of effort, your RABBIT-3 will be generating complete, detailed user information.

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RABBIT-3

Daily Disk Catalog

- generates disk accounting information for each user.
- Security Tracer Record
- provides step by step security information of user activities.
- Disk Critical Alert
- alerts operator that free disk is below efficient levels.

the file violation attempt, and the offending job rolled out or killed. Under RABBIT-4 maximum security, the RSTS operating system may be disabled until the intrusion is cleared.

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For more information on the RABBIT system contact

Telephone: (404) 955-2553 TWX: 810-766-2256

RABBIT-4

RABBIT-4 creates and maintains a complete log of file access and violation attempts. Included in the log is: date, time, job, program, project-programmer and keyboard. The log of file accesses is available at all times to assigned security management. A recap report of violations and access attempts may be run at will.

RABBIT.4 is written in PDP-11 macro to ensure optimized performance. The logged data is in ASCII stream format for prompt reporting. RABBIT.4 is controllable through OPSER and has an optional system crash recovery capability.

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CIRCLE 21 ON READER CARD

as

USING THE EXAMPLE

In order to use this common area, we will create a program that references the resident common area and the overlaid segments. The only requirements are that:

• the correct segment is overlaid (by calling it) prior to reference,

- in BASIC+2, the MAP name that is used must:
 - use the same name as the zero length segment in the root of the resident common,
 - use the FILL variable to offset individual record definitions into the overlaid segment.

STEP 1

The sample data description portion of a program may look like this:

Points to notice:	Sample BASIC+2 Program			
MAP names must be +->100 identical	Map (OVRDAT) ! Same for all areas Contrl.Name\$(2≸) = 6≸ , Contrl.InUse≸(2≸) , Contrl.Status≸(2≸)			
+->200	Map (OVRDAT) Slot.Screen\$ = 1000%			
	, Slot.Status% , Slot.Etc%			

Now the code part of the program is almost normal, but notice the calls to the NAMEd segments to force the remapping of the OVRDAT map.

1000 ! Main - line

...

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1100) ! Access Control Area
CALLs remap to new +->	Call CTRLOO ! Remap to control area
segment	For $I_{1}^{*} = 0_{1}^{*}$ to 2_{1}^{*}
	If Contrl.InUse%(I%)=0%
	Then
	User $3 = 13 + 13$
	Contrl.InUse%(I%)=-1%
	Contrl. Name\$(I%) =
	Next I%
	NEXC 18
1	1 torong a terrestion plat
	! Access a transaction slot
+->	Call SLOTO1 If User# = 1%
+->	Call SLOTO2 If User% = 2%
+->	Call SLOTO3 if User% = 3%
	If (Slot.Status% and 1%) <> 0%
	Then

STEP 2

Compile your program and create a command file that has an additional line in it (or add the additional line if you use the dialogue mode of TKB):

At this point, you should get several diagnostic messages from the task builder. There will be one for each different map statement in the source modules. The message informs you that the PSECT reference conflicts with the PSECT in the library OVRDAT. This is to be expected, because the library was built using the PSECT named OVRDAT with a zero length associated with it, and you have just said that it is 8192 bytes long. Now what happens is that you get your task built properly, because TKB uses the starting address of the resident common as the starting address of your PSECT and as a result your variables are now stored in the resident common. If you want further details feel free to give me a call.

ELEPHANT TRAPS

There are several areas that you may term "elephant traps," because they are large holes that you can fall into and never be heard from again. They are the following:

- Remember it's shared, someone can get your MAPs,
- "Sophomores" (semaphores) are essential,
- Beware of single threaded bottlenecks,
- Check overlay segment length carefully,
- You lose chunks of address space (APRs).

SHARED DATA

In the usual programming environment, you know that when you put information into a variable, it doesn't change until you put something else in it. When you share data through resident common areas, remember that all variables in the common area can be changed by another program when your back is turned. Therefore, the protocol for using resident common must be carefully defined prior to implementation. Two possible forms are:

- use of flags in a control area
- use of slots that are unique to one user/transaction.

USE OF SEMAPHORES

You should consider carefully the correct use of semaphores for one program to be able to signal another program about significant events. RSTS/E does not have many system services for allowing different jobs to synchronize their activities. There is Send/Receive (with sleep) or forced data into psuedo keyboards that will wake up a sleeping job. Otherwise, you have to design a flag setting and sleep n seconds arrangement between jobs. One thing to be VERY wary of is SUSPENDing and RESUMEing jobs. It is easy to have both jobs SUSPENDed at the same time.

SINGLE THREADED BOTTLENECKS

DO NOT accidently create another FIP. A single threaded server is a program that will process one request to completion before processing the next request on the queue. When there are only two or three users of disk intensive server that is single threaded, you will see rapid system degradation. Therefore, design any single threaded server with a careful eye to the amount of disk activity that it must perform to finish a request.

OVERLAY SEGMENT LENGTH

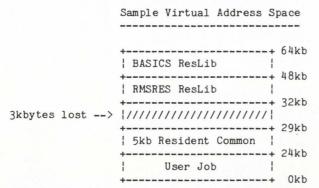
One problem in using overlaid segments with BASIC + 2 is that you can not be sure that your MAP does not exceed the actual length of an overlay segment. This is because of the deliberate introduction of a possible conflict that TKB warns you about. The choices are two, either:

• Make ALL overlay segments reserve 8.kbytes, which will rapidly inflate the physical memory requirements of a resident common region, sometimes even double, or

• Verify the length of a PSECT in your program's MAP against the length of an overlay segment in the resident common area.

LOSS OF ADDRESS SPACE

The PDP-11 allows you to divide your virtual address space into eight pages. Each page may be from 64 bytes to 8192 bytes in size. Each page also starts at an 8kbyte boundary. This feature is used by RSTS/E to separately load and control run time systems, resident libraries and programs that are all part of the same job. However, if you use only a portion of one page (e.g., a 5kbyte resident common area) you lose access to the remaining portion of the virtual addresses of that page (e.g., 3kbytes are unavailable). For example,

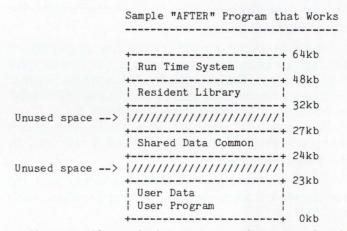


This implies that you have to design the resident common area balancing the advantage of shareable data against the possible loss of virtual address space.

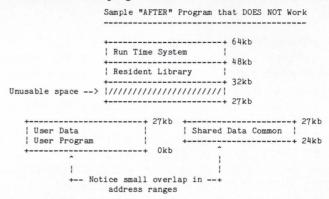
Sample "BEFORE" Program

	++	64kb
	Run Time System	
	++ Resident Library	48kb
	++	32kb
Unused address>	1//////////////////////////////////////	
Space	++	??kb
	User Data	
	User Program	
	++	Okb

This program will only be able to use a resident common if combined program size and non-shared data is less than the current available memory size minus the size (rounded up to 8kbytes) of the resident common area. For example, if the program is 26kbytes prior to splitting it into 3kbytes of shared data and 23kbytes words of program and non-shared data, then you could do it. It would look like this:



However, if you had a program that was already 30kbytes and you want to share 3kbytes of data you are left with 27kbytes of program and non-shared data. If your program currently uses 32kbytes worth of resident libraries (i.e., BASICS and RMSRES), you will have to resort to overlaying RMS (not a good thing) or use some other shoe-horning if you will insist on trying to use the resident common area.



On the other hand, if it allows you to cleanly break your program into two or three smaller pieces, maybe it is not a bad idea after all.

FINAL NOTE

JUST BE VERY CAREFUL ABOUT ALL THIS GOOD STUFF. YOU SHOULD LEARN HOW TO USE IT BEFORE YOU COMMIT YOURSELF TOO FAR. IT IS VERY USEFUL, BUT EVERYTHING HAS ITS COST.



A Column For The Advanced RSTS/E User

Steven L. Edwards, Software Techniques, Inc.

MONITORING FREE DISK SPACE

For production systems, the worst thing that can happen is that the operating system crash. The second worst thing that can happen is that you run out of disk space. This invariably happens five minutes before your all night batch job completes, or after your data entry operators have spent five hours entering transactions into a temporary file. The consequences are rarely convenient.

The ideal solution would be for the operating system to stall that job until the system manager can free up some more disk space. The ideal solution is beyond the scope of this column. However, we can set up a job that will monitor the free disk space, and do something about it.

The program presented here allows the system manager to establish a 'yellow' limit and a 'red' limit for each disk device on the system.

When the free space falls below the 'yellow' limit, the program sends a message to the operator's service console via OPSER.

When the free space falls below the 'red' limit, the program sends a message to the operator's service console. In addition, the program then suspends all jobs, and sets all active terminals to 'SPEED O' to prevent unintended typeahead. The program then suspends itself. When the system manager frees up sufficient disk space, s/he can resume this program which will restore all of the jobs it tampered with.

Please note the following:

1. This program has only been tried on two systems, both of which had two RMO2 type disks. Be prepared to fix any bugs you may uncover.

2. Entering a value of 0, or just typing < RETURN> to a limit question effectively disables that limit.

3. This program assumes job 1 is ERRCPY and job 2 is OPSER. It does not suspend jobs 1 and 2.

4. This program assumes that you do not want to set KBO to 'SPEED O' in case a 'red' limit is exceeded.

5. If you have malicious users, do not establish 'red' limits for any disk they can create files on.

6. If your programs extend files that have clustersizes greater than the pack clustersize, you must ensure that there are free clusters of the size required above the limits set.

7. This program will not catch directory extension failures.

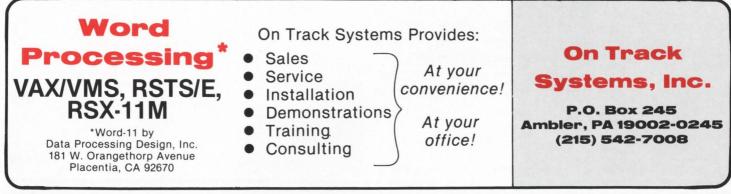
1	Title:	FREMON	
	Version:	Y0101A	
	Edit date:	02-MAY-1983	
	Written by:	STEVEN L. EDW	ARDS
	Date:	17-JAN-1983	
	Package:	In-House	
	Description:	MONITOR FREE	DISK SPACE
11	Copyright (C. Software Tech Los Alamitos, Title to and in Software 1	nniques , CA 90720 d ownership of th	e software shall at all times remain
			ocument is subject to change without onstrued as a commitment by Software
			ed and Software Techniques has no his time, unless stated elsewhere in
01	Modif	`ication History	
	Ver/Edit	Date	Reason (Who)
11	X0101A	17-JAN-1983	Initial conception. (SLE)
	X0102A	02-MAY-1983	CHANGE THE SIZE OF THE LIMIT ARRAY CHANGE METHOD OF DETERMINING WHEN ' SEND A MESSAGE TO OPSER. MINOR CODE CLEAN-UPS. (SLE)
	Y0101A	02-MAY-1983	PROMOTE TO GROUP SUPPORT. (SLE)
00!	Progr	am Description	
011	THIS SYSTEM.	PROGRAM MONITORS	THE FREE SPACE ON ALL DISKS ON THE

when the free space falls below the 'yellow' limit, the program sends a message to the operator's service console via opser.

WHEN THE FREE SPACE FALLS BELOW THE 'RED' LIMIT, THE PROGRAM SENDS A MESSAGE TO THE OPERATOR'S SERVICE CONSOLE VIA OPSER. THE PROGRAM THEN SUSPENDS ALL JOBS AND SETS ALL ACTIVE TERMINALS TO SPEED O TO PREVENT UNINTENDED TYPE-AHEAD. WHEN THE OPERATOR FREES UP SUFFICIENT DISK SPACE, S/HE CAN RESUME THIS FROGRAM WHICH WILL RESTORE ALL OF THE JCBS IT TAMPERED WITH.

5001		
1	Compile time variables	
1		
501	.DEFINE .NAME\$ = "Fremon"	
1	.DEFINE .VERSION\$ = "Y01.01A"	1Y0101A
1	.DEFINE .CHAN.KB% = 1%	
1		
1	Program Name.	
1	Program Version.	
!	Channel number for terminal I/O.	
900!		
1	Dimension Declaration	
1		
1 901	-929 local dimension declarations	
! 930	-949 library dimension declarations	
! 950	-979 MAP statements	
901	DIM	
	PRIORITY%(63%)	
	.RED LIMIT(80%)	1X0102A
	TERMINAL SPEED\$(128\$)	
	YELLOW_LIMIT(80%)	1X0102A
1		
1	JOB PRIORITY ARRAY.	
1	RED LIMIT ALARM ARRAY.	
1	TERMINAL SPEED ARRAY (RECEIVE, TRANSMIT).	
1	YELLOW LIMIT ALARM ARRAY.	

902	MA P	(UUTE1) !	DLOCK DOD NONTROD WADLING DADW T		2010	UNTCNT_POINTER\$ = UNTCNT\$
		I ,DATA	BLOCK FOR MONITOR TABLES PART I		>	FOR DEVICE_INDEX% = 0% TO (DEVOKB% - 2%) STEP 2% FOR UNIT_NUMBER% = 0% TO PEEK(DEVCNT% + DEVICE_INDEX%)
		JOBX2_FILL\$! JOB NUMBER TIMES 2 / FILL BYTE			UNIT_STATUS# = PEEK(UNTCNT_POINTER#)
			I MAX KB NUMBER / MAX JOB NUMBER I TABLE OF MAX UNIT NUMBER		1	GOTO 2020
		, DEV PTR%	I TABLE OF POINTERS TO DEVICE DDB'S			IF ((UNIT_STATUS# < 0%) OR ((UNIT_STATUS# AND 4096%) <> 0%))
		,FILL%	! FILLER		\	DSK\$ = CVT\$\$(SWAP\$(PEEK(DEVNAM\$ + DEVICE_INDEX\$)))
		,JOBTEL% ,FILL\$ = 6%	! JOB TABLE ! FILLER		1.1.1.1.1.1.1.1	+ NUM1\$(UNIT_NUMBER\$) + ":"
		.UNTCNT#	I TABLE OF UNIT STATUS / OPEN COUNT		<u>`</u>	SUB_SCRIPT\$ = (DEVICE_INDEX\$ * 4\$) + UNIT_NUMBER\$ PRINT #.CHAN.KB\$, "Enter YELLOW limit for "
		,SATCTL%	! TABLE OF FREE BLOCK COUNTS (LSW)			+ DSK\$ + SP;
		,FILL% ,SATCTM%	! FILLER ! TABLE OF FREE BLOCK COUNTS (MSW)		`	INPUT #.CHAN.KB\$, YELLOW_LIMIT(SUB_SCRIPT\$) PRINT #.CHAN.KB\$, "Enter RED limit for "
		,FILL\$ = 4%	! FILLER		`	+ DSK\$ + SP;
	MAD	(UUTE1) !	1		1	<pre>INPUT #.CHAN.KB\$, RED_LIMIT(SUB_SCRIPT\$)</pre>
(MAP		THE MONITOR TABLES I DATA BLOCK		2020	UNTCNT_POINTER# = UNTCNT_POINTER# + 2#
			ABLE TO ACCESS ENTIRE THING AT ONCE		1	NEXT UNIT_NUMBER#
		! UUTB1\$ - 30\$! UUTB1\$ = SYS(CHR\$(6\$) + CHR\$(-3\$))		``````````````````````````````````````	NEXT DEVICE_INDEX#
		001014 - 500	I		1	MAKE SURE WE ARE USING THE CORRECT POINTER VALUES.
1	MAP	(UUTB2) !	DLOCK FOR MONITOR TARLES DART IT		1	LOOP THROUGH THE MONITOR'S TABLES OF DEVICES.
		I DATA	BLOCK FOR MONITOR TABLES PART II		!	LOOP THROUGH ALL UNIT NUMBERS. SKIP DEVICES THAT ARE NOT MOUNTED OR ARE NFS.
		FILL\$ = 4%	! FILLER		i	CONSTRUCT THE DEVICE NAME.
		, DEVNAM% , FILL%	! DEVICE NAME TABLE ! CSR TABLE		!	CALCULATE THE PROPER SUBSCRIPT.
		, DEVOKB%	! NUMBER OF DISKS TIMES 2		1	ENTER THE YELLOW LIMIT FOR THIS DEVICE. ENTER THE RED LIMIT FOR THIS DEVICE.
		,FILL\$ = 10% ,LOGNAM%	I FILLER I TABLE OF SYSTEM LOGICAL NAMES		1	BUMP THE POINTER.
		,FILL\$ = 8%	I FILLER		!	TRY THE NEXT UNIT NUMBER. TRY THE NEXT DEVICE.
			1			
1	MA P	(UUTB2) ! ! REMAP	THE MONITOR TABLES PART II DATA BLOCK		2999	PRINT #.CHAN.KB%; CR + LF + "Detaching"
			ABLE TO ACCESS ENTIRE THING AT ONCE			+ CR + LF
		! UUTB2\$ = 30%	! UUTB2\$ = SYS(CHR\$(6\$) + CHR\$(-12\$))			+ CR + LF
		001020 - 300	1 00102¢ = 010(010¢(0¢) + 0100¢(-12¢))			+ CR + LF + CR + LF
0001					\	CLOSE #. CHAN. KB%
9991 1		Start of Initia	lization			$TEMP_0 \$ = SYS(CHR\$(6\$) + CHR\$(7\$))$
!					i	TELL THE USER WHAT WE ARE ABOUT TO DO.
1000	ONERRO	R GOTO 19000			1	CLOSE THE TERMINAL CHANNEL. DETACH FROM THE TERMINAL.
1						DETROIT FROM THE TENTANE.
1	Set st	andard error trap	•		3000!	MATH BROCRAW LOOP
1010	PRINT		RSION\$ + HT + "Software Techniques"		1	MAIN PROGRAM LOOP
		+ CR + LF + "Fr UNLESS EO%	ee disk space monitor." + CR + LF			LOONAL DOTHERD LOONAL
1					3010	LOGNAM_POINTER% = LOGNAM% SATCTL_POINTER% = SATCTL%
1	Print	standard header o	n 'Run' entry.		1	SATCTM_POINTER# = SATCTM#
1030	UUTB1\$	= SYS(CHR\$(6%) +	CHR\$(-3%))		,	UNTCNT_POINTER\$ = UNTCNT\$ FOR DEVICE_INDEX\$ = 0\$ TO (DEVOKB\$ - 2\$) STEP 2\$
1		= SYS(CHR\$(6%) +			1	FOR UNIT_NUMBER% = 0% TO PEEK(DEVCNT% + DEVICE_INDEX%)
		% = CNT_KB_MAXCNT % = SWAP%(CNT KB	MAXCNT\$() AND 255\$		`	UNIT_STATUS# = PEEK(UNTCNT_POINTER#) GOTO 3020
1	OUR_JO	B_NUMBER% = (JOBX	2_FILL% AND 255%) / 2%			IF ((UNIT_STATUS# < 0%)
1	NULL 9	\$ = STRING\$(9%, 0	٤)	IX0102A		OR ((UNIT_STATUS# AND 4096#) <> 0#))
1	NULL_1	0\$ = STRING\$(10%,		IX0102A	`	<pre>DSK\$ = CVT\$\$(SWAP\$(PEEK(DEVNAM\$ + DEVICE_INDEX\$))) + NUM1\$(UNIT_NUMBER\$) + ":"</pre>
	PASS_I	NTERVALS = 30% COUNTS, YELLOW 1	\$, YELLOW_2\$ = 0\$	1X0102A 1X0102A		+ RAD\$(PEEK(LOGNAM_POINTER\$))
i					1	+ RAD\$(PEEK(LOGNAM_POINTER\$ + 2\$)) SUB_SCRIPT\$ = (DEVICE_INDEX\$ * 4\$) + UNIT_NUMBER\$
1		NITOR TABLES - PA			Ň	<pre>FREE_BLOCKS = PEEK(SATCTL_POINTER\$)</pre>
1		NITOR TABLES - PA E NUMBER OF TERMI	NALS ON THIS SYSTEM.		1	FREE_BLOCKS = FREE_BLOCKS + 65536.0 IF (FREE_BLOCKS < 0.0)
1		E MAXIMUM JOB NUM	BER.		Υ	<pre>FREE_BLOCKS = FREE_BLOCKS + 65536.0 * PEEK(SATCTM_POINTER\$)</pre>
1	GET OU	R JOB NUMBER.			1	IF FREE_BLOCKS < RED_LIMIT(SUB_SCRIPT\$) THEN GOSUB 5000
!	A FEW			X0102A		ELSE IF FREE_BLOCKS < YELLOW_LIMIT(SUB_SCRIPT%)
1		MORE NULLS. LIZE OUR PASS INT	ERVAL.	X0102A X0102A		THEN YELLOW_COUNT\$ = YELLOW_COUNT\$ + 1\$ GOSUB 4000
1		LIZE THE YELLOW M		X0102A	`	
1	Define	various variable	9		3020	LOGNAM_POINTER\$ = LOGNAM_POINTER\$ + 10\$ UNTCNT POINTER\$ = UNTCNT POINTER\$ + 2\$
						SATCTL_POINTER% = SATCTL_POINTER% + 2%
1110	OPEN "		PUT AS FILE #.CHAN.KB%		1	SATCTM_POINTER% = SATCTM_POINTER% + 2%
1		, MODE 0%			1	NEXT UNIT_NUMBER\$ NEXT DEVICE_INDEX\$
i	Open th	he terminal.			1	
2000!					!	MAKE SURE WE ARE USING THE CORRECT POINTER VALUES. LOOP THROUGH THE MONITOR'S TABLES OF DEVICES.
1		Start of MAIN			1	LOOP THROUGH ALL UNIT NUMBERS.
1					1	SKIP DEVICES THAT ARE NOT MOUNTED OR ARE NFS.



CIRCLE 13 ON READER CARD

	TRY THE	CONSTRUCT THE DEVICE NAME. CALCULATE THE PROPER SUBSCRIPT. GET THE NUMBER OF FREE BLOCKS ON THIS UNIT. CHECK FOR A RED LIMIT VIOLATION. CHECK FOR A YELLOW LIMIT VIOLATION. BUMP THE VARIOUS POINTERS. TRY THE NEXT UNIT NUMBER.	
3030	IF THEN ELSE	YELLOW_COUNT% YELLOW_1% = YELLOW_1% + 1% YELLOM_1% , YELLOW_2% = 0%	1X0102A 1X0102A 1X0102A
! ! !	IF THEN ELSE	WE GOT ANY YELLOW LIMIT VIOLATIONS, INCREMENT THE "DON'T BOTHER OPSER" FLAG, CLEAR THE FLAGS.	X0102A X0102A X0102A X0102A
3999 \ \		COUNT\$ = 0\$ ASS_INTERVAL\$ OO	1X0102A 1X0102A
!		OR A LITTLE WHILE. BACK TO TRY AGAIN.	
4000! ! !		YELLOW LIMIT EXCEEDED	
4010	IF THEN	YELLOW_2\$ = 0\$ MESSAGE\$ = "Yellow limit exceeded on " + DSK\$ + BEL + BEL	IX0102A IX0102A IX0102A
``	ELSE	GOSUB 10100 YELLOW_2\$ = YELLOW_1\$ YELLOW_2\$ = YELLOW_2\$ - 1\$	IX0102A IX0102A IX0102A
1	SEND A	MESSAGE (VIA OPSER) IF WE HAVEN'T LATELY.	X0102A X0102A
4999	RETURN		
1	BACK TO	MAINLINE CODE.	
5000! ! !		RED LIMIT EXCEEDED	
501C	MESSAGE GOSUB 1	E\$ = "Rec limit exceeded on " + DSK\$ + EEL + EEL COICO	
1	SEND A	MESSAGE (VIA OPSER).	
5020 \ \	TERMINA	IY\$(INDEX\$) = -1\$ FOR INDEX\$ = 1\$ TO MAXCNT\$ AL_SPEED\$(INDEX\$) = 0\$ FOR INDEX\$ = 1\$ TO CNT_KB\$ @BER\$ = 3\$	
! ! !	INITIAL	IZE THE PRIORITY ARBAY. IZE THE TERMINAL SPEED ARRAY. NT JOB 3 (ASSUMING JOB 1 IS ERRCPY AND JOB 2 IS OPSER).	
5030 \	TEMF_0% GOTC 50	:= PEEK(JOBTEL\$ + (JOB_NUMBER\$ * 2\$)) 40 F TEMP_0\$ = 0\$	
X.	GOTC 50	50	
X X		IF TEMP_0\$ = -1\$ Y\$(JOD_NUMBER\$) = PEEK(TEMP_0\$ + 28\$) AND 255\$ = SYS(CHR\$(6\$) + CHR\$(-13\$) + CHR\$(JOE_NUMBER\$) + CHR\$(255\$) + CHR\$(-12\$) URLESS (JOE_NUMBER\$ = OUR_JOE_NUMBER\$)	
	SKIP TH QUIT IF SAVE TH	: AEDRESS OF THE JOB DATA BLOCK FOR THE JOB. HE JOB IF II DOES NOT EXIST (0). WE ARE AT THE END OF THE JOB TABLE (-1). HE PRIORITY OF THE JOB.	
5040 \	JOB_NUM GOTO 50	⊞ER≴ = JOB_NUMBER≴ + 1≴ 30	
i	BUMP TH	E JOB NUMBER.	
5050 \ \ \ \		= PEEK(DEVPTR\$ + DEVOKB\$) EX\$ = 1\$ TO CNT_KB\$ TEMP_1\$ = PEEK(PEEK(TEMP_0\$ + (INDEX\$ * 2\$)) + 2\$) JOB_NUMBER\$ = TEMP_1\$ AND 255\$ GOTO 5060 IF ((JOB_NUMBER\$ = 0\$)	
1		OR ((JOB_NUMBER\$ AND 1\$) = 1\$)) UNIT_NUMBER\$ = SWAP\$(TEMP_1\$) AND 255\$ TEMP_0\$ = SYS(CHR\$(6\$) + CHR\$(-4\$) + CHR\$(UNIT_NUMBER\$) + CHR\$(19\$))	
1		$TEMP_0 \$ = SYS(CHR\$(6\$) + CHR\$(16\$) + CHR\$(0\$) + CHR\$(UNIT_NUMBER\$))$	
1		<pre>TERMINAL_SPEED\$(UNIT_NUMBER\$%) = (ASCI1(MID(TEMP_0\$, 14% OR SWAP\$(ASCI1(MID(TEMP_0\$, 16%, 1\$)))) TEMP_0\$ = SYS(CHR\$(6\$) + CHR\$(16\$) + CHR\$(0\$)</pre>	, 1≴))
		+ CHR\$(UNIT_NUMBER\$) + NULL_9\$ + CHR\$(1\$) + CHR\$(0\$) + CHR\$(1\$))	!X0102A

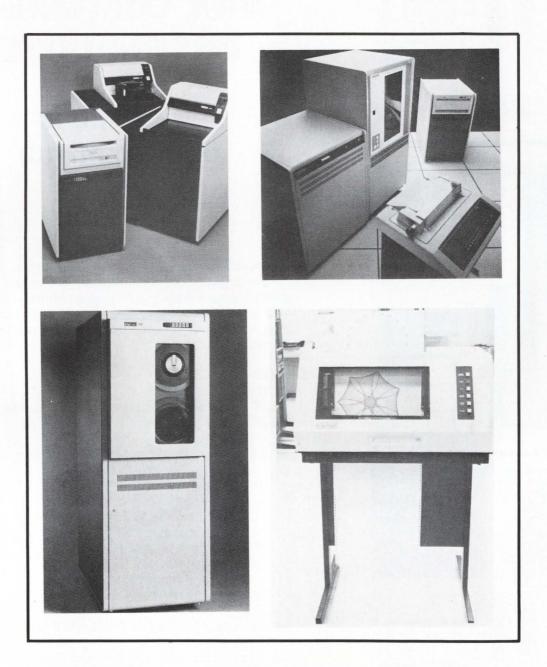
5060	NEXT IN	DE X%	
	OF ALL USER'S	ROUGH THE TERMINAL DEVICE DATA BLOCKS, SETTING THE SPEED TERMINALS OWNED BY A JOB TO ZERO. THIS PREVENTS THE FROM FILLING THEIR TYPE-AHEAD BUFFERS FULL OF CHARACTERS EY REALLY DON'T WANT. NOTE THAT WE ALSO SKIP KEO:.	
500		<pre>= SYS(CHR\$(6\$) + CHR\$(-13\$) + CHR\$(OUR_JOB_NUMBER\$) + CHR\$(255\$) + CHR\$(-128\$)) T_NUMBER\$ = 1\$ TO CNT_KB\$ GOTO 5510 GOTO 5510 IF (TEMP_0\$ = 0\$) TEMP_0\$ = SYS(CHR\$(6\$) + CHR\$(16\$) + CHR\$(0\$) + CHR\$(UNIT_NUMBER\$) + NULL_9\$ + CHR\$(TEMP_0\$ AND 255\$) + CHR\$(0\$) + CHR\$(SWAP\$(TEMP_0\$) AND 255\$)) TEMP_0\$ = SYS(CHR\$(6\$) + CHR\$(-4\$) + CHR\$(UNIT_NUMBER\$) + CHR\$(CMR\$(-4\$) + CHR\$(UNIT_NUMBER\$) + CHR\$(17\$))</pre>	1X0102A
510	NEXT UN FOR JOB	IT_NUMBER\$ _NUMBER\$ = 1\$ TO MAXCNT\$ TEMP_0\$ = PRIORITY\$(JOB_NUMBER\$) TEMP_0\$ = SYS(CHR\$(6\$) + CHR\$(-13\$) + CHR\$(JOB_NUMBER\$) + CHR\$(255\$) + CHR\$(TEMP_0\$)) UNLESS ((TEMP_0\$ = -1\$)	
	NEXT JO	OR (JOB_NUMBER\$ = OUR_JOB_NUMBER\$)) B_NUMBER\$	
		OURSELVES. EVERYONE TO THEIR INITIAL VALUES.	
999	RETURN		
	BACK TO	MAINLINE CODE.	
999	GOTO 32	700	
	Prevent	fall through errors.	
0000!	TT CT CHC		
00001		Local Subroutines	
0100!		SEND & MESSAGE TO OPSER	
0110	IF THEN	LEN(MESSAGE\$) <= 19\$ TEMP_0\$ = SYS(CHR\$(6\$) + CHR\$(22\$) + CHR\$(-1\$) + CHR\$(J\$) + "OPSER" + SP + NULL_10\$ + CHR\$(LEN(MESSAGE\$) + 1\$) + MESSAGE\$)	!X0102A
	ELSE	<pre>TEMP_0\$ = SYS(CHR\$(6\$) + CHR\$(22\$) + CHR\$(-1\$) + CHR\$(0\$) + "OPSER" + SP + NULL_10\$ + CHR\$(255\$) + LEFT(MESSAGE\$, 19\$)) MESSAGE\$ = RIGHT(MESSAGE\$, 20\$)</pre>	1X0102A
		GOTO 10110	
	SEND THE	E DATA TO 'OPSER'.	
0199	RETURN		
	BACK TO	MAINLINE CODE.	
9000!		Error Handler	
9004	IF THEN	ERR = 4\$ IF ERL = 10110\$ THEN SLEEP 1\$ RESUME 10110	
	NO ROOM	FOR USER ON DEVICE WHILE TALKING TO OPSER.	
9011		ERR = 11% RESUME 32700	
	End of f	file on device.	
9032		ERR = 32\$ IF ERL = 10110\$ THEN SLEEP 1\$ RESUME 10110	
	NO BUFFE	ER SPACE AVAILABLE WHILE TALKING TO OPSER.	
9999	ONERROR	GOTO 0	
	Give up.		
2700!		Completion Routines	
2710	CLOSE #1	7	
	Close al	l channels.	
	END		

The VAX-SCENE

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June 1983



INSIDE:

Using VAX Command Procedures (or Your File is My Command)



USING VAX COMMAND PROCEDURES (or, Your File is My Command) By Bob Meyer

Greetings, again, VAX and/or MACRO fans . . .

Being just about the laziest person on the planet, I recently discovered the joys of writing command files on the VAX. Wow! The command files under DCL are really something. For those of you who thought command files looked like this:

MAC FILE = FILE TKB FILE = FILE UT SEND KBXX: OH, MACRO PERSON, YOUR TASK-BUILD IS DONE . . .

you're in for a treat. Some of it's many features are:

String symbol definition Numeric symbols I/O to and from the terminal If-Then statements I/O to and from disk files Parameter substitution Build in lexical functions Gotos and labels

and lots more that I haven't even learned yet.

The first use I had for writing command files was a simple list of assembly and link instructions for a small program:

MAC BOB LINK BOB RUN BOB

RSTS/E INTERNALS MANUAL

The RSTS community has been clamoring for years for a book that details the inner workings of RSTS/E. Well, clamor no more. Michael Mayfield of Northwest Digital Software, and M Systems, the publisher of The RSTS Professional and The DEC Professional Magazines, have teamed up to produce the RSTS/E Monitor Internals Manual.

This manual describes the internal workings and data structures of the RSTS/E monitor. It also notes differences in the internal structures between version 7.1 and earlier versions of the monitor. Future updates will include changes for new versions of the monitor.

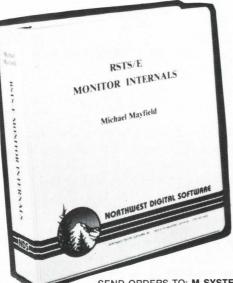
Information is available for all levels of users:

- Gain a basic understanding of the workings of the monitor for optimizing system performance.
- Information on disk structures allows recovery of data from corrupted disk packs.
- Special uses of runtime systems and resident libraries allow complex applications to be developed without degrading system performance.
- Write your own custom device drivers for that "foreign" device you need to add but thought you couldn't.

CONTENTS:

Chapter 1 describes the structures used by the monitor that are resident on disk. These include the directory structure, disk allocation tables, Save Image Library (SIL) formats, bootstrap formats and bad block mapping.

Chapter 2 describes the tables used within the monitor to control system resources and provide program services. These tables provide job, memory, file and device control, as well as program services such as interjob communication.



Chapter 3 contains information on writing and installing a custom device driver. It describes the entry points and information the driver must provide to the monitor as well as the subroutines and macros the monitor provides for the driver.

Chapter 4 contains information that enhances information already orovided by Digital on writing custom resident libraries and runtime systems. It concentrates mainly on non-standard uses of resident libraries and runtime systems to increase system performance and functionality.

Appendix A provides six quick reference foldout charts:

- The directory structure.
- The monitor tables.
- Fixed memory locations and common data structures.
- Monitor subroutines.
- Device driver entry points.
- Device driver macros.

Appendix B provides examples of the peek sequences required to access most of the monitor tables. It also contains an example program that uses many of the monitor tables to display a job and open files status.

Appendix C provides an example device driver.

Appendix D provides an example runtime system that doubles as a menu system for restricting specified users to a menu of options.



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Soon I learned that the same command file could be used for other assemblies as well, simply by replacing the file names with a parameter symbol: (the dollar signs in the command files are required, and seem to help DCL determine data from commands)

\$ MAC 'P1' \$ LINK 'P1' \$ RUN 'P1'

Where P1 is the first parameter I type after activating the command file: (assuming the cmd file is named 'ASM.COM')

@ASM BOB

DCL will then replace all occurrences of 'P1' with the string 'BOB' as it executes the file. Up to eight parameters may be specified, separated by spaces. Using single quotes around a string specify that the string should be substituted with something else; in the case of P1 thru P8, the strings are replaced with parameters from the user's command line activating the command file. An example of substitution would be:

\$INQUIRE FILE "Enter file to assemble" \$MAC 'FILE' \$LINK 'FILE' \$RUN 'FILE'

When this command file is executed, DCL will prompt the user with the string specified in the INQUIRE command, and take input from the terminal (actually, from SYS\$IN-PUT). That input will be assigned to the symbol 'FILE':

@ASM Enter file to assemble: BOB \$ MAC BOB \$ LINK BOB \$ RUN BOB

We can avoid seeing all this on the screen by adding the line:

\$SET NO VERIFY

to the command file. (This is often used in LOGIN.COM) Comments can be easily inserted in command files:

- \$!\$! command file to assemble & link programs
- \$!
- \$ SET NO VERIFY
- \$ INQUIRE FILE 'File'
- \$ MACRO FILE
- \$ LINK FILE
- \$ EXIT

The manuals recommend typing out commands in full, in an effort to make them more readable.

Several string operators are available; for example:

\$ IF P1 .EQS. "" THEN GOTO GET—INPUT
\$ FILE: = 'P1'
\$ GOTO DOIT
\$ GET—INPUT:
\$ INQUIRE FILE 'File'
\$ DOIT:
\$ MACRO FILE
\$ LINK FILE
\$ EXIT

In this case, if parameter 1 (P1) was not specified, the command processor will transfer control to the line following the label GETINPUT:. If, however, a string was supplied when invoking the command file, the symbol 'FILE' will be assigned the string, and the command file will branch around the INQUIRE statement and begin the assembly and link process.

Some other string operators are:

.GES. String greater than or equal to .GTS. Greater than .LES. Less than or equal to .LTS. Less than .NES. Not equal to

Several math operators are also at your disposal:

.EQ. Equal to

- .GE. Greater than or equal to
- .GT. Greater than
- .LE. Less than or equal to
- .LT. Less than
- .NE. Not equal

Also the logicals: .OR., .AND., .NOT. and faithful $+,-,^*,/$.

As an example of the arithmetic operators, let's assume we need to execute a program, a predetermined number of times:

- \$ COUNT = 0
- \$ RUN:
- \$ RUN DOODAH.EXE
- COUNT = COUNT + 1
- \$ IF COUNT .GE. 5 THEN GOTO EXIT
- \$ GOTO RUN
- \$ EXIT:
- \$ EXIT

String symbols can be used like RSTS CCL's on a userdefined level. For example, the RSTS user 'migrating' to VAX might add something to his/her LOGIN.COM file (which gets executed every time you log in) like this:

\$ SY: = SHOW SYSTEM

The ':=' is a string equate; whenever the string 'SY' is typed to DCL, it is replaced with 'SHOW SYSTEM' and the user gets the VMS equivlant of a RSTS systat.

Some other examples of practical use of the command language were pointed out to me by my comrade Bernie Velivis, system manager of Squibb's Molecular Modeling Graphics Lab in Princeton, NJ. Some of these include:

TE*CO:= = \$SYS\$SYSTEM:TECO.EXE/VT/HOLD

The '*' (I can't spell asterisk ...) is similar to the '-' (I CAN spell dash) in RSTS's CCL structure; it acts as an ?abbreviation point in the command string, allowing the user to type TE, TEC, or TECO to invoke TECO. The '\$' before SYS\$SYS-TEM: tells VMS to execute the image 'TECO.EXE' and pass to it the parameters '/VT/ HOLD' as well as any parameters the user might have typed (such as a file name to edit).

Another interesting feature is the ability to perform I/O to and from disk files. Note the following command file:

> \$ OPEN INPUT 'P1'
> \$ READ:
> \$ READ
> INPUT REC/END
> = DONE
> \$ WRITE SYS\$OUT PUT REC
> \$ GOTO READ
> \$ DONE:
> \$ CLOSE INPUT
> \$ WRITE SYS\$OUTPUT '

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CIRCLE 51 ON READER CARD

SYS\$OUTPUT "[End of file]" \$ EXIT First we open the file specified ('P1

First we open the file specified ('P1') when calling the file. Each time we encounter the READ statement, we get one line from the input file. The input record is assigned to

The '-' at the end of line 2 is used as a line continuation marker (which is helpful when your in verbose mode). The lexical functions used in the previous example, "F\$logical("TT")" and "F\$time()", are used to return the current system time and terminal device name. Lexical functions available are:

We then close the input file, display a message saying we've reached end of file, and exit. (Reminds me of the good ol' days of BASIC PLUS . . .)

the symbol 'REC'. Then we write the symbol to the terminal

(SYS\$OUTPUT), and loop. When we hit eof on the input file,

we take the error path specified in the read statement (/END = DONE) and control is transferred to the label DONE:.

The following command file gives an example of writing a file from DCL. This command file is executed by LGICMD.COM (the command file in SYS\$MANAGER: that is executed for every user who logs in) to track user logins by time and terminal.

> \$OPEN/ WRITE OUT SYS\$ MANAGER: LOGINS.LIS **\$WRITE OUT** "F\$LOGICAL ("SYS\$ LOGIN") SUCCESS-FULLY LOGGED IN INTO '-\$'F\$LOGICAL ("tt")' AT "F\$TIME()" **\$CLOSE OUT** \$EXIT: \$EXIT

The OPEN/ WRITE statement will create a new file if the specified file does not exist or append to a file if it does.

FUNCTION

F\$CVSI(Bit-position,width,integ)

F\$CVUI(Bit-position,width,integ)

F\$DIRECTORY()

F\$EXTRACT(offset, length, string)

F\$LENGTH(string)

F\$LOCATE(substring, string)

F\$LOGICAL(logical-name)

F\$MESSAGE(code)

F\$MODE()

F\$PROCESS

F\$TIME()

F\$USER()

F\$VERIFY()

ACTION

Signed value extracted from the specified integer, converted to an ASCII literal.

Unsigned value extracted from the specified integer, converted to an ASCII literal.

Current default directory.

substing starting at 'offset' of length 'length'.

Length of specified string.

Returns the offset of the substring within string indicated; or, the length of the string in the substring is not found.

Returns the logical translation of a logical name.

Message text associated with the specified status code. Useful for displaying error messages.

One of the character strings INTERACTIVE or BATCH.

Current process name string.

Current date and time of day in the format: dd-mmm-yyy hh:mm:ss.cc.

Current uic ([g,m]

A numeric value of 1 if verification is set on; a numeric value of 0 if verification is set off.

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CIRCLE 170 ON READER CARD

ACCTNG.B2S SYSTEM-WIDE RESOURCE ACCOUNTING SYSTEM FOR RSTS/E

By Philip Hunt, OL.LF.B.P., 6400 E. Broad St., Columbus, OH 43213

DESCRIPTION

ACCTNG is a system program that will remove the RSTS imposed restrictions of the amount of CPU/KCT and other resource time that may be used under any one account without 'overflowing' the internal RSTS accounting structures. ACCTNG will keep track of all system usage in an indexed data file that is updated nightly with all usage that day. Since this file is kept in BP2 'double-precision' numbers, data will not be lost.

USAGE

ACCTNG is run from any terminal; it will find its own data file, and initialize one if one doesn't currently exist. It creates its data file in the account in which ACCTNG resides.

Example:

Ready

RUN \$ACCTNG ACCTNG — V1.02 RSTS V7.2-04 N OLFBP 11/70

A)dd, L)ist, C)lear or E)xit:

At this point there are four options:

1) Add current statistics to file. This is usually done every night but may be done more often without damage.

2) List all or any account in a money-type format (see enclosed samples) that shows all usage FOR THE YEAR for CPU,KCT,Device Time, Connect-time. It also shows other items of interest. An account that is deleted, will be marked with a '(D)' next to the PPN and no updates are then done until recreated. This brings along an explanation of the last two columns, UPDATE shows how many times the account was found during nightly update and LUD shows the last time the account was found. You may direct the report output to any file or device. If selective mode is requested, a project, programmer number will be requested by ACCTNG.

3) Clear (Priviledged) — This CLEARS the YEARLY FILE after forcing a List operation; it will save the current file under ACCTNG.OLD and create a new ACCTNG.SYS file.

4) Exit the program.

LIST Example:

A)dd, L)ist, C)lear	or E)xit: L					
Output to <kb:> ?</kb:>						
Selective <n> ? Y</n>	ES					
Project, Programmer	? 1,2					
PPN	CPU-TIME	KCTS	DEVICE	CONNECT	UPDATE	LUD
[001,002]	8:02.5	56048	0	1:44	2 10-	Jan-83
Project, Programmer	? 1,10					
PPN	CPU-TIME	KCTS	DEVICE	CONNECT	UPDATE	LUD
[001,010]	5:33.6	59666	46	1:42	2 10-	Jan-83
Project, Programmer	? ^Z					
Ready						
ADD (Nightly update						
RUN \$ACCTNG						
ACCTNG - V1.02	RSTS V7.2	2-04 N OLF	BP 11/7	0		
A)dd, L)ist, C)lear Accounts Updated	or E)xit: A					
Updated Accounts: New Accounts:						

I have also included a sample report showing a complete system dump of all accounts.

DO NOT FORGET, THIS PROGRAM MUST BE COMPILED IN DOUBLE PRECISION BP2.

Compilation Example:

BP2 OLD ACCTNG.B2S BUILD/IND COM/DOUBLE/OBJ TKB @ACCTNG

PIP [PPN]ACCTNG.TSK<232>=ACCTNG.TSK

WHERE: [PPN] IS WHERE YOU WOULD LIKE ACCTNG TO RESIDE

	1		
	! PROGRAM:		
	! AUTHOR :	PJH OLFBP	
	! VERSION:	V07	
	! EDIT :		
	! DATE :	072981	
	1		
	! THIS PROGRAM WILL DO		
	1) A)DD VALUES FROM FILE IN DOUBLE PR	SYSTEM TABLES TO ACCTNG	
		VALUES FROM MY DATA FILE	
	TO OUTPUT SPECIFI		
		ILE (DOES A LIST OF ALL	
	I FIRST)		
	1		
	!*****NOTE: THIS PROGRAM MUS		
	DOUBLE PRE		
	1		
1000	DIM F1P\$(30%)	× ×	
1000	ON ERROR GOTO 19000		
	on minor oore , yeee		
1005	MAP (PPNREC)		
1005	PPN\$ = 6%.		
	CPU.		
	KCT.		
	CNNCT.		
	DEVICE,		
	LAST.UPD,		
	UPDT,		
	LUD\$ = 9%	<u>\</u>	

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		DTREC\$	TE\$=9%, ME\$=8%, E\$=9%, E\$=8%, T,				
1010	VERSION\$		= "1"			Ν.	
	EDT\$	TONA	= "02"	(04)		1	
	INSTALLAT R.PACKAGE		= ERT\$ = SYS((0%) CHR\$(12%))		1	
	CHANGE		R. PA	CKAGE\$ TO I	FIP%	N 1	
	PACKAGE.F PACKAGE.F PACKAGE.F	PROG%	= FIP% = "["+	(5%) NUM1\$(PACK)	AGE.PROJ%)+	\ ICAN F	IS SO WE IND WHERE WE ME FROM
	PACKAGE.F	SPEC\$	"]" = "_"+	CHR\$(FIP%(2		1 /	
				PACKAGE. FSI	(25%))+":"+	\ 1	
	PRINT "AC PRINT	CCTNG -	V";VER	SION\$;".";H	EDT\$,INSTALLAT	TION\$ \	
	OPEN "KB: OPEN PACK C	AGE.FS ORGANIZ	PEC\$+"A ATION I MODIFY,	CCTNG.SYS" NDEXED FIXE ALLOW NONE		, FILE ∉1%,	
	F CLOSE #1	MAP PPN PRIMARY	KEY PP	N\$ NODUPLIC	CATES	١	
1011	OPEN PACK	DRGANIZ	ATION I	CCTNG.SYS" NDEXED FIXE ALLOW NONE			
	٢	AP PPN	REC,	N\$ NODUPLIC		1	
	GET #1%,	KEY #0	% EQ "0	00000"		1	
	F.DATE\$ =	INIT.	DATE\$+"	n		1	
	F.TIME\$ = U.DATE\$ =	UPD.D	TIME\$+" ATE\$+""			,	
	U.TIME\$ =	UPD.T	IME \$+""			1	
	CUR.ACCT= SYSTEM.AC GOTO 1020	CTS =		Т		ì	
1012			AD FORM	AT, SO REIN	TI TIN		
	GOSUB 401 GOTO 1011					`	
1013	INO FILE GOSUB 401 PRINT		INIT O	NLY		1	
1020				lear or E).	xit: ";	X	
		EFT (CVT	\$\$(CMD\$,-1%),1%)		1	
	CMD% = IN GOTO 1020 ON CMD% () IF CM	D% = 0%		0,32700,5000,6	5000 \	
	GOTO 3270						
2000		, CLEAR	ALL UP	DATE FLAGS)		
	RESET.FL. FILE.ACC					1	
	FIND #1%	, KEY#0	% GT "O	00000"		1	
	WHILE -19	% GET #1%				,	
	1	FILE.AC LAST.UP	CT=FILE D=0%	C. ACCT+1		\ \	
2002	NEXT	UPDATE	W 1%			X	
2005	INDX% = SYSTEM.A					~	
	WHILE -1	¥.	INDX%+1%			>	
		INFO\$=5	CHR	\$(RESET.FLA \$(0%)+CHR\$(R\$(SWAP\$(INDX G\$)+CHR\$(RESE 0\$)+CHR\$(1\$)+	T.FLAG%)+	å
		SYSTEM	ACCT-S	STEM. ACCT+	STRING\$(0%,21	x))/	
		CHANGE	INFO\$ 7	TO FIPS		N	
			FIP\$(8%) FIP\$(7%) 0000			1	
2010			1%, KEY =UPDTD%→	#0% EQ PPN ⊧1%	1\$	Ν	
2015		GOTO 20)50				
2020		KCT=0				Ň	
		CPU=0 CNNCT=0)			1	
		DEVICE:	= 0			N	
			6=NEWADI	D\$+1\$		1	
2050		PUT #19		0% EQ PPN\$			
2050					DTD#/ 10#		
2060		CPU.HIC	GH=(FIP9	FIP%(14%)+ %(22%) AND		1	
!HIGH	IS HIGH 6				HIGH#65536.)	1	

	DEVICE.TEMP=256. #FIP%(20%)+FIP%(19%)	1	
	CNNCT.TEMP =256.*FIP\$(16\$)+FIP\$(15\$)	1	
urrou to	KCT.TEMP =256.#FIP%(18%)+FIP%(17%) S LOW ORDER 10 BITS OF FIP(21,22)	1	
THIGH IS		AND 04111	
	KCT.HIGH =FIP\$(21\$)+(256.*(FIP\$(22\$) KCT.TEMP =KCT.TEMP+(65536.*KCT.HIGH)		
	KCT=KCT+KCT.TEMP	1	
	CPU=CPU+CPU.TEMP	\	
	DEVICE=DEVICE+DEVICE.TEMP	1	
	CNNCT = CNNCT+CNNCT. TEMP	1	
	UPDT=UPDT+1	\	
	LAST. UPD=-1%	1	
	LUD\$ = DATE\$(0\$)	1	
	UPDATE #1%	1	
	NEXT		
2500	PRINT "Accounts Updated"	X	
	GET #1%, KEY#0% EQ "000000"	1	
	UPD.DATE\$=DATE\$(0%)	\	
	UPD.TIME\$=TIME\$(0%)	\	
	SYS.ACCT = SYSTEM.ACCT	\	
	ACCT.CNT = FILE.ACCT+NEWADD%	\	
	CUR.ACCT = ACCT.CNT	1	
	UPDATE #1%	X	
	PRINT	1	
	PRINT "Updated Accounts: ";UPDTD% PRINT "New Accounts: ";NEWADD%	1	
	PRINT "New Accounts: ":NEWADDS	1	
	RETURN		
3000	ILIST FROM CURRENT FILE		
2	ZERO, MSG \$=""	\	
	PROG\$ = 0\$	1	
	PROJ% = 0%	1	
	PPN \$= STR ING \$(0%,6%)	X	
	IF CLEAR.FLAG\$ THEN ZERO.MSG\$=" and INITIALIZATION" SELECTIVE\$ = "N"	1	
3005	LINPUT "Output to <kb:> ";OUTP\$</kb:>	1	
	OUTP\$="_KB:" IF OUTP\$=""	1	
301€	GOTO 3017 IF CLEAF. FLAG%	1	
3010	INPUT "Selective (N) ";ANS¢	1	
	SELECTIVE SELECTIVE SHOW , MANY	1	
	SELECTIVE\$="N" SELECTIVE\$="Y" IF LEFT(ANS\$,1%)="Y"		

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CIRCLE 178 ON READER CARD

3015	IF SELECTIVE\$="Y" THEN INPUT "Project, Programmer ";FROJ≴,PRO GOSUB 10000 GOTO 3015 IF PPN\$="000000" TOTLINE≸=0\$ FIND 015, KEY 00≸ EQ PPN\$	XGX /	6 6 6 6 8		<pre>PRINT #3%, "Acctng Dump";ZERO.MSG\$;" of Snapshot File on "; DATE\$(0%);" at ";TIME\$(0%);TAB(70%);"Page ";PGE% \ PRINT #3%, "File Created: ";F.DATE\$;" at ";F.TIME\$,\ PRINT #3%, "File Accounts: ";CUR.ACCT \ PRINT #3%, "File Updated: ";U.DATE\$;" at ";U.TIME\$,\</pre>
3017	OPEN OUTP\$ FOR OUTPUT AS FILE 3%		å		PRINT #3%, "System Accounts: ";SYSTEM.ACCTS
3020	GET #1 GOTO 3020 IF PPN\$="0000000" GOSUB 11000 IF SELECTIVE\$="Y" THEN 3015 ELSE 3020	1	12 ه ه ه	010	PRINT #3%, " PPN CPU-TIME KCTS ";\ PRINT #3%, " DEVICE CONNECT UPDATE LUD"\ TOTLINE%1% \
3030	IEOF ON LIST FROM CLEAR/INIT FUNCTION RETURN		& & 15	000	RETURN ! TIME - HHHH:MM FROM TT (BEING MINUTES,NOT SECONDS
4000	ICLEAR CURRENT FILE INPUT "Are you sure <no> ";AANS\$ GOTO 4005 IF AANS\$="YES"</no>	2	& & &		TT=TT=60.\ IGET INTO SECONDS HOURS.FLAG#=-1#\ GOTO 16010
4005	PRINT "CLEAR aborted!!" CLEAR.FLAG\$ = -1\$		& 16 &	000	! T I M E - H H H H : M M : S S . S ! RETURNS A STRING SS\$ FROM TT
4010	GOSUB 3000 PRINT "Initializing data file"				HOURS.FLAG#=0#
4010	CLOSE #15 kill package.fspec\$+"ACCTNG.OLD" goto 4012		& 16 & &	010	SS&="" \TT≸=0\$ \TT\$≤TT/3600. \IFT1≸≤TT/3600. TT1\$<>0\$ THEN TT=3600.®TT1\$
4011	Ino file to kill print "No old file to be killed, continuing" goto 4012	λ	& & &		GOSUB 16040 SS\$=SS\$#"":" ! CLEAR STRING TO BE RETURNED AND INITIALIZE ! T1\$= # OF HOURS - SUBTRACT FROM TOTAL IF NON-ZERO
4012	NAME PACKAGE.FSPEC\$+"ACCTNG.SYS" AS PACKAGE.FSPEC\$+"ACCTNG.OLD" PRINT "Old file: ";PACKAGE.FSPEC\$+"ACCTNG.SYS" +" named to ";PACKAGE.FSPEC\$+"ACCTNG.OL goto 4015	\ \	& & &	020	TT1\$=TT/60\$ \ IF (TT1\$ OR TT\$)<>0\$ THEN TT=TT-60\$*TT1\$
4013	Ino file to rename PRINT "No current file, continuing"	1	& & &		<pre>\ RETURN IF HOURS.FLAG\$!ONLY WANT HH:MM, SO EXIT \ SS\$=SS\$+":" ! NOW T1\$= 0 OF MINUTES - SUBTRACT FROM TOTAL IF NON-ZERO AND ! TACK ON TO STRING</pre>
1015	goto 4015			030	TT1\$=TT ! PREPARE TO GET SECONDS.TENTHS OF SEC.
4015	OPEN PACKAGE.FSPEC\$*"ACCING.SYS" FOR OUTPUT AS ORGANIZATION INDEXED FIXED, FILESIZE 250\$, ACCESS MODIFY, ALLOW NONE,	FILE #1%,	& 160 & &		TT\$=NUM\$(TT1\$+TT\$*100\$) \ SS\$=SS\$+MID(TT\$,2\$+TT\$,LEN(TT\$)-2\$-TT\$) \ TT\$=1\$
	MAP PPNREC, PRIMARY KEY PPN\$ NODUPLICATES PPN\$="000000"	1	& &		I TACK THIS PART ONTO THE STRING
	INIT.DATE\$= DATE\$(0\$) INIT.TIME\$= TIME\$(0\$)	X	& 160 & &	050	RETURN
	UPD.DATE\$ = "NONE" UPD.TIME\$ = "NONE" SYS.ACCT = O\$ ACCT.CNT = O\$ PUT #15		& 190 & & & &		<pre>IERROR PROCESSING IF ERR-5≸ AND ERL=1010 THEN PRINT "?ACCTNG- Data file not found-INITIALIZING" \ RESUME 1013</pre>
	PRINT "File Cleared!!!" RETURN	1	ል & 190 &	010	IF ERR=155% AND ERL=3015 THEN PRINT "%Account not found" \ RESUME 3015
5000	IRESERVED RETURN		& 190 &	20	IF (ERR=11\$ OR ERR=155\$) AND (ERL=3020) THEN
6000	IRESERVED Return		å		PRINT #35, CHR\$(12\$); ////////////////////////////////////
10000	<pre>IGIVEN PROJ\$,PROG\$, WILL GENERATE PPN\$ 6 CHARS : PP1\$ = NUM1\$(PROJ\$) PP2\$ = NUM1\$(PROG\$) PP1\$ = RIGHT("000",LEN(PP1\$)+1\$)+PP1\$</pre>	LONG	& 190 & &	25	IF (ERR=155\$) AND (ERL=1011) THEN PRINT "\$No Date record found" RESUME 1012
	PP2\$ = RIGHT("000",LEN(PP2\$)+1\$)+PP2\$ PPN\$ = PP1\$ + PP2\$ RETURN	`	Ł	30	IF ERL=3017 THEN PRINT "?Bad output specification" \ RESUME 3000
11000	IDISPLAY A LINE FROM FILE GOSUB 12000	\	& 190 &	40	IF ERR=5≸ AND ERL=2005 THEN RESUME 2500
	PRINT #3%, "["+LEFT(PPN\$,3%)+","+RIGHT(PPN\$,4%) PRINT #3%," (D)"; IF LAST.UPD>=0% TT=CPU/10	+"]"; \ \ \	& 190 &	50	IF ERR=155% AND ERL=2010 THEN RESUME 2020
	TTTS=INT((TT-INT(TT))*10%) GOSUB 16000 SS\$=SS\$+"."+CHR\$(48%+TTTS)		& & 190 &	060	IF (ERR=155% OR ERR=11%) AND ERL=2000 THEN RESUME 2005
	SS\$=FNSP\$(SS\$,16%) PRINT #3%, TAB(13%);SS\$; SS\$ = NUM1\$(KCT)	1	& 190	70	if erl=4010 then resume 4011 Ino old file to kill-ok
	SS\$ = FNSP\$(SS\$,12\$) PRINT #3\$, TAB(33\$);SS\$;		å 190	80	if erl=4012 then resume 4013 Ino current file-ok
	TT=DEVICE GOSUB 15000		& 199 &	80	IF ERR=11% THEN RESUME 32700
	SS\$=FNSP\$(SS\$,9\$) PRINT #3\$, TAB(41\$);SS\$; TT=CNNCT GOSUB 15000		ል ል 199 ል	90	PRINT "?ACCTNG - ";ERT\$(ERR);" (";NUM1\$(ERL);")-FATAL"\ RESUME 32700
	SS\$=FNSP\$(SS\$,9%) PRINT #3%, TAB(53%);SS\$;	1	& 327 &		PRINT CLOSE #1%,#2%
	PRINT #3%, TAB(65%);FNSP\$(NUM1\$(UPDT),4%); PRINT #3%, " ";LUD\$; PRINT #3%		& 327 &	10	DEF FNSP\$(AR\$,P\$)=SPACE\$(P\$-LEN(AR\$))+AR\$
	PRINT 03% IF SELECTIVE\$="Y" LCOUNT%=LCOUNT%+1% TOTLINE%=0% IF LCOUNT%>57% RETURN		4 327 4 4 4	67	END
12000	IDISPLAY TITLE RETURN IF TOTLINE%	1	& Acc	tng D	ump of Snapshot File on 07-Jan-83 at 05:11 PM Page 1
	GOTO 12010 IF SELECTIVE\$="Y" PGE%=PGE%+1%	1	&		ated: 14-Jan-82 at 10:62 AM File Accounts: 373
	PRINT #3%, CHR\$(12%);	`			ated: 06-Jan-83 at 08:54 PM System Accounts: 305

&

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PPN	CPU-TIME	KCTS	DEVICE	CONNECT	UPDATE LUD	
[000,001]	0.0	0	0 DEVICE	0	70 06-Jan-83	
[001,000]	10:03.0	151976	47	1:39	70 06-Jan-83	
	48:42:52.7	3584947	28	2:57	70 06-Jan-83	
[001,001]	73:09:04.5	38855575	2401:14	1621:49	70 06-Jan-83	SO THAT'S WHAT
[001,002]	1:08:02.1	742939	30	9:28	70 06-Jan-83	JO I HAIJ WHAI
[001,003]		713129	2:03	10:14	70 06-Jan-83	
[001,004]	1:35:35.0	15523246	183:18	253:54	70 06-Jan-83	
[001,005]	21:59:59.1	15523240	103:10	255.54	70 06-Jan-83	AN ERROR LOG
[001,006]	0.0		5	2:37	70 06-Jan-83	A TIN LAKUK LUG
[001,007]	6:22.2	24932	0	2:31	70 06-Jan-83	
[001,008]	14.6	1286			70 06-Jan-83	
[001,009]	13:27.9	64593	1:45	2:27	70 06-Jan-83	() LOOKS LIKE!
[001,010]	65:13:36.4	41108124	254:09	1498:10		LINE:
[001,011]	1:06.5	8862	0	24	70 06-Jan-83	
[001,012]	0.9	45	0	0	70 06-Jan-83	
[001,013]	0.0	0	0	0	70 06-Jan-83	
[001,014]	0.4	9	0	0	70 06-Jan-83	
[001,015]	45.5	6956	8	4	70 06-Jan-83	
[001,016]	0.0	0	0	0	70 06-Jan-83	(1) (3)
[001,017]	0.0	0	0	0	70 06-Jan-83	
[001,018]	21:17.9	287039	49	3:10	70 06-Jan-83	
[001,019]	0.0	0	0	0	70 06-Jan-83	A IN Start Change
[001,020]	7.5	1555	0	5	70 06-Jan-83	Let LEPPen
[001,021]	75:37:35.1	42391838	256:27	827:23	70 06-Jan-83	K I I CONTROL
[001,022]	1.6	149	0	0	70 06-Jan-83	
[001,023]	0.0	0	0	0	70 06-Jan-83	Y THE Y
[001,024]	4:56.7	67142	57	1:06	70 06-Jan-83	
[001,025]	33:20:00.7	18459087	176:39	798:41	69 06-Jan-83	
[001,026]	12:09:08.8	5554523	54:30	454:50	49 06-Jan-83	
[001,030]	0.0	0	0	0	70 06-Jan-83	Real - // I
[001,031]	6:13.0	77003	29	1:31	47 06-Jan-83	
[001,032]	41.2	2863	0	28	30 06-Jan-83	
[001,033] (D)	0.0	0	0	0	17 05-Sep-82	
[001,035]	57:03.6	791168	2:03	52:33	31 06-Jan-83	#00007 (***03)
[001,038]	20:08:44.6	13882483	93:07	118:38	70 06-Jan-83	
[001,040]	4:41:15.1	3424433	21:12	119:29	31 06-Jan-83	10004 10003
[001,041] (D)	0.0	0	0	0	17 02-Jun-82	
[001,045]	1:16:18.6	816248	7:55	60:24	31 06-Jan-83	[2003] [2003]
[001,050]	66:01:24.7	41539474	201:12	1610:15	70 06-Jan-83	Loco7 10001
[001,051]	13:04:47.1	6859340	9:44	233:22	67 06-Jan-83	/300/ kood
[001,052]	1:12:50.1	946693	1:41	37:06	70 06-Jan-83	
[001,055]	0.0	0	0	0	70 06-Jan-83	
[001,056]	0.0	0	0	0	70 06-Jan-83	
[001,060]	7:16.2	36739	5	6:17	70 06-Jan-83	
[001,060] [001,061] (D)	0.0	0139	õ	0.11	17 07-Jul-82	
[001,061] (D)	5:18.3	57181	6	1:01	48 06-Jan-83	EMILY DAHLSTROM
	37:29.3	548774	1:48	5:49		
[001,063]	26:23.0	296 994	1:29	7:32		- //
[001,064]	2:43:41.5	126 1838	2:59	13:42		
[001,065]		9083616	69:04	136:44	30 06-Jan-83	
[001,070]	10:47:47.5	9003616	09:04	130:44	70 06-Jan-83	
[001,076]	30.8	230149	0	2:56	52 06-Jan-83	
[001,077]	14:22.5	230149	1	2:50	J2 00-0411-03	



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RSTS/E MULTITERMINAL INPUT/OUTPUT SERVICE

By Michael H. Koplitz, Computer Integrated Management Systems

The information in this document is believed to be accurate and correct, however the author assumes no liability for any errors which may appear in this document.

The multiterminal input/output service allows for the development of software which controls many terminals from one input/output channel. This feature is useful for the development of user interfaces, data entry programs, and other RSTS/E system programs. Several terminals can be doing the same task but only ONE job is in use. The efficiency of multiterminal service can be seen in computer shops where there is a high volume of data input.

SYSGEN OPTION

To be able to use multiterminal input/output service, the system manager must specify during the SYSGEN dialogue that multiterminal service is desired. The SYSGEN question and long description are:

An optional feature of the RSTS/E terminal service allows one job to interact with several terminals through special forms of the Record I/O GET and PUT statements. This feature is useful in applications where the same basic function is performed on several terminals and a separate job for each is undesirable or at least inefficient. Would you like to include this feature (YES OR NO)?

Multi-terminal service? #Y# Answer Y or linefeed to this question.

THE MASTER TERMINAL

To use multiterminal input/output service a master terminal must be selected. This master terminal will be opened on one of the twelve RSTS/E channels. The OSC should not be used as the master terminal. This can interfere with OPSER, if OPSER if running.

OPEN "KB:" AS FILE #1%

This OPEN statement associates the current keyboard with channel one. Any channel, one through twelve can be used. Also a keyboard other than the current keyboard, can be used for the master terminal.

OPEN "KB10:" AS FILE #2%

This OPEN statement associates keyboard ten with channel two. Keyboard ten becomes the master keyboard.

Keyboards which are to be included in multiterminal input/output services must be assigned to the job.

OUTPUT FOR MULTITERMINAL SERVICES

The PUT or PRINT statements are used to transmit data to the terminals under the multiterminal service. The PUT statement is used in combination with the FIELD statement to transmit data to the multiterminal service. PUT #5%, RECORD 32767% + 1% + T%, COUNT N% The RECORD option must be 32767% + 1% to indicate to RSTS/E that the output is to be for multiterminal services. The variable T% in the RECORD count is a specific keyboard unit number to transmit to. Therefore, all of the terminals assigned to the job can be transmitted to by using RECORD clause 32767% + 1%, or transmission can be to a single terminal by using 32767% + 1% + T%. The COUNT option is used to transmit a specific amount of the buffer to the terminal.

The error, ?NOT A VALID DEVICE (ERR = 6), can occur if the terminal being addressed is not the master keyboard nor one of the slave terminals. A slave terminal is one which is assigned to the job, and is not indicated in the OPEN statement of the master terminal. Another possible error is, ?I/O CHANNEL NOT OPEN (ERR = 9), which indicates that the terminal in question is not assigned to the job.

The PRINT or PRINT-USING statement can be used with multiterminal input/output services. The RECORD clause is used to designate multiterminal input/output services.

PRINT #5%, RECORD 32767% + 1% + T%, A\$

PRINT #5%, RECORD 32767% + T% + 1% + T%, USING "III", A\$;

The FIELD statement is NOT used if the PRINT or PRINT-USING statements are used.

To transmit binary data add 4096% to the RECORD clause.

INPUT FROM MULTITERMINAL SERVICES

Input from multiterminal services can be requested from a specific keyboard or from any of the terminals that might have some data available. The type of input is specified in the RECORD clause of the GET statement.

To get input from a specific keyboard the following statement is used.

GET #5%, RECORD 32767% + 1% + T%

The variable T% is the unit number of the terminal from which data is to be accepted. A FIELD statement is then used so that the program can read the data. The variable RECOUNT contains the length of the buffer.

FIELD #5%, RECOUNT AS A\$

The information from keyboard T% is contained in variable A\$. If there is not any data available, the error, ?DATA ER-ROR ON DEVICE (ERR = 13), is generated. The system WILL NOT stall (KB state) on the GET statement.

To get data from any terminal on the multiterminal service the following GET statement is used.

GET #5%, RECORD 32767% + 1% + 16384% + 5%Where S% has several meanings:

Value: 0% Meaning:

The GET statement waits until input is available from any of the terminals on the service. The program stalls until data is received. After data is received the error, ?DATA ERROR ON DEVICE (ERR = 13), may be generated under some cir-

cumstances (race condition with [†]C). A race condition can occur when two jobs are accessing the same data.

Value: 1 to 255 Meaning:

The GET statement waits up to S% seconds for input from any terminal. If no data is received in S% second the error, ?DATA ERROR ON DEVICE, is generated.

Value: 8192% Meaning:

The GET statement immediately generates the error, ?DATA ERROR ON DEVICE, if no data is available.

A 1Z entered at either the master or slave terminal generates the error, ?END OF FILE ON DEVICE.

When the GET statement is executed and data is available, the first character of the buffer is the unit number of the terminal where the data was generated. To convert this value into a Software Tools for RSTS/E

to him.

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unit number use the ASCII function.

LIMIT.BAS — AN EXAMPLE OF MULTITERMINAL SERVICES

LIMIT.BAS allows the system manager to prevent specific terminals from logging into the system. The program uses multiterminal input/output services to acthe "/S" is not used ALL terminals that are not in use at the moment will be limited.

It is HIGHLY recommended that if the time set for LIMIT.BAS to start is in the future, that the "/S" switch NOT BE USED. If LIMIT.BAS is to start in the future it will detach. If LIMIT.BAS is to start immediately the prompt LIM > will appear.

minals. Second is the stop time; when should LIMIT.BAS stop limiting the terminals. LIMIT.BAS will sleep until the start time is reached. It will send a message to KBO: indicating its start and stop times.

complish this task. All terminals on the system can be

"limited." LIMIT.BAS has special features built in to allow a

user to free up the terminal if the LIMIT password is known

start time, when should LIMIT.BAS start limiting the ter-

LIMIT.BAS has two time parameters. The first is the

LIMIT.BAS allows for a password to be established which when entered on a terminal will cause LIMIT.BAS to free up that terminal for operator use. The password MUST be six characters in length. The password feature need not be used. A < CR> entered to the password question will set up this situation.

A message for the session is prompted for. LIMIT.BAS will display the inputted message each time a user enters something which is not the password on a "limited" terminal. A switch exists to this message. This switch "/S" allows for the entering of individual terminals that are to be limited. If

Note that there is a ten second time limit on entering a command to LIMIT.BAS. The time limit was set up to insure that any terminals that are "limited" are serviced in some reasonable amount of time. If the program completely stalled at the LIM > prompt, then terminals might not be serviced for very long periods of time.

LIMIT.BAS will first attempt to open the master terminal. LIMIT.BAS uses KB40: as the master. This can be adjusted by changing line 07. J% contains the master terminal number. If the master terminal is not available an error is generated and LIMIT.BAS will ask for a different terminal to be the master terminal.

There might be some terminals which are not assigned to LIMIT.BAS because they are in use. LIMIT.BAS does contain some commands which will force the terminal to become assigned to it. There are commands which cause LIMIT.BAS to free up a terminal. Also assigning terminals which are not in use can be done. The commands are entered at the LIM > prompt. The commands are:

Command	Meaning
ASSIGN	Assigns a terminal to LIMIT.
CHANGE	Changes a LIMIT feature.
DEASSIGN	Deassigns a terminal from LIMIT.
DETACH	Detaches the job. Attaches to the job if
	other commands are desired.
END	Ends LIMIT.
FREE	Lists the free terminals on the system,
	those not controlled by LIMIT.
HELP	Displays HELP message.
LIST	Lists all assigned terminals.
SEIZE	Siezes a terminal which is in use.
VIEW	Lists all of LIMIT's parameters.

It is best to detach LIMIT.BAS; it runs more efficiently!

REFERENCES

RSTS/E Programming Manual RSTS/E System Generation Manual (for SYSGEN dialogue)

```
......
01
       THIS PROGRAM ALLOWS SYSTEM MANAGER TO PREVENT CERTAIN TERMINALS
   11
   1*
       FROM LOGGING INTO THE SYSTEM.
      THIS PROGRAM WAS WRITTEN BY M H KOPLITZ
   1*
   1
              COMPUTER INTEGRATED MANAGEMENT SYSTEMS
   REVISION #1 (TO VERSION 2) M H KOPLITZ OCT 2, 1981
   1
   1*
              PURPOSE: DO NOT LET ANY TIME BE ENTERED EXCEPT
                      XX:XX.
   05
       DIM A(60%).B(60%)
07
       J% = 56%
                                           INUMBER OF TERMINALS
       M$ = "KB" + NUM1$(J%) + ":"
                                           IMASTER KB NUMBER
08
                    V 2.1
10
       PRINT "LIMIT
                  LIMIT systems use"
15
       ON ERROR GOTO 31000
                                           !ERROR ROUTINE
       X$ = SYS(CHR$(6$) + CHR$(-7$))
                                           ISET ^C TRAP
16
17
       OPEN M$ AS FILE #12%
                                           IOPEN THE MASTER FILE FOR
                                            MULTI-TERMINAL SERVICES
20
       PRINT
                                           IPRINT HEADER
       INPUT "Time to start <NOW>";A$
                                           IASK FOR TIME START
30
40
       IF A$ = "" THEN A$ = CVT$$(TIME$(0),-1%)
GOTO 50 IF LEN(A$) = 5%
       PRINT "?Time must be in the form XX:XX"
                                          IF NULL USE CURRENT TIME
I ELSE CHECK LEGALITY.
       GOTO 30
       INPUT "Time to end";B$
IF B$ = "" THEN 32767
50
                                           IASK FOR TIME STOP
```

```
GOTO 50
                                                     IIF NULL STOP, ELSE CHECK
 1
                                                     I TO SEE IF VALID.
         INPUT "Password for session"; P9$
63
         PRINT "$No password for session"
IF P9$ = ""
P9$ = "" IF P9$ = ""
 ١
         GOTO 70 IF P9$ = "~"
         PRINT "%Password must be 6 characters";
                       length"
                    in
                 IF LEN(P9$) <> 6
 1
         GOTO 63
                  IF P9$ = "" OR LEN(P9$) <> 6%
                                                    IENTER PASSWORD
                                                     FOR SESSION
70
         PRINT "Message for the session";
         INPUT LINE C$
         C$ = CVT$$(C$,4%)
         C1$ = CVT$$(C$,32$ + 4$)
         C$ = LEFT(C$, INSTR(1%, C$, "/")-1%)
IF INSTR(1%, C$, "/") <> 0%
         C$ = C$ + CHR$(10) + CHR$(13)
                                                     IMESSAGE TO BE SENT
71
         IF A$ > TIME$(0) THEN GOSUB 1500
                                                    ITTME TO START NOT YET DETACH
72
         IF A$ > TIME$(0) THEN SLEEP 30%
                  L95 = 15
                 GOTO 72
                                                     ITIME TO START NOT HERE YET
75
         X$ = SYS(CHR$(6$) + CHR$(9$) + CHR$(0)) IGET KB NUMBER
         T_{3} = (ASCIT(MTD(X_{3}, 2, 1))/2)
                                                    IT% HAS KE NUMBER
77
78
         X$ = SYS(CHR$(6$) + CHR$(-5$) + CHR$(0)
                 + "$System will be limited"
+ " until: "
                  + B$ + CHR$(13%)
                  + CHR$(10))
                                                    ISEND MESSAGE TO KBO:
         GOTO 7000
79
                 IF RIGHT(C1$,
INSTR(1$,C1$,"/") + 1$)
                             20
                 AND INSTR(1%,C1$,"/") <>0
80
         FOR X = 9 TO J%-1%
                                                    TALL LEGAL TERMINALS
85
                 IF T% = X THEN B(X) = 0%
                                                     ISKIP CURRENT TERMINAL
                  GOTO 120
90
                  X$ = SYS(CHR$(6%) + CHR$(10%)
                    + STRING$(20%.0%)
                     + "KB" + CHR$(X)
                     + CHR$(255%))
                                                    LASSIGN THE KB
                 X$ = SYS(CHR$(6$) + CHR$(-5$))
100
                          + CHR$(X)
                          + "$System will be
+ "limited until: "
                          + B$ + CHR$(13%)
+ CHR$(10))
                                                    ISEND SYSTEM LIMITED PHRASE
                                                    IFLAG TERMINAL AS ASSIGNED
110
                 A(X) = X
120
         NEVT V
                                                    ICONTINUE ASSIGNING TERMINALS
                                                    IIN DETACHED CONDITION
125
        GOTO 217 IF L9% = 1%
        WAIT 10
                                                    WAIT 10 SECONDS
130
                                                    IREADY FOR LIMIT COMMAND INPUT
        PRINT "LIM>":
140
                                                    IGET THE COMMAND
150
        INPUT D$
155
        D$ = CVT$$(D$,32$ + 4$ + 2$ + 8$)
                                                    ICONVERT ENTRY TO USABLE DATA
         IF D$ = "END" THEN 32766
                                                    IIF END GOTO END
160
         IF D$ = "HELP" THEN GOSUB 1000
                                                    IGOSUB HELP PARAGRAPH
170
180
        IF D$ = "DEASSIGN" THEN GOSUB 1100
                                                    IDEASSIGN SECTION
190
        TE D$ = "ASSIGN" THEN GOSUB 1200
                                                    LASSIGN SECTION
        IF D$ = "SEIZE" THEN GOSUB 8000
                                                    ISEIZE A TERMINAL IN USE
195
                                                    ICHANGE LIMIT PARAMETERS
197
        IF D$ = "CHANGE" THEN GOSUB 9000
        IF D$ = "LIST" THEN GOSUB 1300
                                                    ILIST ALL ASSIGNED DEVICES
200
                                                    ILIST ALL UNASSIGNED DEVICES
210
        IF D$ = "FREE" THEN GOSUB 1400
        IF D$ = "DETACH" THEN GOSUB 1500
                                                    IDETACH OUT
215
217
        GOSUB 2000
                                                    ITIME CHECK
        IF D$ = "VIEW" THEN GOSUB 9500
                                                    IVIEW CERTAIN DATA
218
                                                    ISEE IF ANYONE TALKED TO US
        GOSUB 5000
219
        X = SYS(CHR$(6%) + CHR$(9%) + CHR$(0)) !GET KB NUMBER
220
223
        IF ASCII(MID(X$,2%,1%)) > 80%
                 THEN SLEEP 1%
                 GOTO 217
                                                    IDON'T ASK OUESTION IF DET.
1
                                                    IBACK TO LIM QUESTION
230
        GOTO 140
```

GOTO 63 IF LEN(B\$) = 5% PRINT "?Time must be in the form XX:XX"

60

	1
~	I* PRINT HELP MESSAGES
1	1
1005	PRINT IHELP MESSAGE PRINT "ASSIGN";TAB(30);
1015	"Assign a KB to Limit" PRINT "CHANGE";TAB(30);
1020	"Change a Limit feature" PRINT "DEASSIGN";TAB(30);
1030	"Deassign a KB from Limit" PRINT "DETACH";TAB(30);
1040	"Detach the job" PRINT "END";TAB(30);
1050	"End Limit now" PRINT "FREE";TAB(30);
1060	"List free terminals" PRINT "HELP";TAB(30);
1070	"This help message" PRINT "LIST";TAB(30);
1075	"List all assigned KB" PRINT "SEIZE";TAB(30); "Seize a terrinel which is inves"
1077	"Seize a terminal which is inuse" PRINT "VIEW";TAB(30); "View limit parameters"
1080	RETURN IFINSHED
1100	1.
1	I DEASSIGN A KB:
1	1
1105 1107	INPUT "Deassign KB:";N\$!DEASSIGN A KB:
\ 1110	GOTO 1105 IF N# = 0#
	+ STRING\$(20\$,0\$) + "KB" + CHR\$(N\$) + CHR\$(255\$))
1115	IDEASSIGN A DEVICE X = SYS(CHR$(6$) + CHR$(-5$) + CHR$(N$)$
	+ "\$Terminal now available" + CHR\$(13\$) + CHR\$(10))
1120 1130 1140	
1200	
1	I* ASSIGN KB:
1]#]***********************************
1203 1204	
\ 1205	GOTO 1203 IF N% = 0%
	THEN PRINT "\$Device already assigned"
`	GOTO 1250 IERROR MESSAGE
1207 \ 1208	T\$ = (ASCII(MID(X\$,2,1))/2) IT\$ HAS KB NUMBER
\	assigned" IF T≸ = N≸ GOTO 1250 IF T≸ = N≸ ICAN'T ASSIGN TERMINAL
	AT WHICH LIMIT IS CURRENTLY RUNNING
1210	+ STRING\$(20%,0%)
	+ "KB" + CHR\$(N\$) + CHR\$(255\$)) IASSIGN THE KB CODE
1215	<pre>X\$ = SYS(CHR\$(6\$) + CHR\$(-5\$) + CHR\$(N\$) + "\$System will be limited" + " until: "</pre>
	+ B\$ + CHR\$(13\$) + CHR\$(10)) ISEND SYSTEM LIMITED PHRASE
1220 1250	A(N%) = N% PLACE TKB IN ASSIGNED TABLE
1	I* LIST ROUTINE
1	1
1310	FOR X = 9 TO JX ILOOP THRU KB:
1320 1330	PRINT "KB";X;": "; IF A(X) > 0 !PRINT IF ASSIGNED
	PRINT
	1
1400	1 FREE ROUTINE
1	
1410	
1420	PRINT "KB";X;": "; IF A(X) = 0 !PRINT IF UNASSIGNED
1435	PRINT
	1
1,	I DETACH THE LIMIT JOB
	1

1510 PRINT "Detaching....." X\$ = SYS(CHR\$(6%) + CHR\$(7%)) PRINT DETACHING ON KB 1520 ITHE DETACH CODE 1530 RETURN . TIME CHECK 2010 IF B\$ < TIME\$(0) THEN 32766 ITIME TO STOP 2030 RETURN **!FINISHED** 5000 I* MULTI-TERMINAL SERVICES FIELD #12%, 7% AS 29\$ 5010 LSET 29\$ = "" GET 412\$, RECORD 32767\$ + 1\$ + 16384\$ + 8192\$!SEE IF ANYTHING OUT THERE IF MID(29\$ + "",2\$,6\$) = P9\$ THEN 6000 !INPUTTED PASSWORD PRINT #12\$, RECORD 32767\$ + 1\$ + ASCII(LEFT(29\$ + "",1\$)), C\$ C\$ C\$ C\$ 5020 5030 5040 5050 IRETURN IS DONE FROM ERROR R. 5060 RETURN !* PART OF MULTI-TERMINAL SERVICES, !* DEASSIGN TERMINAL X\$ = SYS(CHR\$(6%) + CHR\$(11%) 6010 + STRING\$(20\$,0\$) + "KB" + LEFT(29\$ + "",1\$) + CHR\$(255\$)) IDEASSIGN A DEVICE 6015 X\$ = SYS(CHR\$(6%) + CHR\$(-5%) + LEFT(Z9\$ + "",1%) + "%Terminal is now available" + CHR\$(13%) + CHR\$(10))ISEND SYSTEM LIMITED PHRASE IDELETE FROM ASSIGNED TO TABLE A(N%) = 0% B(N%) = N% 6020 6030 6040 IADD TO UNASSIGNED TABLE GOTO 5010 7000 I* SELECTED TERMINALS ONLY. PRINT "Enter terminal numbers"; 7010 " (end with <cr>)" INPUT "KB number";N% INPUT TERMINAL NUMBER 7020 INPUT *KE number , ... GOTO 125 IF N\$ = 0\$ X\$ = SYS(CHR\$(6\$) + CHR\$(9\$) + CHR\$(0)) IGET KE NUMBER X\$ = SYS(CHR\$(6\$) + CHR\$(9\$) + CHR\$(0)) IGET KE NUMBER IT\$ HAS KE NUMBER 7030 ITERMINATE IF ZERO 7040 7050 7060 IDON'T ASSIGN CURRENT TERMINAL 7070 GOSUB 1205 IGOTO ASSIGN ROUTINE AND ASSIGN TERMINAL 7080 GOTO 7020 IKEEP PROCESSING 8000 I* SEIZE A TERMINAL INPUT "Terminal to seize";N\$ PRINT "\$Console can not be seized" IF N\$ = 0\$ GOTO 8010 IF N\$ = 0\$ PRINT "\$Terminal already limited" IF A(N\$) <> 0\$ GOTO 8010 IF A(N\$) <> 0\$ GOTO 8010 IF A(N\$) <> 0\$ CHEC(CHEC(64) - CHEC(04) + CHEC 8010 IINPUT THE TERMINAL NUMBER 8020 IERROR, CONSOLE TERMINAL 8030 ITERMINAL ALREADY ASSIGNED X\$ = SYS(CHR\$(6\$) + CHR\$(9\$) + CHR\$(0)) T\$ = (ASCII(MID(X\$,2,1))/2) 8040 IT% HAS KE NUMBER PRINT "%Current terminal can not "; "be assigned" IF T% = N% GOTO 8100 IF T% = X% 8042 ICAN'T ASSIGN TERMINAL AT WHICH LIMIT IS CURRENTLY RUNNING X\$ = SYS(CHR\$(6\$) + CHR\$(-4\$) + CHR\$(N\$) + CHR\$(3\$)) FOR X = 1 TO 3 8044 IFORCE 3 °C TO TERMINAL X\$ = SYS(CHR\$(6\$) + CHR\$(-4\$))8050 + CHR\$(N\$) + "BYEF" + CHR\$(13) + CHR\$(10)) IFORCE THE BYEF 8060 INIT COUNTER FOR ATTEMPTS U9% = 0% TO SEIZE TERMINAL X\$ = SYS(CHR\$(6\$) + CHR\$(10\$) + STRING\$(20\$,0\$) 8070 + "KB" + CHR\$(N%) + CHR\$(255%)) ASSIGN THE KB CODE 8080 X\$ = SYS(CHR\$(6\$) + CHR\$(-5\$) + CHR\$(N\$)
+ "\$System will be"
+ " limited until: "
+ B\$ + CHR\$(13\$) + CHR\$(10)) ISEND SYSTEM LIMITED PHRASE IPLACE TKB IN ASSIGNED TABLE IFUNCTION FINISHED 80 90 $A(N_{s}) = N_{s}$ 8100 RETURN

1

1

1

1

9000 \	1	**
1	<pre>!* CHANGE PARAMETER ROUTINES !*</pre>	
1	1 *************************************	**
9010 9020	De - OIE TE OIE () MM	INEW TIME TO STOP INULL MEANS NO CHANGE
9030 9040	INPUT "Password change";Q1\$	INEW PASSWORD
1	PRINT "%Password must be 6 characters";	
1	" long" IF LEN(Q1\$) <> 6 GOTO 9030 IF LEN(Q1\$) <> 6	
~	P9\$ = Q1\$	PASSWORD VERIFIED AND CHANGED
9050 9060 9070	INPUT "Message change";Q1\$ C\$ = Q1\$ IF Q1\$ <> "" RETURN	INEW MESSAGE INULL MEANS NO CHANGE IFUNCTION FINISHED
1	!* !* VIEW ROUTINE	
9510	PRINT "Time to stop: ";B\$	PRINT TIME TO STOP
9520 9530	PRINT "Password is: ";P9\$	PRINT PASSWORD
9540		FUNCTION FINISHED
1.		***
,	I* I* ERROR ROUTINE	
1	ii	•••
31005	IF ERR = 28 THEN RESUME 32760	1°C TYPED
31007	IF ERR = 11 AND ERL > 130 THEN RESUME 219	
١.	THEN RESUME 219	1°Z TYPED AND BEYOND
31010	IF ERL = 17 THEN RESUME 31900	LINE 140 IN CODE IMASTER TERMINAL NOT AVAILABLE
31100		IDEASSIGN FAILED
1	RESUME 217 IF ERL = 1210	!KB WAIT EXHAUSTED
51120	THEN PRINT "\$Device not";	
١	" available" RESUME 1250	
31125 31130	TE ERI. = 8070 THEN RESUME 31800	ITERMINAL SERVICE ERROR ISEIZE DID NOT WORK
31140 31190	IF ERL = 31930 THEN RESUME 31910	IERROR ON NEW MASTER TERMINAL
31200		IUNKNOWN ERROR-TERMINATE
31800 \	1	***
1	I SIEZE TERMINAL ERROR SECTION	
1	1	***
31810	U9% = U9% + 1	IINCREMENT COUNTER
31820 31830		ISLEEP ONE SECOND ITRY 10 TIMES TO GET IT
	PRINT "\$Terminal can not be seized"; " - try again"	
31850		IGOTO RETURN STATEMENT

1	I MASTER TERMINAL NOT AVAILABLE	
1	1	
31910	PRINT "?Master terminal not available"	
	PRINT "Input another terminal to be"; " the master";	
1	INPUT K%	IGET NEW MASTER TERMINAL
`		NUMBER
31930 31940	OPEN K\$ AS FILE #12% GOTO 20	IOPEN THE MASTER TERMINAL IRETURN PROCESSING
32760		
1	!# !# EOJ	
~	i	
32766	X\$ = SYS(CHR\$(6\$) + CHR\$(12\$))	IDEASSIGN ALL DEVICES
32767	X\$ = SYS(CHR\$(6\$) + CHR\$(-5\$) + CHR\$(X)	
	+ " \$LIMIT going offlineSYSTEM back " + "to normal"	
	+ CHR\$(13) + CHR\$(10\$)) IF A(X) <> 0 FOR X = 9 TO J\$	
1	X\$ = SYS(CHR\$(6%) + CHR\$(-5%) + CHR\$(0%)	
	+ "\$LIMIT going offlineSYSTEM back" + " to normal"	
	+ CHR\$(13) + CHR\$(10%))	
`	X\$ = SYS(CHR\$(6\$) + CHR\$(9\$) + CHR\$(0\$))	
1	X1\$ = SYS(CHR\$(6\$) + CHR\$(8\$) + CHR\$(PEEK(518)/2)	
	+ SPACE\$(23) + CHR\$(0%) + CHR\$(255%)) IF ASCII(MID(X\$,2,1)) > 80%	
١.		ALL DONE AND OVER, DETERMINE WHETHER
		TO KILL THE JOB OFF
		OR NOT THEN DO IF NECESSARY
*		

A MENU MANAGEMENT PROGRAM

By R. David Broom Omega Optical 13515 N. Stemmons Dallas, TX 75234 (214) 241-4141

INTRODUCTION

Every site provides some way of informing its personnel what options are available to them. This is usually accomplished through a menu display. Once the list of options is displayed, the operator keys in a numeric entry or a mnemonic name for the option desired and the menu program chains to the program associated with that option.

OBJECTIVES and **GOALS**

"MENU" is a menu management program where mnemonic names are associated with the option desired. In the development of "MENU" four goals were set:

1) Lower the number of keystrokes entered by an operator,

2) Lower the number of characters outputted to a terminal,

3) "Hide" the fact that the operator is executing a system program like PIP, DTR, ATPK, et cetera,

4) Allow the operators to do as much as possible on one line of entry. That is, pass user-selected parameters.

Implementing these four goals will give the appearance that the operator is doing more with less work. For example, if we have a program that cancels a customer's order, then under this type of menu management program the operator keys in the following:

ME CAN ACCOUNT.INVOICE

where ME means menu

CAN means cancel account number, invoice number

If we did not allow the passing of parameters, the operator would have typed the same number of keystrokes, but not on the same line. He/she would also have waited for the menu display, and the display of the cancellation program. Some general goals where added:

1) Provide the ability to page within a menu,

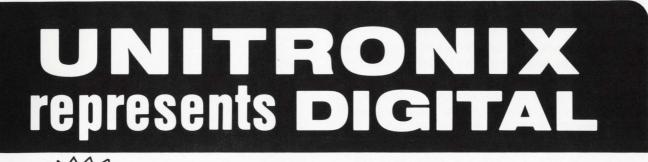
2) Call one menu while in another menu,

3) Provide operator security checks,

4) Pass user parameters by CCL and chaining,

5) Develop a menu that is easy to modify,

6) Display options according to terminal type.





TERMINALS VT100 LA12 VT101 LA34 VT102 LA50 VT125 LA100

VT131 LA120

CPUs PDP 11/23, 11/44 & VAX

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IMPLEMENTATION

To accomplish some of these goals a formatted text file was developed. This text file contains the mnemonic names of options, their description, what action is to be performed, program name or CCL or command file or DTR procedure or menu file to use with the action, line number to chain to and what terminal type can access this option. This type of text file is easy to modify and maintain. Figure 1 is an example of such a text file and Figure 2 is the display associated with it.

Security checks were implemented by using another action code: "O". The ways of doing a security check are infinite. One way is to create an operator file where each operator has a unique ID, and flag what secure options an operator can access. This is currently the type of system we are using.

Terminal type is checked because we have a mixture of VT52s and VT100s. Certain tasks were written to take advantage of the VT100's display options.

An interesting method of implementing this menu system is to apply the feature patch 21.3.5 which modifies the RSX run time system to chain to "MENU". An article with a related idea appears in the December issue of RSTS PRO, "A User Written Keyboard Monitor," by Ken Harris. A few modifications have to be made to "MENU" for this to work, such as the removal of the wait for entry.

Another approach is the feature patch 10.12.4 to modify "LOGIN" to chain to a specific program and add a CCL for "MENU". This is the approach we took. The entire menu is always displayed when our operators first log-in. Once an operator is logged in there is no need to redisplay the menu unless he/she is unsure of the options. To enter an option without displaying the menu listing, the operator enters a CCL for menu followed by an option, followed by parameters if they are allowed.

CONCLUSION

This type of menu system seems to work effectively. "MENU" accomplished our major goals and as a side benefit, operator training time has markedly decreased.

APPENDIX

Digital Equipment Corporation, "Invoking a Menu Program RSX.RTS Feature Patch," RSTS/E V7.2 Maintenance Notebook, (June 1982), Seq 21.3.5.

Digital Equipment Corporation, "Login Can Chain to a Specific Program - LOGIN Feature Patch," RSTS/E V7.2 Maintenance Notebook, (June 1982), Seq 10.12.4.

Harris, Ken, "A User Written Keyboard Monitor," RSTS Professional, (December 1982), Vol. 4, Num. 6

```
Customer Service Menu Selection ...

1;ACCOUNT;A/R account inquiry;C;INQUIR.BAC;0

;BYE;Log off the system;L;BYE/Y;0

;CLEAR;Clear all temporary files;T;CLEAR.PIP;0

;INQUIRY;General laboratory inquiry ;L;INQ;0

;LIST;Listing of current transactions for an account;P;ACCOUNT-TRANS;0

;MIN;Main menu for Lab;S;$MAIN.MEN;0

;NIGHT;Produce nightly reports;A;NIGHT.LOG=NIGHT.ATP/DET;0
```

FIGURE 1

1	Customer	Service Inquiry Menu Selection
!		
[ACCOUNT	- A/R account inquiry
]	BYE	- Log off the system
]	CLEAR	- Clear all temporary files
]	INQUIRY	- General laboratory inquiry by Job/Tray/Account
[LIST	- Listing of current transactions for an account
]	MAIN	- Main menu for Lab
[NIGHT	- Produce nightly reports

OPTION>

FIGURE 2

	AM: MENU.B2S
	ON: 1.1
	TE: 03-Nov-82
	OR: R. David Broom
51	TE: Omega Optical, Dallas,Tx

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•••••	
<prog< td=""><td>RAM DESCRIPTION></td></prog<>	RAM DESCRIPTION>
	Menu management program.
	PURPOSE: to provide control and entry for multi-
	menu system.
	USE AND INSTALLATION:
	1. Install CCL for this program (@ 30000)
	CCL name can be any appropriate name.
	2. This program chains and/or CCLs on filespecs
	provided in the control file MENU.TXT.
	3. Build MENU.TXT using the format listed below.
	MENU.TXT must reside in the user account.
	The first line is a title for the menu and cannot
	exceeded 60 characters. The following lines are menu
	description lines.
Menu :	line construct:
T;OPT.	IONNAME; TEXT; ACTIONCODE; PROG. NAME; LINENO
	WHERE:
	T=1 OPTION valid for VT100 only.
	T=5 OPTION valid for VT52 only.
	T=H OPTION valid for hardcopy only.
	T=" " OPTION valid for any terminal type
	OPTIONNAME
	up to 12 character name for menu option. TEXT

optional descriptive text (up to 60 characters).

	ACTIONCODE		8			
		e ATPK command file PROG.NAME	4			
		to PROG.NAME at LINENO passing	&			
		ther text in core common.	&			
		e DTR command file PROG.NAME	&			
	L = execut	e CCL using PROG.NAME as CCl and	&			
	passir	ng any other text as switches.	&			
		IEVE procedure at PROG.NAME	&			
	S = Switch	to new menu control file at PROG.	NAME &			
	T = Execut	e a PIP command file	\$			
	PROG.NAME		å			
	Execute na	ame (program name or CCL)	&			
	LINENO		&			
	Line numbe	er used in chaining	å			
All ite	All items are separated by ";" &					
NOTES:	NOTES:					
			&			
MENU w:	111 page. Up to 15 :	items are displayed per screen.	å			
This ca	an be modified by in	ncreasing the arrays and changing	&			
MENU.L.	MIT%		å			
			&			
***********	****************	*****	****&			
01**********	*****************		****& &			
	COMMON DE	CLARATION	&			
COMMON	M% (30%)	! Sys call array for termina	& 1 &			
controll	,KB.CHN%	! Keyboard channel	&			
	MENU.CHN%	! Menu channel	&			
	. PAUS. TIME%	! Time limit on user input	&			
	MENU.LIMIT%	! Limit of menu items to dis	play&			
	, OPTION%	! Selected option	&			
	.OPT.LEN%	! Length of option entry	&			
	, N%	! Menu item counter	&			
	, VTCHK%	! Terminal specification	å			
	A MAXIMUM A A COUNTY					

! Line number used in chaining

! Text to display

! Restart menu display

! Title of menu

,TERM.FLAG\$=1%,OPTION\$(15%)=12% ! Type of terminal for display ! Option names ,ACTION\$(15%)=1% ! Action to perform , PROG. NAME\$(15%)=26% ! Execute name for program MAIN PROGRAM ON ERROR GOTO 19000 KB.CHN% = 1% MENU. CHN% = 2% PAUS.TIME%=60% MENU.LIMIT#=15% MEIN.MENU\$="MENU.TXT" ! Default menu, TH: OPEN "_KB:" AS FILE KB.CHN\$,MODE 256\$ X\$=SYS(CHR\$(11\$)+CHR\$(KB.CHN\$)) ! Clear type ahead ! Default menu, THIS ACCOUNT GOSUB 14000 ! Test for VT100/VT52/HARDCOPY OPTIONS =0% IF CCL% GOSUB 10000 ! Read MENU GOSUB 12000 IF OPTION% ! Execute if possible THEN 1010 CCL%, OPTION%, OPT.LEN%=0% OPT\$="" GOSUB 10000 ! Initialize MENU GOSUB 11000 GOSUB 12000 IF OPTION% ! Interpret user request ! Execute if possible & &

10000!	***************************************
:	
1	LOCAL SUBROUTINES
!	
1	Initialize the menu settings
1	
	OPEN MAIN.MENU\$ FOR INPUT AS FILE MENU.CHN%
1	LINPUT #MENU.CHN%,TITLE\$
	! Open menu file and get title
10010	N% = 1%
1	IF CCL%=0%
	THEN X%=FN.CLEAR.SCREEN%
1	PRINT TAB(3%); STRING\$(73%,61%)
\	PRINT "! ";TITLE\$;TAB(79%);"!"
1	PRINT "!";TAB(79%);"!"
	! If no CCL entry then clear
	! screen and display title
	· corteen und diopray exert
10020	LINPUT #MENU.CHN%, MENU.LINE\$
1	TERM. FLAG\$=LEFT\$(MENU.LINE\$,1%)
1	GOTO 10020 IF TERM.FLAG\$="H" AND VTCHK\$<>1\$
	! This is a hardcopy option
\	GOTO 10020 IF TERM. FLAG\$="5" AND VTCHK\$<>2\$
	! This is a VT52 option
1	GOTO 10020 IF TERM. FLAG \$="1" AND VTCHK\$<>3%
`	
	! This is a VT100 option

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GOTO 1010

, VTCHK% ,LINENO%(15%)

. TEXT\$=60%

,TITLE\$=60%

SC.POS\$=POS(MENU.LINE\$,";",1\$) SC.POS1\$=POS(MENU.LINE\$,";",SC.POS\$+1\$) OPTION\$(N\$)=SEG\$(MENU.LINE\$,SC.POS\$+1\$,SC.POS1\$-1\$) 2 ! Option name SC. POS%=POS(MENU.LINE\$, "; ", SC. POS1\$+1\$) TEXT\$=SEG\$(MENU.LINE\$, SC. POS1\$+1\$, SC. POS\$-1\$) ! Option description SC. POS1\$=POS(MENU.LINE\$, "; ", SC. POS\$+1\$) ACTION\$(N\$)=SEG\$(MENU.LINE\$, SC.POS\$+1\$, SC.POS1\$-1\$) ! Action to perform for option 2 SC. POS\$ = POS(MENU.LINE\$, "; ", SC. POS1\$+1\$) PROG. NAME\$(N\$) = SEG\$(MENU.LINE\$, SC. POS1\$+1\$, SC. POS\$-1\$) ! Program or command file to use& LINENO\$(N\$)=VAL(RIGHT\$(MENU.LINE\$,SC.POS\$+1\$)) & ! Chaining point for program OPTION\$=N\$ IF OPT\$=LEFT\$(OPTION\$(N\$),OPT.LEN\$) AND CCL\$ 2 & RETURN IF OPTIONS I If a CCL entry and option ! matches then save option value& and get out. GOTO 10020 IF CCL.\$! If a CCL entry do not display & ! option line PRINT "[";OPTION\$(N\$);" - ";TEXT\$;TAB(79\$);"]" N%=N%+1% GOTO 10020 IF N%<=MENU.LIMIT% PRINT PRINT "[To PAGE hit [RETURN] ... "; TAB(79%); "]" 2 10030 RETURN IF CCL. I Should never get here on CCL \$! unless no match is found 2 PRINT PRINT " Option> "; r WAIT PAUS.TIMES LINPUT #KB. CHNS. OPT\$ OPT\$=EDIT\$(OPT\$,32\$) I Make it upper case GOTO 10010 IF OPT\$="" AND N\$>MENU.LIMIT\$! Page for other options & 10040 CLOSE #MENU. CHNS ٤ RETURN Determine if user input is executable X%=POS(OPT\$,SP,1%) USER.DATA\$=EDIT\$(RIGHT\$(OPT\$,X\$+1\$),8\$+128\$+32\$) IF X\$! Save rest of entry line OPT\$=LEFT\$(OPT\$,X%-1%) IF X% & OPT.LEN%=LEN(OPT\$) RETURN IF OPT.LEN%<3% ! Input must be at least 3 chars& RETURN IF LEFT\$(OPT\$, I\$)=LEFT\$(OPTION\$(OPTION\$), I\$) FOR OPTION\$=1\$ TO N\$ 2 FOR IS=OPT. LENS TO 3% STEP -1% & ! Loop through all options ! trying to match user input 2 ! with an option & 11099 OPTIONS =0% ! No match, so return back with & ! a bad request RETURN Process MENU function code A] C] Execute ATPK command file PROG.NAME Chain to PROG.NAME at LINENO 1. 2. Execute DTR command file PROG.NAME Execute CCL PROG.NAME with USER.DATA\$ D 3. I. 1 4. Execute DATATRIEVE procedure PROG.NAME P 5. S 1 6. Switch to MENU PROG.NAME Т Perform a PIP command file 1 7. X\$=FN. PUT. CORE. COMMON\$("") ! Clear core common ON POS("ACDL PST", ACTION\$(OPTION\$), 1\$) GOSUB 12100, I[A]] 12150. 1[C 12200 . I[D]] I[L]I 12250, I[P]! 12300, 12350 . I[S]] 12400 I[T]I 12010 RETURN PROG.NAME\$="PER "+PROG.NAME\$(OPTICN\$) 12100 X%=FN.CLEAR.SCREEN% X\$=FN.CCL\$(PROG.NAME\$) ! Clear screen ! Ferform syscall for CCL ! Execute ATPK file PROG.NAME 2 RETURN ! Clear screen X%=FN.CLEAR.SCREEN% 12150 X\$=FN.PUT.CORE.COMMON\$(USER.DATA\$) IF LEN(USER.DATA\$) ! Put DATA in core common & CHAIN PROG. NAME \$ (OPTION\$) LINE LINENO\$ (OPTION\$) ! Chair to PROG.NAME at LINENO RETURN & PROG. NAME \$= "DTR @"+FROG. NAME \$(OPTION\$) 12200 ! Clear screen ! Perform syscall for CCL X%=FN.CLEAR.SCREEN% X\$=FN.CCL\$(PROG.NAME\$) ! Execute DATATRIEVE command file PROG.NAME å RETURN PROG.NAME\$=PROG.NAME\$(OPTION≸) PROG.NAME\$=PROG.NAME\$+" "+USER.DATA\$ IF LEN(USER.DATA\$) 12250

X%=FN.CLEAR.SCREEN% I Clear screen X\$=FN.CCL\$(PROG.NAME\$) ! Perform syscall for CCL
! Execute CCL PROG.NAME + user input RETURN & 12300 PROG. NAME \$= "DTR : "+PROG. NAME \$(OPTION\$) ME\$(UFILUM;) I Clear screen I perform syscall for CCL X%=FN. CLEAR. SCREEN% X\$=FN.CCL\$(PROG.NAME\$) RETURN ! Do DATATRIEVE procedure PROG. NAME 12350 MAIN.MENU\$=PROG.NAME\$(OPTION\$) ! Get next menu RETURN 12400 PROG.NAME\$="PIP @"+PROG.NAME\$(OPTION\$) X%=FN.CLEAR.SCREEN% Clear screen X\$=FN.CCL\$(PROG.NAME\$) ! Perform syscall for CCL RETURN Determine Terminal Characteristics CHANGE SYS(CHR\$(6\$)+CHR\$(16\$)+CHR\$(0\$)+CHR\$(255\$)+ STRING\$(24%,0%)) TO M% ! Get terminal characteristics VTCHK% = 1% Hardcopy terminal RETURN IF M%(11%)=128% NO SCOPE set VTCHK%=2% VT52 RETURN IF M%(9%)=128% NO XON set PRINT #KB.CHN%, CHR\$(155%); "[c"; ! Ask for VT100 I.D. WAIT 1% ! Wait only 1 second GET #KB. CHN% X%=RECOUNT MOVE FROM #KB.CHN%, X\$=X% VTCHK\$=3\$ IF POS(X\$, "[?1;", 1\$) ! This should work for VT100 and CIT-101 14020 RETURN LOCAL FUNCTIONS DEF FN.CLEAR.SCREENS PRINT CHR\$(155%)+"[2J"+CHR\$(155%)+"[0:0H": IF VTCHK#=3% PRINT CHR\$(155\$)+"H"+CHR\$(155\$)+"J"; IF VTCHK\$=2\$ PRINT FNEND 15100 DEF FN. PUT. CORE. COMMON \$(D\$)=SYS(CHR\$(8\$)+D\$) ! Put DATA D\$ in core common 15200 DEF FN.CCL\$(D\$)=SYS(CHR\$(14%)+D\$) I Do CCL STANDARD ERROR HANDLING RESUME 32767 IF ERL=10030 ! User is quitting on time-out or CTRL/Z RESUME 14020 IF ERL=14000 Time-out waiting on VT100 I.D. call RESUME 11099 IF ERL=11000 ! Subscript out of range on a CCl entry RESUME 10030 IF ERL=1020 AND ERR=11% ! End of file on read RESUME 10020 IF ERL=1020 AND (ERR=52% OR ERR=55%) 1 ! Line number specification is illegal number ! or subscript out of range IF ERR=5% 19010 PRINT "That option does not exist" IF ERL=12150 THEN 1 **RESUME 32767** ! Tried to chain to a non-existent program ! Menu file does not exist. Just act like ! nothing happen. PRINT "Error: ";ERR;" ";ERT\$(ERR);" at line ";ERL 19900 **RESUME 32767** CCL ENTRY POINT ON ERROR GOTO 19000 OPT\$=SYS(CHR\$(7%)) CCL\$=0\$ USER . DATAS=" X\$=POS(OPT\$, SP, 1%) OPT\$=RIGHT\$(OPT\$, X%+1%) GOTO 1000 IF LEN(OPT\$)<3% ! Assume not a CCL entry if ! input is less than 3 chars X\$=POS(OPT\$,SP,1\$) USER.DATA\$=RIGHT\$(OPT\$,X\$+1\$) IF X\$ OPT\$=LEFT\$(OPT\$.X%-1%) IF X% OPT.LEN%=LEN(OPT\$) CCL% = (OPT. LEN%>=3%) GOTO 1000 ! Pull-off any user parameters ٤ ! and save them & 32767 END &

NEW PRODUCTS

ROSS OFFERS STOCK OPTION SOFTWARE SYSTEM

A stock option software system that manages and tracks all employee stock option plan information is now available from Ross Systems. The new accounting and tracking system meets all current legislative requirements for incentive stock options (ISO) and provides complete information for public, tax and internal reporting of stock option programs.

Through its database design, Ross' Stock Option System gives users complete control over their option plan to accommodate changing regulations. As an interactive system, it tracks all activity relating to the granting, exercising, cancellation and repricing of employee stock options. Additionally, the system meets Federal Incentive Stock Option (ISO) accounting and reporting requirements. The Ross System is available to run on Digital Equipment Corporation's PDP-11 Series and VAX series minicomputers under RSTS/E and VMS operating systems. It provides a selection of standard reports, accommodates specific inquiries about participants and can interface with both company personnel and compensation systems. It can accept information regarding employee status or location on a daily basis and is compatible with a number of communications formats.

The Ross Stock Option System is based on Ross' interactive INTAC database management system for organizing and reporting business information. Ross Systems, Inc. provides financial decision support, applications and database software for financial managers in manufacturing, business services and banking. Software is available for Digital Equipment Corporation minicomputers or through Ross' worldwide timesharing network. For information, contact Jack Brown at Ross Systems, 1860 Embarcadero Road, Palo Alto, CA 94043, (415) 865-1100.

DIGITAL ANNOUNCES DECmail/RSTS

In 1860, the century's most modern method of delivering the mail galloped into the daily lives of many of the people of this country. Relays of swift horses and unflinching men — the Pony Express — raced from St. Joseph, Missouri to Sacramento, California in record time.

Later came Wells Fargo coaches and four. They sped mail, cargo and passengers ever faster. Trains and planes followed. But even they were not the ultimate. Because now, in 1983, an entirely new and incredibly faster method of delivering messages has just been announced by Digital.

This new method enables RSTS/E users to receive and send messages to business associates down the hall, across town, around the country and overseas — elecronically. To individuals or entire groups all at once. To file and retrieve and edit these messages almost without effort at enormous savings in cost and time. It's called DECmail/RSTS and it's here now.

DECmail/RSTS was sired, reared and broken in just for you. It was designed by the original RSTS/E development team — the computer gang that knows the RSTS/E system intimately. Their experience and expertise has made DECmail/RSTS a downright superior product.

What does DECmail do?

Everything an electronic message system should. With DECmail, you can create, edit, send and process messages. The message could be a simple memo: "Your preliminary figures for quarter just ended seem too good to be true. Please check and verify at once."

Or a complicated specification completed overnight and forwarded for first morning review: "Report rechecked and confirmed. Here's the detailed breakdown for immediate action."

DECmail also stores messages it receives while you are out. Before you begin your work in the morning, for example, it brings you up-to-date. "You have 3 new messages." And it shows you their essence and whom they're from. There's more.

As for filing systems, DECmail will organize your material for you easily, efficiently, electronically. It stores, searches for and retrieves messages held in your files.

If this sounds as easy as falling off a horse, it is. DECmail designers knew RSTS/E software had to be sophisticated enough for advanced users and yet simple and uncomplicated for computer novices. DECmail is both.

For beginners, easy-to-learn commands and flexible options make learning a cinch. If there's trouble along the way, the on-line help facility can solve problems fast. For advanced users, DECmail offers many features including the flexibility to support many project teams by sharing specified folders. This way, everyone keeps abreast of the latest data. Managers, project leaders, researchers, administrators — all communicate quickly, easily and more effectively with DECmail.

Who can use DECmail/RSTS Version 1.0?

DECmail runs on RSTS/E Version 7.2 and 8.0. No update will be needed for DECmail/RSTS until after RSTS/E V8.0 is released. Thanks to DECmail update kits, you won't be saddled with an obsolete software product down the line.

Specifically, DECmail/RSTS allows single-node use on a single RSTS/E system. When used with DECnet/E, it provides multi-node use. Multiple standalone RSTS/E systems running DECmail/RSTS can be connected in a network, allowing users on one system to exchange mail with users on other systems. VAX/VMS systems running the VMSmail facility also can exchange messages with systems running DECmail/RSTS through DECnet.

DECmail/RSTS can be used from any RSTS/E command terminal. Supported terminals include word processing systems (WS78, WS278 and DECmate) running CX protocol.

All this is available now for only \$3,000 per system (only \$1500 per additional license). And just as important, DECmail can increase productivity and improve your working environment. It is a new way for people in business to communicate with each other.

Order your DECmail/RSTS software

now. As a gesture of appreciation for acting promptly, Digital will send you an authentic western belt buckle as a gift. Only one per pardner, please. Call toll free (800) DEC-INFO any working day between 8:30 a.m. and 5:00 p.m. EST.

NORTHWEST DIGITAL SOFTWARE ANNOUNCES RPM — RSTS

RPM is a new performance optimization package from Michael Mayfield and Northwest Digital Software. RPM can drastically improve total system performance by identifying problem areas in system usage, the programs that caused the problems and the problem areas within each program. RPM is the only product that can do all this.

System tuning with RPM uses a step by step "cookbook" approach. No knowledge of system tuning or monitor internals is required. Problem areas are identified using an automatic procedure which provides a report describing, in plain English, not numbers, where the system performance can be improved. It even makes suggestions for improvement.

On-line plotting, histograms and other reports can be used to further identify problem areas. Extended monitor data collection allows plotting of information not normally available, such as seek distance, disk usage, and cache, memory, file processor (FIP), and small buffer utilization. The programs causing the problems are then identified and can be examined in extreme detail. Detailed examination of a program includes CPU usage, a count of I/O requests and disk overhead by channel and a count of monitor calls and disk overhead by call. For information, contact Northwest Digital Software, Inc., Box 2-743, Spring Valley Road, Newport, WA 99156, (509) 447-2620.

REPORTING SYSTEM AVAILABLE FOR VAX

Enterprise Technology Corporation has announced the release of the Generalized Reporting System (GRS) for Digital Equipment Corporation VAX-11 computers. GRS is a userfriendly report writer, query and audit package designed for use with a wide variety of file formats and database systems, and is a low-cost alternative to Digital's Datatrieve facility.

GRS is compatible with the applications packages from most major software vendors, as well as with systems developed by end users. It includes a powerful English-like query language which can be used by both programmers and non-programmers. With GRS, management personnel can meet their own needs for ad-hoc information retrieval, and systems personnel can vastly decrease the amount of time spent responding to user requests for reports. GRS is also ideal for use by auditors.

The GRS software for VAX systems is supplied under a perpetual license for a one-time fee of \$6,500.00 which includes full documentation and one year of support. More information can be obtained by contacting Enterprise Technology Corporation at 305 Madison Avenue, New York, NY 10165, (212) 972-1860.

SPSORT FOR RSTS/E

The System Performance House, Inc. is pleased to announce a new high-performance RSTS/E sort called SPSORT.

- SPSORT can sort a file onto itself using no scratch space except for a 5 block command file.

- SPSORT can sort a file from one device onto another file on another device.

- SPSORT sorts a wide variety of data types including all Basic-Plus CVT types, all 8-, 16-, 32- and 64-bit signed and unsigned data types, FORTRAN IV-Plus 1*4, DIBOL decimal, packed, and EBCDIC. In addition, the user may specify any two byte-collating sequences desired.

- SPSORT will sort any type of fixed length record file whose record length is 512 bytes or less and whose total size is one billion bytes or less. There is no limit on record count.

- SPSORT is extremely fast. In one test, SPSORT sorted 100,000 64-byte records with 23 byte keys in 16 minutes on an 11/24 with RK07 disk drives.

 SPSORT has an extremely versatile calling format enabling the user to chain into and out of it in a number of ways.

Available with SPSORT are "frontend" modules which emulate OMSI SORT-1 Plus from Oregon Software, FSORT3 from E. G. and H., and DEC's SORTG/SORTM without change in user software!

SPSORT, complete with emulators, is attractively priced. The System Performance House will provide a superior sort at a lower price than its competitors. Fixed priced unlimited licenses are available to authorized DIGITAL distributors and OEM's at exceptionally low rates. For specifications and pricing information, contact the System Performance House, Inc., 5522 Loch More Court, Dublin, OH 43017, 614-265-7788.

VSELECT FOR VAX FROM EG&H

Evans Griffiths & Hart, Inc. has announced the release of a new product, VSELECT, a high speed machinelanguage record extractor for the VAX under VMS. The package handles RMS fixed-record-length sequential, relative, and indexed files (prologue 2 only).

VSELECT employs user-specified selection criteria to extract from one or more input files those records needed for specific report or application. Optionally, the extracted records are modified by the deletion, insertion, or rearrangement of fields. The records are then written to an output file that may be sorted by VSORT, EGH's fast machine-language sort package. VSORT may be invoked directly through VSELECT.

Selection keys for VSELECT include ASCII, signed and unsigned integer, floating point (single, double, giant, or huge), packed decimal, and the relative file record occupancy flag.

VSELECT supports a number of sophisticated options, among which are: (1) hooks for user-coded selection criteria and output field transformations; (2) conditional field copying; (3) the conditional generation of multiple output records from a single input record; (4) the generation of backpointers into the input file(s) —



You are invited to London International Press Centre

> on Tuesday 28th June 1983 From 10 am to 6 pm

to attend a seminar presented by

Carl Marbach Dave Mallery Al Cini

The Team from DEC Professional, RSTS Professional and Personal and Professional Magazines in the U.S.A.

Price: \pounds 50, if payment received before 6th June 1983 and \pounds 60 thereafter

> and includes: Morning Coffee & Biscuits Three-Course Lunch with Wine Afternoon Tea & Biscuits

For more information:

Digital Equipment Computer Users Society 2.0. Box 53, Reading, RG2 0TW Telephone: Reading (0734) 387725 Telex: 848327/8 including RFAs for RMS files; and (5) the insertion of string literals in output records.

VSELECT is callable from any of the VAX native mode languages that support the standard DEC procedure calling conventions. In addition, the package is supplied with a selfcontained interactive interface that can accept partial or complete command files.

A single-CPU license for VSELECT is \$1500.00, discounted to \$500.00 when the package is purchased with VSORT. OEM discounts are available. Contact Evans Griffiths & Hart, Inc., 55 Waltham St. Lexington, MA 02173. Tel: (617) 861-0670.

STAR PLAN INTRODCES MARCO LANGUAGE UTILITY

Star Plan Data Processing of Milwaukee, Wisconsin is pleased to announce MPR, a multipurpose macro language utility. MPR is a system engineering tool designed primarily as a program language preprocessor. The basic technical principles involved are the reduction in volume of the code body, and the addition of a level of abstraction above the source language. This decreases the cost of maintenance and increases the flexibility and portability of the application system.

Users of the utility are able to work it with anything that is kept in text files. Examples include BASIC+2, COBOL, DIBOL, TECO, BASIC+, RPG, C, RNO, PASCAL, LISP, ATPK and batch control files, and documentation.

MPR includes features for interactive applications and is frequently used to tailor control files according to the user's specific requirements.

The product's single most common application is sharing data definitions used by many programs in a system. For example, if you put record definitions for an application into external files and ".INCLUDE" them in the programs in the application, the entire application can be rebuilt with the new definitions after changing only one file.

With individual programs, "common" areas can be defined in one place for a program with many external subroutines. This can save considerable time when compared to the conventional technique of editing each module separately to bring them up to date.

Another practical application for MPR is source-level libraries of program code. By using MPR on standard system routines, programmers can have the advantages of standard libraries without the memory and runtime cost of automatically bringing in unnecessary code.

OEMs often need to create several versions of a software product with minor, but pervasive differences. Using MPR, it is possible to have a single body of source code for a class of very similar applications, which expands differently for different clients or operating systems. It is also possible to create an integrated software package and unbund'e various features by either including or excluding the relevant code when creating a client's system.

MPR is a highly versatile software tool that will benefit large scale application programming jobs. It has proven its adaptability to a wide range of specific in-house methodologies.

MPR is priced at \$875.00 for a single-CPU binary license and is available for all of Digital Equipment Corporation's PDP-11 series minicomputers using RSTS/E, and will soon be released under VAX VMS and RSX-11M+. Further information about MPR can be obtained from Star Plan Data Processing, 2040 W. Wisconsin Avenue, Suite #354, Milwaukee, WI 53233, ATTN: Noah M. Dixon, (414) 933-0800.

BASIC+, BASIC+2, DIBOL, RSTS/E, VAX VMS, and RSX11 are trademarks of Digital Equipment Corp.

MPR is a trademark of Star Plan Data Processing, Inc.

FREE LITERATURE FROM CROSS

Cross Information Company is offering a series of articles entitled: Advanced Business Communications: Applications for Computer Tele/conferencing. They are designed to help businesses and educational institutions develop and implement productive and cost-effective computer tele/conferencing on VAX systems.

The CIC series offers practical advice on planning, cost justification, and strategies for using computer tele/conferencing. CIC has organized these articles into specific applications:

Association management

• Improving software development and programmer productivity

• Managing the sales/marketing department

• Educational and training applications

• Uses in quality/study circles

• Telecommunications management

nanagement

Corporate strategic planning
General corporate management, and

u I G

• Legal office and education. Requestors are asked to indentify which articles are desired. There is no charge for these articles. There are available from Cross Information Company, Corporate Communications, 934 Pearl Mall — Suite B, Boulder, Colorado 80302-5181, 303-499-8888.

EDT MANUAL FROM COMPUTEREASE

DEC's EDT editor, available for its PDP-11 and VAX product lines, is a flexible, powerful editor. However, the complexity of the EDT manual has intimidated many of DEC's less experienced users, programmers as well as clerical workers. Many users learn as much as they need to know for their basic functions and then end their learning process, as far as EDT goes. As a solution to this problem, ComputerEase Publishing announces the availability of their latest publication, "Understanding EDT."

Recognizing that a good manual takes technical skill as well as good writing ability, ComuterEase teamed a programmer with an experienced technical writer to produce the manual. To further ensure the technical accuracy of the manual, all composition was performed using the EDT Editor itself.

"Understanding EDT" is especially

worthwhile for novices; it begins by explaining how to use a CRT and the basic concepts of using a text processor. The command section of the manual progresses in difficulty from the basic commands to the more advanced ones. In presenting the material this way, the aim is bring the user along gradually, so that he or she gains more capability with each succeeding chapter. "Understanding EDT" covers all three editing modes: keypad, nokeypad, and line mode, with the emphasis on keypad mode. In addition, the manual discusses strategies in using the editor and provides examples for applying it in text processing and programming environments.

The "Understanding EDT" manual contains more than 150 pages of instructional material. It is available from ComputerEase Publishing for \$49.95 per manual in quantities of nine or less, and for \$34.95 each in quantities of ten or greater. To place an order or for more information, contact Director of Marketing, ComputerEase Publishing, P.O. Box 390272, Mountain View, CA 94039.

PRODUCT UPDATES

SORTC FOR RSTS FROM JBM GROUP

The JBM Group announces that its SORTC Compounding Sort V2.5 is now available for PDP-11s using the RSTS/E operating system. This disk sort, written entirely in MARCO-11, is one of the many software products in JBM's MIDWARE line.

Originally written as a high-speed, direct replacement for DEC's SORTG/ SORTM DIBOL sort package, it is now expanded to handle DMS-500 Type 1 and Type 2 files. This permits use in BASIC-Plus and BASIC-Plus-2 shops that use native mode RSTS block I/O.

SORTC operates in standalone mode or is available in subroutine form to be called from within a DIBOL, MACRO-11 or BASIC-Plus-2 program. The compounding feature permits summing of selected numeric fields, consolidating data records and producing totals without the need for a post-processing routine. This reduces the number of times data records are shuffled about during sorting thus limiting both CPU usage and disk requirements.

SORTC may be fine-tuned by the user to balance system resource loading. Complete documentation includes numerous examples of option use.

SORTC is available from stock on either 800 or 1600 bpi 9-track tape. The standalone version, including basic sort mechanics, is priced at \$1,400.00. The optional subroutine capability is an additional \$500.00. If purchased together the price is \$1,700.00. For further information or to order SORTC, contact: The JBM Group, Inc. Dept. 22B, 332 West Church Road, King of Prussia, PA 19406, TWX 510-660-3999, TEL 215-337-3138.

PORTRAIT DISPLAY OPTION FOR AMBASSADOR

Ann Arbor has announced a portrait display configuration of the Ann Arbor Ambassador for text editing applications. This configuration uses a vertical screen to provide a full-page display in an $8-\frac{1}{2} \times 11$ format.

The Portrait Display Ambassador, unlike other full-page terminals, allows the user to select the display size most appropriate for the application. If small characters become tiring, the user can "zoom" part of the text off-screen to increase character size, yet instantly "zoom" it back to a full-page display for context or review and final editing. This multiple-format capability also permits use of existing programs written for conventional 24-line terminals while the user adds full-page programs to his or her library.

The Ann Arbor Ambassador is especially suited to text editing because it is one of the few interactive terminals available that can be made to locally rearrange data on the screen. In addition, its editing commands are controlled by hardware line pointers rather than firmware, which speeds up their execution and virtually eliminates the need for pad characters. Its ANSI-standard coding permits the use of decimal parameters in the edit commands so the application software can insert or delete multiple lines or characters without repetitively sending the same command. A meta shift is provided for use with editors, such as EMACS, that use the parity bit for data transmission.

Other features of the Ambassador include detachable keyboard with sixty programmable key levels, a split-screen operation with userdefinable scrolling regions, multiple page/window capability, printer output, and both block-mode and KSR operation. Options include DEC mode for VT100 software compatibility and a tilt/swivel stand.

The Portrait Display Ambassador is available from Ann Arbor distributors or from the factory direct. The suggested retail price is \$1795; delivery is from stock to 30 days. For further information, contact Kathryn Straith, Ann Arbor Terminals, Inc., 6175 Jackson Road, Ann Arbor, MI, 313/663-8000.

COMBOARD/3780 NOW AVAILABLE FOR VAX, PDP-11S

Software Results Corporation is announcing the addition of the ability to communicate via 3780 protocol to the COMBOARD_{TM} series of frontend processors for DEC VAX and PDP-11 systems.

The COMBOARD/3780 is a complete hardware/software system designed to function under the VMS, RXS-111M/M+ and UNIX operating systems. It provides an RJE link to IBM or any other mainframe supporting the standard 3780 protocol. It performs the communications processing at speeds up to 56 kbps. The off-loading of this processing from the DEC CPU allows the DEC machines to respond rapidly, processing user applications more efficiently.

The COMBOARD/3780 offers the ability to route files directly from the central system to end-user disk areas. The end-user is automatically notified

when files are placed. This saves the users time and provides security for the received data. Files can be previewed before valuable printing time is used. Graphics can be previewed before being spooled to plotting devices.

The COMBOARD/3780 also has the ability to act as a central system. This feature allows the DEC systems to communicate with PC's or any RJE station able to communicate via 3780 protocol.

End users make more cost effective use of both the central system and DEC machines because of the ease of use and efficiencies of the COM-BOARD.

COMBOARD/3780 is available in three models supporting transfer rates from 1200 to 56,000 bps. Prices range from \$9,800 to \$17,900 including hardware and software. For information contact Software Results Corp., 2887 Silver Drive, Columbus, Ohio 43211 (614/267-2203).

> HYBRICON OFFERS UNIBUS, Q-BUS BOARDS OFF-THE-SHELF

Hybricon Corporation offers wire

wrappable socket boards "off-theshelf." These boards are available with socket pins installed (4-format) or for use with DIP sockets (2-format). Hybricon's dual size 2-DE2-VHF can hold 52-16 pin DIPS, and costs only \$80.00 in unit guantities. The guad size, 2-DE4B-VHF, can handle 111-16 pin DIPS at a unit cost of \$140.70. Loaded with socket pins, the 4-DE2-VHF sells for \$198.60 and the 4-DE4B-VHF for \$358.30. Hybricon's patented pattern, VHF ground planes and VHF accessories enables you to mix analog and digital logic as well as high speeds (in excess of 250 MHz). These boards are directly compatible with Digital Equipment Corporation's hardware. Extender boards related accessories, wire wrapping services and custom design are all offered by Hybricon Corporation, 410 Great Road, Littleton, Massachusetts 01460, (617) 486-3174.

INMAC HAS COMPACT GENDER CHANGER

A self-contained, compact gender changer available from Inmac makes it easy and economical to recon-

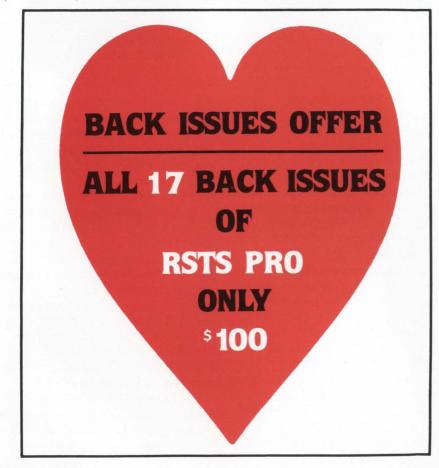


figure office terminals by joining cables whose genders conflict.

Inmac offers the gender changers as an inexpensive alternative to replacing or modifying EIA 232 cables. The unit — which will fit into a shirt pocket — consists of two EIA 232 25-pin connectors placed back to back. Both male-to-male and female-to-female connectors are available, with all 25 pins connected.

The gender changers are marketed by the company's Data Communications Division, and are guaranteed for one year. Data Communications also manufactures standard and Clear Signal_{TM} custom cables for a wide variety of micro and minicomputers and peripheral devices.

Inmac offers 24 hour shipment of phone and mail orders from regional sales and distribution centers throughout the United States. The company also maintains distribution facilities in Europe. For a free copy of Inmac's 100-page, four-color catalog listing over 2,000 mini and computer accessories, contact: Ron Becht, Product Marketing Manager, INMAC, Department 118, 2465 Augustine Drive, Santa Clara, CA 95051, 408-727-1970.

SOFTOOL RELEASES CHANGE CONTROL ENVIRONMENT

Softool Corporation is pleased to announce the release of its Change Control (CC) Environment.

CC automates the management of software changes and their documentation. It also controls who can make changes and to what components, thereby eliminating many unworkable manual procedures. CC can handle anything: source code, object code, test data, documents, . . . and it can deal with any language: FORTRAN, COBOL, PASCAL, ADA[®] JOVIAL, Assembly . . .

CC provides comprehensive features that include: automatic reconstruction of previous versions, problem tracking, difference reports, management reports, access control, archiving, compression, encryption, automatic recovery, and more. CC is an interactive tool with an easy-touse interface, on-line help, and interactive tutorials. CC is a stand-alone component of SOFTOOL, which is an integrated set of tools. CC offers complete support for source code management. CC is upgradable to the SOFTOOL Change and Configuration Control (CCC_{TM}) Environment which, in addition to change control, also offers full support for configuration control. A Programming Environment is also available.

A permanent license for CC is

\$12,000 with significant multiple copy discounts and a fifty percent credit toward any subsequent upgrade to CCC. For further details, contact Softool Corporation at 340 South Kellogg Avenue, Goleta, CA 93117. Telephone: (805) 964-0560.

SEMINARS, MEETINGS, SHOWS

PRECISION VISUALS SPONSORS FREE GRAPHICS SEMINARS

Precision Visuals, Inc., is currently sponsoring a free half-day graphics seminar in cities throughout the United States and Canada for end users, programmers, analysts, technical managers, and others who design and implement applications systems employing computer graphics.

Entitled "Understanding and Applying Today's Computer Graphics," the PVI seminar will be given in more than 40 cities this year, and is open to any interested graphics user by simply contacting PVI in advance to reserve a place.

According to Don Van Dyken,

PVI's director of marketing, the seminar will emphasize the potential for increased productivity offered by computer graphics applications in the fields of business, industry, government, education, and enginering. Participants will be given demonstrations of such diverse applications as business graphics, computer-aided design, contouring, demographics, 3-D surfaces, and electronics design.

In addition to surveying graphics applications and trends, the seminar will present options available in computer graphics software and hardware, and define the issues a potential buyer must consider when shopping for a graphics package. The advantages of a device-independent package will be explored as part of a discussion of efforts currently underway by standards organizations to establish a device-independent graphics standard.

Precision Visuals, based in Boulder, Colorado, is a leading supplier of graphics software tools offering device, machine, and even application independence. The company's DI-3000 package has been installed at more than 400 sites worldwide.

According to Van Dyken, seating is limited at the seminars, and those interested in attending should make reservations as soon as possible. For dates and locations of the seminars and registration information, contact Robbie Becker at Precision Visuals, 6260 Lookout Road, Boulder, CO 80301, (303) 530-9000.

PEOPLE, PLACES, THINGS

NCCS, EG&H INTEGRATE VSORT

North County Computer Services, Inc. of Escondido, California and Evans Griffiths & Hart, Inc. of Lexington, Massachusetts announced that they have agreed to integrate OEM versions of the software packages VSORT and VSELECT, developed by Evans Griffiths & Hart, into NCCS's new USER-11V database package for Digital Equipment Corporation's VAX computer. USER-11V is the VAX version of USER-11, a high-performance database management system for RSTS/E and CTS-500 operating environments, which is currently in use at over 200 installations. It is expected that the increased sort speed provided by VSORT and the record extraction and reformatting capabilities provided by VSELECT will significantly improve the performance and increase the capabilities of USER-11V's SORT/SELECT module.

The USER-11V database management system is organized in functional modules that are quickly combined and modified to implement many different types of applications. In addition to SORT/SELECT, there are database, report, data entry and modification, transfer, and program generation modules. USER-11V is fully dictionary-driven to simplify the maintenance of applications as they change over time.

VSORT, a very fast sort package for the VAX was released by Evans Griffiths & Hart in November 1982. It sorts fixed-length records from sequential or relative files from three to seven times faster than DEC's SORT-11, and also requires less disk space. VSELECT is a high speed package that extracts and optionally reformats records that meet usersupplied selection criteria. The package was released in March 1983. For further information contact: North County Computer Services, 2235 Meyer Avenue, Escondido, California 92025. Telephone: 619-745-6006, or Evans Griffiths and Hart, Inc., 55 Waltham Street, Lexington, Massachusetts 02173. Telephone: 617-861-0670.

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key Systems. 214-934-0031 Disposing RK05-K cartridges. Lawrence

University, Box 599, Appleton, WI 54912, (414) 735-6571.

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Horace D. Stephens

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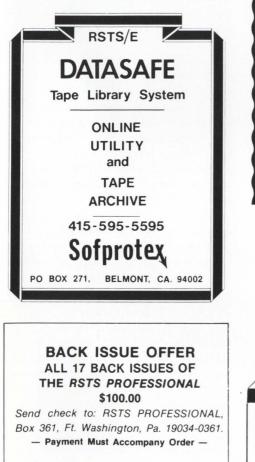
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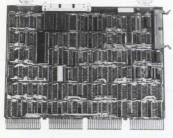


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... continued from page 81

SYSTEM INDUSTRIES AND MASSTOR JOIN FORCES

System Industries and MASSTOR Systems Corporation recently announced an Agreement providing for the joint marketing and development of products.

This Agreement enables System Industries to offer the DEC VAX mini and supermini-computer user storage capacities ranging from 160 megabytes through 440 giga bytes per subsystem. The products will allow large main frames such as CDC, IBM, HONEYWELL and UNIVAC real time access to very large on-line data bases in addition to the DEC VAX systems. MASSTOR's M860 Mass Storage System allows storage capacities ranging from 55 giga bytes to 440 giga bytes per system. The System Industries/MASSTOR local area networking architecture will provide interconnection of up to 64 computers over distances of one mile and longer distances with commercially available telecommunications services. Reliable multimegabit data communcations at speeds up to 50 million bits per second is provided.

MASSTOR designs, produces, markets and services very largecapacity (55 gigabytes or more online) data storage systems and high speed (approximately 50 megabits per second) general purpose local area and remotely-coupled data networks.

System Industries designs, produces, markets and services large capacity disk systems and high performance tape systems for DEC and Data General computers. Now in its 14th year, System Industries has over 30,000 data storage systems installed worldwide. The firm is publicly held, traded OTC (NASDAQ symbol — SYSM). For information, contact System Industries, Inc., 408/942-1212.

EMULEX ANNOUNCES AGREEMENT WITH SYSTIME COMPUTERS

Emulex Corporation recently an-

nounced that it has reached agreement with Systime Computers Limited (of Leeds, England) to supply selected Emulex products for use in Systime's activity as one of Britain's largest computer manufacturers.

"Systime has agreed to use Emulex's disk and disk/tape controllers and Emulex's communications controllers and multiplexers for its PDP-11 and VAX-11 systems," reports James F. Martin, Emulex's Vice President, International Sales. "The products involved are Emulex's SC21, SC71, SC750, and V-Master/780 disk controllers, TC11" tape controllers, and CS11 communications multiplexers."

"We're delighted that Systime has selected Emulex products after an extensive evaluation process," stated Mr. Martin.

The SC21 is a single-board, fullyembedded, software transparent disk controller that Systime will use in its PDP-11/44 applications. In VAX-11/730 applications, the SC21/V will be used, which consists of a hardware/software package developed and supported by Emulex. The SC71 is, where requested, a fully embedded, software transparent disk controller designed for use with the PDP-11/70.

The SC750 is a single-board, fullyembedded disk controller designed specifically for the VAX-11/750 CMI bus. It allows the user to mix up to four disk drives of varying capacities and media types, while maintaining full software transparency. The SC750 emulates the DEC RH750 Massbus Adapter with RM03, RM05, RM80, and RP06 drives.

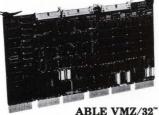
The V-Master/780 is a four-board, fully-embedded disk or disk/tape controller designed specifically for the VAX-11/780 SBI bus. While maintaining full software transparency, it allows the user to mix up to eight disk drives or four disk drives and four tape drives — all of varying capacities and media types.

Emulex, headquartered in Costa Mesa, California, designs, manufactures, and markets disk and tape controllers, communications controllers and multiplexers, and mass storage peripheral subsystems for mini- and micro- computers manufactured by Digital Equipment Corporation.

If you're in the market for communications modules, make the ABLE connection now. And join the thousands who already have.

We are known as the innovators. Most of our products are industry "firsts" which become popular quickly, then settle into a stage of steady long-term acceptance. These four DEC-compatible, communications devices fit the pattern perfectly. They are ABLE originals. They achieved instant success worldwide. They provide top performance. And they are very reliable. Read on to find the one for you.

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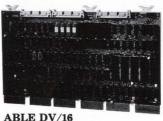
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needs for DMA communications requirements, serves UNIBUS systems equally well, and beats them all for MTBF, throughput and products soon to come.

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16-line DV11 replacement

of 4 or 8 lines with modem control and full system software compatibility. It takes less than half the space of a DV11 and uses word transfer instead of byte DMA to gain a 2 to 1 speed advantage or permit operation in half the bandwidth required for data transfers.

Q-BUS DMA.

The Q/DH is an asynchronous controller which makes DH-class performance possible on PDP-11/23 and LSI-11/23 Q-BUS systems. It connects the standard Q-BUS to as many as 16 async lines with DMA output capabil-



ABLE Q/DH^{**} 8 or 16-line DH/DM for Q-BUS

ities and allows optimum Q-BUS utilization. Features include software compatibility with RSTS/E and RSX operating systems, large input silo, modem control on all lines.

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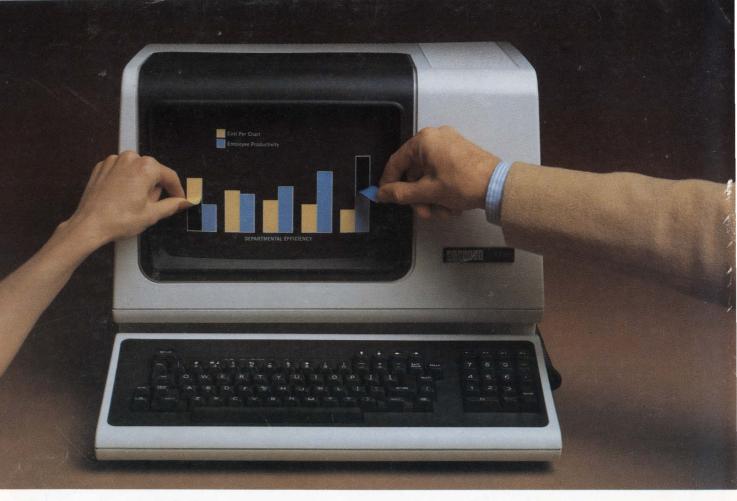
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